

NATURAL TRANSFORMATIONS (continued)

Remark:

Is this category $\mathcal{C} \cdot \rightarrow \cdot \rightarrow \cdot \rightarrow \cdot$ equivalent to \mathcal{C}^{op} $\cdot \leftarrow \cdot \leftarrow \cdot \leftarrow \cdot$? Yes! But: given $\mathcal{D} \cdot \rightarrow \cdot \rightarrow \cdot \rightarrow \cdot \dots$ and $\mathcal{D}^{op} \cdot \leftarrow \cdot \leftarrow \cdot \leftarrow \cdot \dots$ (∞ objects) they are NOT equivalent, since \mathcal{D} has an initial object and \mathcal{D}^{op} has no initial objects.

The numerical category is equivalent to its opposite (map n to n and $A \in \mathbb{R}^{m \times n}$ to $A^T \in \mathbb{R}^{n \times m}$), but Set is NOT equivalent to Set^{op} , and $\text{Vect}(\mathbb{R})$ is NOT equivalent to $\text{Vect}(\mathbb{R})^{op}$. For example, in Set every morphism into the initial object \emptyset is an isomorphism (the only map into \emptyset is its identity morphism $\text{id}_{\emptyset}: \emptyset \rightarrow \emptyset$) but in Set^{op} that is not the case: in Set^{op} , every morphism out of the terminal object $\{x\}$ is not necessarily an isomorphism. Set^{op} is equivalent to the completely atomized algebras

examples:

1) length:

$$\left. \begin{array}{l} \text{List}(X) \rightarrow \mathbb{N} \\ [x_1, \dots, x_n] \mapsto n \end{array} \right\} \text{ is a natural transformation } \eta: \text{List} \Rightarrow K_{\mathbb{N}}$$

where $K_{\mathbb{N}}$ is the constant functor:

$$\begin{aligned} K_{\mathbb{N}}: \text{Set} &\rightarrow \text{Set} \\ X &\mapsto \mathbb{N} \\ f &\mapsto \text{id}_{\mathbb{N}} \end{aligned}$$

2) $\det(\cdot)$ is a natural transformation:

$$\eta_{\mathbb{F}}: \mathbb{F}^{n \times n} \rightarrow \mathbb{F} \quad \Rightarrow \quad \eta: H \Rightarrow L \quad \text{where } H, L \text{ are as follows:}$$

$$A \mapsto \det(A)$$

$$\begin{aligned} H: \text{Field} &\rightarrow \text{Mon}, & L: \text{Field} &\rightarrow \text{Mon} \\ \mathbb{F} &\mapsto \mathbb{F}^{n \times n}, & \mathbb{F} &\mapsto (\mathbb{F}, 1, \cdot) \end{aligned}$$

Indeed, given $\varphi: \mathbb{F} \rightarrow \mathbb{G}$ field homomorphism, we have:

$$\det \circ (H(\varphi)) = \det \circ \varphi = \varphi \circ \det$$

3) $\left. \begin{array}{l} X \longrightarrow X \times X \\ x \longmapsto (x, x) \end{array} \right\}$ is a natural transformation $\eta: \text{Id}_{\text{Set}} \Rightarrow \text{Pair}$
 where:

$$\begin{array}{l} \text{Pair}: \text{Set} \longrightarrow \text{Set} \quad , \quad f \times f: (x, y) \longmapsto (f(x), f(y)) \\ X \longmapsto X \times X \\ f \longmapsto f \times f \end{array}$$

4) $\eta: \text{Id}_{\text{Set}} \Rightarrow \mathcal{P}$ is a natural transformation:

$$\begin{array}{l} \eta_x: X \longrightarrow \mathcal{P}(X) \\ x \longmapsto \{x\}, \phi, X \end{array}$$

even if η_x is as follows:

$$\begin{array}{l} \eta_x: X \longrightarrow \mathcal{P}(X) \\ x \longmapsto \begin{cases} \{x\} & \text{CH} \\ X & \neg \text{CH} \end{cases} \quad \text{CH} := \begin{array}{l} \text{Continuum} \\ \text{Hypothesis} \end{array} \end{array}$$

η is still a natural transformation.

Instead, if η_x is as follows:

$$\begin{array}{l} \eta_x: X \longrightarrow \mathcal{P}(X) \\ x \longmapsto \begin{cases} \{x\} & X = \mathbb{N} \\ X & \text{otherwise} \end{cases} \end{array}$$

then η is not a natural transformation anymore
