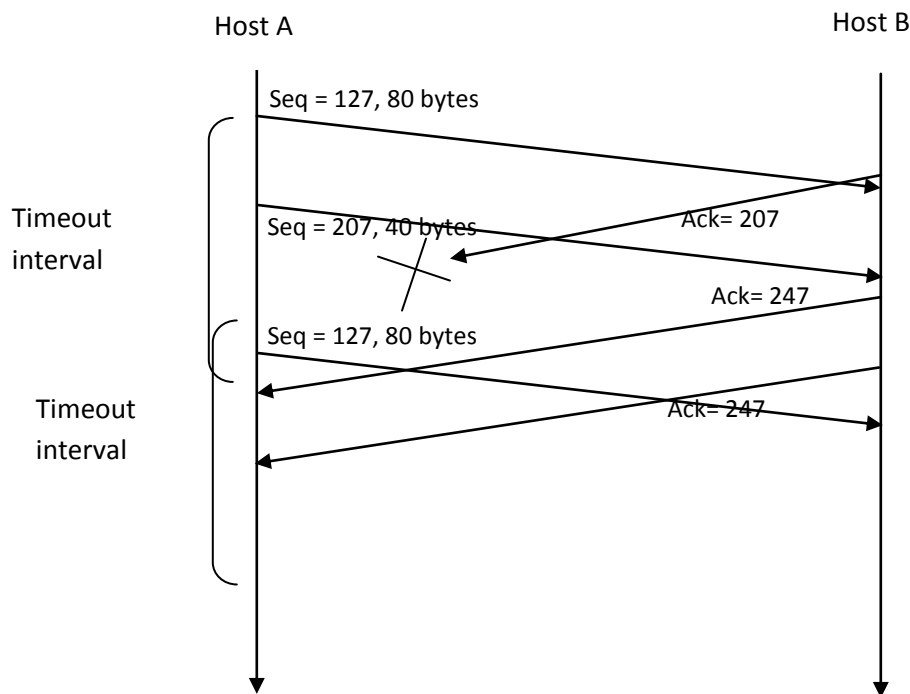


Solutions to Homework #5

Problem 27

- a) In the second segment from Host A to B, the sequence number is 207, source port number is 302 and destination port number is 80.
- b) If the first segment arrives before the second, in the acknowledgement of the first arriving segment, the acknowledgement number is 207, the source port number is 80 and the destination port number is 302.
- c) If the second segment arrives before the first segment, in the acknowledgement of the first arriving segment, the acknowledgement number is 127, indicating that it is still waiting for bytes 127 and onwards.
- d)



Problem 28

Since the link capacity is only 100 Mbps, so host A's sending rate can be at most 100Mbps. Still, host A sends data into the receive buffer faster than Host B can remove data from the buffer. The receive buffer fills up at a rate of roughly 40Mbps. When the buffer is full, Host B signals to Host A to stop sending data by setting $RcvWindow = 0$. Host A then stops sending until it receives a TCP segment with $RcvWindow > 0$. Host A will thus repeatedly stop and start sending as a function of the $RcvWindow$ values it receives from Host B. On average, the long-term rate at which Host A sends data to Host B as part of this connection is no more than 60Mbps.

Problem 40

- a) TCP slowstart is operating in the intervals [1,6] and [23,26]
- b) TCP congestion avoidance is operating in the intervals [6,16] and [17,22]
- c) After the 16th transmission round, packet loss is recognized by a triple duplicate ACK. If there was a timeout, the congestion window size would have dropped to 1.
- d) After the 22nd transmission round, segment loss is detected due to timeout, and hence the congestion window size is set to 1.
- e) The threshold is initially 32, since it is at this window size that slow start stops and congestion avoidance begins.
- f) The threshold is set to half the value of the congestion window when packet loss is detected. When loss is detected during transmission round 16, the congestion window size is 42. Hence the threshold is 21 during the 18th transmission round.
- g) The threshold is set to half the value of the congestion window when packet loss is detected. When loss is detected during transmission round 22, the congestion window size is 29. Hence the threshold is 14 (taking lower floor of 14.5) during the 24th transmission round.
- h) During the 1st transmission round, packet 1 is sent; packet 2-3 are sent in the 2nd transmission round; packets 4-7 are sent in the 3rd transmission round; packets 8-15 are sent in the 4th transmission round; packets 16-31 are sent in the 5th transmission round; packets 32-63 are sent in the 6th transmission round; packets 64 – 96 are sent in the 7th transmission round. Thus packet 70 is sent in the 7th transmission round.
- i) The threshold will be set to half the current value of the congestion window (8) when the loss occurred and congestion window will be set to the new threshold value + 3 MSS . Thus the new values of the threshold and window will be 4 and 7 respectively.
- j) threshold is 21, and congestion window size is 1.
round 17, 1 packet; round 18, 2 packets; round 19, 4 packets; round 20, 8 packets; round 21, 16 packets; round 22, 21 packets. So, the total number is 52

Problem 44

- a) It takes 1 RTT to increase CongWin to 6 MSS; 2 RTTs to increase to 7 MSS; 3 RTTs to increase to 8 MSS; 4 RTTs to increase to 9 MSS; 5 RTTs to increase to 10 MSS; 6 RTTs to increase to 11 MSS; and 7 RTTs to increase to 12MSS.
- b) In the first RTT 5 MSS was sent; in the second RTT 6 MSS was sent; in the third RTT 7 MSS was sent; in the fourth RTT 8 MSS was sent; in the fifth RTT, 9 MSS was sent; and in the sixth RTT, 10 MSS was sent. Thus, up to time 6 RTT, $5+6+7+8+9+10 = 45$ MSS were sent (and acknowledged). Thus, we can say that the average throughput up to time 6 RTT was $(45 \text{ MSS})/(6 \text{ RTT}) = 7.5 \text{ MSS/RTT}$.

Problem 46

- a) Let W denote the max window size measured in segments. Then, $W \cdot \text{MSS} / \text{RTT} = 10 \text{Mbps}$, as packets will be dropped if the maximum sending rate exceeds link capacity. Thus, we have $W \cdot 1500 \cdot 8 / 0.15 = 10 \cdot 10^6$, then W is about 125 segments.
- b) As congestion window size varies from $W/2$ to W , then the average window size is $0.75W = 94$ (ceiling of 93.75) segments. Average throughput is $94 \cdot 1500 \cdot 8 / 0.15 = 7.52 \text{Mbps}$.
- c) $94/2 \cdot 0.15 = 7.05$ seconds, as the number of RTTs (that this TCP connections needs in order to increase its window size from $W/2$ to W) is given by $W/2$. Recall the window size increases by one in each RTT.

Problem 48

- a) Let W denote the max window size. Then, $W \cdot \text{MSS} / \text{RTT} = 10 \text{Gbps}$, as packets will be dropped if maximum sending rate reaches link capacity. Thus, we have $W \cdot 1500 \cdot 8 / 0.15 = 10 \cdot 10^9$, then $W = 125000$ segments.
- b) As congestion window size varies from $W/2$ to W , then the average window size is $0.75W = 93750$ segments. Average throughput is $93750 \cdot 1500 \cdot 8 / 0.1 = 7.5 \text{Gbps}$.
- c) $93750/2 \cdot 0.15 / 60 = 117$ minutes. In order to speed up the window increase process, we can increase the window size by a much larger value, instead of increasing window size only by one in each RTT. Some protocols are proposed to solve this problem, such as ScalableTCP or HighSpeed TCP.