NASM: data and bss Beyond the podcast

ICS312 Machine-Level and Systems Programming

Henri Casanova (henric@hawaii.edu)



NASM Program Structure

data segment

bss segment

text segment

declaration of initialized data declaration of uninitialized data

statically allocated data that is allocated for the duration of program execution

code



```
      var1
      dd
      -9

      var2
      db
      "dcba"

      a1
      times 3
      dw
      011b

      wt
      db
      0A3h, 0
```

- What is the layout and the content of the data memory segment on a LITTLE ENDIAN machine?
 - Byte per byte, in hex

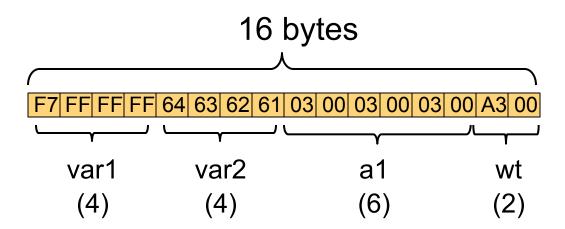
.

```
      var1
      dd
      -9

      var2
      db
      "dcba"

      a1
      times 3
      dw
      011b

      wt
      db
      0A3h.
      0
```





```
x1 dw -22
msg db "ba", 0
array times 2 dd 0230
wt dw 0.011b.020o
```

- What is the layout and the content of the data memory segment on a BIG ENDIAN machine?
 - Byte per byte, in hex

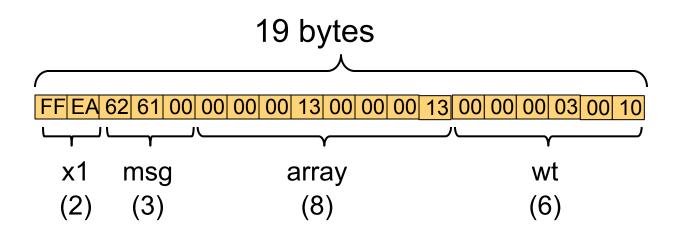
.

```
      x1
      dw
      -22

      msg
      db
      "ba", 0

      array
      times 2
      dd
      0230

      wt
      dw
      0.011b.020o
```





More practice?

Of course we can easily come up with tones of practice examples...

Should we do one more right now?

Homework Assignment #2

- Homework Assignment #2 is posted
- With what we've seen you can get started on Exercise #1
- For Exercise #2 we need to go further a bit...
 - We might do it all today



Our first instructions

- At this point we need to introduce a few assembly instructions
 - adding integers
 - subtracting integers
 - moving data between registers / memory locations / constants

Different kinds of operands

- Assembly instructions can have operands, and it's important to know what kind of operands are possible
- Register: specifies one of the registers
 - □ add eax, ebx
 - □ means eax = eax + ebx
- Memory: specifies an address in memory.
 - □ add eax, [ebx]
 - □ means eax = eax + content of memory at address ebx
- Immediate: specifies a fixed value (i.e., a number)
 - □ add eax, 2
 - □ means eax = eax + 2
- Implied: not actually encoded in the instruction
 - □ inc eax
 - □ means eax = eax + 1

Additions, subtractions

Additions

```
\square add eax, 4 ; eax = eax + 4
```

- □ add al, ah ; al = al + ah
- Subtractions

```
□ sub bx, 10 ; bx = bx - 10
```

- □ sub ebx, edi ; ebx = ebx edi
- Increment, Decrement

```
□ inc ecx ; ecx++ (a 4-byte operation)
```

□ dec dl ; dl-- (a 1-byte operation)

The move instruction

- This instruction moves data from one location to another mov dest, src
- Destination goes first, and the source goes second
- At most one of the operands can be a memory operand
 - □ mov eax, [ebx] ; OK
 - mov [eax], ebx ; OK
 - □ mov [eax], [ebx] ; NOT OK
- Both operands must be exactly the same size
 - □ For instance, AX cannot be stored into BL
- Examples:
 - □ mov ax, ebx ; **NOT OK**
 - □ mov bx, ax ; OK
- This type of "exceptions to the common case" make programming languages difficult to learn and assembly may be the worst offender
 - By contrast, Lisp is known for being very consistent (ICS313)

Use of Labels

- It is important to constantly be aware that when using a label in a program, the label is a pointer, not a value
- Therefore, a common use of the label in the code is as a memory operand, in between square brackets '[' ']'
- mov AL, [L1]
 - Copy the value at address L1 into register AL
- Question: how does the assembler know how many bits to move?
- Answer: it's up to the programmer to do the right thing, that is load into appropriately sized registers
 - In the above example, since AL is a 1-byte register, then 1 byte is moved
- LABELS HAVE NO TYPE!
 - So although it's tempting to think of them as variables, they are much more limited: just pointers to a byte somewhere

Moving to/from a register

- Say we have the following data segment
 L db 0F0h, 0F1h, 0F2h, 0F3h
- Example: mov AL, [L]
 - Will copy 1 byte from memory into AL
- Example: mov [L], AX
 - □ Will copy the 2 bytes in AX to memory (overwrite values)
- Example: mov [L], EAX
 - Will copy the 4 bytes in AX to memory (overwrite values)
- What bytes are written where depends on Little vs. Big Endian...

Mov and Little Endian

Consider the following data segment

L1 db 0AAh, 0BBh, 0CCh, 0DDh

L2 dd 0AABBCCDDh

- The instruction: mov eax, [L1] puts DDCCBBAA into eax
 - □ Note that we're loading 4x1 bytes as a 4-byte quantity
- The instruction: mov eax, [L2] puts AABBCCDD into eax!!!
 - Meaning that the memory content was DDCCBBAA
- When declaring a value in the data segment, that value is declared as it would be appearing in registers when loaded "whole"
 - □ It would be confusing to write numbers in little endian in the program
 - So all numerical values you write are in register-order not memory-order

Example

Data segment (little endian):

```
L1 db 0AAh, 0BBh
L2 dw 0CCDDh
L3 db 0EEh, 0FFh
```

Program:

```
mov eax, [L2]
mov ax, [L3]
mov [L1], eax
```

What's the final memory content?



Solution

Data segment:

L1 db 0AAh, 0BBh

L2 dw 0CCDDh

L3 db 0EEh, 0FFh

L1 L2 L3

AA BB DD CC EE FF

Solution

```
L1 L2 L3

AA BB DD CC EE FF
```

```
mov eax, [L2] ; eax = FF EE CC DD
```

mov [L1], eax ; content at L1 is EE FF EE FF

L1 L2 L3

EE FF EE FF EE FF

Final memory content

Moving immediate values

- Consider the instruction: mov [L], 1
- The assembler will give us an error: "operation size not specified"!
- This is because the assembler has no idea whether we mean for "1" to be 01h, 0001h, 00000001h, etc.
 - Labels have no type (they're NOT variables)
- Therefore the assembler must provide us with a way to specify the size of immediate operands
- mov dword [L], 1
 - 4-byte double-word
- Size specifiers: byte, word, dword, qword (and tword)

Size Specifier Examples

■ mov [L1], 1 ; Error

mov byte [L1], 1 ; 1 byte

mov word [L1], 1 ; 2 bytes

mov dword [L1], 1; 4 bytes

mov [L1], eax ; 4 bytes

mov [L1], ax ; 2 bytes

mov [L1], al ; 1 byte

mov eax, [L1] ; 4 bytes

■ mov ax, [L1] ; 2 bytes

mov ax, 12 ; 2 bytes

Brackets or no Brackets

- mov eax, [L]
 - Puts the content at address L into eax
 - □ Puts 32 bits of content, because eax is a 32-bit register
- mov eax, L
 - Puts the address L into eax
 - □ Puts the 32-bit address L into eax
- mov ebx, [eax]
 - □ Puts the content at address eax (= L) into ebx
- inc eax
 - Increase eax by one
- mov ebx, [eax]
 - □ Puts the content at address eax (= L + 1) into ebx

Example

```
first db 00h, 04Fh, 012h, 0A4h second dw 165 third db "adf"
```

```
mov eax, first
inc eax
mov ebx, [eax]
mov [second], ebx
mov byte [third], 11o
```

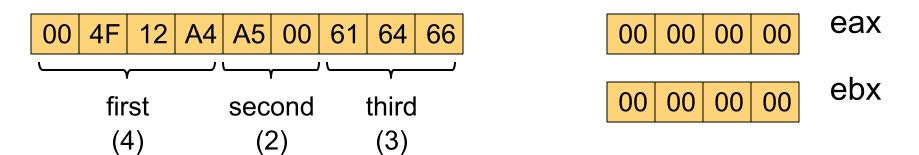
What is the content of the data segment after the code executes on a **Little Endian** Machine?

.

Example

```
first db 00h, 04Fh, 012h, 0A4h second dw 165 third db "adf"
```

```
mov eax, first
inc eax
mov ebx, [eax]
mov [second], ebx
mov byte [third], 11o
```

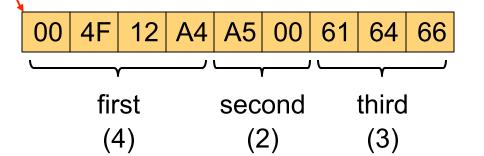


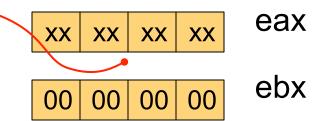


```
first db 00h, 04Fh, 012h, 0A4h second dw 165 third db "adf"
```

```
mov eax, first
inc eax
mov ebx, [eax]
mov [second], ebx
mov byte [third], 11o
```

Put an **address** into eax (this works because our addresses are 32-bit and thus fit into 4-byte registers, just like any other 4-byte values!)



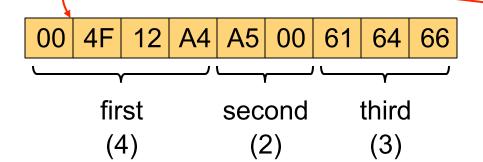


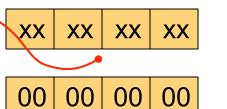


```
first db 00h, 04Fh, 012h, 0A4h second dw 165 third db "adf"
```

```
mov eax, first
inc eax
mov ebx, [eax]
mov [second], ebx
mov byte [third], 11o
```

Increment that address by 1, thus now pointing to the next byte





ebx

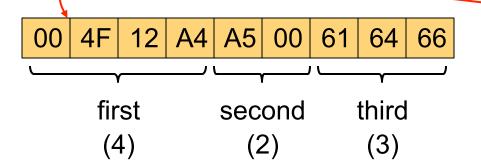
eax

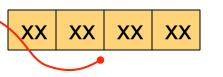


```
first db 00h, 04Fh, 012h, 0A4h second dw 165 third db "adf"
```

```
mov eax, first
inc eax
mov ebx, [eax]
mov [second], ebx
mov byte [third], 11o
```

Put the 4 bytes at that address into ebx (note the Little Endian)





ebx

eax

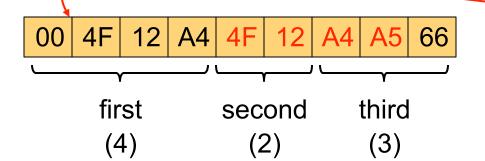
A5 A4 12 4F



```
first db 00h, 04Fh, 012h, 0A4h second dw 165 third db "adf"
```

```
mov eax, first
inc eax
mov ebx, [eax]
mov [second], ebx
mov byte [third], 11o
```

Copy 4 bytes to memory at address **second**



eax

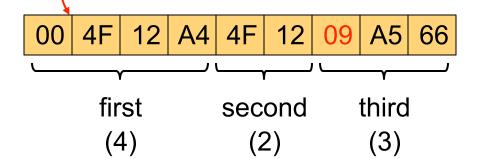
A5 A4 12 4F

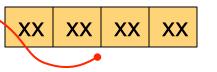
ebx

```
first db 00h, 04Fh, 012h, 0A4h second dw 165 third db "adf"
```

```
mov eax, first
inc eax
mov ebx, [eax]
mov [second], ebx
mov byte [third], 11o
```

Write 1 byte at address third





A5 A4 12 4F

eax

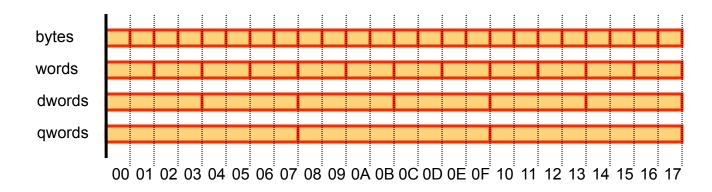
ebx

Assembly is Dangerous

- The previous example is really a terrible program
- But it's a good demonstration of why the assembly programmer must be really careful
- For instance, we were able to store 4 bytes into a 2byte label, thus overwriting the first 2 characters of a string that merely happened to be stored in memory next to that 2-byte label
 - again: LABELS ARE NOT VARIABLES AT ALL
- Playing such tricks can lead to very clever programs that do things that would be impossible (or very cumbersome) to do with many high-level programming language (e.g., in Java)
- But you really must know what you're doing

x86 Assembly is Dangerous

- Another dangerous thing we did in our assembly program was the use of unaligned memory accesses
 - We stored a 4-byte quantity at some address
 - We incremented the address by 1
 - We read a 4-byte quantity from the incremented address!
 - This really removes all notion of a structured memory (it's only bytes)
- Some architectures only allow aligned accesses
 - Accessing an X-byte quantity can only be done for an address that's a multiple of X!



Practice #3

Consider the following program

```
      var1
      dd
      179

      var2
      db
      0A3h, 017h, 012h

      var3
      db
      "bca"
```

```
mov eax, var1
add eax, 3
mov ebx, [eax]
add ebx, 5
mov [var1], ebx
```

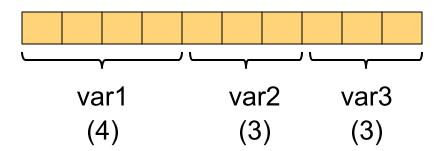
What is the final layout of the data segment starting at address var1 on a Little Endian Machine?

```
      var1
      dd
      179

      var2
      db
      0A3h, 017h, 012h

      var3
      db
      "bca"
```

```
mov eax, var1
add eax, 3
mov ebx, [eax]
add ebx, 5
mov [var1], ebx
```



```
      var1
      dd
      179

      var2
      db
      0A3h, 017h, 012h

      var3
      db
      "bca"
```

```
mov eax, var1
add eax, 3
mov ebx, [eax]
add ebx, 5
mov [var1], ebx
```

.

```
      var1
      dd
      179

      var2
      db
      0A3h, 017h, 012h

      var3
      db
      "bca"
```

```
mov eax, var1
add eax, 3
mov ebx, [eax]
add ebx, 5
mov [var1], ebx
```

```
B3 00 00 00 A3 17 12 62 63 61

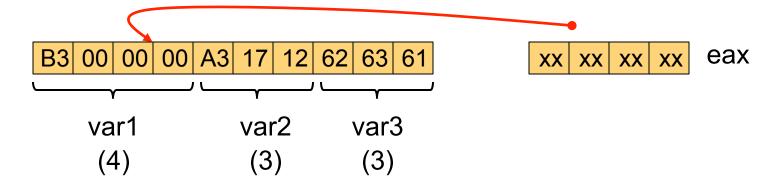
var1 var2 var3
(4) (3) (3)
```

```
      var1
      dd
      179

      var2
      db
      0A3h, 017h, 012h

      var3
      db
      "bca"
```

```
mov eax, var1
add eax, 3
mov ebx, [eax]
add ebx, 5
mov [var1], ebx
```



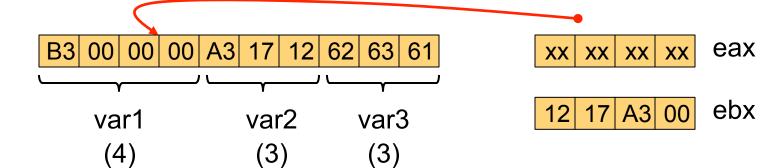
M.

```
      var1
      dd
      179

      var2
      db
      0A3h, 017h, 012h

      var3
      db
      "bca"
```

```
mov eax, var1
add eax, 3
mov ebx, [eax]
add ebx, 5
mov [var1], ebx
```



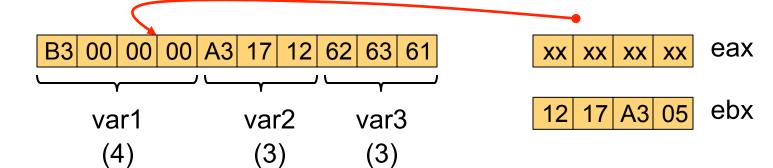
ж.

```
        var1
        dd
        179

        var2
        db
        0A3h, 017h, 012h

        var3
        db
        "bca"
```

```
mov eax, var1
add eax, 3
mov ebx, [eax]
add ebx, 5
mov [var1], ebx
```



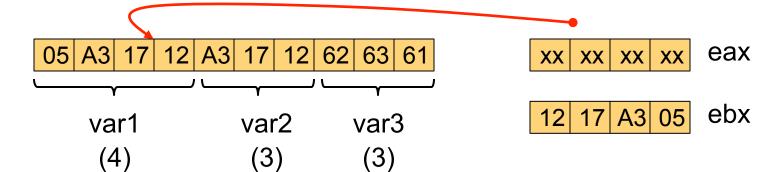
.

```
      var1
      dd
      179

      var2
      db
      0A3h, 017h, 012h

      var3
      db
      "bca"
```

```
mov eax, var1
add eax, 3
mov ebx, [eax]
add ebx, 5
mov [var1], ebx
```



Practice #4

Consider the following program

```
      var1
      db
      "b","ca",0

      var2
      db
      3,0,0,0

      var3
      times 2 dw
      012h
```

```
mov eax, var3
mov ebx, var1
sub eax, 4
add ebx, [eax]
mov dword [ebx], 42
```

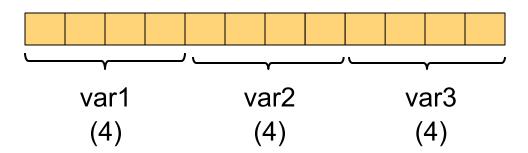
What is the final layout of memory starting at address var1 on a Little Endian Machine?

```
      var1
      db
      "b","ca",0

      var2
      times
      db
      3,0,0,0

      var3
      times
      2 dw
      012h
```

```
mov eax, var3
mov ebx, var1
sub eax, 4
add ebx, [eax]
mov dword [ebx], 42
```



```
      var1
      db
      "b","ca",0

      var2
      times
      db
      3,0,0,0

      var3
      times
      2 dw
      012h
```

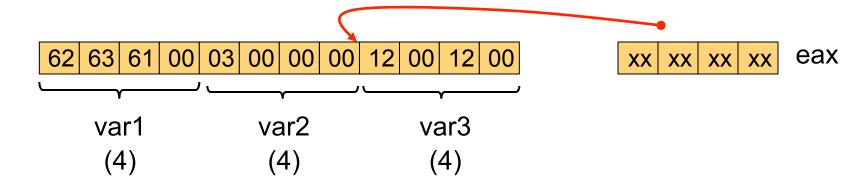
```
mov eax, var3
mov ebx, var1
sub eax, 4
add ebx, [eax]
mov dword [ebx], 42
```

```
      var1
      db
      "b","ca",0

      var2
      times
      db
      3,0,0,0

      var3
      times
      2 dw
      012h
```

```
mov eax, var3
mov ebx, var1
sub eax, 4
add ebx, [eax]
mov dword [ebx], 42
```

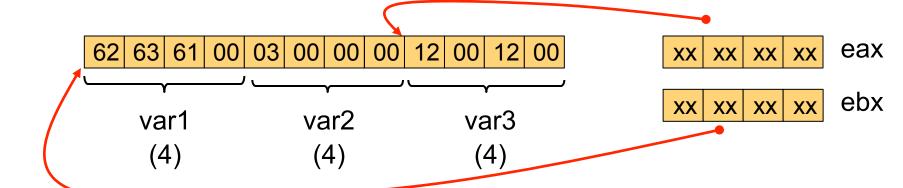


```
      var1
      db
      "b","ca",0

      var2
      times
      db
      3,0,0,0

      var3
      times
      2 dw
      012h
```

```
mov eax, var3
mov ebx, var1
sub eax, 4
add ebx, [eax]
mov dword [ebx], 42
```

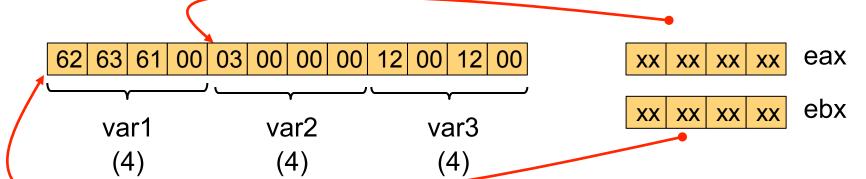


```
      var1
      db
      "b","ca",0

      var2
      times
      db
      3,0,0,0

      var3
      times
      2 dw
      012h
```

```
mov eax, var3
mov ebx, var1
sub eax, 4
add ebx, [eax]
mov dword [ebx], 42
```

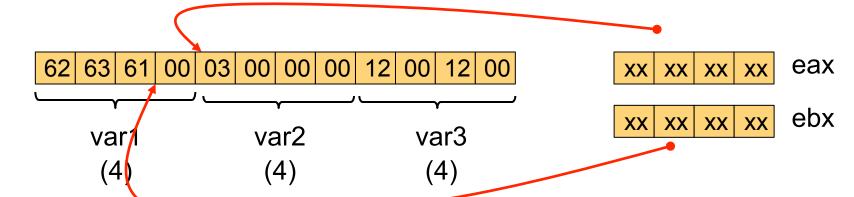


```
      var1
      db
      "b","ca",0

      var2
      times
      db
      3,0,0,0

      var3
      times
      2 dw
      012h
```

```
mov eax, var3
mov ebx, var1
sub eax, 4
add ebx, [eax]
mov dword [ebx], 42
```

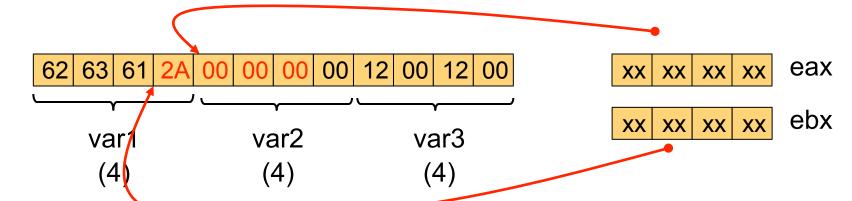


```
      var1
      db
      "b","ca",0

      var2
      times
      db
      3,0,0,0

      var3
      times
      2 dw
      012h
```

```
mov eax, var3
mov ebx, var1
sub eax, 4
add ebx, [eax]
mov dword [ebx], 42
```





Conclusion

- Should we make up other practices now or are we good?
 - Somebody wants to propose one?
- We can now do Homework #2
- Next lecture we'll go through a practice quiz (Zoom poll) on this module...