1 Inference rules

1.1 Formation of contexts, types, and their elements

MLTT rule:

$$\frac{\Gamma, x : A \vdash B(x) \text{ type}}{\Gamma \vdash A \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \ \vdash A \equiv B \text{ type}}{\Gamma \ \vdash A \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \vdash A \equiv B \text{ type}}{\Gamma \vdash B \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \vdash a \colon\! A}{\Gamma \vdash A \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \vdash a \equiv b : A}{\Gamma \vdash a : A}$$

MLTT rule:

$$\frac{\Gamma \vdash a \equiv b : A}{\Gamma \vdash b : A}$$

1.2 Judgemental equality is an equivalence relation

MLTT rule:

$$\frac{\Gamma \vdash A \text{ type}}{\Gamma \vdash A \equiv A \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \ \vdash A \equiv B \text{ type}}{\Gamma \ \vdash B \equiv A \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \vdash A \equiv B \text{ type} \qquad \Gamma \vdash B \equiv C \text{ type}}{\Gamma \vdash A \equiv C \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \vdash a : A}{\Gamma \vdash a \equiv a : A}$$

MLTT rule:

$$\frac{\Gamma \ \vdash \ a \equiv b : A}{\Gamma \ \vdash \ b \equiv a : A}$$

MLTT rule:

$$\frac{\Gamma \vdash a \equiv b : A \qquad \Gamma \vdash b \equiv c : A}{\Gamma \vdash a \equiv c : A}$$

1.3 Variable conversion

MLTT rule:

$$\frac{\Gamma \vdash A \equiv A' \text{ type} \qquad \Gamma, x : A, \Delta \vdash B(x) \text{ type}}{\Gamma, x : A', \Delta \vdash B(x) \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \vdash A \equiv A' \text{ type} \qquad \Gamma, x : A, \Delta \vdash \mathcal{J}}{\Gamma, x : A', \Delta \vdash \mathcal{J}}$$

1.4 Substitution

Metasyntactic notation (used in many places):

t[a/x] := t with all occurences of x replaced by a

MLTT rule:

$$\frac{\Gamma \vdash a : A \qquad \Gamma, \, x : A, \, \Delta \vdash \mathcal{J}}{\Gamma, \, \Delta[a/x] \vdash \mathcal{J}[a/x]} \, S$$

MLTT rule:

$$\frac{\Gamma \vdash a \equiv a' : A \qquad \Gamma, \ x : A, \ \Delta \vdash B \text{ type}}{\Gamma, \ \Delta[a/x] \vdash B[a/x] \equiv B[a'/x] \text{ type}}$$

MLTT rule:

$$\frac{\Gamma \vdash a \equiv a' \colon A \qquad \Gamma, \ x \colon A, \ \Delta \vdash b \colon B}{\Gamma, \ \Delta[a/x] \ \vdash \ b[a/x] \equiv b[a'/x] \colon B[a/x]}$$

1.5 Weakening

MLTT rule:

$$\frac{\Gamma \vdash A \text{ type} \qquad \Gamma, \Delta \vdash \mathcal{J}}{\Gamma, x : A, \Delta \vdash \mathcal{J}} W$$

MLTT rule:

$$\frac{\Gamma \vdash A \text{ type} \qquad \Gamma \vdash B \text{ type}}{\Gamma, x : A \vdash B \text{ type}} W$$

1.6 Generic elements

MLTT rule:

$$\frac{\Gamma \vdash A \text{ type}}{\Gamma, x : A \vdash x : A} \delta$$

1.7 Manipulating variables

Derived rule:

$$\frac{\Gamma, x: A, \Delta \vdash \mathcal{J}}{\Gamma, x': A, \Delta[x'/x] \vdash \mathcal{J}[x'/x]} x'/x$$

Derived rule:

$$\frac{\Gamma \ \vdash \ B \ \text{type} \qquad \Gamma, \ x \colon\! A, \ y \colon\! B, \ \Delta \ \vdash \ \mathcal{J}}{\Gamma, \ y \colon\! B, \ x \colon\! A, \ \Delta \ \vdash \ \mathcal{J}}$$

$\mathbf{2}$ Π types

2.1 Dependent functions

2.1.1 Π formation

MLTT rule:

$$\frac{\Gamma, x: A \vdash B(x) \text{ type}}{\Gamma \vdash \Pi_{(x:A)}B(x) \text{ type}} \Pi$$

MLTT rule:

$$\frac{\Gamma \vdash A \equiv A' \text{ type} \qquad \Gamma, \ x \colon A \vdash B(x) \equiv B'(x) \text{ type}}{\Gamma \vdash \Pi_{(x \colon A)} B(x) \equiv \Pi_{(x \colon A')} B'(x) \text{ type}} \ \Pi\text{-eq}$$

2.1.2 Π introduction

MLTT rule:

$$\frac{\Gamma, x : A \vdash b(x) : B(x)}{\Gamma \vdash \lambda x . b(x) : \Pi_{(x : A)} B(x)} \lambda$$

MLTT rule:

$$\frac{\Gamma, x : A \vdash b(x) \equiv b'(x) : B(x)}{\Gamma \vdash \lambda x . b(x) \equiv \lambda x . b'(x) : \Pi_{(x : A)} B(x)} \lambda - \text{eq}$$

2.1.3 Π elimination

MLTT rule:

$$\frac{\Gamma \vdash f : \Pi_{(x:A)}B(x)}{\Gamma, x:A \vdash f(x) : B(x)} \text{ ev}$$

MLTT rule:

$$\frac{\Gamma \vdash f \equiv f' : \Pi_{(x:A)} B(x)}{\Gamma, \ x : A \vdash f(x) \equiv f'(x) : B(x)} \text{ ev-eq}$$

2.1.4 Π computation

MLTT rule:

$$\frac{\Gamma, \ x \colon\! A \ \vdash \ b(x) \colon\! B(x)}{\Gamma, \ x \colon\! A \ \vdash \ (\lambda y . b(y))(x) \equiv b(x) \colon\! B(x)} \ \beta$$

MLTT rule:

$$\frac{\Gamma \vdash f : \Pi_{(x:A)}B(x)}{\Gamma \vdash \lambda x. f(x) \equiv f : \Pi_{(x:A)}B(x)} \eta$$

2.2 Ordinary functions

2.2.1 Basic

Defined rule:

$$\frac{\Gamma \, \vdash A \; \mathrm{type} \quad \Gamma \, \vdash B \; \mathrm{type}}{\Gamma \, \vdash A \, {\rightarrow} B \; \mathrm{type}} \, \rightarrow \,$$

Defined rule:

$$\frac{\Gamma \vdash A \text{ type} \qquad \Gamma \vdash B \text{ type}}{\Gamma \vdash A \to B \equiv \Pi_{(x:A)} B \text{ type}}$$

Derived rule:

$$\frac{\Gamma \ \vdash B \text{ type} \qquad \Gamma, \ x \colon\! A \ \vdash \ b(x) \colon\! B(x)}{\Gamma \ \vdash \ \lambda x \ldotp b(x) \colon\! A \to\! B} \ \lambda$$

Derived rule:

$$\frac{\Gamma \vdash f : A \to B}{\Gamma, x : A \vdash f(x) : B} \text{ ev}$$

Derived rule:

$$\frac{\Gamma \vdash B \text{ type} \quad \Gamma, x : A \vdash b(x) : B}{\Gamma, x : A \vdash (\lambda y . b(y))(x) \equiv b(x) : B} \beta$$

Derived rule:

$$\frac{\Gamma \vdash f : A \to B}{\Gamma \vdash \lambda x . f(x) \equiv f : A \to B} \eta$$

2.2.2 Identity

Defined rule:

$$\frac{\Gamma \vdash A \text{ type}}{\Gamma \vdash \text{id}_A : A \to A}$$

Defined rule:

$$\frac{\Gamma \vdash A \text{ type} \qquad \Gamma \vdash B \text{ type}}{\Gamma \vdash \text{id}_A \equiv \lambda x. x : A \rightarrow A}$$

2.2.3 Composition

Defined rule:

$$\frac{\Gamma \vdash A \text{ type} \quad \Gamma \vdash B \text{ type} \quad \Gamma \vdash C \text{ type}}{\Gamma \vdash \text{comp} : (B \to C) \to (A \to B) \to (A \to C)}$$

Defined rule:

$$\frac{\Gamma \vdash A \text{ type} \quad \Gamma \vdash B \text{ type} \quad \Gamma \vdash C \text{ type}}{\Gamma \vdash \text{comp} \equiv \lambda g. \lambda f. \lambda x. g(f(x)) : (B \to C) \to (A \to B) \to (A \to C)}$$

Remark:

$$g \circ f \equiv \text{comp}(g, f) \equiv \lambda x. g(f(x))$$

Derived rule:

$$\frac{\Gamma \ \vdash f : A \mathop{\rightarrow} B \qquad \Gamma \ \vdash g : B \mathop{\rightarrow} C \qquad \Gamma \ \vdash h : C \mathop{\rightarrow} D}{\Gamma \ \vdash (h \circ g) \circ f \ \equiv \ h \circ (g \circ f) : A \mathop{\rightarrow} D}$$

3 The natural numbers

3.1 Basic

Defined rule:

$$\frac{1}{1+1}$$
 N type N-form

Defined rule:

$$\overline{\vdash 0_{\mathbb{N}} : \mathbb{N}}$$

Defined rule:

$$\vdash$$
 succ_N : N

3.2 Induction principle

Metasyntactic shortcut (only for this document):

$$\widehat{P}$$
 := Γ , $n : \mathbb{N} \vdash P(n)$ type

Defined rule:

$$\frac{\textcircled{P} \qquad \Gamma \, \vdash \, p_0 : P(0_{\mathbb{N}}) \qquad \Gamma \, \vdash \, p_S : \Pi_{(n:\mathbb{N})} \left(P(n) \to P(\operatorname{succ}_{\mathbb{N}}(n)) \right)}{\Gamma \, \vdash \, \operatorname{ind}_{\mathbb{N}}(p_0, p_S) : \Pi_{(n:\mathbb{N})} P(n)} \, \, \mathbb{N} \text{-ind}$$

Defined rule:

$$\frac{ \textcircled{P} \qquad \Gamma \vdash p_0 : P(0_{\mathbb{N}}) \qquad \Gamma \vdash p_S : \Pi_{(n:\mathbb{N})} \left(P(n) \rightarrow P(\operatorname{succ}_{\mathbb{N}}(n)) \right) }{ \Gamma \vdash \operatorname{ind}_{\mathbb{N}}(p_0, p_S, 0_{\mathbb{N}}) \equiv p_0 : P(0_{\mathbb{N}}) }$$

Defined rule:

$$\frac{\bigcirc \quad \Gamma \vdash p_0 : P(0_{\mathbb{N}}) \quad \Gamma \vdash p_S : \Pi_{(n:\mathbb{N})} \left(P(n) \to P(\operatorname{succ}_{\mathbb{N}}(n)) \right)}{\Gamma, \ n : \mathbb{N} \vdash \operatorname{ind}_{\mathbb{N}}(p_0, p_S, \operatorname{succ}_{\mathbb{N}}(n)) \equiv p_S(n, \operatorname{ind}_{\mathbb{N}}(p_0, p_S, n)) : P(\operatorname{succ}_{\mathbb{N}}(n))}$$

3.3 Addition

Defined rule:

$$\overline{m: \mathbb{N} \vdash \operatorname{add-zero}_{\mathbb{N}}(m): \mathbb{N}}$$

Defined rule:

$$\overline{m: \mathbb{N} \vdash \text{add-zero}_{\mathbb{N}}(m) \equiv m: \mathbb{N}}$$

Defined rule:

$$\overline{m: \mathbb{N} \vdash \text{add-succ}_{\mathbb{N}}(m): \mathbb{N} \to \mathbb{N} \to \mathbb{N}}$$

Defined rule:

$$\overline{m: \mathbb{N} \vdash \text{add-succ}_{\mathbb{N}}(m) \equiv \lambda n. \text{succ}_{\mathbb{N}} : \mathbb{N} \to \mathbb{N} \to \mathbb{N}}$$

Defined rule:

$$\overline{m: \mathbb{N} \vdash \operatorname{add}_{\mathbb{N}}(m): \mathbb{N} \to \mathbb{N}}$$

Defined rule:

$$m: \mathbb{N} \vdash \operatorname{add}_{\mathbb{N}}(m) \equiv \operatorname{ind}_{\mathbb{N}}(\operatorname{add-zero}_{\mathbb{N}}(m), \operatorname{add-succ}_{\mathbb{N}}(m)) : \mathbb{N} \to \mathbb{N}$$

Remark:

$$m + 0_{\mathbb{N}} \equiv \operatorname{add}_{\mathbb{N}}(m, 0_{\mathbb{N}}) \equiv m$$

Remark:

$$m + \operatorname{succ}_{\mathbb{N}}(n) \equiv \operatorname{add}_{\mathbb{N}}(m, \operatorname{succ}_{\mathbb{N}}(n)) \equiv \operatorname{succ}_{\mathbb{N}}(m+n)$$