

# **Final Review**

## **ICS312 Machine-Level and Systems Programming**

Henri Casanova ([henric@hawaii.edu](mailto:henric@hawaii.edu))

# What to Study for the Final?

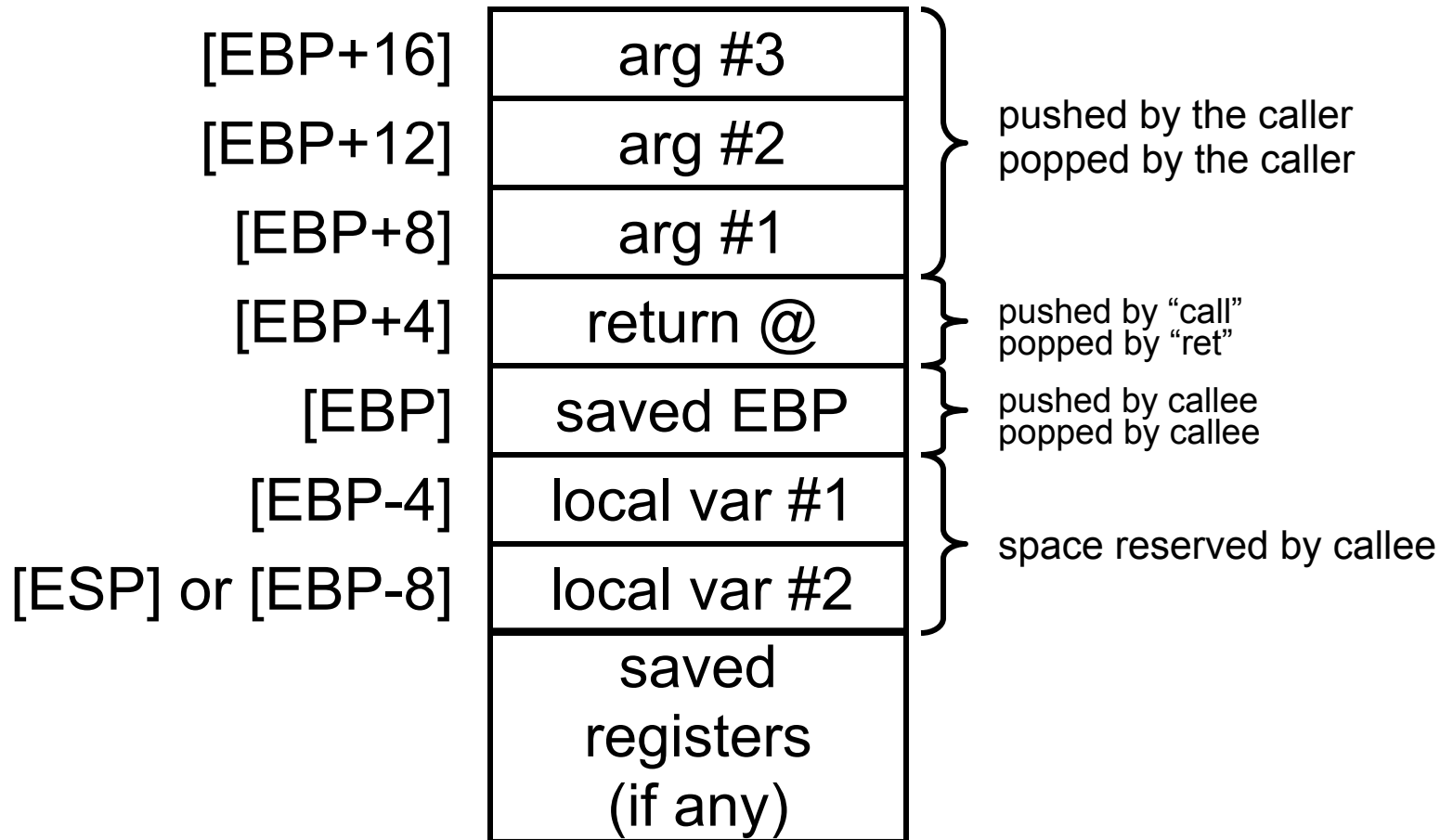
- Open-notes/computer, Laulima (just like the midterm)
- Non-cumulative: only post-midterm content
  - But of course you must remember pre-midterm assembly programming (e.g., when you implement a function)
- Types of questions
  - Mostly Problem-like questions
    - Write a piece of code to do something
    - What does this piece of code do?
    - What does the stack look like?
- What to study
  - All lecture notes
  - Homework assignment answers, past quizzes



# **“Obvious” Possible Questions**

- Translate some C code into assembly
- Translate from assembly code into C
- Subprograms
  - Write a function with arguments and parameters
  - Draw a picture of the stack at some point in a program
  - Reason about how many things are on the stack
- Compiler
  - General knowledge
  - Write simple regular expressions / FSAs / CFGs

# The Main Stack Structure





# The Stack

- Should we go back and look at the subprogram pencil-and-paper homework assignment?
- Or make up some example on the fly?



# Write a Function that ...

- Write a function that takes 2 addresses as arguments and swaps their content, without destroying any register
  - watch out for indirections...
  - `void f(int *x, int *y);`

# Write a Function that ...

- Write a function that takes 2 addresses as arguments and swaps their content, without destroying any register

```
f:    push    ebp
      mov     ebp, esp
      pusha
      mov     eax, [ebp+8]      ; eax = 1st address
      mov     ebx, [eax]       ; ebx = content at 1st address
      mov     ecx, [ebp+12]    ; ecx = 2nd address
      mov     edx, [ecx]       ; edx = content at 2nd address
      mov     [eax], edx       ; store edx to 1st address
      mov     [ecx], ebx       ; store ebx to 2nd address
      popa
      pop     ebp
      ret
```



# Translate this Function...

```
int f(int x, int y) {  
    int z;  
    z = x + y;  
    return z - 2;  
}
```



# Translate this Function...

```
int f(int x, int y) {  
    int z;  
    z = x + y;  
    return z - 2;  
}
```

```
f:  
push    ebp  
mov     ebp, esp  
sub     esp, 4  
  
mov     eax, [ebp+8]  
add     eax, [ebp+12]  
mov     [ebp-4], eax  
mov     eax, [ebp-4]  
sub     eax, 2  
  
add     esp, 4  
pop     ebp  
ret
```

# Same with popa

```
int f(int x, int y) {  
    int z;  
    z = x + y;  
    return z - 2;  
}
```

```
f:  
push    ebp  
mov     ebp, esp  
sub     esp, 4  
pusha    ; after local variables  
  
mov     eax, [ebp+8]  
add     eax, [ebp+12]  
mov     [ebp-4], eax  
mov     eax, [ebp-4]  
sub     eax, 2  
mov     [somewhere], eax  
  
popa  
mov     eax, [somewhere]  
add     esp, 4  
pop     ebp  
ret
```

# Stack picture

f(4,3);

....

```
void f(int a, int b) {
```

```
    g(a+1, b-1);
```

```
}
```

```
void g(int x, int y) {
```

```
    int z = x+1;
```

```
    z = z * y;
```

```
    // SHOW STACK
```

```
}
```

# Stack picture

f(4,3);

....

```
void f(int a, int b) {
```

```
    g(a+1, b-1);
```

```
}
```

```
void g(int x, int y) {
```

```
    int z = x+1;
```

```
    z = z * y;
```

```
    // SHOW STACK
```

```
}
```

3
4
ret@ to main
saved EBP
2
5
ret@ to f
saved EBP
z = 6 12

# The Big Picture

## High-level code

```
char *tmpfilename;  
int num_schedulers=0;  
int num_request_submitters=0;  
int i,j;  
  
if (!(f = fopen(filename,"r"))) {  
    xbt_assert(0,"Cannot open file %s",filename);  
}  
while(fgets(buffer,256,f) {  
    if (strcmp(buffer,"SCHEDULER",9))  
        num_schedulers++;  
    if (strcmp(buffer,"REQUESTSUBMITTER",16))  
        num_request_submitters++;  
}  
fclose(f);  
tmpfilename = strdup("/tmp/jobsimulator_
```

**ASSEMBLER**

## Machine Code (object files)

```
010000101010110110  
10  
10  
10  
101  
101  
00  
01  
111  
000  
010  
000  
0100001010101001  
000101010101000101  
111100001010101001  
00010101011101011  
010000000010000100  
000010001000100011
```

## RUNNING PROGRAM

**LOADER**

## Machine Code (executable)

**COMPILER**

## Assembly code

```
sll $t3, $t1, 2  
add $t3, $s0, $t3  
sll $t4, $t0, 2  
add $t4, $s0, $t4  
lw $t5, 0($t3)  
lw $t6, 0($t4)  
slt $t2, $t5, $t6  
beq $t2, $zero, endif  
add $t0, $t1, $zero  
sll $t4, $t0, 2  
add $t4, $s0, $t4  
lw $t5, 0($t3)  
lw $t6, 0($t4)  
slt $t2, $t5, $t6  
beq $t2, $zero, endif
```

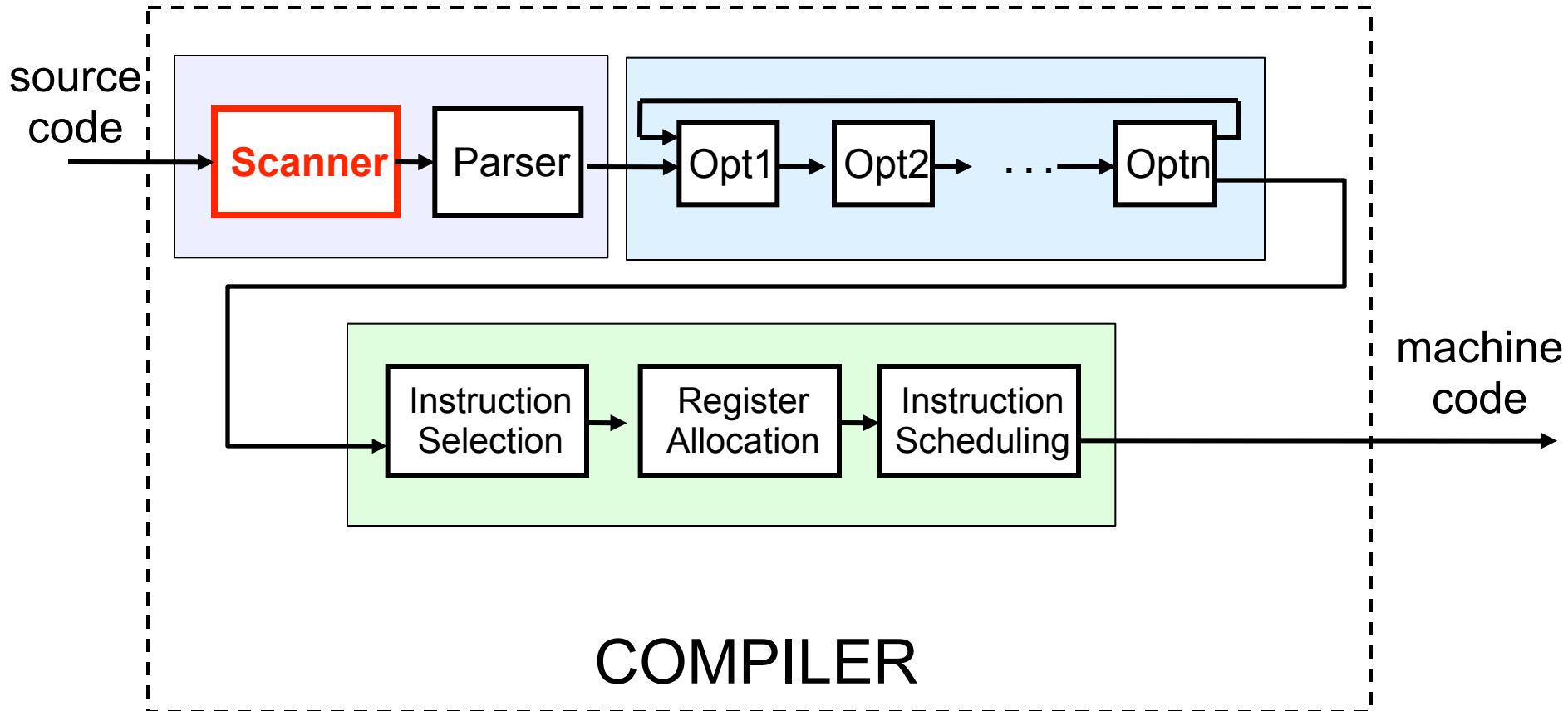
**Hand-written  
Assembly code**

```
sll $t3, $t1, 2  
add $t3, $s0, $t3  
sll $t4, $t0, 2  
add $t4, $s0, $t4  
lw $t5, 0($t3)  
lw $t6, 0($t4)  
slt $t2, $t5, $t6  
beq $t2, $zero, endif
```

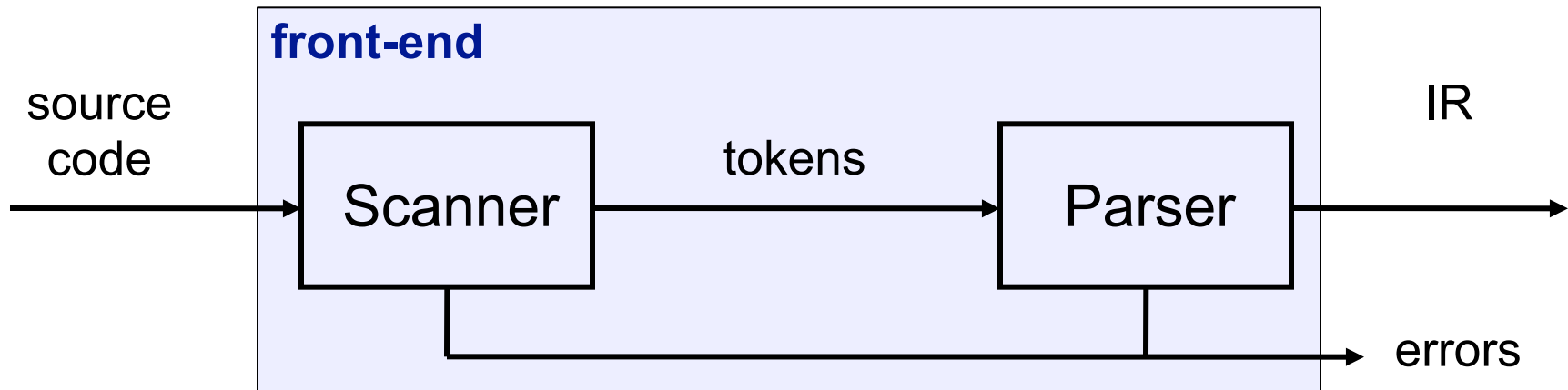
**LINKER**

```
010000101010110110  
101010101111010101  
101001010101010001  
101010101010100101  
111100001010101001  
000101010111101011  
010000000010000100  
000010001000100011  
101010101011101110  
101010101010010000  
000010101110101111  
001010101011111111  
11111111111101010  
010101111110110101  
110101010101010101  
111110101010101010
```

# The Compiler's Big Picture



# The front-end's big picture



- The scanner:

- uses regular expressions (implemented as finite automata) to identify correct “words”

- The parser:

- uses a context-free grammar to identify correct “sentences”

# A few regular expressions

- What is this pattern (in english)?:  $a^*b^+a^*$
- What is the pattern for all sentences over the  $\{0,1\}$  alphabet that start with a 0, alternate with a 1 and a 0 an arbitrary number of times, and end with any number of 0?



# A few regular expressions

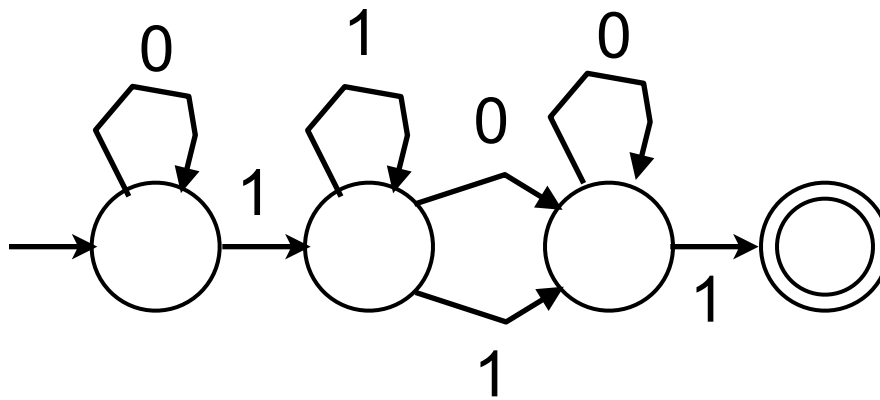
- What is this pattern (in english)?:  $a^*b^+a^*$ 
  - The set of strings over alphabet  $\{a,b\}$  that begin with 0 or more a's, then have 1 or more b's, and then end with 0 or more a's.
  - Examples: aaaaabbbbba, ba, b, aba, abbaa
- What is the pattern for all sentences over the  $\{0,1\}$  alphabet that start with a 0, alternate with a 1 and a 0 an arbitrary number of times, and end with any number of 0's?
  - $0(10)^*0^*$

# RE to NFA

- Draw the NFA for:  $0^*1^+(0|1)0^*1$

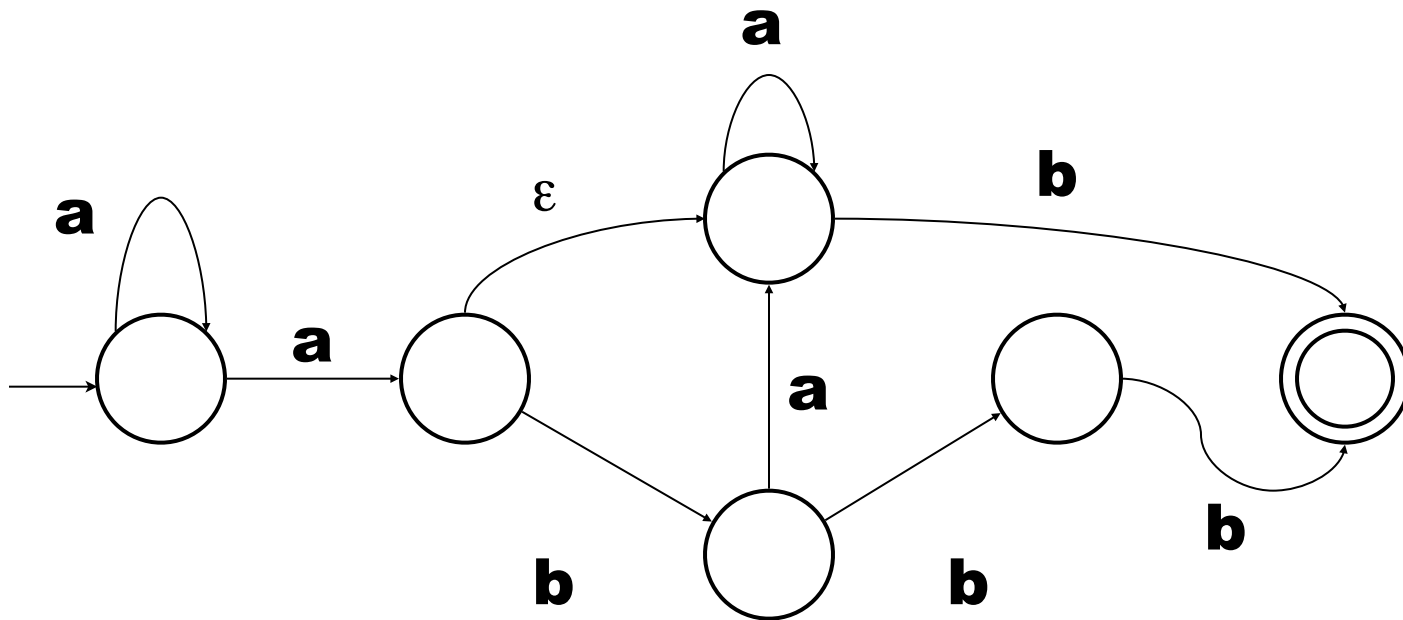
# RE to NFA

- Draw the NFA for:  $0^*1^+(0|1)0^*1$

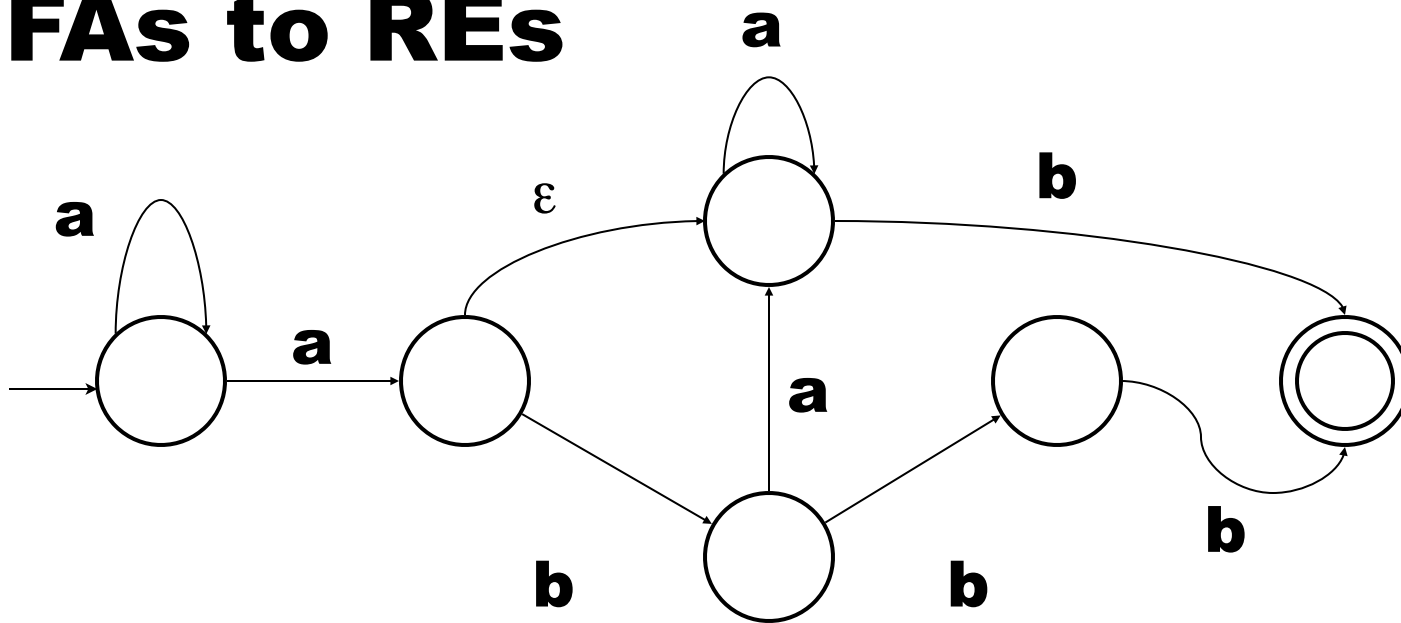


# NFAs to REs

- Write the RE for this NFA



# NFAs to REs



$$\begin{aligned}\text{RE} &= a^+a^*b \mid a^+ba^+b \mid a^+bbb \\ &= a^+b \mid a^+ba^+b \mid a^+bbb \\ &= a^+b(\epsilon \mid a^+b \mid bb)\end{aligned}$$



# Write a CFG

- Write the CFG over the alphabet  $\{x,y\}$  that describes all possible strings

# Write a CFG

- Write the CFG over the alphabet  $\{x,y\}$  that describes all possible strings
  - $S \rightarrow \varepsilon$
  - $S \rightarrow x S$
  - $S \rightarrow y S$

# Write a CFG

- Write the CFG over the alphabet  $\{a,b\}$  that describes all strings that are palindromes



# Write a CFG

- Write the CFG over the alphabet  $\{a,b\}$  that describes all strings that are palindromes

- $S \rightarrow \varepsilon \mid a \mid b$

- $S \rightarrow a S a$

- $S \rightarrow b S b$

# Write a CFG

- Write the CFG over the alphabet  $\{0,1\}$  that describes all strings that end with a sequence of  $n$  0's and  $n$  1's (for any  $n > 0$ )
  - example: 010110100010101**00001111**
  - example: **00000001111111**

# Write a CFG

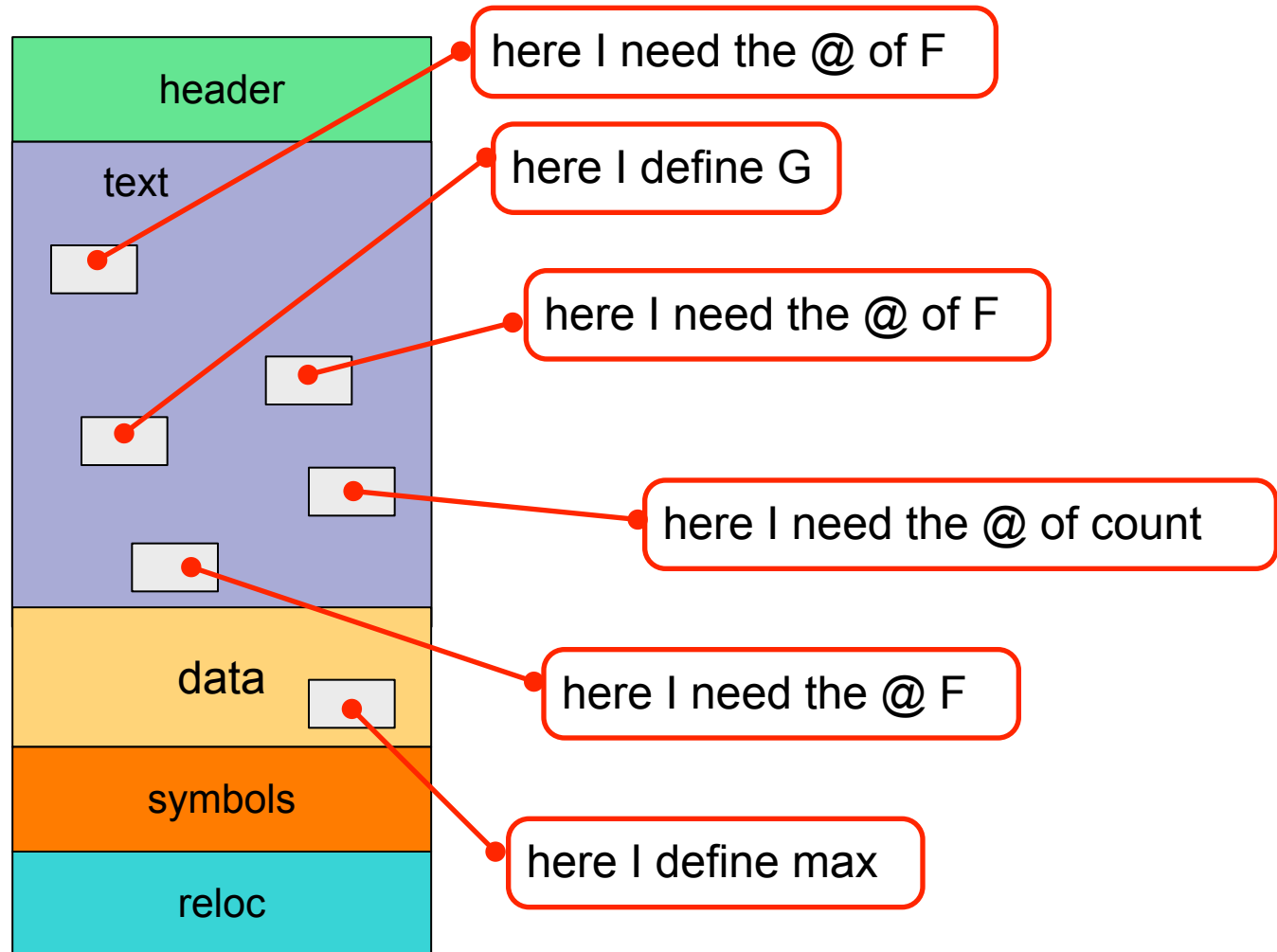
- Write the CFG over the alphabet  $\{0,1\}$  that describes all strings that end with a sequence of  $n$  0's and  $n$  1's (for any  $n > 0$ )

- $S \rightarrow A B$

- $A \rightarrow 0 A \mid 1 A \mid \varepsilon$

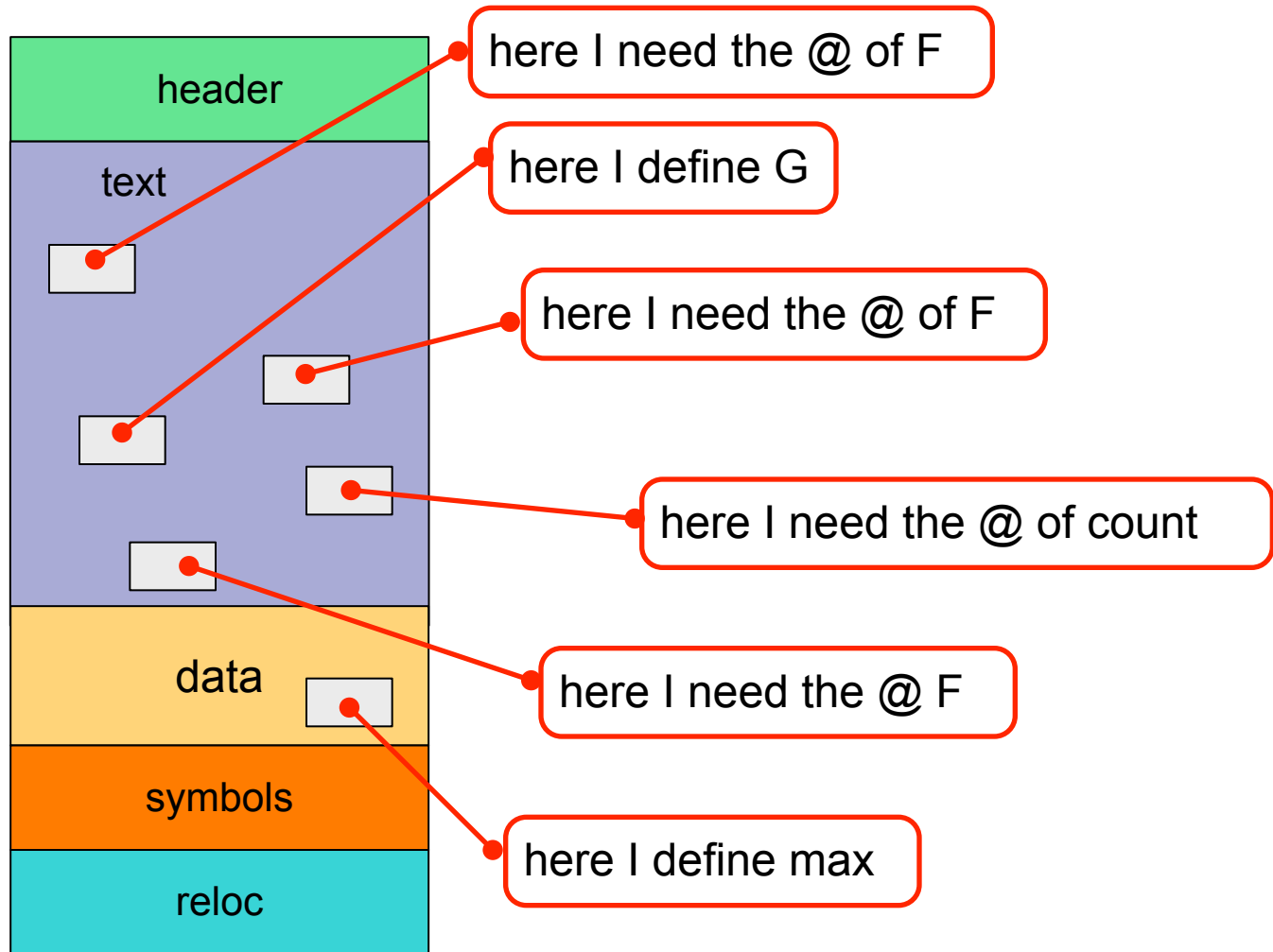
- $B \rightarrow 0 B 1 \mid 01$

# Object files from the assembler



# Object files from the assembler

how many entries in the symbol table?  
how many entries in the relocation table?





# The end

- Our final is Thursday December 17th at 2:15PM
  - Scheduled for 2 hours
  - Should be doable in much less