

iptables reference

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iptables.h

#define's

IPT__x

IPT_KEYOFFSET the offset to the actual key in the byte array

IPT_KEYLEN(k) the length of byte array **k**

IPT_KEYPTR(k) pointer to the location of the key in the byte array **k**

IP4__x

IP4_KEYLEN the length of the byte array to hold an IPv4 binary key

IP4_MAXMASK the maximum number of bits in an IPv4 binary mask

IP4_PFXSTRLEN buffer size for an ipv4 prefix string (addr/len)

IP6__x

IP6_KEYLEN the length of the byte array to hold an IPv6 binary key

IP6_MAXMASK the maximum number of bits in an IPv6 binary mask

IP6_PFXSTRLEN buffer size for an ipv6 prefix string (addr/len)

STR__x

STR_IS_IP4(s) detects an ipv4 address by an in-string **'-'**-character

STR_IS_IP6(s) detects an ipv6 address by an in-string **'.'**-character

KEY__x

KEY_IS_IP4(k) checks for an ipv4 binary key by checking the byte array's length

KEY_IS_IP6(k) checks for an ipv6 binary key by checking the byte array's length

KEY_AF_FAM(k) returns **AF_INET(6)** or **AF_UNSPEC** based on binary key length

KEY_LEN_FAM(*af*) returns byte array length for given AF family.

MAX_*x*

MAX_BINKEY buffer size to hold both ipv4/ipv6 binary keys

MAX_STRKEY buffer size to hold both ipv4/ipv6 prefix/len strings

AF_*x*

AF_UNKNOWN(*f*) true if *f* is not AF_INET or AF_INET6

RDX_*x*

RDX_ISLEAF(*rn*) true if radix node *rn* is a LEAF node

RDX_ISINTERNAL(*rn*) true if radix node *rn* is an INTERNAL node

RDX_ISROOT(*rn*) true if radix node *rn* is a ROOT node (LE-mark, TOP, RE-mark)

MSK_ISROOT(*rm*) true if radix mask node *rm* is a ROOT node (LE,TOP,RE)

RDX_ISRCHILD(*rn*) true if radix node *rn* is its parent right child

RDX_MAX_KEYLEN the maximum length for a byte array.

RDX FLAG

IPTF_DELETE additional radix node flag to indicate node was deleted

When an iteration is happening, radix nodes are flagged for deletion rather than being removed immediately. All iterators run with an upvalue that has a garbage collector which will eventually remove radix nodes flagged for deletion when it is safe to do so.

The flag is set in **rn_flags** in a radix node, alongside the flags defined by **radix.h**

RDX node types

A table has a stack which allows for pushing arbitrary data combined with a type identifier for that data. The stack is only used by the radix iterator which is read only and allows for graphing a radix tree. See 'ipt2dot.lua'

Radix node type include:

- **TRDX_NONE** – indicates an unknown type
- **TRDX_NODE_HEAD** – a radix_node_head
- **TRDX_HEAD** – a radix head, a struct inside a radix node head

- ‘TRDX_NODE’ – a single radix node
- TRDX_MASK_HEAD – the head of the mask tree
- TRDX_MASK – a radix mask node

Structures

entry_t

The type `entry_t` has 2 members:

- `rn[2]`, an array of two radix nodes: a leaf & an internal node.
- `void *value`, which points to user data.

The radix tree stores/retrieves pointers to **radix leaf nodes** using binary keys. So a user data structure must begin with an array of two radix nodes: a leaf node for storing the binary key, associated with the user data, and an internal node to actually insert the leaf node into the tree.

Retrieving the data is done by matching, either exact or using a longest prefix match, against a binary key which yields a pointer to a leaf node. That pointer is then recast to `entry_t *` in order to access the user data associated with the matched binary key in the tree via the `value` pointer.

purge_t

The type `purge_t` has the following members:

- `struct radix_node_head *head`, the head of a radix tree
- `purge_f_t *purge`, a callback function pointer; to free user data
- `void *args`, an opaque pointer to be interpreted by `purge`

This structure is used to relay contextual arguments to the user callback function upon deletion time. The different levels of memory ownership include:

- `radix.c`, which owns:
 - the ‘mask’ tree completely, but
 - only the `radix_node_head` for the ‘key’-tree, not the leafs
- `iptable.c`, which owns:
 - the radix nodes that are part of `entry_t`, and
 - the binary key which it derives from `const char *` strings
- `user.c`, she owns:
 - the memory pointed to by the `void *value` pointer in `entry_t`

So when `user.c` calls `iptable.c`’s `tbl_create`, it supplies a *purge* function whose signature is:

```
typedef void purge_f_t(void *pargs, void *value);
```

When `user.c` calls `iptable.c`'s `tbl_del`, it supplies the prefix string to be deleted and a contextual argument for its `purge` callback function (`pargs` here). `tbl_del` then:

- derives a binary from the prefix given
- calls `radix.c`'s `rnhdeladdr` to remove the binary key, (if succesful, two radix nodes are returned (ie an `entry_t`))
- frees the `entry_t`'s memory and the binary key memory, and finally
- calls back the `purge` function supplying it with both the
 - `void *pargs`, the contextual argument for this deletion, and
 - `void *value`, from the `entry_t` that was deleted.

It may seem a bit convoluted, but it allows the user's `purge` callback to examine a request to free user memory in context.

`stackElm_t`

A stack element has members:

- `int type`, denotes the type of this element
- `void *elm`, an opaque element pointer
- `struct stackElm_t *next`, pointer to the next stack element

An iptable contains a stack, which is only used by `rdx_firstnode` and `rdx_nextnode` to perform a preorder (node, left, right) traversal and node processing of all nodes of all types of all trees used in the table. I.e. it traverses both the 'key'-tree as well as the 'mask'-tree and yields all nodes of all types used in either tree.

This makes it possible to produce a dot-file of a radix tree and create an image using graphviz's dot tool.

`table_t`

An iptable has the following members:

- `struct radix_node_head *head4`, the head of the ipv4 radix tree
- `struct radix_node_head *head6`, the head of the ipv6 radix tree
- `size_t count4`, the number of ipv4 prefixes present in the ipv4 tree
- `size_t count6`, the number of ipv6 prefixes present in the ipv6 tree
- `purge_f_t *purge`, user callback for freeing user data
- `int itr_lock`, indicates the presence of active iterators
- `stackElm_t *top`, the stack to iterate across all radix nodes in all trees
- `size_t size`, the current size of the of the stack

Two separate radix trees are used to store ipv4 resp. ipv6 binary keys. Table operations detect the type of prefix used and access the corresponding tree. Each time a prefix is added or deleted, the tree's counter is updated. The `purge`

function pointer is supplied upon tree creation time and is called whenever user data can be freed, see `purge_t`.

The `itr_lock` is actually a Lua specific feature to track the presence of any currently active tree iterators (there are a few). This allows for postponed radix node removal while some iterator is still traversing one of the trees.

Finally, the `*top` and `size` exist in order to be able to graph the tree(s).

iptables.c

`max_mask`

```
uint8_t max_mask[RDX_MAX_KEYLEN] = {-1, .., -1};
```

`max_mask` serves as an AF family's default mask in case one is needed, but not supplied. Thus missing masks are interpreted as host masks.

Key functions

`key_alloc`

```
uint8_t *key_alloc(int af);
```

Allocate space for a binary key of AF family type `af` and return its pointer. Supported `af` types include:

- `AF_INET`, for ipv4 protocol family
- `AF_INET6`, for the ipv6 protocol family

`key_copy`

```
uint8_t *key_copy(uint8_t *src);
```

Copies binary key `src` into newly allocated space and returns its pointer. The `src` key must have a valid first length byte. $0 \leq KEY_LEN(src) \leq MAX_KEYLEN$. Returns `NULL` on failure.

`key_bystr`

```
uint8_t *key_bystr(uint8_t *dst, int *mlen, int *af, const char *s);
```

Store string `s`' binary key in `dst`. Returns `NULL` on failure. Also sets `mlen` and `af`. `mlen=-1` when no mask was supplied. Assumes `dst` size `MAX_BINKEY`, which fits both ipv4/ipv6.

key_byfit

```
uint8_t *key_byfit(uint8_t *m, uint8_t *a, uint8_t *b)
```

Set *m* to the largest possible mask for key *a* such that:

- *a*'s network address is still *a* itself, and
- *a*'s broadcast address is less than, or equal to, *b* address

Note: the function assumes $a \leq b$.

key_bylen

```
uint8_t *key_bylen(uint8_t *binkey, int mlen, int af);
```

Create a mask for given *af* family and mask length *mlen*.

key_bypair

```
uint8_t * key_bypair(uint8_t *a, const void *b, const void *m);
```

Set key *a* such that *a*/*m* and *b*/*m* are a pair that fit in *key*/*m*-1. Assumes masks are contiguous.

key_masklen

```
int key_masklen(void *key);
```

Count the number of consecutive 1-bits, starting with the msb first.

key_tostr

```
const char *key_tostr(char *dst, void *src);
```

src is byte array; 1st byte is usually total length of the array.

- iptable.c keys/masks are `uint8_t *`'s (unsigned)
- radix.c keys/masks are `char *`'s. (may be signed; sys dependent)
- radix.c keys/masks's KEYLEN may deviate: for masks it may indicate the total of non-zero bytes i/t array instead of its total length.

key_tostr_full

```
const char *key_tostr_full(char *dst, void *src);
```

Return the full string for a key without shortending contiguous zero's for ipv6 keys. If the src represents a mask, it may be shorter than than the protocol's actual key-length, so supply trailing zeros as well. See the remarks for `key_tostr`. Only has effect for ipv6 keys. Embedded ipv4 addresses are printed as hex digits, not as integers.

key_bynum

```
int key_bynum(void *key, size_t number);
```

Create key from number. Returns NULL on failure, key otherwise.

key_incr

```
uint8_t *key_incr(uint8_t *key, size_t num);
```

Increment `key` with `num`. Returns `key` on success, NULL on failure (e.g. when wrapping around the available address space) which usually means the resulting `key` value is meaningless.

key_decr

```
uint8_t *key_decr(uint8_t *key, size_t num);
```

Decrement `key` with `num`. Returns `key` on success, NULL on failure (e.g. when wrapping around the available address space) which usually means the resulting `key` value is meaningless.

key_invert

```
int key_invert(void *key);
```

inverts a key, usefull for a mask. Assumes the LEN-byte indicates how many bytes must be (and can be safely) inverted.

key_reverse

```
int key_reverse(void *key);
```

reverse the bytes of a key. For ipv6, the nibbles are also swapped. Assumes the LEN-byte indicates how many bytes must be (and can be safely) reversed.

key_network

```
int key_network(void *key, void *mask);
```

key_broadcast

```
int key_broadcast(void *key, void *mask);
```

set key to broadcast address, using mask

- 1 on success, 0 on failure
- mask LEN ≤ key LEN, ‘missing’ mask bytes are taken to be 0x00 (a radix tree artifact).

key_cmp

```
int key_cmp(void *a, void *b);
```

Returns -1 if a<b, 0 if a==b, 1 if a>b; or -2 on errors

key_isin

```
int key_isin(void *a, void *b, void *m);
```

return 1 iff a/m includes b, 0 otherwise note:

- also means b/m includes a
- any radix keys/masks may have short(er) KEYLEN’s than usual

radix node functions

__dumprn‘

```
void _dumprn(const char *s, struct radix_node *rn);
```

dump radix node characteristics to stderr

rdx_flush

```
int rdx_flush(struct radix_node *rn, void *args);
```

free user controlled resources.

Called by walktree, rdx_flush:

- frees the key and, if applicable,

- uses the purge function to allow user controlled resources to be freed. The purge function is supplied at tree creation time.

Note: `rdx_flush` is called from `walktree` and only on *LEAF* nodes, so the `rn` pointer is cast to pointer to `entry_t`. As a `walktree_f_t`, it always returns the value 0 to indicate success, otherwise the walkabout would stop.

`rdx_firstleaf`

```
struct radix_node *rdx_firstleaf(struct radix_head *rh);
```

Find first non-ROOT leaf of in the prefix or mask tree Return NULL if none found.

Note: the /0 mask is never stored in the mask tree even if stored explicitly using prefix/0. Hence, the /0 mask won't be found by this function.

`rdx_nxtleaf`

```
struct radix_node *rdx_nextleaf(struct radix_node *rn);
```

Given a `radix_node` (INTERNAL or LEAF), find the next leaf in the tree. Note that the caller must check the `IPTF_DELETE` flag herself. Reason is to allow upper logic to be performed across both deleted and normal leaves, such as actually deleting the flagged nodes..

`pairleaf`

```
struct radix_node *pairleaf(struct radix_node *oth);
```

Find and return the leaf whose key forms a pair with the given leaf node such that both fit in an enclosing supernet with the given leaf's masklength-1. Note:

- 0/0 is the 'ultimate' supernet which combines 0.0.0.0/1 and 128.0.0.0/1

`rdx_firstnode`

```
int rdx_firstnode(table_t *t, int af_fam);
```

initialize the stack with the radix head nodes

`rdx_nextnode`

```
int rdx_nextnode(table_t *t, int *type, void **ptr);
```

pop top and set type & node, push its progeny

- stackpush will ignore NULL pointers, so its safe to push those return 1 on success with type,ptr set return 0 on failure (eg stack exhausted or unknown type)

table functions

tbl_create

```
table_t *tbl_create(purge_f_t *fp):
```

Create a new iptable with 2 radix trees.

tbl_walk

```
int tbl_walk(table_t *t, walktree_f_t *f, void *fargs);
```

run f(args, leaf) on leafs in IPv4 tree and IPv6 tree

tbl_destroy

```
int tbl_destroy(table_t **t, void *pargs);
```

Destroy table, free all resources owned by table and user

- return 1 on success, 0 on failure

tbl_get

```
entry_t *tbl_get(table_t *t, const char *s);
```

Get an exact match for addr/mask prefix.

tbl_set

```
int tbl_set(table_t *t, const char *s, void *v, void *pargs);
```

tbl_del

```
int tbl_del(table_t *t, const char *s, void *pargs);
```

tbl_lpm

```
entry_t *tbl_lpm(table_t *t, const char *s);
```

`tbl_lsm`

```
struct radix_node *tbl_lsm(struct radix_node *rn);
```

Given a leaf, find a less specific leaf or fail

- used by `tbl_lpm` in case the match is flagged for deletion

`tbl_stackpush`

```
int tbl_stackpush(table_t *t, int type, void *elm);
```

`tbl_stackpop`

```
int tbl_stackpop(table_t *t);
```