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1  ┌────────────────────────── MODULE TCS ───────────────────────────┐
  See DISC'2018: Multi-Shot Distributed Transaction Commit
5  EXTENDS Naturals, Integers, FiniteSets, Sequences, Functions, TLC,
6          FiniteSetsExt
7  └──────────────────────────────────────────────────────────────────┘

8  CONSTANTS
9      Key,          the set of keys, ranged over by  $k \in Key$ 
10     Tid,          the set of transaction identifiers, ranged over by  $t \in Tid$ 
11     RSet,         RSet[ $t$ ]: the read set of  $t \in Tid$ 
12     WSet,         WSet[ $t$ ]: the write set of  $t \in Tid$ 
13     CVer,         CVer[ $t$ ]: the commit version of  $t \in Tid$ 
14     Shard,        the set of shards, ranged over by  $s \in Shard$ 
15     Coord,        Coord[ $t$ ]: the coordinator of  $t \in Tid$ 
16     KeySharding  KeySharding[ $k$ ]: the shard that holds  $k \in Key$ 

18  NotTid  $\triangleq$  CHOOSE  $t : t \notin Tid$ 

20  Ver  $\triangleq$  0 .. Cardinality(Tid) with a distinguished minimum version 0
21  Slot  $\triangleq$  0 .. Cardinality(Tid) - 1

23  TKey( $t$ )  $\triangleq$  WSet[ $t$ ]  $\cup$  { $kv[1] : kv \in RSet[t]$ }
24  TSharding( $t$ )  $\triangleq$  {KeySharding[ $k$ ] :  $k \in TKey(t)$ }

26  ASSUME
27       $\wedge RSet \in [Tid \rightarrow SUBSET (Key \times Ver)]$ 
28       $\wedge \forall t \in Tid: RSet[t] \setminus * TODO: \text{one version per object}$ 
29       $\wedge WSet \in [Tid \rightarrow SUBSET Key]$ 
30       $\wedge \setminus * TODO: \text{"no blind update" assumption}$ 
31       $\wedge CVer \in [Tid \rightarrow Ver]$ 
32       $\wedge \setminus * TODO: \text{higher than any of the versions read}$ 
33       $\wedge Coord \in [Tid \rightarrow Shard]$ 
34       $\wedge KeySharding \in [Key \rightarrow Shard]$ 
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36  VARIABLES
37      next,        next[ $s$ ]  $\in Z$  points to the last filled slot
38      txn,         txn[ $s$ ][ $i$ ] is the transaction (identifier) to certify in the  $i$ -th slot
39      vote,        vote[ $s$ ][ $i$ ] is the vote for txn[ $s$ ][ $i$ ]
40      dec,         dec[ $s$ ][ $i$ ] is the decision for txn[ $s$ ][ $i$ ]
41      phase,       phase[ $s$ ][ $i$ ] is the phase for txn[ $s$ ][ $i$ ]
42      msg,         the set of messages in transit
43      submitted   the set of  $t \in Tid$  that have been submitted to TCS

45  sVars  $\triangleq$   $\langle next, txn, vote, dec, phase \rangle$ 
46  vars  $\triangleq$   $\langle next, txn, vote, dec, phase, msg, submitted \rangle$ 
47  └──────────────────────────────────────────────────────────────────┘
48  Message  $\triangleq$  [type : { "PREPARE" },  $t : Tid, s : Shard$ ]

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49    $\cup [type : \{\text{"PREPARE\_ACK"}\}, s : Shard, n : Int, t : Tid, v : \{\text{"COMMIT"}, \text{"ABORT"}\}]$ 
50    $\cup [type : \{\text{"DECISION"}\}, p : Int, d : \{\text{"COMMIT"}, \text{"ABORT"}\}, s : Shard]$ 

52    $Send(m) \triangleq msg' = msg \cup m$ 
53    $Delete(m) \triangleq msg' = msg \setminus m$ 
54    $SendAndDelete(sm, dm) \triangleq msg' = (msg \cup sm) \setminus dm$ 

55 |-----|
56    $TypeOK \triangleq$ 
57    $\wedge next \in [Shard \rightarrow Int]$ 
58    $\wedge txn \in [Shard \rightarrow [Slot \rightarrow Tid \cup \{NotTid\}]]$ 
59    $\wedge vote \in [Shard \rightarrow [Slot \rightarrow \{\text{"COMMIT"}, \text{"ABORT"}, \text{"NULL"}\}]]$ 
60    $\wedge dec \in [Shard \rightarrow [Slot \rightarrow \{\text{"COMMIT"}, \text{"ABORT"}, \text{"NULL"}\}]]$ 
61    $\wedge phase \in [Shard \rightarrow [Slot \rightarrow \{\text{"START"}, \text{"PREPARED"}, \text{"DECIDED"}\}]]$ 
62    $\wedge msg \subseteq Message$ 
63    $\wedge submitted \subseteq Tid$ 

64 |-----|
65    $Init \triangleq$ 
66    $\wedge next = [s \in Shard \mapsto -1]$ 
67    $\wedge txn = [s \in Shard \mapsto [i \in Slot \mapsto NotTid]]$ 
68    $\wedge vote = [s \in Shard \mapsto [i \in Slot \mapsto \text{"NULL"}]]$ 
69    $\wedge dec = [s \in Shard \mapsto [i \in Slot \mapsto \text{"NULL"}]]$ 
70    $\wedge phase = [s \in Shard \mapsto [i \in Slot \mapsto \text{"START"}]]$ 
71    $\wedge msg = \{\}$ 
72    $\wedge submitted = \{\}$ 

73 |-----|
74    $ComputeVote(t, s, n) \triangleq \text{"ABORT"} \quad \text{TODO}$ 
75    $ComputeDecision(vs) \triangleq$ 
76    $\text{IF } \exists v \in vs : v = \text{"ABORT"} \text{ THEN } \text{"ABORT"} \text{ ELSE } \text{"COMMIT"}$ 

77 |-----|
78    $Certify(t) \triangleq \text{Certify } t \in Tid$ 
79    $\wedge t \in Tid \setminus submitted$ 
80    $\wedge Send([type : \{\text{"PREPARE"}\}, t : \{t\}, s : TSharding(t))$ 
81    $\wedge submitted' = submitted \cup \{t\}$ 
82    $\wedge \text{UNCHANGED } sVars$ 

84    $Prepare(t, s) \triangleq \text{Prepare } t \in Tid \text{ on } s \in Shard \text{ when receive "PREPARE}(t)" \text{ message}$ 
85    $\wedge \exists m \in msg :$ 
86    $\wedge m = [type \mapsto \text{"PREPARE"}, t \mapsto t, s \mapsto s]$ 
87    $\wedge next' = [next \text{ EXCEPT } !s = @ + 1]$ 
88    $\wedge txn' = [txn \text{ EXCEPT } !s[txn'[s]] = t]$ 
89    $\wedge vote' = [vote \text{ EXCEPT } !s[next'[s]] = ComputeVote(t, s, next'[s])] \quad \text{TODO}$ 
90    $\wedge phase' = [phase \text{ EXCEPT } !s[next'[s]] = \text{"PREPARED"}]$ 
91    $\wedge SendAndDelete(\{[type \mapsto \text{"PREPARE\_ACK"},$ 
92    $s \mapsto s,$ 
93    $n \mapsto next'[s],$ 

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94          $t \mapsto t,$ 
95          $v \mapsto \text{vote}'[s][\text{next}'[s]]],$ 
96          $\{m\}$ 
97      $\wedge \text{UNCHANGED } \langle \text{dec}, \text{submitted} \rangle$ 

99  $\text{PrepareAck}(t, s) \triangleq$  PrepareAck for  $t \in \text{Tid}$  on shard  $s \in \text{Shard}$  when receive all "PREPARE_ACK" messages for  $t$ 
100      $\wedge s = \text{Coord}[t]$ 
101      $\wedge \text{LET } ms \triangleq \{m \in \text{msg} : m.\text{type} = \text{"PREPARE\_ACK"} \wedge m.t = t\}$ 
102          $vs \triangleq \{m.v : m \in ms\}$ 
103          $ss \triangleq \{m.s : m \in ms\}$ 
104     IN  $\wedge ss = \text{TSharding}(t)$ 
105          $\wedge \text{SendAndDelete}(\{[type \mapsto \text{"DECISION"},$ 
106              $p \mapsto \text{ChooseUnique}(ms, \text{LAMBDA } m : m.s = \text{shard}).n,$ 
107              $d \mapsto \text{ComputeDecision}(vs),$ 
108              $s \mapsto \text{shard}] : \text{shard} \in ss\},$ 
109              $ms)$ 
110      $\wedge \text{UNCHANGED } \langle s\text{Vars}, \text{submitted} \rangle$ 

112  $\text{Decision}(s) \triangleq$  Decide on shard  $s \in \text{Shard}$  when receive a "DECISION" message
113      $\wedge \exists m \in \text{msg} :$ 
114          $\wedge m.\text{type} = \text{"DECISION"}$ 
115          $\wedge m.s = s$ 
116          $\wedge \text{dec}' = [\text{dec} \text{ EXCEPT } ![s][m.p] = m.d]$ 
117          $\wedge \text{phase}' = [\text{phase} \text{ EXCEPT } ![s][m.p] = \text{"DECIDED"}]$ 
118          $\wedge \text{Delete}(\{m\})$ 
119      $\wedge \text{UNCHANGED } \langle \text{next}, \text{txn}, \text{vote}, \text{submitted} \rangle$ 
120 |-----|
121  $\text{Next} \triangleq$ 
122      $\vee \exists t \in \text{Tid} : \text{Certify}(t)$ 
123      $\vee \exists t \in \text{Tid}, s \in \text{Shard} :$ 
124          $\vee \text{Prepare}(t, s)$ 
125          $\vee \text{PrepareAck}(t, s)$ 
126      $\vee \exists s \in \text{Shard} :$ 
127          $\vee \text{Decision}(s)$ 

129  $\text{Spec} \triangleq \text{Init} \wedge \square[\text{Next}]_{\text{vars}}$ 
130 |-----|

\ * Modification History
\ * Last modified Sun Jun 13 17:13:35 CST 2021 by hengxin
\ * Created Sat Jun 12 21:01:57 CST 2021 by hengxin

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