



TLA+

Distributed Systems

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توصیف رسمی

نمایش ریاضی سیستم

$/, *, -, +$

● جبر

$\wedge \vee \neg \Rightarrow \equiv$

● منطق

● مجموعه ها

\cap intersection \cup union \subseteq subset \setminus set difference

Propositional Logic

\wedge conjunction (and)

\vee disjunction (or)

\neg negation (not)

\Rightarrow implication (implies)

\equiv equivalence (is equivalent to)

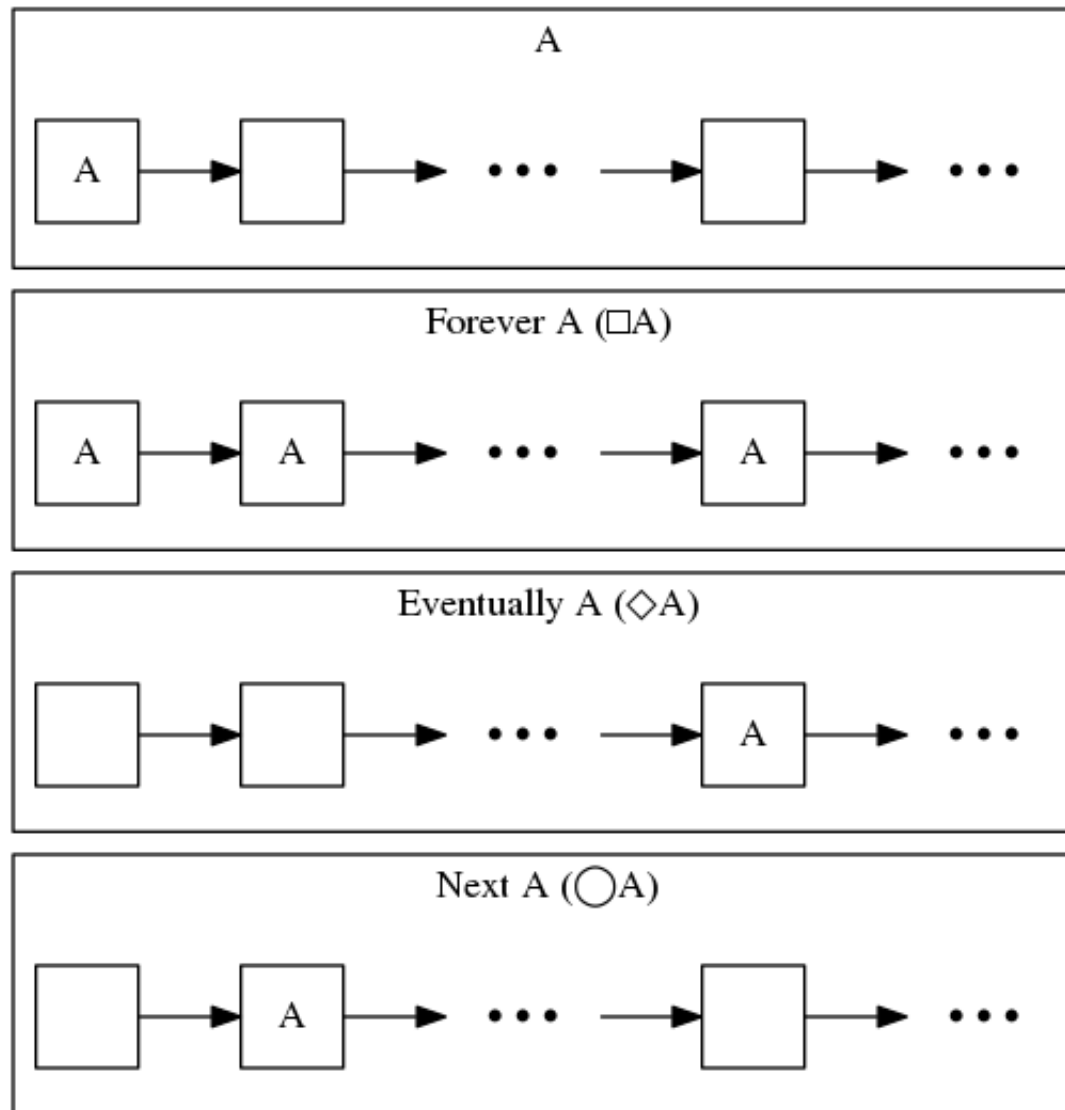
Predicate Logic

Predicate logic extends propositional logic with the two quantifiers:

\forall universal quantification (for all)

\exists existential quantification (there exists)

Temporal Logic



Specifying a simple clock

To specify the hour clock, we describe **all its possible behaviors**.

We write an **initial predicate** that species the possible initial values of hr , and a next-state relation that species how the value of hr can change in any step.

$$HCini \triangleq hr \in \{1, \dots, 12\}$$

... is informal.

$$HCnext \triangleq hr' = \text{IF } hr \neq 12 \text{ THEN } hr + 1 \text{ ELSE } 1$$

The temporal formula $\Box F$ asserts that formula F is always true.

In particular, $\Box HCnext$ is the assertion that $HCnext$ is true for every step in the behavior.

Weather station

$$\begin{aligned} \begin{bmatrix} hr & = & 11 \\ tmp & = & 23.5 \end{bmatrix} &\rightarrow \begin{bmatrix} hr & = & 12 \\ tmp & = & 23.5 \end{bmatrix} \rightarrow \begin{bmatrix} hr & = & 12 \\ tmp & = & 23.4 \end{bmatrix} \rightarrow \\ &\begin{bmatrix} hr & = & 12 \\ tmp & = & 23.3 \end{bmatrix} \rightarrow \begin{bmatrix} hr & = & 1 \\ tmp & = & 23.3 \end{bmatrix} \rightarrow \dots \end{aligned}$$

$$HCini \wedge \Box HCnext$$

$$HCini \wedge \Box (HCnext \vee (hr' = hr))$$

$$HCini \wedge \Box [HCnext]_{hr}$$

TLA+

- Reserved words that appear in small upper-case letters (like EXTENDS) are written in ASCII with ordinary upper-case letters.
- When possible, symbols are represented pictorially in ASCII—for example, \square is typed as `[]` and \neq as `#`. (You can also type \neq as `/=`.)
- When there is no good ASCII representation, T_EX notation is used—for example, \in is typed as `\in`. The major exception is \triangleq , which is typed as `==`.

Hour Clock

MODULE *HourClock*

EXTENDS *Naturals*

VARIABLE *hr*

$HCini \triangleq hr \in (1 \dots 12)$

$HCnxt \triangleq hr' = \text{IF } hr \neq 12 \text{ THEN } hr + 1 \text{ ELSE } 1$

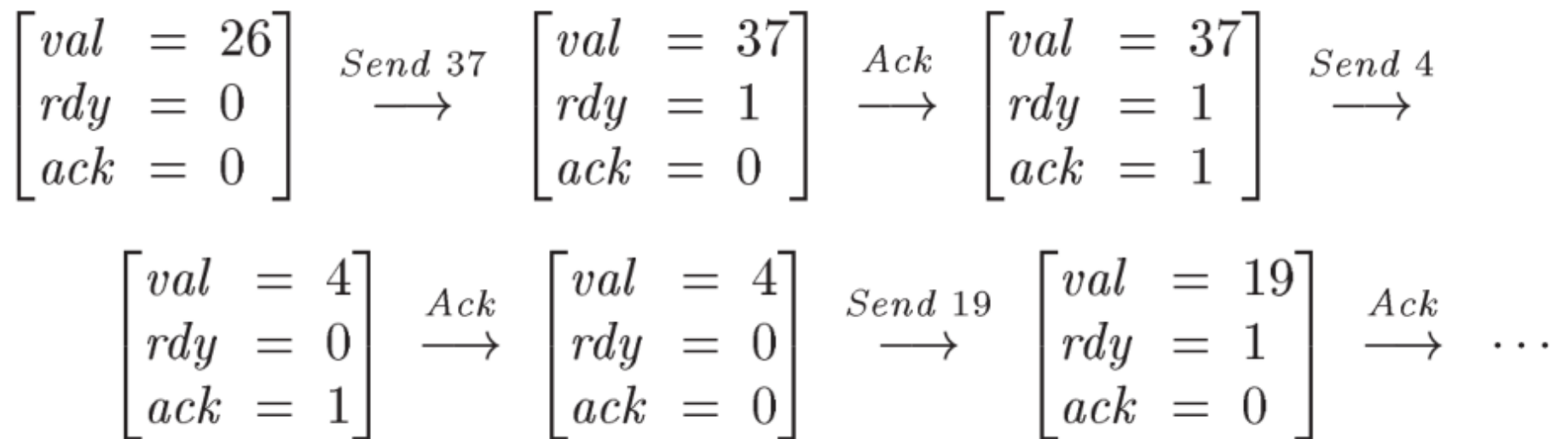
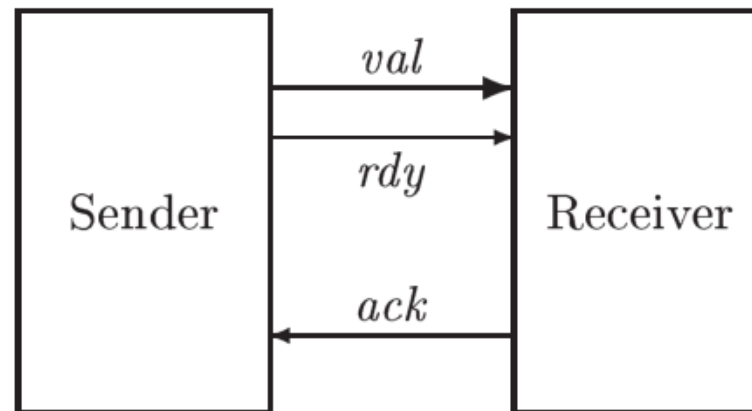
$HC \triangleq HCini \wedge \Box[HCnxt]_{hr}$

THEOREM $HC \Rightarrow \Box HCini$

Hour Clock

```
----- MODULE HourClock -----  
EXTENDS Naturals  
VARIABLE hr  
HCini == hr \in (1 .. 12)  
HCnxt == hr' = IF hr # 12 THEN hr + 1 ELSE 1  
HC == HCini /\ [] [HCnxt]_hr  
-----  
THEOREM HC => [] HCini  
=====
```

An Asynchronous Interface



EXTENDS *Naturals*

CONSTANT *Data*

VARIABLES *val, rdy, ack*

$$\begin{aligned} \textit{TypeInvariant} \triangleq & \wedge \textit{val} \in \textit{Data} \\ & \wedge \textit{rdy} \in \{0, 1\} \\ & \wedge \textit{ack} \in \{0, 1\} \end{aligned}$$

$$\begin{aligned} \textit{Init} \triangleq & \wedge \textit{val} \in \textit{Data} \\ & \wedge \textit{rdy} \in \{0, 1\} \\ & \wedge \textit{ack} = \textit{rdy} \end{aligned}$$

$$\begin{aligned} \textit{Send} \triangleq & \wedge \textit{rdy} = \textit{ack} \\ & \wedge \textit{val}' \in \textit{Data} \\ & \wedge \textit{rdy}' = 1 - \textit{rdy} \\ & \wedge \text{UNCHANGED } \textit{ack} \end{aligned}$$

$$\begin{aligned} \textit{Rcv} \triangleq & \wedge \textit{rdy} \neq \textit{ack} \\ & \wedge \textit{ack}' = 1 - \textit{ack} \\ & \wedge \text{UNCHANGED } \langle \textit{val}, \textit{rdy} \rangle \end{aligned}$$

$$\textit{Next} \triangleq \textit{Send} \vee \textit{Rcv}$$

$$\textit{Spec} \triangleq \textit{Init} \wedge \Box[\textit{Next}]_{\langle \textit{val}, \textit{rdy}, \textit{ack} \rangle}$$

THEOREM $\textit{Spec} \Rightarrow \Box \textit{TypeInvariant}$

$$\begin{bmatrix} big & = 0 \\ small & = 0 \end{bmatrix}$$

The big jug is filled from the faucet.

↓

$$\begin{bmatrix} big & = 5 \\ small & = 0 \end{bmatrix}$$

The small jug is filled from the big one.

↓

$$\begin{bmatrix} big & = 2 \\ small & = 3 \end{bmatrix}$$

The small jug is emptied (onto the ground).

↓

$$\begin{bmatrix} big & = 2 \\ small & = 0 \end{bmatrix}$$

A little thought reveals that there are three kinds of steps in a behavior:

- Filling a jug.
- Emptying a jug.
- Pouring from one jug to the other. There are two cases:
 - This empties the first jug.
 - This fills the second jug, possibly leaving water in the first jug.

EXTENDS *Integers*

VARIABLES *big, small*

$$\begin{aligned} Init \quad \triangleq \quad & \wedge big = 0 \\ & \wedge small = 0 \end{aligned}$$
$$\begin{aligned} Next \quad \triangleq \quad & \vee FillSmall \\ & \vee FillBig \\ & \vee EmptySmall \\ & \vee EmptyBig \\ & \vee SmallToBig \\ & \vee BigToSmall \end{aligned}$$

$$FillSmall \triangleq small' = 3$$

$$\left[\begin{array}{l} big = 2 \\ small = 1 \end{array} \right] \rightarrow \left[\begin{array}{l} big = 2 \\ small = 3 \end{array} \right]$$

$$\left[\begin{array}{l} big = 2 \\ small = 1 \end{array} \right] \rightarrow \left[\begin{array}{l} big = \sqrt{42} \\ small = 3 \end{array} \right]$$

$$FillSmall \triangleq \begin{array}{l} \wedge small' = 3 \\ \wedge big' = big \end{array}$$

$$\begin{aligned}
 \textit{FillBig} \quad &\triangleq \quad \wedge \textit{big}' = 5 \\
 &\quad \wedge \textit{small}' = \textit{small}
 \end{aligned}$$

$$\begin{aligned}
 \textit{EmptySmall} \quad &\triangleq \quad \wedge \textit{small}' = 0 \\
 &\quad \wedge \textit{big}' = \textit{big}
 \end{aligned}$$

$$\begin{aligned}
 \textit{EmptyBig} \quad &\triangleq \quad \wedge \textit{big}' = 0 \\
 &\quad \wedge \textit{small}' = \textit{small}
 \end{aligned}$$

$$\begin{aligned}
\textit{SmallToBig} \triangleq & \quad \vee \wedge \textit{big} + \textit{small} > 5 \\
& \quad \wedge \textit{big}' = 5 \\
& \quad \wedge \textit{small}' = \textit{small} - (5 - \textit{big}) \\
& \vee \wedge \textit{big} + \textit{small} \leq 5 \\
& \quad \wedge \textit{big}' = \textit{big} + \textit{small} \\
& \quad \wedge \textit{small}' = 0
\end{aligned}$$

$$Min(m, n) \triangleq \text{IF } m < n \text{ THEN } m \text{ ELSE } n$$

$$\begin{aligned} SmallToBig &\triangleq \wedge big' = Min(big + small, 5) \\ &\quad \wedge small' = small - (Min(big + small, 5) - big) \end{aligned}$$

$$\begin{aligned} SmallToBig &\triangleq \\ \text{LET } poured &\triangleq Min(big + small, 5) - big \\ \text{IN } \quad \wedge big' &= big + poured \\ \quad \wedge small' &= small - poured \end{aligned}$$

$BigToSmall \triangleq$
LET $poured \triangleq Min(big + small, 3) - small$
IN $\wedge big' = big - poured$
 $\wedge small' = small + poured$

$$\begin{aligned} TypeOK \triangleq & \quad \wedge big \in 0 .. 5 \\ & \quad \wedge small \in 0 .. 3 \end{aligned}$$