

The  $\varepsilon$ -Net Theorem

Anne Driemel

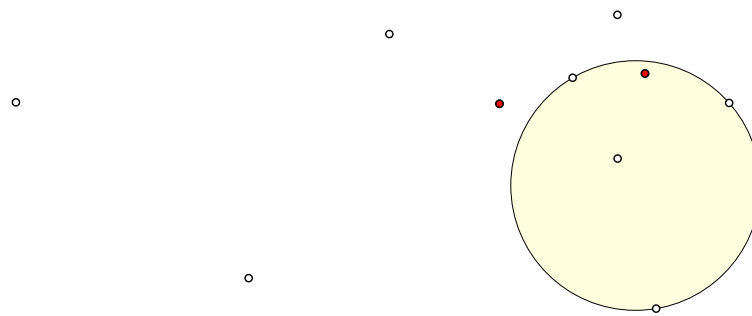
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In the previous lecture we introduced set systems and we showed that the VC-dimension characterizes the growth of subsystems, which we formalized with the growth function  $\Pi_{\mathcal{R}}(m)$ . In this lecture we will introduce the concept of  $\varepsilon$ -nets and show that a small VC-dimension implies the existence of  $\varepsilon$ -nets of size independent of the size of the ground set.

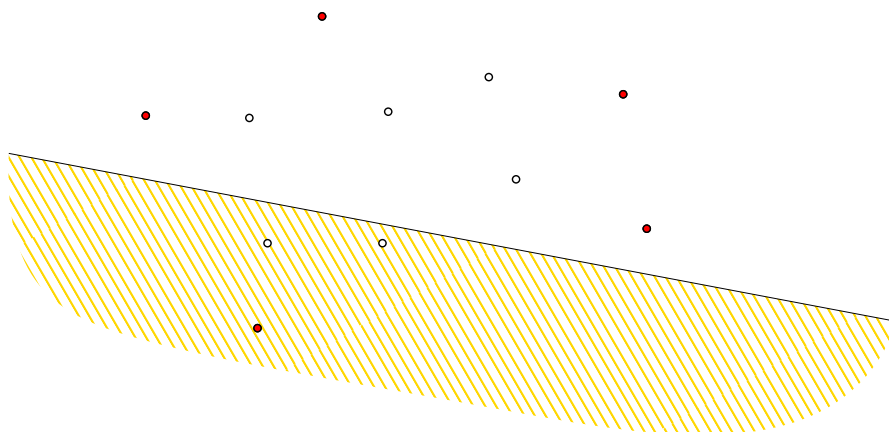
## 1 Definitions

**Definition 21.1** ( $\varepsilon$ -net). Let  $\mathcal{R}$  be a set system with finite ground set  $X$  and let  $\varepsilon \in [0, 1]$  be a real value. A subset  $A \subseteq X$  is an  $\varepsilon$ -net for this set system, if  $A \cap r \neq \emptyset$  for all  $r \in \mathcal{R}$  with  $|r| \geq \varepsilon|X|$ .

**Example 21.2.** The figure below shows a set of 10 points in the plane. Let  $X$  denote this set and consider the set system  $\mathcal{R}$  where every set is defined by a disk  $D$  and contains the subset of  $X$  inside  $D$ . Let  $\varepsilon = \frac{1}{2}$ , an  $\varepsilon$ -net of  $\mathcal{R}$  is shown with red (filled) dots. It can be verified that any disk containing at least 5 points also contains a point of the  $\varepsilon$ -net.



**Example 21.3.** Let  $X$  be a set of  $n$  points in  $\mathbb{R}^d$ . Let  $\mathcal{R}$  be a set system, where every set is defined by a vector  $a \in \mathbb{R}^d$  and a real value  $b \in \mathbb{R}$ , by  $r = \{x \in X \mid \langle a, x \rangle \leq b\}$ . That is,  $\mathcal{R}$  is a set system of halfspaces in  $\mathbb{R}^d$  intersected with a finite groundset. Let  $\varepsilon = \frac{1}{n}$ , then the set of points that lie on the boundary of the convex hull of  $X$  serves as an  $\varepsilon$ -net for this set system.



## 2 The $\varepsilon$ -net theorem

If  $\mathcal{R}$  is a finite set system that consists of all possible subsets of the ground set  $X$ , then no small  $\varepsilon$ -nets can exist for any  $\varepsilon < 1/2$ : no matter what set  $A \subset X$  one chooses, if  $|A| < |X|/2$ , there will be always be a set  $r \in \mathcal{R}$ , namely  $X \setminus A$ , such that  $|r| \geq |X|/2 > \varepsilon|X|$  while  $r \cap A = \emptyset$ . However, if  $\mathcal{R}$  is a finite set system of small VC-dimension, small  $\varepsilon$ -nets do exist:

**Theorem 21.4.** *Let  $\mathcal{R}$  be a finite set system with ground set  $X$  and VC-dimension at most  $d$  and let  $0 < \varepsilon < \frac{1}{2}$  be a parameter, then there exists an  $\varepsilon$ -net of size in  $O(\frac{d}{\varepsilon} \ln \frac{1}{\varepsilon})$  for this set system.*

*Proof.* In our calculations, we will assume  $d \geq 2$  (the result will also hold for  $d = 1$ , since  $O(\frac{1}{\varepsilon} \ln \frac{1}{\varepsilon}) = O(\frac{2}{\varepsilon} \ln \frac{1}{\varepsilon})$ ).

Let  $s = \frac{Cd}{\varepsilon} \ln \frac{1}{\varepsilon}$  for some value of  $C$ , which will be determined later, and assume for simplicity that  $s$  is a natural number. Let  $N$  be a random sample of  $X$  generated by making  $s$  independent draws from  $X$ . We claim that  $N$  is a  $\varepsilon$ -net with positive probability.

For simplicity, assume that all  $r \in \mathcal{R}$  satisfy  $|r| \geq \varepsilon|X|$ . (Otherwise, consider the set system that only contains such sets). Let  $E_0$  be the event that there exists a set  $r \in \mathcal{R}$  with  $N \cap r = \emptyset$  (so  $N$  is not an  $\varepsilon$ -net). We want to bound  $\Pr[E_0]$  from above.

To this end, draw another sample  $M$  in the same way as above. Let  $E_1$  be the event that there exists a set  $r \in \mathcal{R}$  with  $N \cap r = \emptyset$  and  $|M \cap r| \geq k$ , where  $k = \frac{1}{2}\varepsilon s$ . For simplicity assume that  $k$  is a natural number. (Here, we treat  $N$  and  $M$  as multisets, so an element drawn multiple times is also counted the appropriate number of times.)

The remainder of the proof is based on two claims.

**Claim 21.5.**  $\Pr[E_1] \geq \frac{1}{2} \Pr[E_0]$ .

**Claim 21.6.**  $\Pr[E_1] < \frac{1}{2}$ .

Together, the two claims imply  $\Pr[E_0] < 1$ , so  $N$  is an  $\varepsilon$ -net with positive probability.  $\square$

It remains to prove the two claims.

To prove Claim 21.5, we will use the following lemma.

**Lemma 21.7.** *Let  $Y = Y_1 + Y_2 + \dots + Y_s$ , where  $Y_i$  are independent random variables with  $\Pr[Y_i = 1] = \alpha$  and  $\Pr[Y_i = 0] = 1 - \alpha$ . Then*

$$\Pr\left[Y \geq \frac{s \cdot \alpha}{2}\right] \geq \frac{1}{2}$$

*provided that  $s \cdot \alpha \geq 8$ .*

*Proof.* Recall that the variance of a random variable is defined as the mean squared difference with its mean value:

$$\text{Var}(Y) := \mathbf{E}[(Y - \mathbf{E}[Y])^2].$$

For any  $t \geq 0$  the following holds:

$$\text{Var}(Y) = \mathbf{E}[(Y - \mathbf{E}[Y])^2] \geq t^2 \cdot \Pr[(Y - \mathbf{E}[Y])^2 \geq t^2] = t^2 \cdot \Pr[|Y - \mathbf{E}[Y]| \geq t],$$

hence:

$$\Pr[|Y - \mathbf{E}[Y]| \geq t] \leq \frac{\text{Var}(Y)}{t^2} \quad (\text{Chebyshev's inequality})$$

In our case,

$$\mathbf{E}[Y] = \sum_{i=1}^s \mathbf{E}[Y_i] = s \cdot \alpha$$

and, using the fact that the variance of the sum of independent variables is the sum of their variances:

$$\text{Var}(Y) = \sum_{i=1}^s \text{Var}(Y_i) = \sum_{i=1}^s \mathbf{E}[(Y_i - \mathbf{E}[Y_i])^2] = \sum_{i=1}^s (\alpha(1 - \alpha)) \leq s \cdot \alpha$$

So, for  $t = s \cdot \alpha/2$  with  $s \cdot \alpha \geq 8$  we have by Chebyshev's inequality

$$\mathbf{Pr}\left[Y < \frac{s \cdot \alpha}{2}\right] \leq \mathbf{Pr}\left[|Y - \mathbf{E}[Y]| \geq \frac{s \cdot \alpha}{2}\right] \leq \frac{s \cdot \alpha}{(s \cdot \alpha/2)^2} = \frac{4}{s \cdot \alpha} \leq \frac{1}{2}$$

Indeed, the first inequality holds, since

$$Y < \frac{s \cdot \alpha}{2} \implies (s \cdot \alpha) - Y > \frac{s \cdot \alpha}{2} \implies |\mathbf{E}[Y] - Y| > \frac{s \cdot \alpha}{2}$$

□

*Proof of Claim 21.5.* By the definition of conditional probability

$$\mathbf{Pr}[E_1|E_0] = \frac{\mathbf{Pr}[E_1 \cap E_0]}{\mathbf{Pr}[E_0]} = \frac{\mathbf{Pr}[E_1]}{\mathbf{Pr}[E_0]}$$

since  $E_1 \subseteq E_0$ . Therefore, to show  $\mathbf{Pr}[E_1] \geq \frac{1}{2} \cdot \mathbf{Pr}[E_0]$ , it suffices to show  $\mathbf{Pr}[E_1|E_0] \geq \frac{1}{2}$ .

So assume that  $E_0$  occurs. In this case, there exists an  $r \in \mathcal{R}$  with  $N \cap r = \emptyset$ . Let's fix one such set  $r$  and denote it by  $r_N$ .

We have

$$\mathbf{Pr}[E_1|E_0] \geq \mathbf{Pr}[|M \cap r_N| \geq k]$$

by the definition of these events. Note that we have an inequality here since there could be more than one set that is hit by  $M$  many times.

Now define random variables  $Y_1, \dots, Y_s$  with  $Y_i = 1$  if and only if the  $i$ th sample point of  $M$  falls into  $r_N$ . Observe that

$$\mathbf{Pr}[Y_i = 1] = \frac{|r_N|}{|X|} \geq \varepsilon$$

Let  $Y = \sum_{i=1}^s Y_i$  and observe that  $Y = |M \cap r_N|$ .

Lemma 21.7 implies

$$\mathbf{Pr}\left[|M \cap r_N| \geq \frac{s \cdot \alpha}{2}\right] \geq \frac{1}{2}$$

where  $\alpha$  denotes the probability  $\mathbf{Pr}[Y_i = 1]$ . (To satisfy the conditions of Lemma 21.7 we need to choose  $C$  sufficiently large)

Since  $k = \varepsilon s/2$  and since  $\alpha \geq \varepsilon$  we have that

$$|M \cap r_N| \geq \frac{s \cdot \alpha}{2} \implies |M \cap r_N| \geq k$$

Therefore,

$$\mathbf{Pr}[E_1|E_0] \geq \mathbf{Pr}[|M \cap r_N| \geq k] \geq \mathbf{Pr}\left[|M \cap r_N| \geq \frac{s \cdot \alpha}{2}\right] \geq \frac{1}{2}$$

□

*Proof of Claim 21.6.* Instead of choosing  $N$  and  $M$  directly as above, we generate the samples as follows. We first draw a sequence  $A = (z_1, \dots, z_{2s})$  of  $2s$  independent random draws from  $X$ . Then in the second step, we randomly choose  $s$  positions in  $A$  (each subset of  $s$  positions out of  $2s$  having the same probability) and put the elements at the chosen positions into  $N$  and the remaining elements into  $M$ . The resulting distribution of  $N$  and  $M$  is the same as before.

Now, consider  $A$  to be fixed and consider the subsystem  $\mathcal{R}|_A$ . Fix a set  $r \in \mathcal{R}|_A$ . Let  $E_r$  be the event that  $N \cap r = \emptyset$  and  $|M \cap r| \geq k$  (Note that this is like  $E_1$ , but for a fixed set  $r$ ). We distinguish two cases: (i)  $|r| = |A \cap r| < k$ , and (ii)  $|r| = |A \cap r| \geq k$ .

In the first case,  $E_r$  cannot occur, so  $\Pr[E_r] = 0$ .

In the second case,  $N \cap r = \emptyset$  implies  $|M \cap r| \geq k$ , so  $\Pr[E_r] = \Pr[N \cap r = \emptyset]$ . This is the probability that a random sample of  $s$  positions out of  $2s$  avoids the at least  $k$  positions occupied by  $A \cap r$ . Therefore we have

$$\Pr[N \cap r = \emptyset] \leq \left(1 - \frac{k}{2s}\right)^s \leq e^{-(k/2s) \cdot s} = e^{-k/2}$$

where the last inequality follows from the fact that  $1 - x \leq e^{-x}$  for  $x \in [0, 1]$ .

Recall that  $k = \frac{1}{2}\varepsilon s = \frac{1}{2}Cd \ln \frac{1}{\varepsilon}$ , so

$$e^{-k/2} = e^{-\frac{1}{4}Cd \ln(1/\varepsilon)} = \varepsilon^{\frac{1}{4}Cd}$$

Thus we have, for any fixed  $A$  and fixed  $r \in \mathcal{R}|_A$ :

$$\Pr[E_r] \leq \varepsilon^{\frac{1}{4}Cd}.$$

By Theorem 23.9 the number of distinct sets  $r \in \mathcal{R}_A$  is at most

$$\Pi_{\mathcal{R}}(|A|) \leq \left(\frac{e \cdot 2s}{d}\right)^d = \left(\frac{e \cdot 2C}{\varepsilon} \ln \frac{1}{\varepsilon}\right)^d \leq \left(\frac{e \cdot 2C}{\varepsilon^2}\right)^d$$

Thus, for any fixed  $A$  we can use a union bound over all  $r \in \mathcal{R}|_A$  and obtain

$$\Pr[E_1] \leq \sum_{\text{fixed } A} \Pr[E_r] \leq |\mathcal{R}|_A| \cdot \varepsilon^{\frac{1}{4}Cd} \leq \left(\frac{e \cdot 2C}{\varepsilon^2}\right)^d \cdot \varepsilon^{\frac{1}{4}Cd} = \left(e \cdot 2C \cdot \varepsilon^{\frac{1}{4}C-2}\right)^d < \frac{1}{2}$$

if  $C$  is sufficiently large (50 will do),  $d \geq 1$ , and  $\varepsilon \leq \frac{1}{2}$ .

Now, since this bound holds as a worst-case bound for any sample  $A$ , it also holds as a bound on the probability of  $E_1$  over all possible samples  $A$ . In particular, we can define  $F_A$  as the event that the set  $A$  is chosen in the first draw of  $2s$  elements from  $X$  and we can think of the above probability as conditioned on the event  $F_A$  for a fixed set  $A$ . Now, by the law of total probability we have

$$\Pr[E_1] = \sum_{F_A} \Pr[E_1|F_A] \cdot \Pr[F_A] \leq \max_{F_A} (\Pr[E_1|F_A]) \cdot \sum_{F_A} \Pr[F_A] = \max_{F_A} \Pr[E_1|F_A] < \frac{1}{2}.$$

□

### 3 Extension to infinite set systems

**Definition 21.8** ( $\varepsilon$ -net). *Let  $\mathcal{R}$  be a set system with ground set  $X$  and let  $\varepsilon \in [0, 1]$  be a real value. Let  $\mathcal{D}$  be a probability distribution defined on  $X$ . A subset  $A \subseteq X$  is an  $\varepsilon$ -net for this set system, if  $A \cap r \neq \emptyset$  for all  $r \in \mathcal{R}$  with  $\Pr[x \in r] \geq \varepsilon$  for a random sample  $x$  drawn from  $\mathcal{D}$ .*

Note that this is an extension of the definition of  $\varepsilon$ -net that we used above, since for finite set systems we can choose the uniform distribution over  $X$  for  $\mathcal{D}$  and then the two definitions are equivalent. Using this definition and the above proof one can show with minor modifications that the following extended theorem holds true.

**Theorem 21.9.** *Let  $\mathcal{R}$  be a set system with ground set  $X$  and VC-dimension at most  $d$  and let  $0 < \varepsilon < \frac{1}{2}$  be a parameter. Let  $\mathcal{D}$  be a probability distribution defined on  $X$ . There exists an  $\varepsilon$ -net of size in  $O(\frac{d}{\varepsilon} \ln \frac{1}{\varepsilon})$  for this set system.*

### References

- Sarel Har-Peled, Chapter 5 in *Geometric Approximation Algorithms*. AMS Mathematical Surveys and Monographs. 2011.
- Jiří Matoušek, Chapters 10.2 and 10.3 in *Lectures on Discrete Geometry*, Springer Graduate Texts in Mathematics.