Sorting Algorithms

Merge Sort



Overview

Divide and Conquer

Merge Sort



Divide and Conquer



- 1. Base Case, solve the problem directly if it is small enough
- 2. Divide the problem into two or more similar and smaller subproblems
- 3. Recursively solve the subproblems
- 4. Combine solutions to the subproblems

Divide and Conquer - Sort



Problem:

Input: A[left..right] – unsorted array of integers

 Output: A[left..right] – sorted in non-decreasing order

Divide and Conquer - Sort

- 1. Base case
 - at most one element (left ≥ right), return
- **2.** Divide A into two subarrays: FirstPart, SecondPart Two Subproblems:

sort the FirstPart sort the SecondPart

3. Recursively

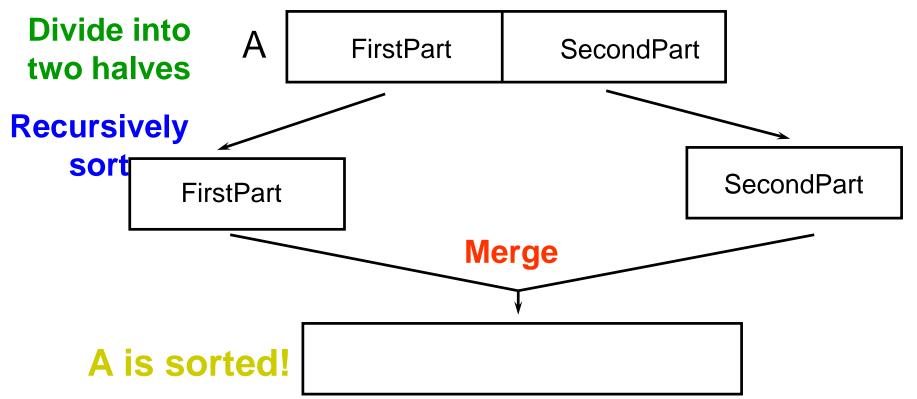
sort FirstPart sort SecondPart

4. Combine sorted FirstPart and sorted SecondPart



Merge Sort: Idea



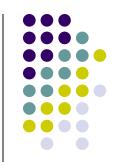


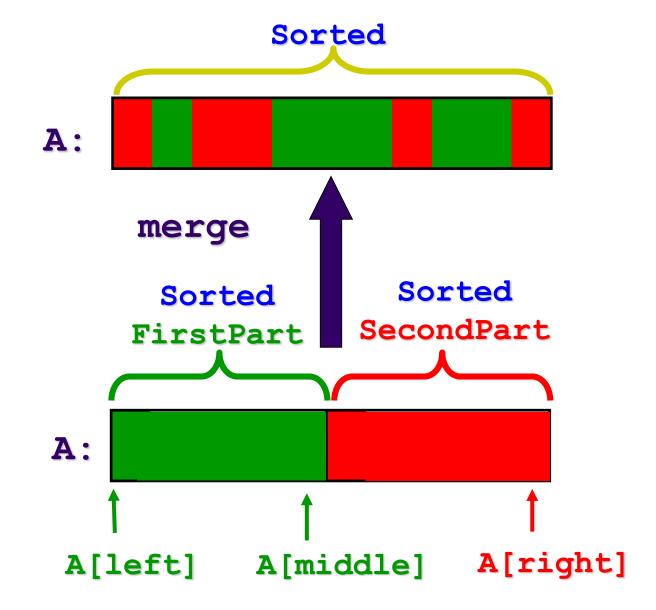
Merge Sort: Algorithm

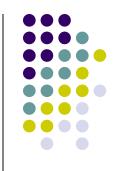


```
Merge-Sort (A, left, right)
      left ≥ right return
 else
      middle \leftarrow b(left+right)/2
                                             Recursive Call
      Merge-Sort(A, left, middle)
      Merge-Sort(A, middle+1, right)
      Merge(A, left, middle, right)
```

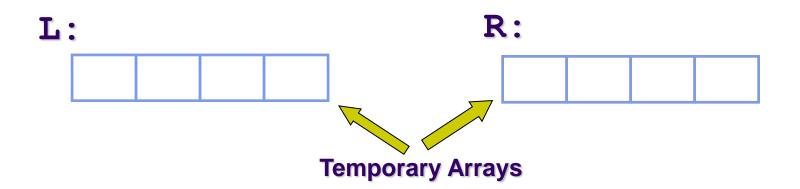
Merge-Sort: Merge

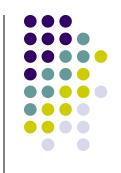


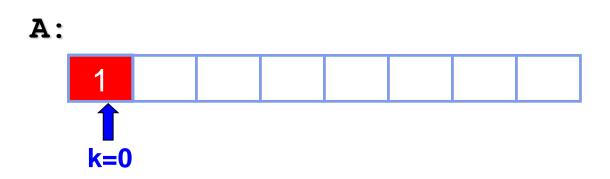


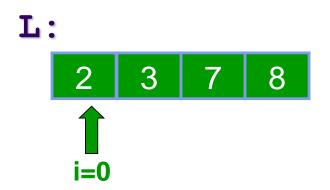


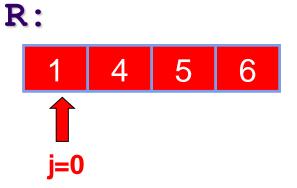


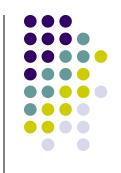












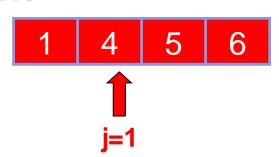


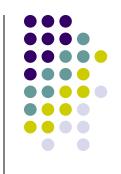


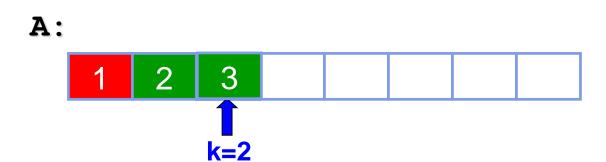
L:



R:



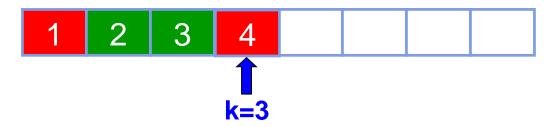




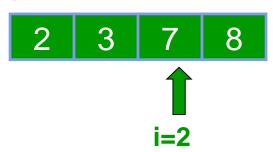




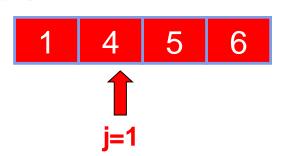


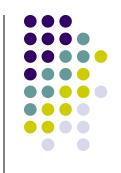


L:

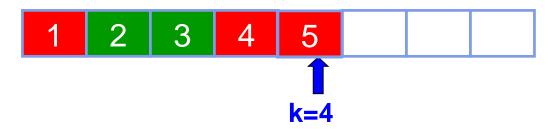


R:





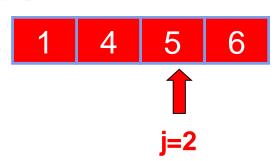


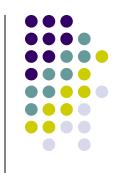


L: 2 3 7



R:

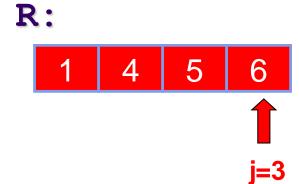


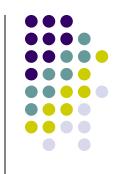






L: 2 3 7 8 i=2





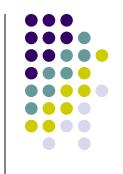




L: 2 3 7 8 i=2







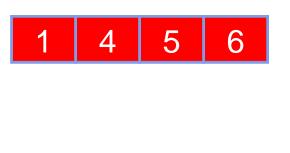




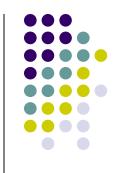
i=3

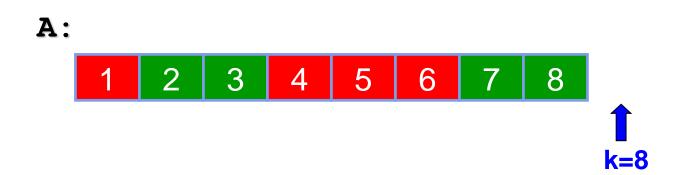
L: 2 3 7 8 1

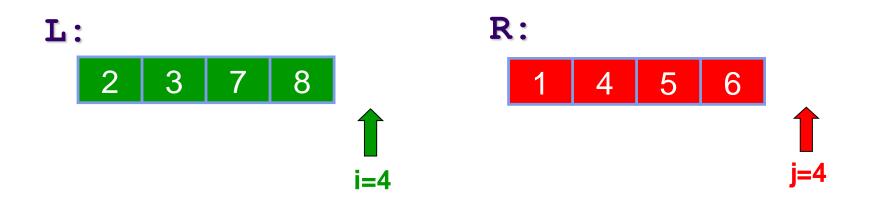




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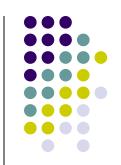


```
Merge(A, left, middle, right)
    n_1 \leftarrow middle - left + 1
2. n_2 \leftarrow right - middle
    create array L[n_1], R[n_2]
3.
    for i \leftarrow 0 to n_1-1 do L[i] \leftarrow A[left +i]
4.
    for j \leftarrow 0 to n_2-1 do R[j] \leftarrow A[middle+j]
5.
6. \mathbf{k} \leftarrow \mathbf{i} \leftarrow \mathbf{j} \leftarrow \mathbf{0}
    while i < n_1 \& j < n_2
7.
           if L[i] < R[j]
8.
                A[k++] \leftarrow L[i++]
9.
    else
10.
                 A[k++] \leftarrow R[j++]
11.
12.
    while i < n_1
     A[k++] \leftarrow L[i++]
13.
14. while j < n_2
15. A[k++] \leftarrow R[j++]
```

 $n = n_1 + n_2$

Space: n

Time: cn for some constant c

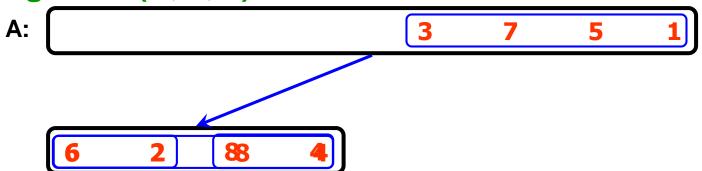


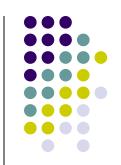
Divide

A: 6 2 8 4 3 3 7 7 5 5 1 1

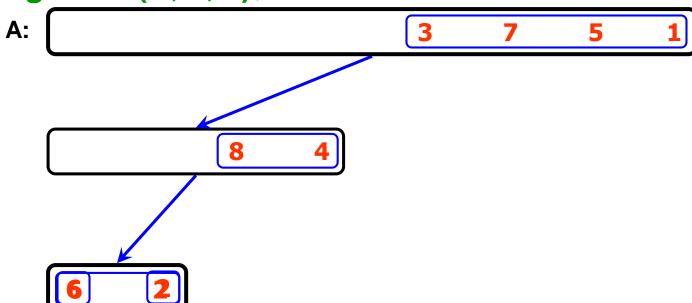


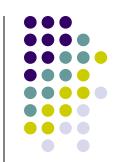
Merge-Sort(A, 0, 3), divide



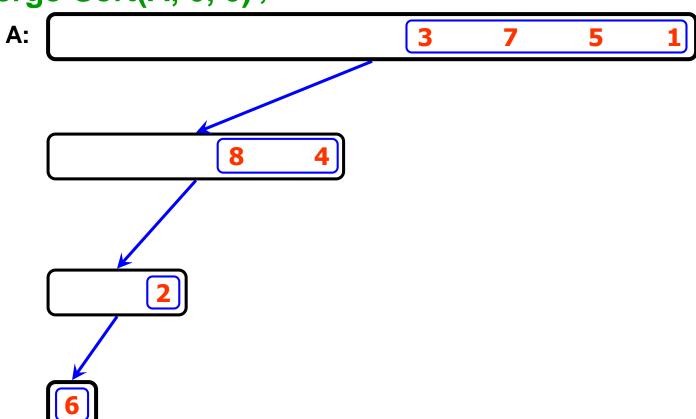


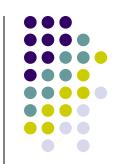
Merge-Sort(A, 0, 1), divide



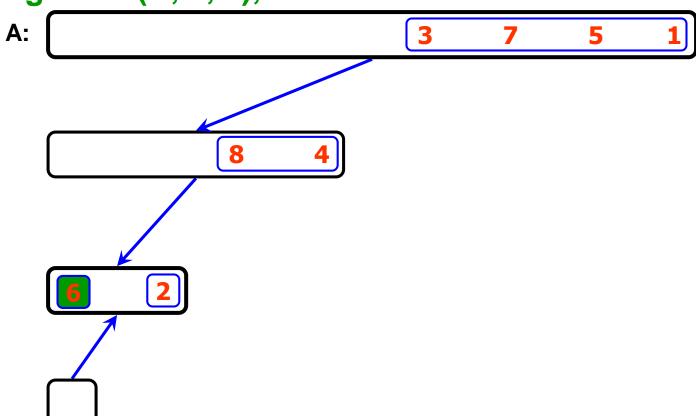


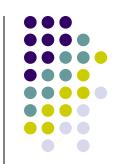
Merge-Sort(A, 0, 0), base case



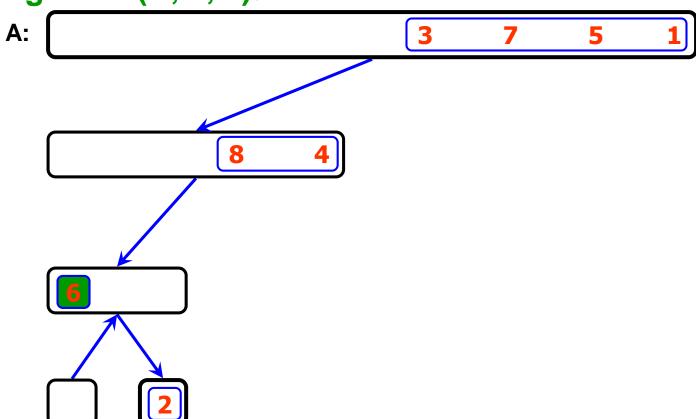


Merge-Sort(A, 0, 0), return



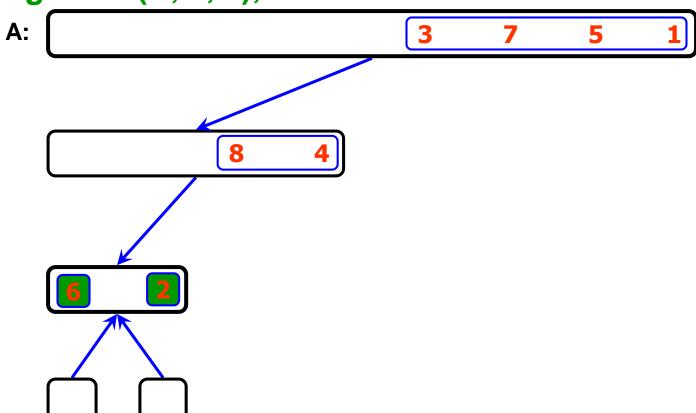


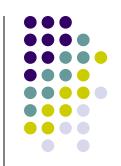
Merge-Sort(A, 1, 1), base case



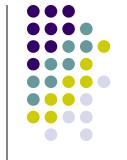


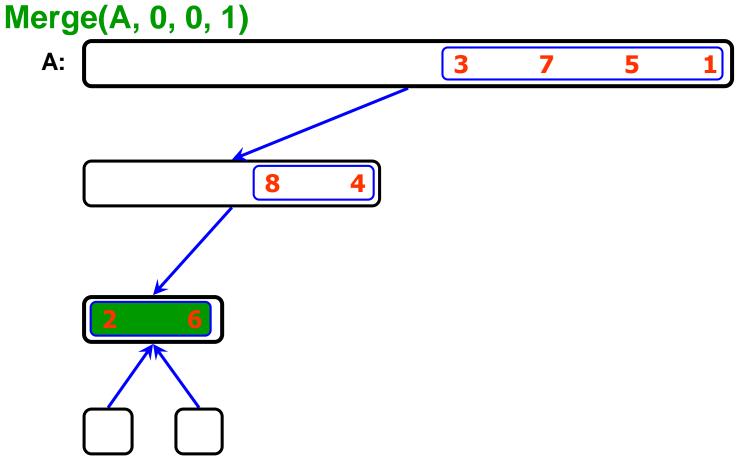
Merge-Sort(A, 1, 1), return



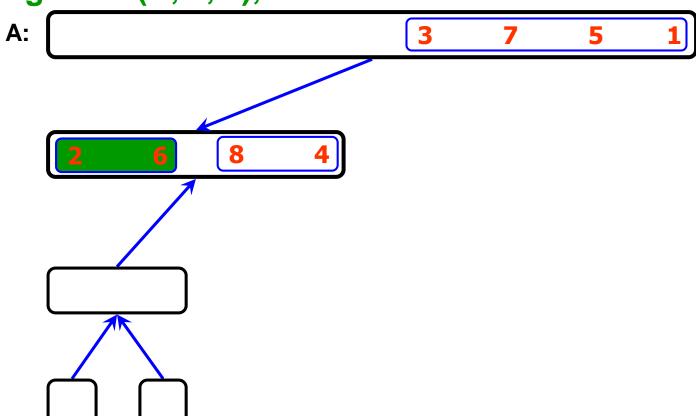


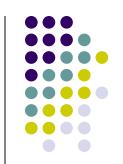




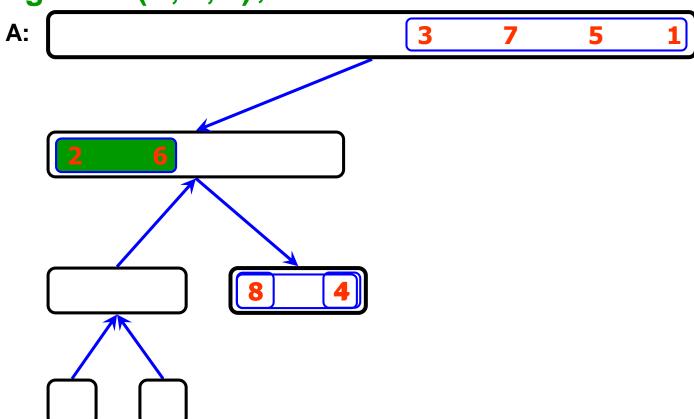


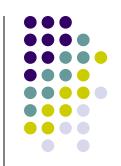
Merge-Sort(A, 0, 1), return



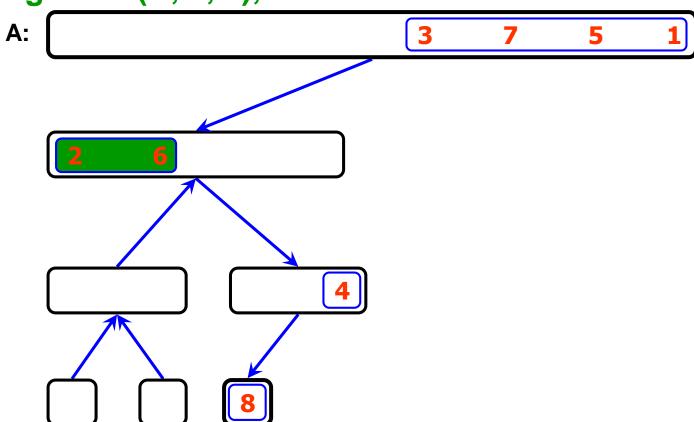


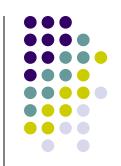
Merge-Sort(A, 2, 3), divide



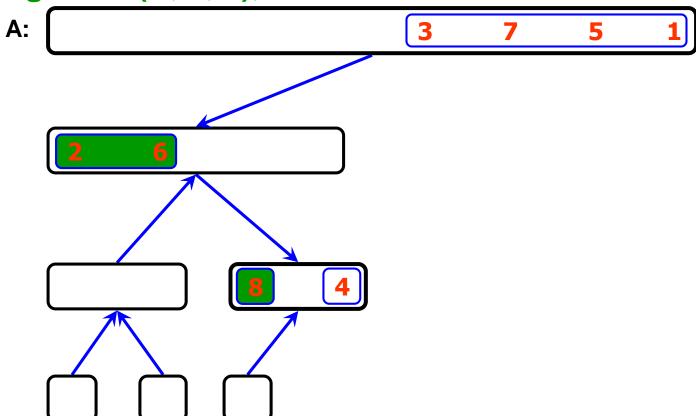


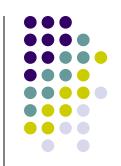
Merge-Sort(A, 2, 2), base case





Merge-Sort(A, 2, 2), return



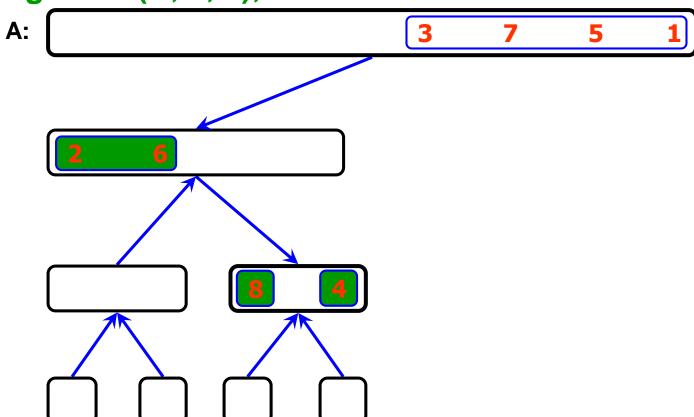


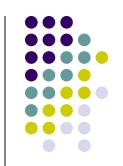
Merge-Sort(A, 3, 3), base case

A: 2 6

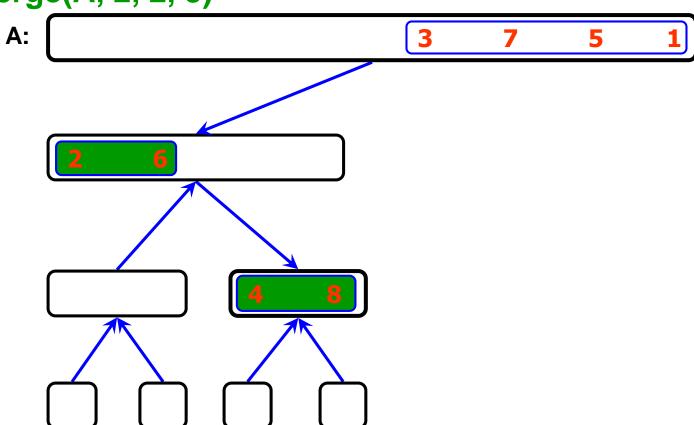


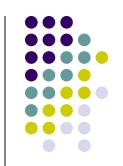
Merge-Sort(A, 3, 3), return



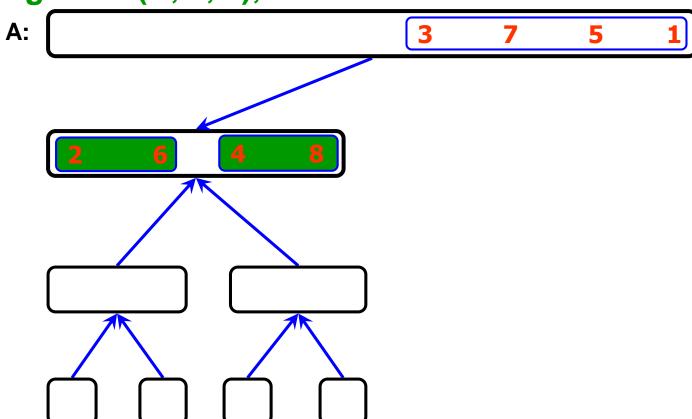


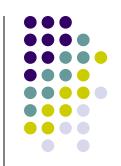
Merge(A, 2, 2, 3)



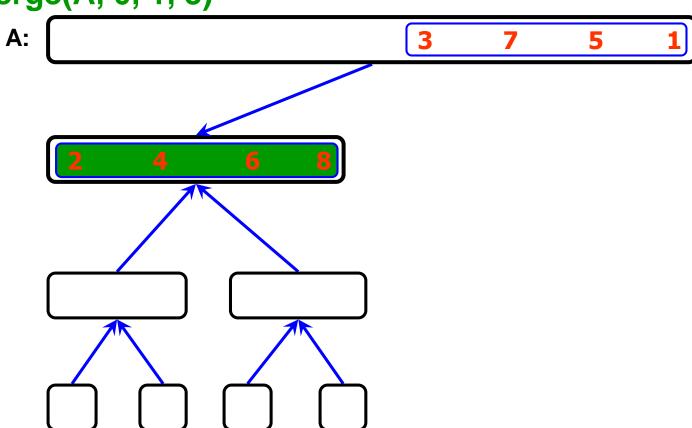


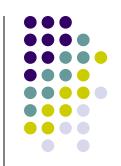
Merge-Sort(A, 2, 3), return



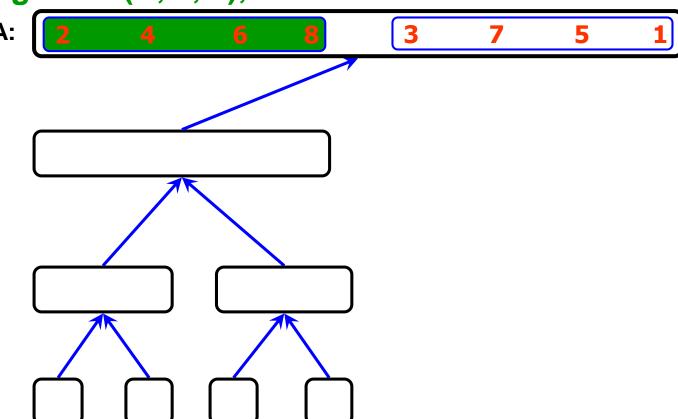


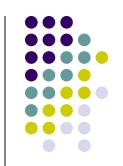
Merge(A, 0, 1, 3)



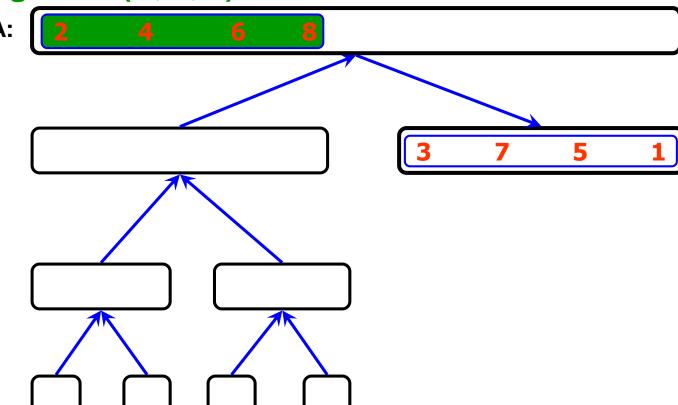


Merge-Sort(A, 0, 3), return





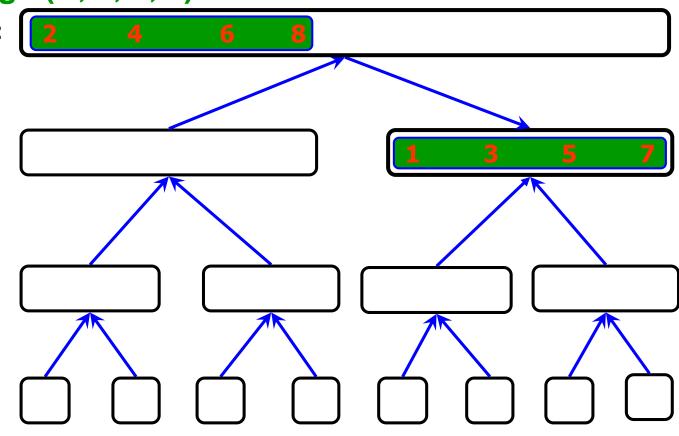
Merge-Sort(A, 4, 7)



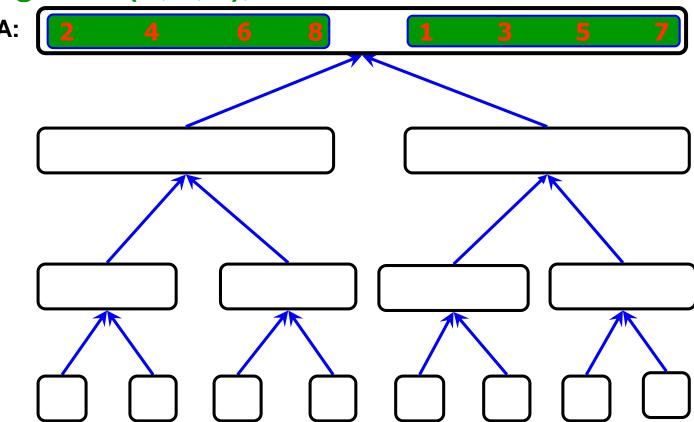


Merge (A, 4, 5, 7)



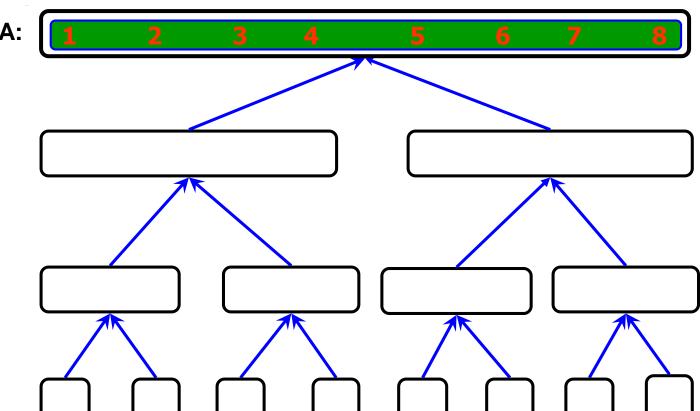


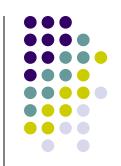
Merge-Sort(A, 4, 7), return





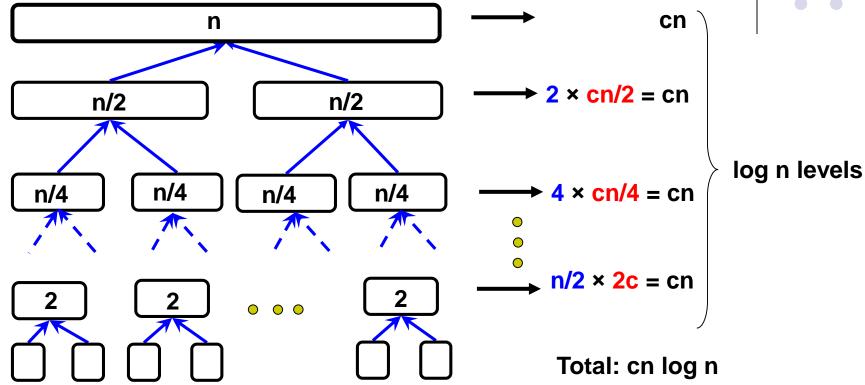
Merge-Sort(A, 0, 7), done!





Merge-Sort Analysis





- Total running time: Θ(nlogn)
- Total Space: Θ (n)

Merge-Sort Summary

Approach: divide and conquer

Time

- Most of the work is in the merging
- Total time: Θ(n log n)

Space:

• $\Theta(n)$, more space than other sorts.

