Primary keys shown.

Foreign keys = FK coded.

Cardinalities in **bold** as **one** to **one** or **one** too **many** in relationships.

Below are:

ERD entities with attributes

Assumptions

Dependencies ->

Relations with PK underlined and coded with FK=.

Relationships with cardinalities in bold

CUSTOMER entity

CUS CODE = PK

CUS LNAME

CUS AREACODE

CUS INITIAL

CUS BALANCE

CUS PHONE

CUS_FNAME

Assume:

- 1. Customer name has only one balance.
- 2. Enforce one phone per customer see notes**
- 3. Not shown here: CUS BALANCE derived from LINE PRICE*LINE UNITS*invoices.

Dependencies:

ERD shows: <u>CUS_CODE</u> —> CUS_FNAME, CUS_INITIAL, CUS_LNAME, CUS_AREACODE,

CUS_PHONE, CUS_BALANCE

transitive: CUS FNAME, CUS INITIAL, CUS LNAME -> CUS AREACODE, CUS PHONE

transitive: CUS_FNAME, CUS_INITIAL, CUS_LNAME —> CUS_BALANCE

Relations:

CUSTOMER NAMES(CUS CODE, CUS FNAME, CUS INITIAL, CUS LNAME)

FK = CUS FNAME, CUS INITIAL, CUS LNAME

CUSTOMER PHONE(CUS FNAME, CUS INITIAL, CUS LNAME, CUS AREACODE, CUS PHONE)

CUSTOMER_BALANCE(CUS FNAME, CUS INITIAL, CUS LNAME, CUS_BALANCE)

Relationships:

CUSTOMER NAMES can have one CUSTOMER PHONE

CUSTOMER PHONE assigned to many CUSTOMER NAMES

CUSTOMER NAMES can have one CUSTOMER BALANCE

CUSTOMER BALANCE assigned one CUSTOMER NAME

CUSTOMER NAMES can have many INVOICE (see next entity's relations)

INVOICE must have one CUSTOMER_NAMES

INVOICE entity

INV NUMBER = PK INV_DATE CUS DATE

Assume:

- 1. CUS DATE is date customer ordered
- 2. INV DATE is date customer order is sourced

Dependencies:

INV NUMBER -> INV_DATE, CUS_DATE

Relation:

INVOICE(INV NUMBER, INV DATE, CUS DATE)

Relationships:

INVOICE can have **many** LINE (see next entity's relations) LINE must have **one** INVOICE

LINE entity

INV NUMBER = PK = invoice number

LINE NUMBER = PK = A line is a line on an invoice.

LINE_UNITS = how many units ordered by customer

P_CODE = product and price code

LINE PRICE = price per unit, not total for the line.

Assume:

No partial dependency: P_CODE fully functionally dependent on INV_NUMBER, LINE_NUMBER LINE_PRICE = P_PRICE - P_DISCOUNT derived based on P_CODE on invoice.

Dependencies:

INV_NUMBER, LINE_NUMBER —> LINE_UNITS, P_CODE, LINE_PRICE transitive: P_CODE —> LINE_PRICE

Relations:

INVOICE_LINE(<u>INV_NUMBER</u>,LINE_NUMBER, LINE_UNITS, P_CODE)

FK = P_CODE

PRICE_PER(<u>P_CODE</u>,LINE_PRICE)

Relationships:

INVOICE_LINE has **one** PRICE_PER
PRICE_PER can have **many** INVOICE_LINE
PRICE_PER has **one** PRODUCT_AT_THIS_PRICE (see next entity's relations)
PRODUCT_AT_THIS_PRICE can have **many** PRICE_PER

PRODUCT entity

P CODE = PK = multidimensional product price code to determine LINE PRICE

P_PRICE = price before discount

P DISCOUNT = discount offered

P_DESCRIPTION = Specific grade; P_DISCOUNT less if P_DESCRIPTION better grade

V CODE = vendor A charges a different price and offer a different grade

P_QOH = quantity on hand; P_DISCOUNT less if P_QOH less

P INDATE = next back order date; later means bigger P DISCOUNT; not a f(P QOH)

P MIN = minimum order size from our warehouse; no impact on PRICE or DISCOUNT

Assume:

Customer specifies product and description (grade) in a PCODE that determine discount and via product description determine base price and vendor code (I considered product description determine both base price and product description simultaneously).

P_DESCRIPTION, P_DISCOUNT=f(P_CODE)

P_PRICE = g(P_DESCRIPTION)

V CODE = f(P DESCRIPTION)

P DISCOUNT=g(V CODE, P QOH, P INDATE)

Dependencies:

P CODE -> P DESCRIPTION, P PRICE, V CODE, P QOH, P INDATE, P DISCOUNT, P MIN

transitive: P DESCRIPTION -> P PRICE

transitive: P DESCRIPTION —> V CODE (in DISCOUNT relation)

transitive: V_CODE, P_QOH, P_INDATE -> P_DISCOUNT

Relations:

PRODUCT AT THIS PRICE(P CODE, P DESCRIPTION, P QOH, P INDATE)

FK = P_DESCRIPTION

FK composite = P_DESCRIPTION, P_QOH, P_INDATE

BASE_PRICE(P_DESCRIPTION, P_PRICE)

DISCOUNT(P DESCRIPTION, P QOH, P INDATE, PDISCOUNT)

FK = P DESCRIPTION

VENDORS(P_DESCRIPTION, V_CODE)

Relationships:

PRODUCT AT THIS PRICE is a factor of one BASE PRICE

BASE_PRICE is a factor yielding many PRODUCT_AT_THIS_PRICE

PRODUCT_AT_THIS_PRICE is a factor of **one** DISCOUNT

DISCOUNT is a factor yielding many PRODUCT AT THIS PRICE

DISCOUNT can be achieved from many VENDORS

VENDORS offer one DISCOUNT

PRODUCT AT THIS PRICE sources from many VENDOR REQUEST (see next entity's relations)

VENDOR REQUEST are for one PRODUCT AT THIS PRICE

VENDOR entity

V CODE

V_ORDER

V NAME

V STATE

V CONTACT

V AREACODE

V PHONE

Assume:

- 1. V CODE is unique combination of V NAME vendor, V ORDER type
- 2. V ORDER is code for order details not documented here: product, speed of delivery, etc
- 3. Each V NAME only has one V CONTACT and V STATE
- 4. Each V CONTACT has one V AREACODE and V PHONE
- 5. A contact can help us order more than once.

Dependencies:

V CODE -> V ORDER, V_NAME, V_STATE, V_CONTACT, V_AREACODE, V_PHONE,,

transitive: V_NAME —> V_STATE, V_CONTACT

transitive: V_CONTACT —> V_AREACODE, V_PHONE

Relations:

VENDOR REQUEST(V CODE, V ORDER, V NAME)

 $FK = V_NAME$

WAREHOUSE(VNAME, V STATE, V CONTACT)

FK = V CONTACT

SALES PERSON(V CONTACT, V AREACODE, V PHONE)

Relationships:

VENDOR_REQUEST is sourced by **one** WAREHOUSE WAREHOUSE can supply **many** VENDOR_REQUEST WAREHOUSE has **many** SALE_PERSON SALES PERSON works at **one** WAREHOUSE

Notes for #2 assumption on CUSTOMER (not essential reading)

Enforce one phone per customer ** see notes

- 1. CUS FNAME, CUS INITIAL, CUS LNAME —> CUS AREACODE, CUS PHONE 1:1
- Allow phone number identify more than one customer (spouse e.g.):
 CUS AREACODE, CUS PHONE does not —> CUS FNAME, CUS INITIAL, CUS LNAME 1:*
- 3. Reason for #2 transitive:

a:* possible for CUS_FNAME, CUS_INITIAL, CUS_LNAME: CUS_AREACODE, CUS_PHONE only if 1:* here

CUS_FNAME, CUS_INITIAL, CUS_LNAME : CUS_CODE

and

CUS_FNAME, CUS_INITIAL, CUS_LNAME : CUS_CODE

which allows 1:* here

CUS_FNAME, CUS_INITIAL, CUS_LNAME: CUS_BALANCE

so don't allow 2 phones per person.

but can allow 2 persons per phone

which allows spouse to open second cust code under same home phone

A->B, C, D where

A = cust code

B = person

C = phone number

D = balance

B must —> C if we want to limit B to one A and limit B to one D

if B does not —> C then two C per B only possible when assign a second A to B

e.g. would be possible if B not —> C

A=1 B=3 C=4

A=2 B=3 C=5

but undesirable when A—>D because if B can have two C via two As then

B can have two D via two As

e.g. would be possible if B not —> C

A=1 B=3 C=4 D=6

A=2 B=3 C=5 D=7

person must —> phone if want to limit person to one cust code and one balance

C not —> B ok if allow C to more than one A and allow C to more than one D

if C not—>B then a second B assigned to C allows requires second A assigned to C and a second D to B.

e.g. would be possible

A=1 B=3 C=5 D=6

A=2 B=4 C=5 D=7

phone can be assigned to more than one person

means a phone assigned to more than one account

means a phone to more than one balance

if e.g. a spouse opens a cust code under same phone number