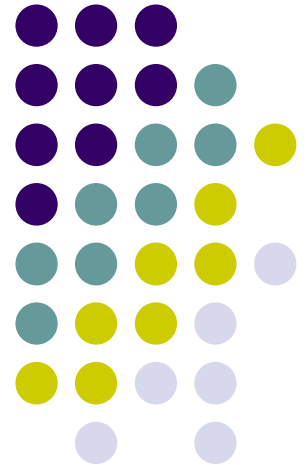


Sorting Algorithms

Quick Sort



Quick Sort



- **Divide:**
 - Pick any element **p** as the **pivot**, e.g, the first element
 - Partition the remaining elements into
 - FirstPart**, which contains all elements $< p$
 - SecondPart**, which contains all elements $\geq p$
- **Recursively sort** the **FirstPart** and **SecondPart**
- **Combine:** no work is necessary since sorting is done in place

Quick Sort



A:

pivot



Partition

FirstPart

SecondPart



Recursive call

Sorted

Sorted

FirstPart

SecondPart



Sorted

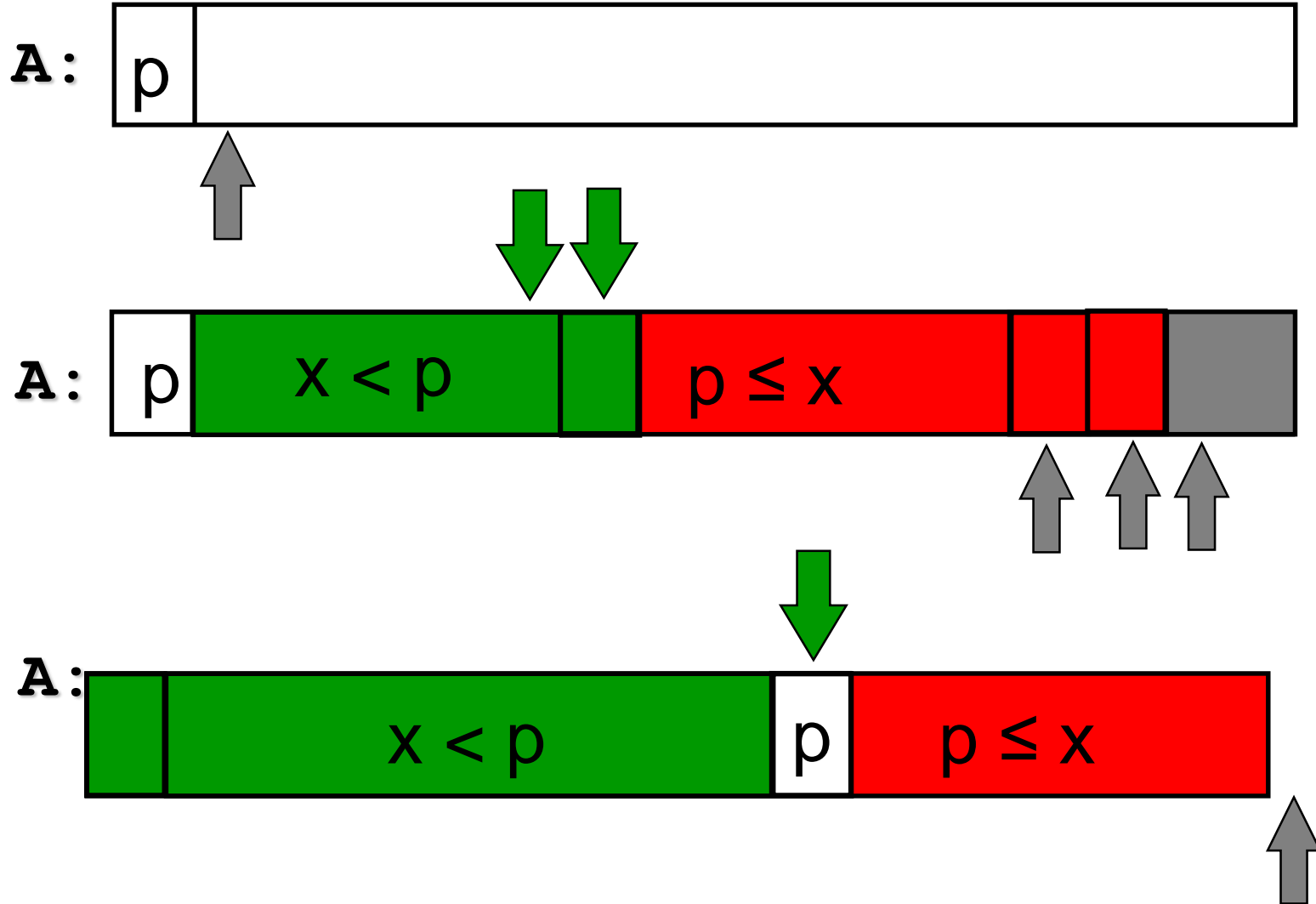
Quick Sort



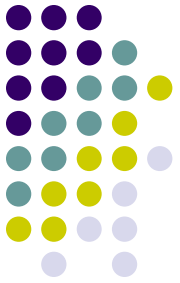
```
Quick-Sort(A, left, right)
  if    left  $\geq$  right  return
  else
    middle  $\leftarrow$  Partition(A, left, right)
    Quick-Sort(A, left, middle-1 )
    Quick-Sort(A, middle+1, right)
  end if
```



Partition

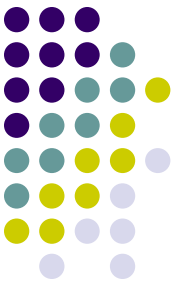


Partition Example

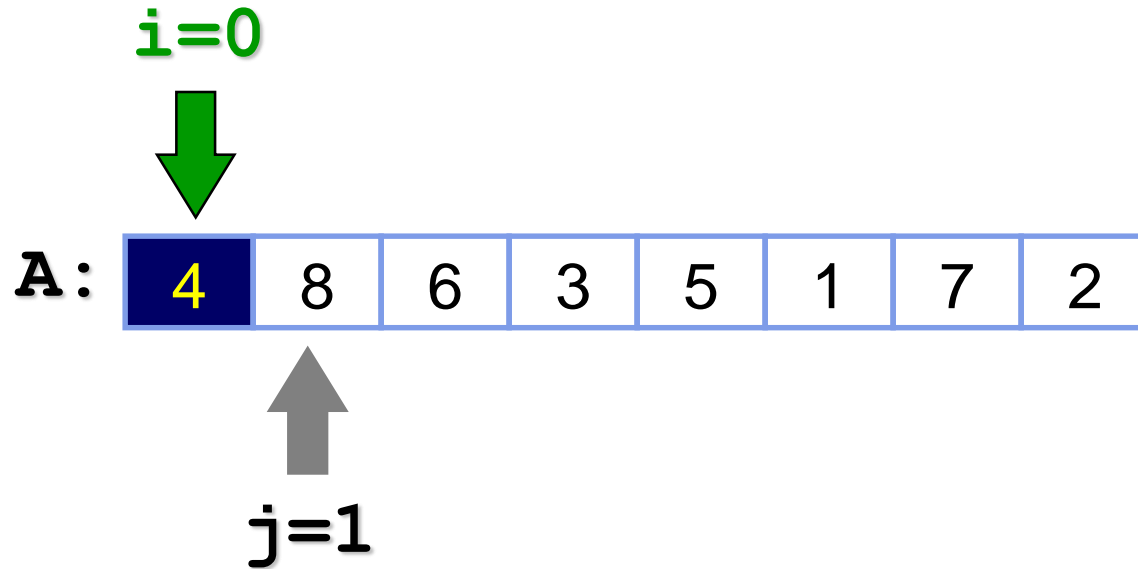


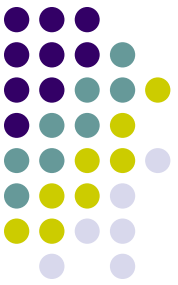
A:

4	8	6	3	5	1	7	2
---	---	---	---	---	---	---	---

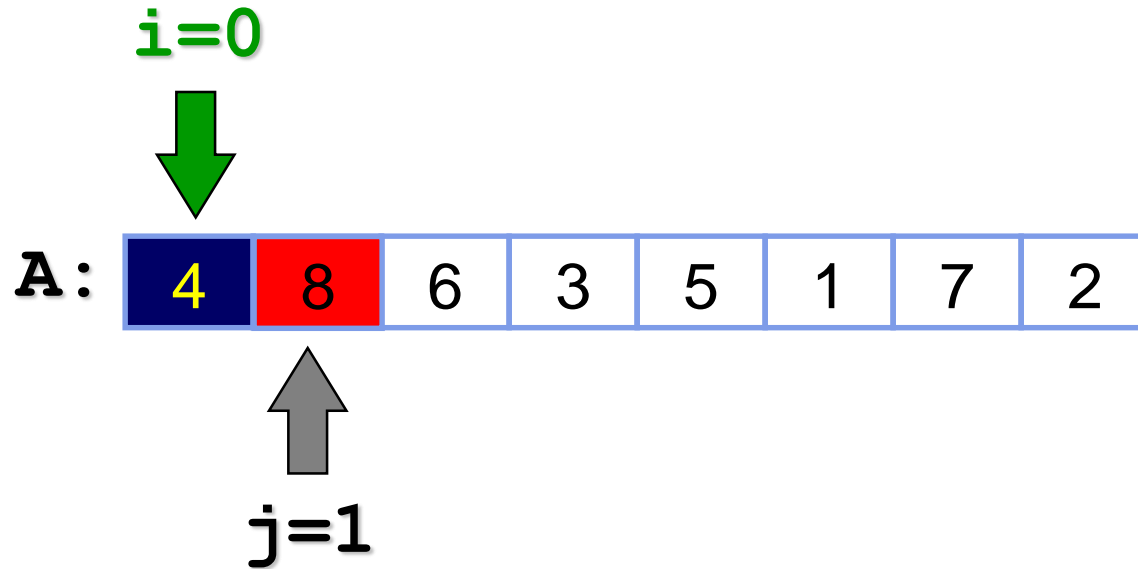


Partition Example



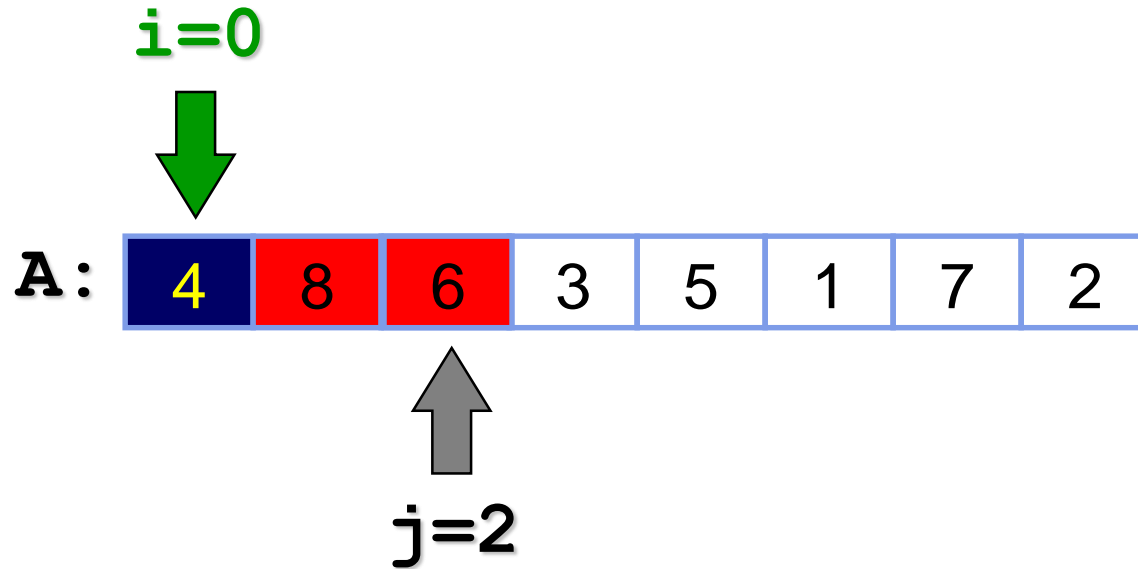


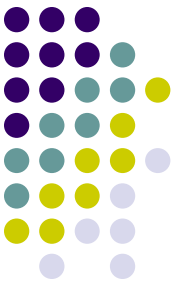
Partition Example



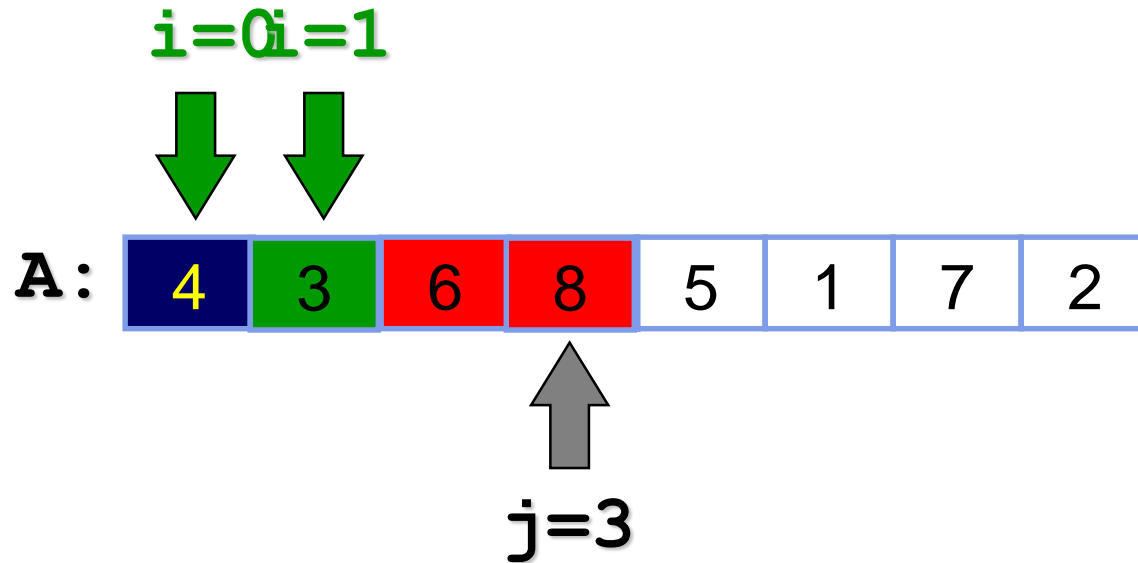


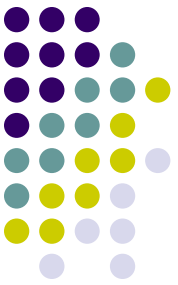
Partition Example



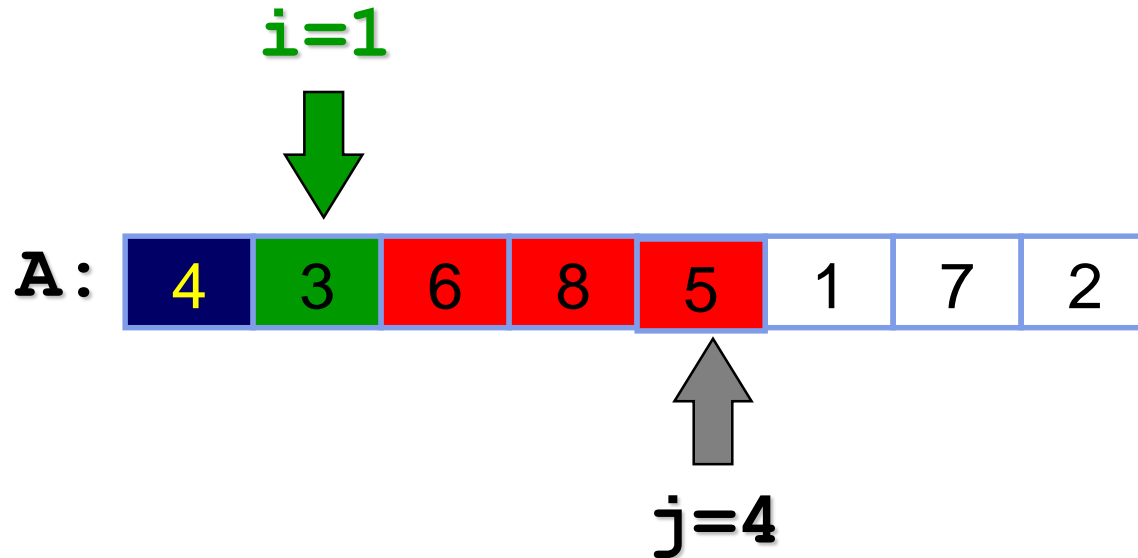


Partition Example



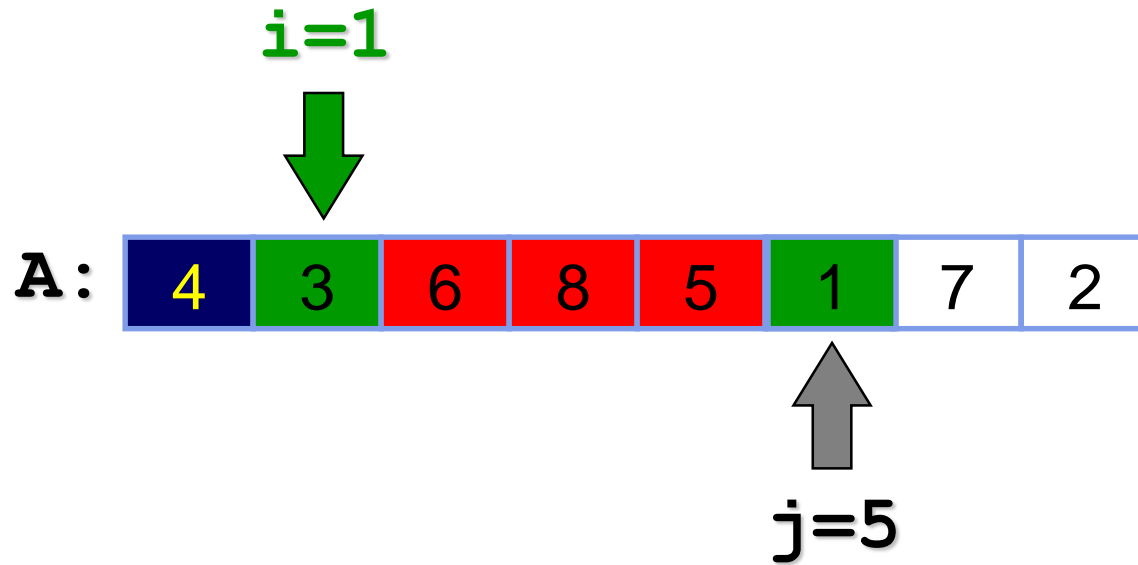


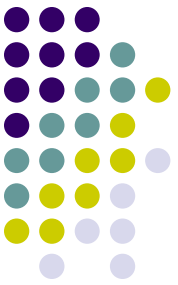
Partition Example



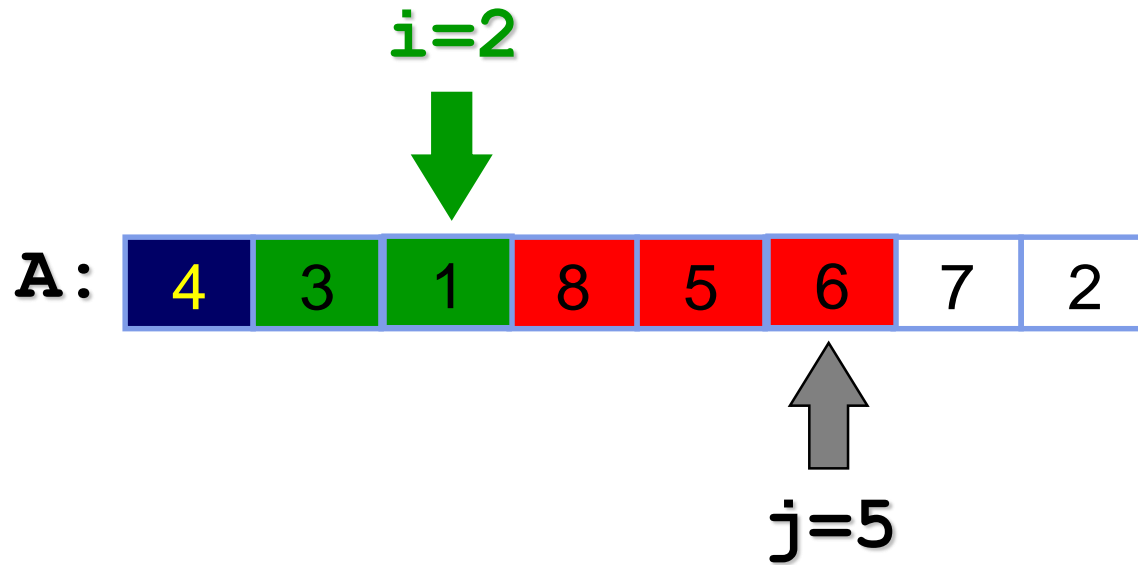


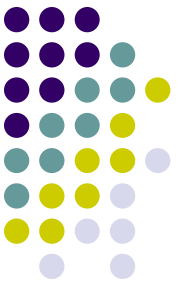
Partition Example



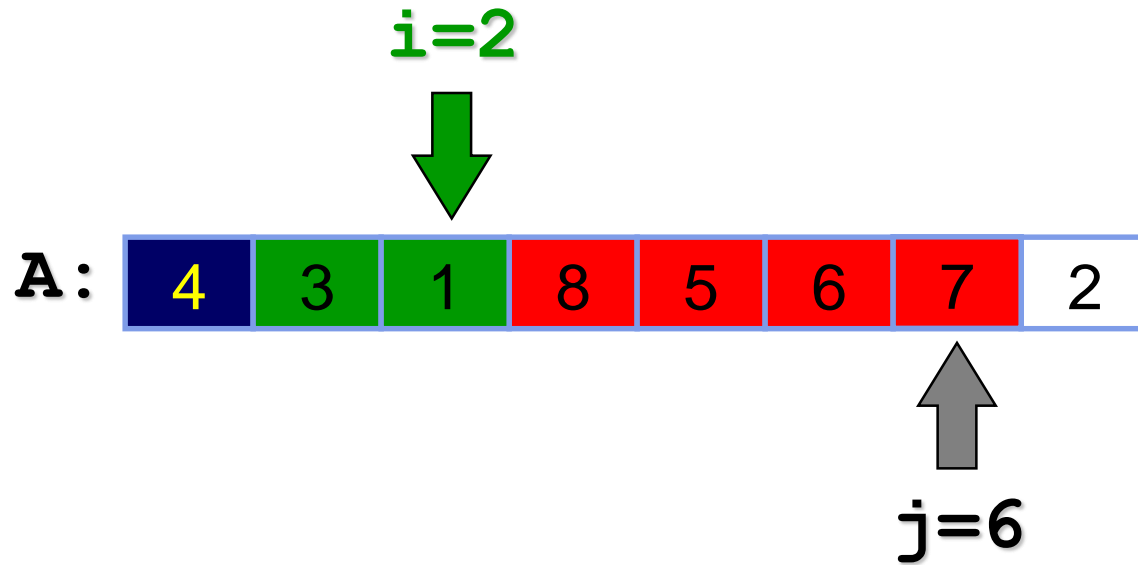


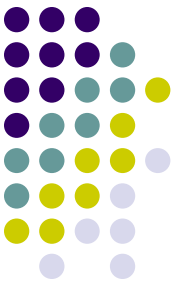
Partition Example



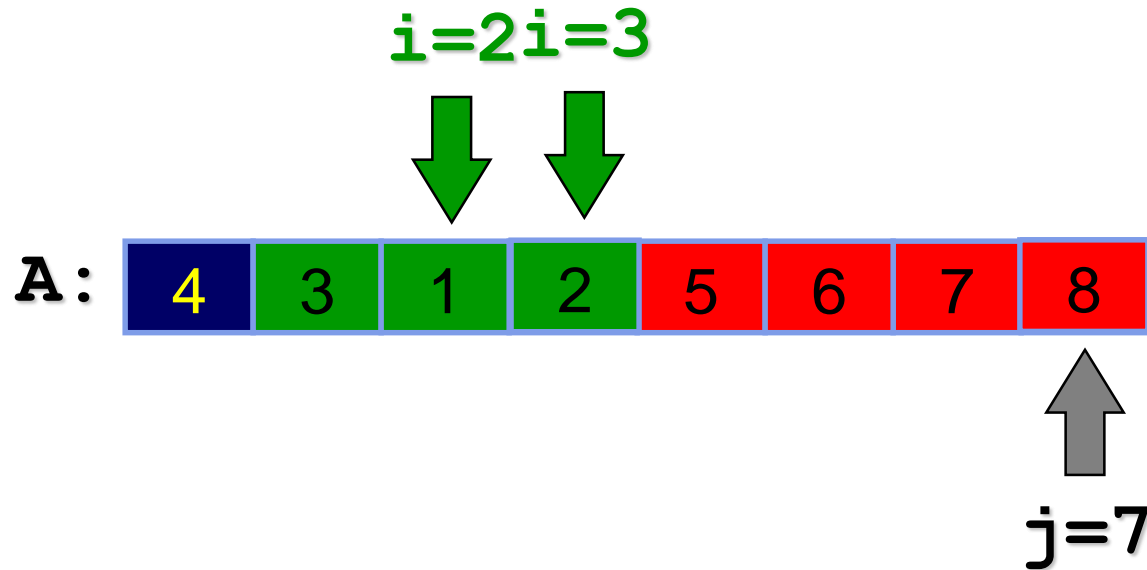


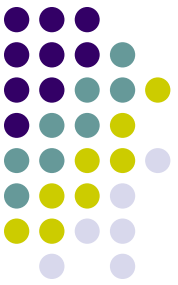
Partition Example



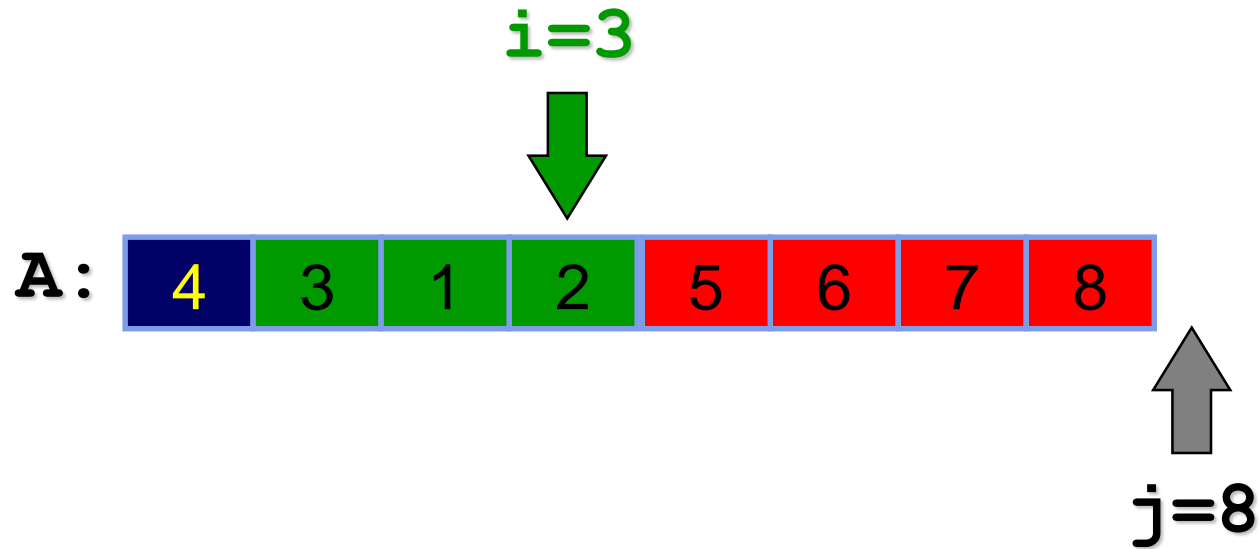


Partition Example



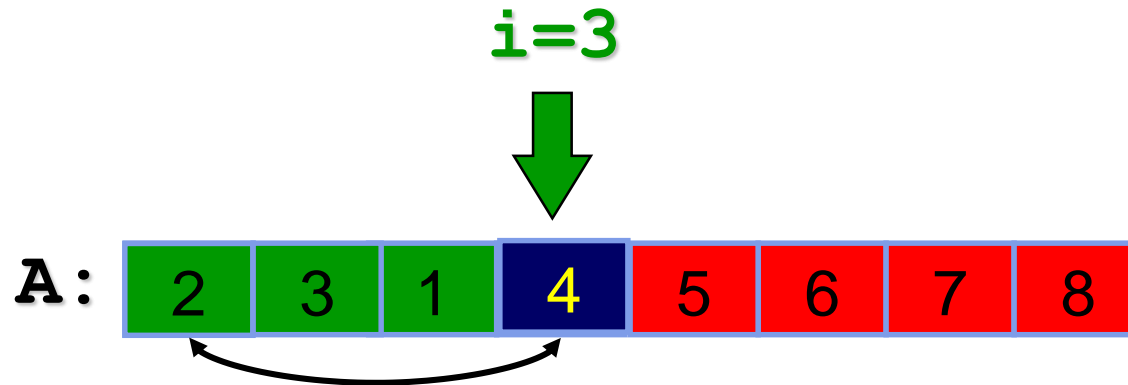


Partition Example



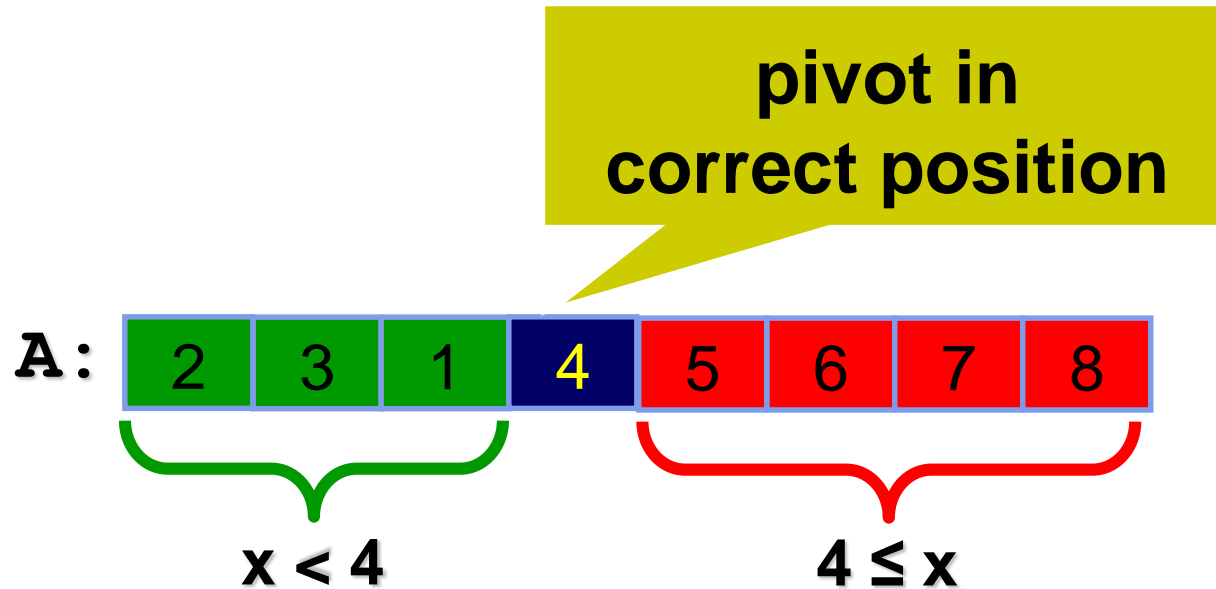


Partition Example





Partition Example





Partition(A, left, right)

```
1.  x ← A[left]
2.  i ← left
3.  for j ← left+1 to right
4.      if A[j] < x then
5.          i ← i + 1
6.          swap(A[i], A[j])
7.      end if
8.  end for j
9.  swap(A[i], A[left])
10. return i
```

n = right – left + 1

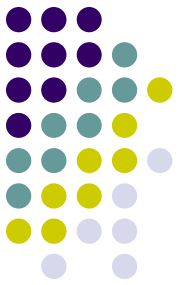
Time: cn for some constant c

Space: constant

Quick-Sort(A, 0, 7)

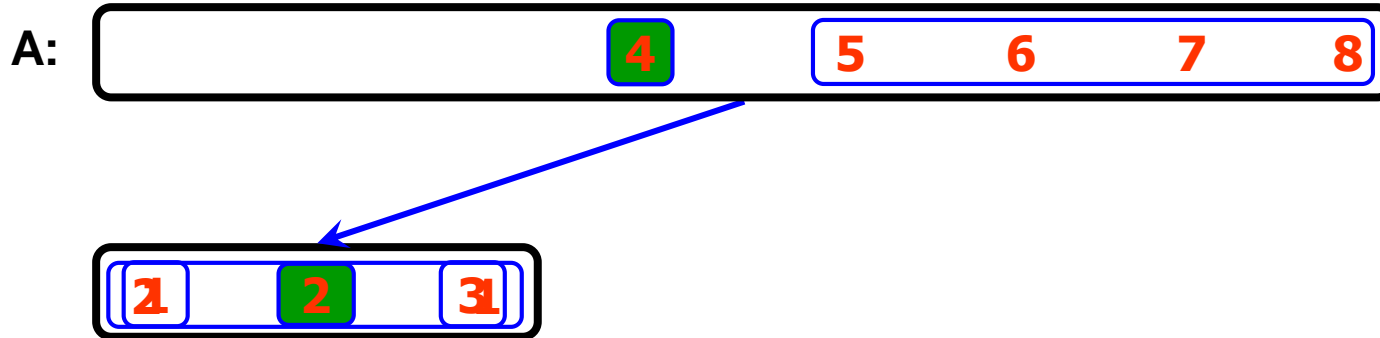
Partition

A:



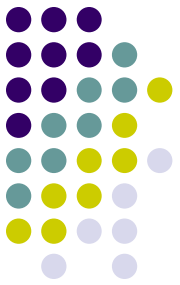
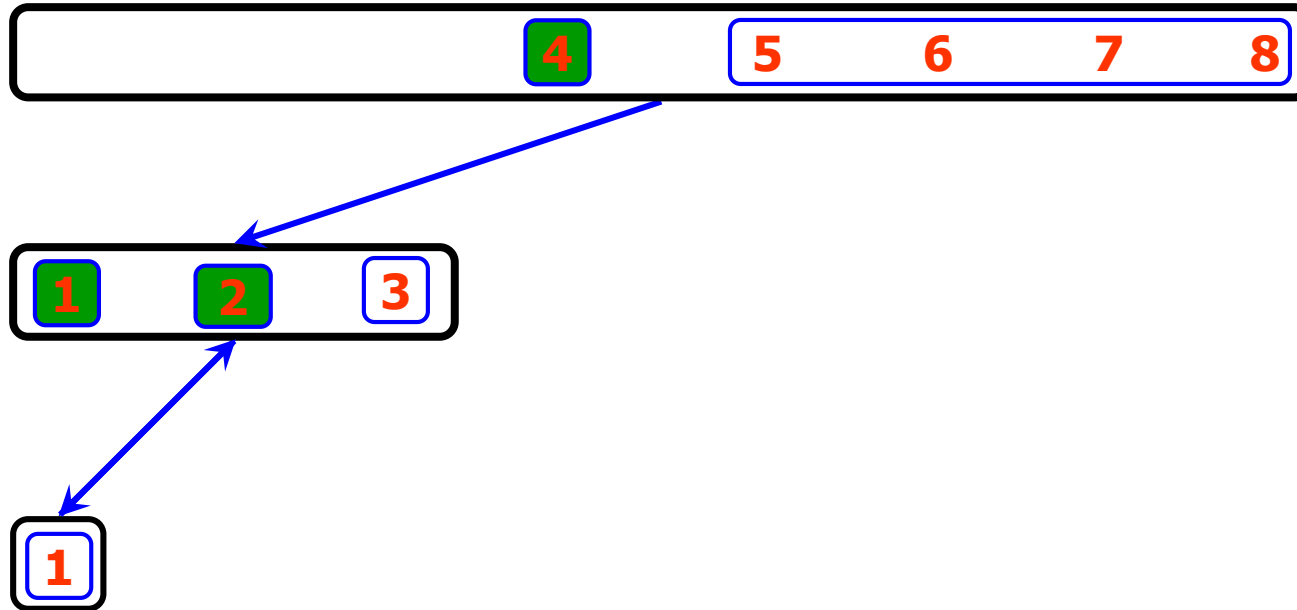
Quick-Sort(A, 0, 7)

Quick-Sort(A, 0, 2), partition



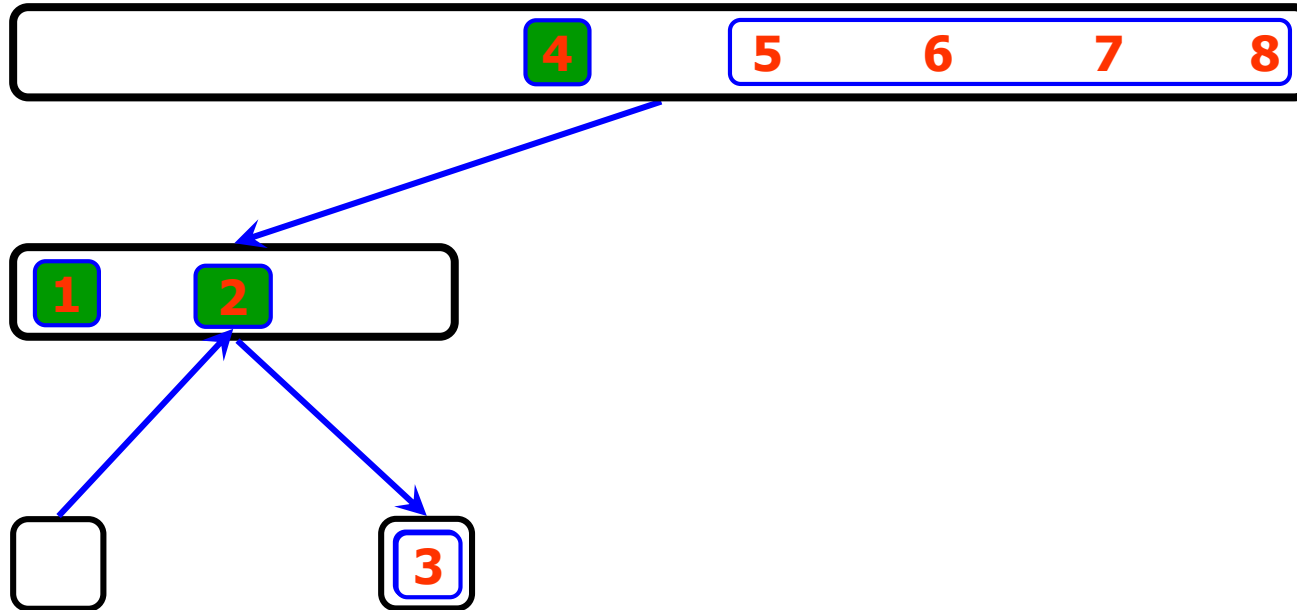
Quick-Sort(A, 0, 7)

Quick-Sort(A, 0, 0), base case



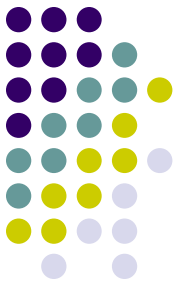
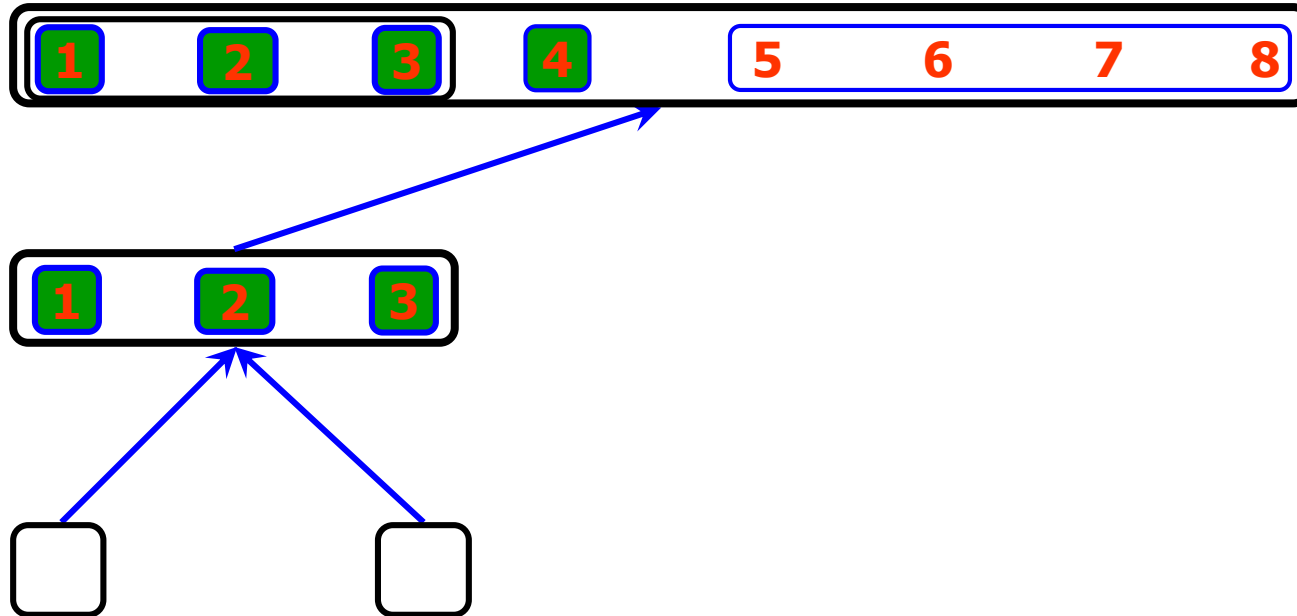
Quick-Sort(A, 0, 7)

Quick-Sort(A, 1, 1), base case



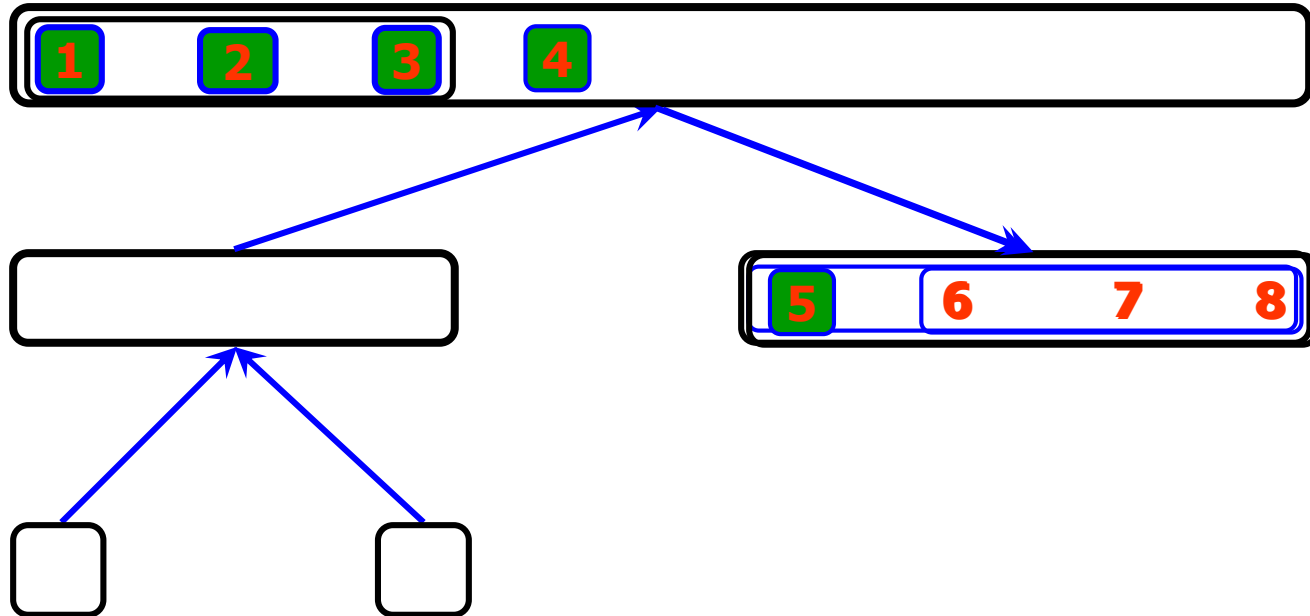
Quick-Sort(A, 0, 7)

Quick-Sort(A, 0, 2), return



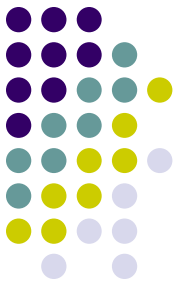
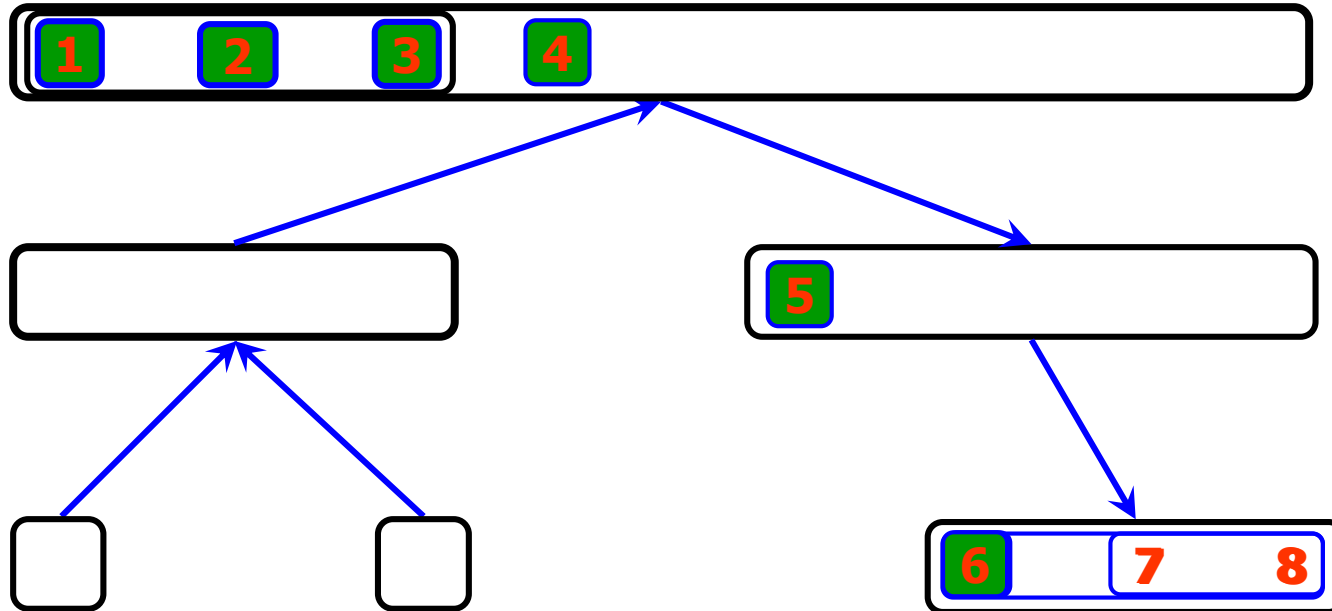
Quick-Sort(A, 0, 7)

Quick-Sort(A, 4, 7), partition



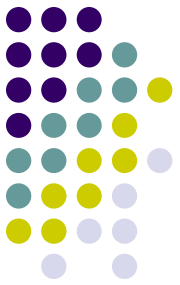
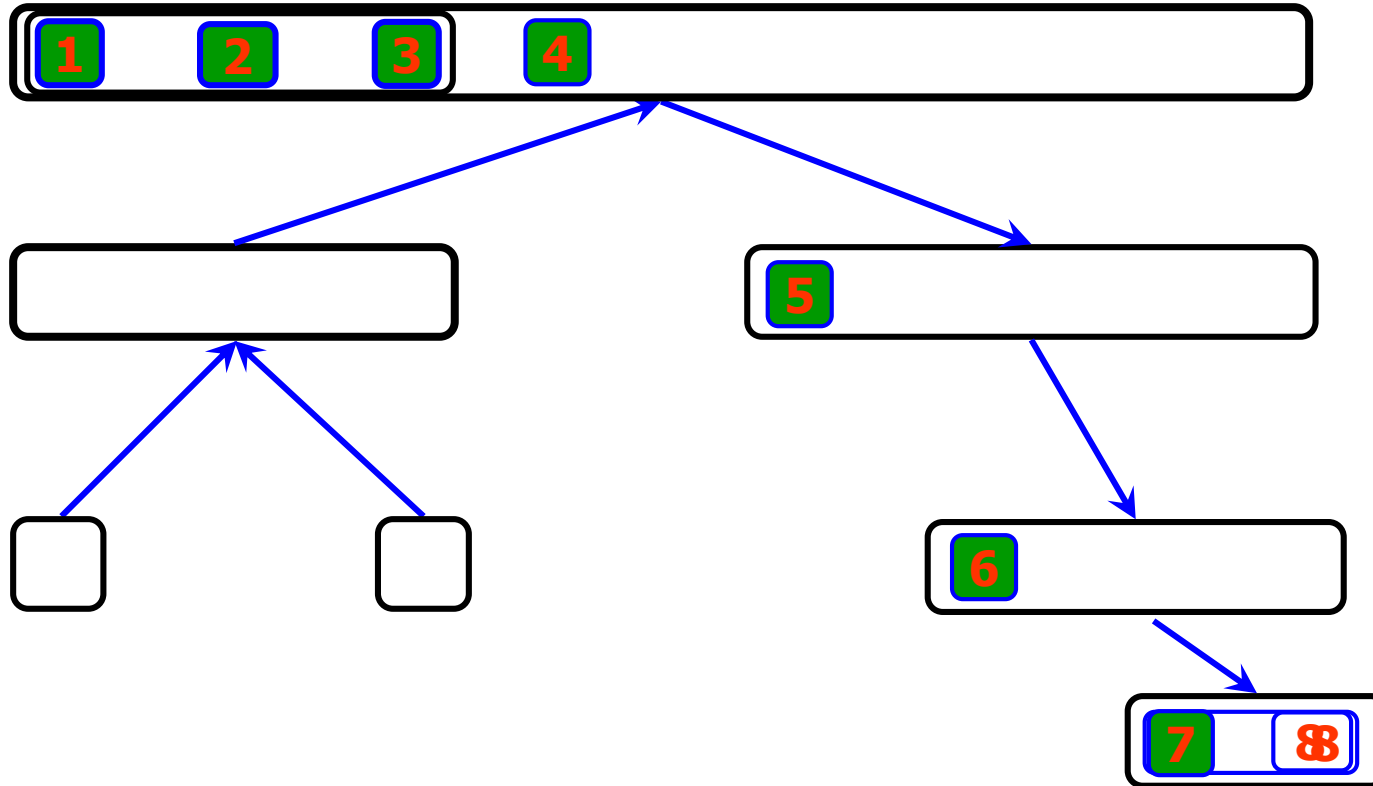
Quick-Sort(A, 0, 7)

Quick-Sort(A, 5, 7), partition



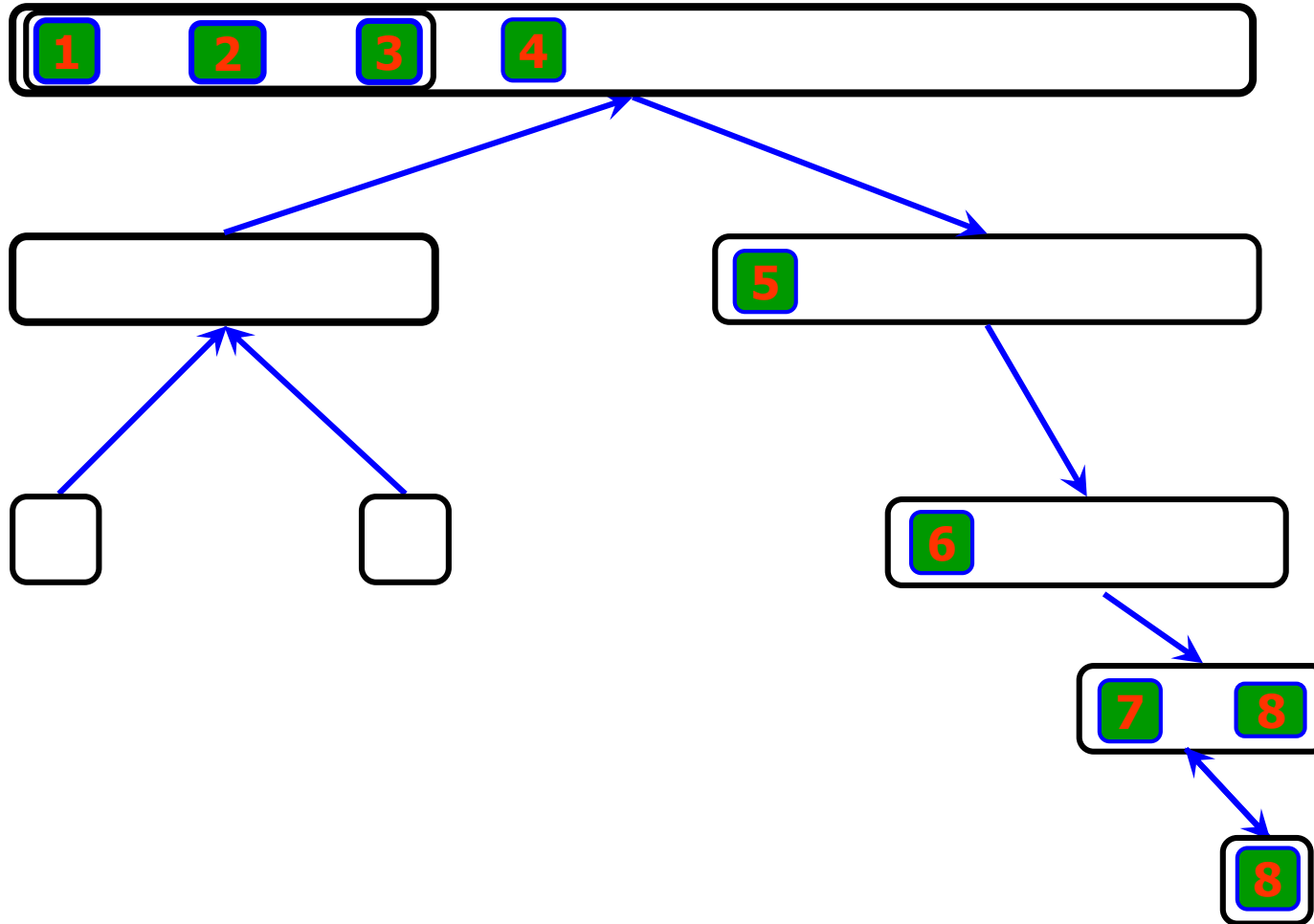
Quick-Sort(A, 0, 7)

Quick-Sort(A, 6, 7), partition



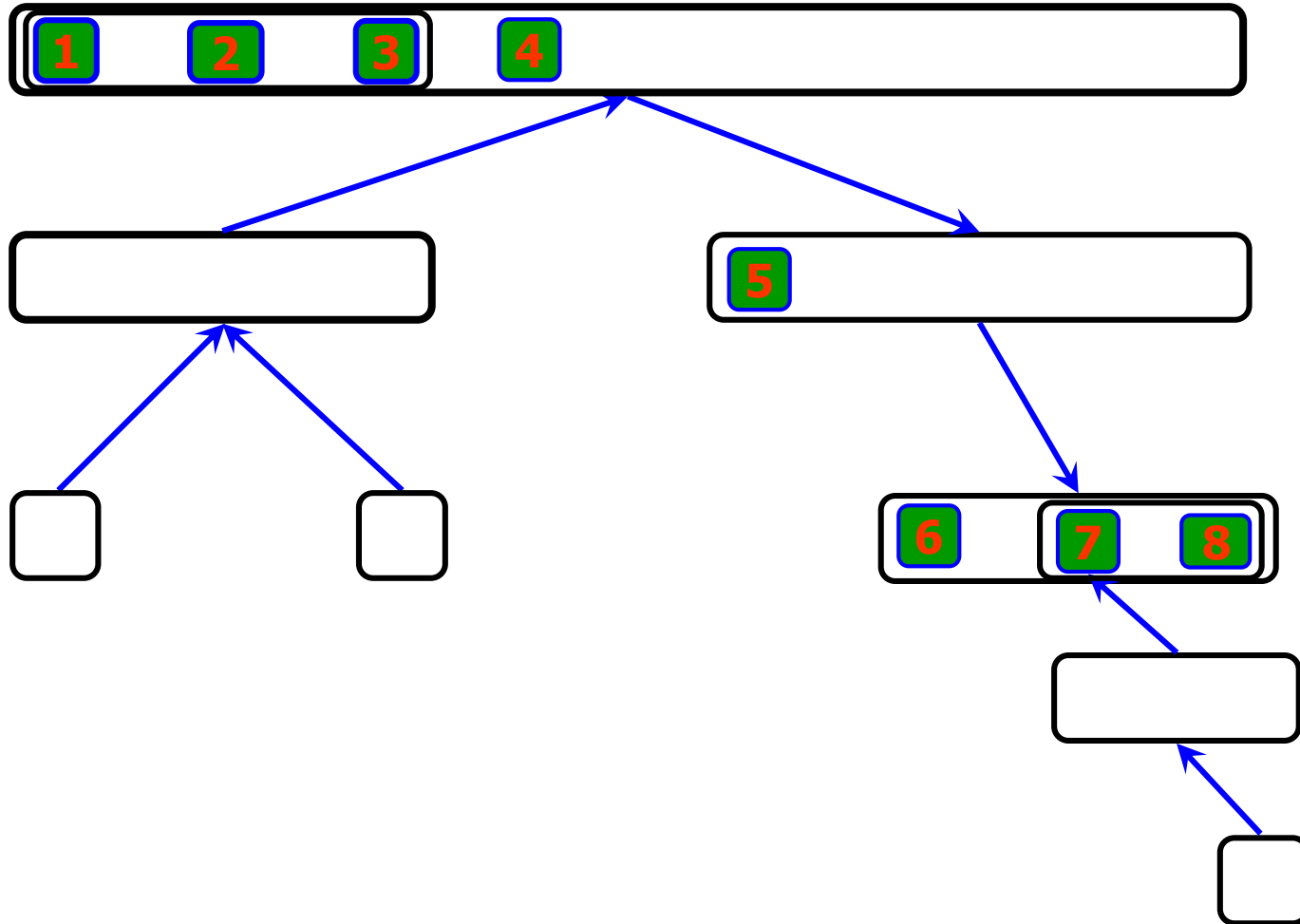
Quick-Sort(A, 0, 7)

Quick-Sort(A, 7, 7), base case



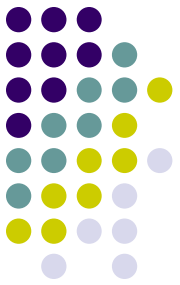
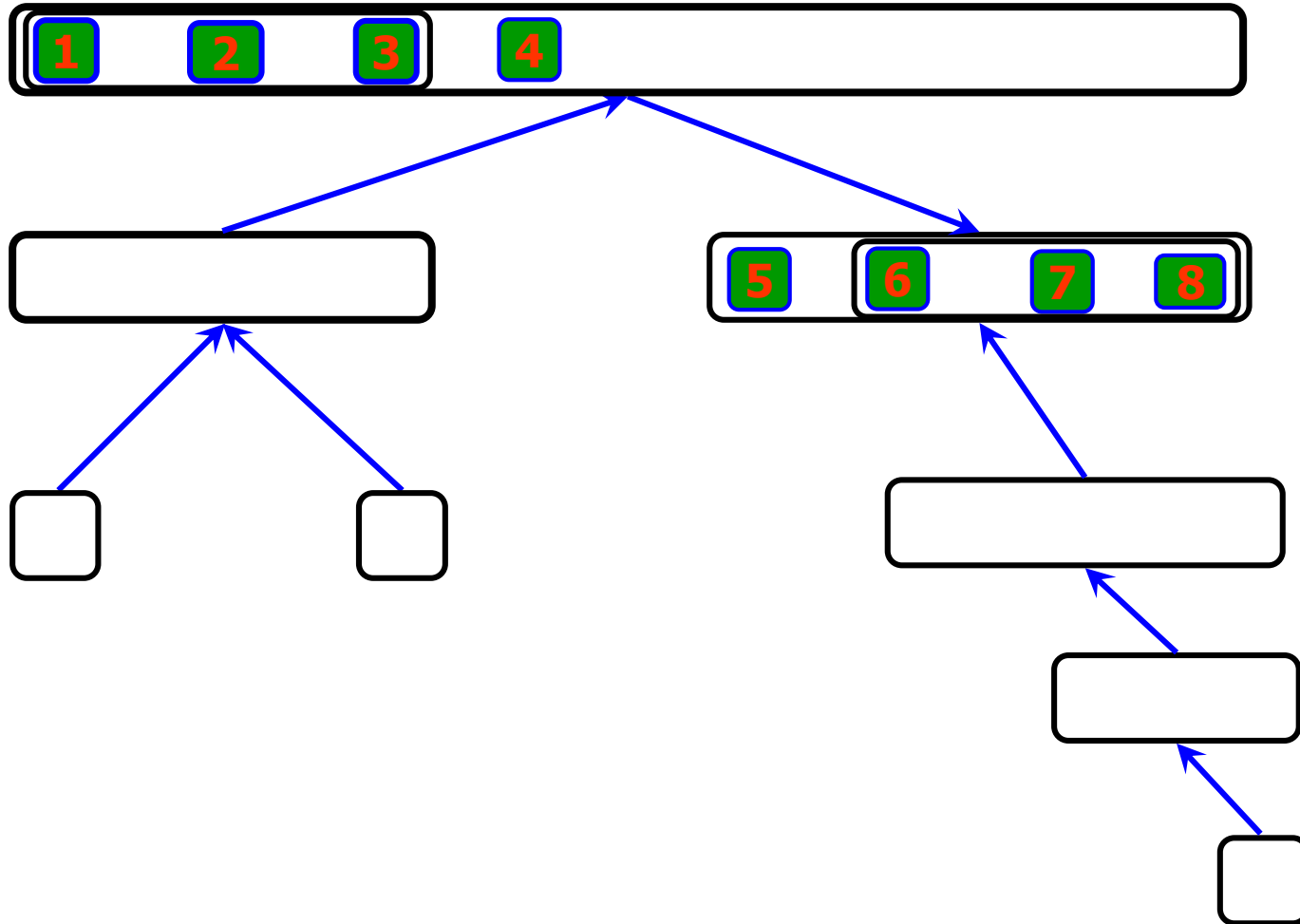
Quick-Sort(A, 0, 7)

Quick-Sort(A, 6, 7) , return



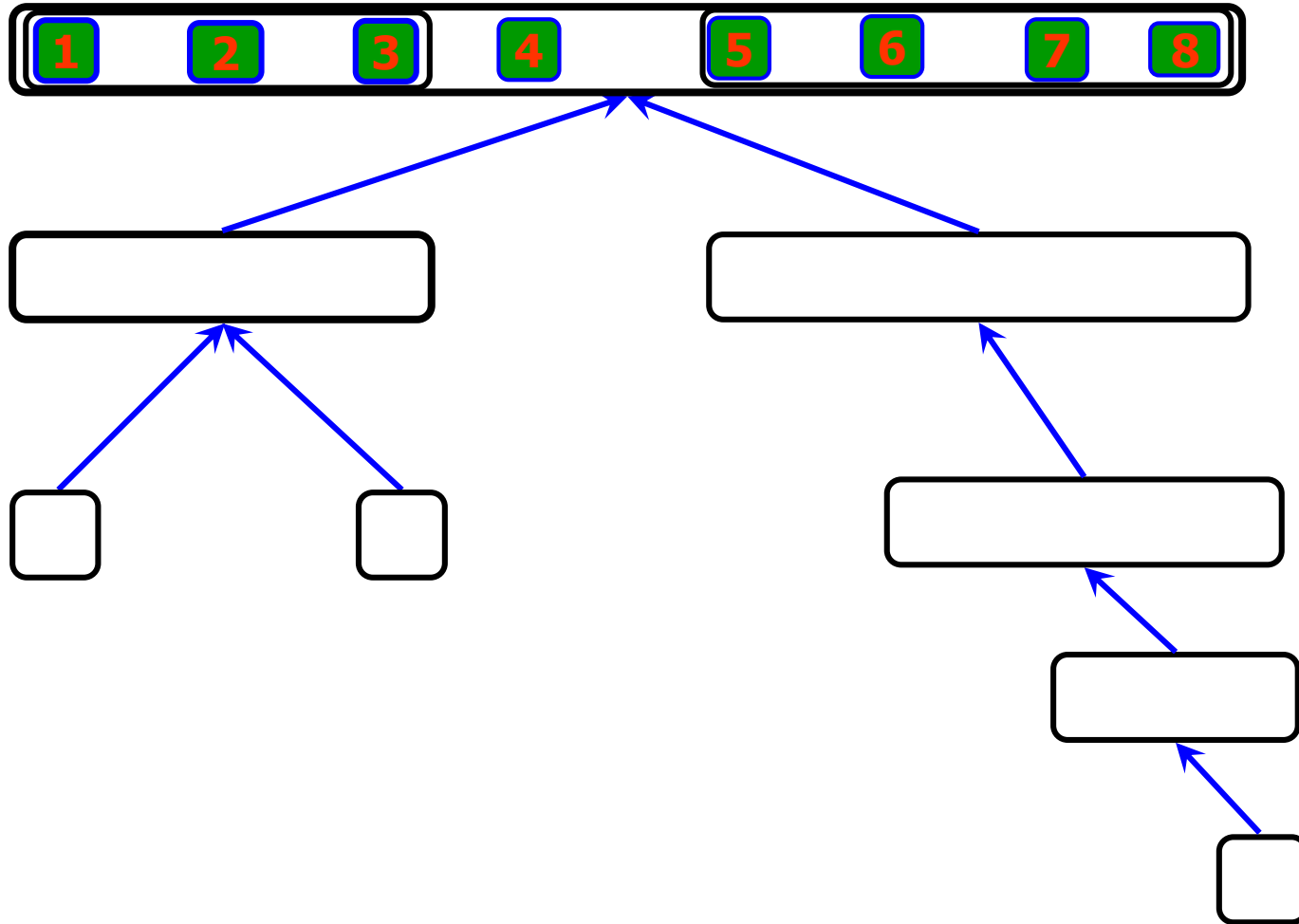
Quick-Sort(A, 0, 7)

Quick-Sort(A, 5, 7) , return



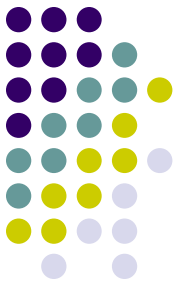
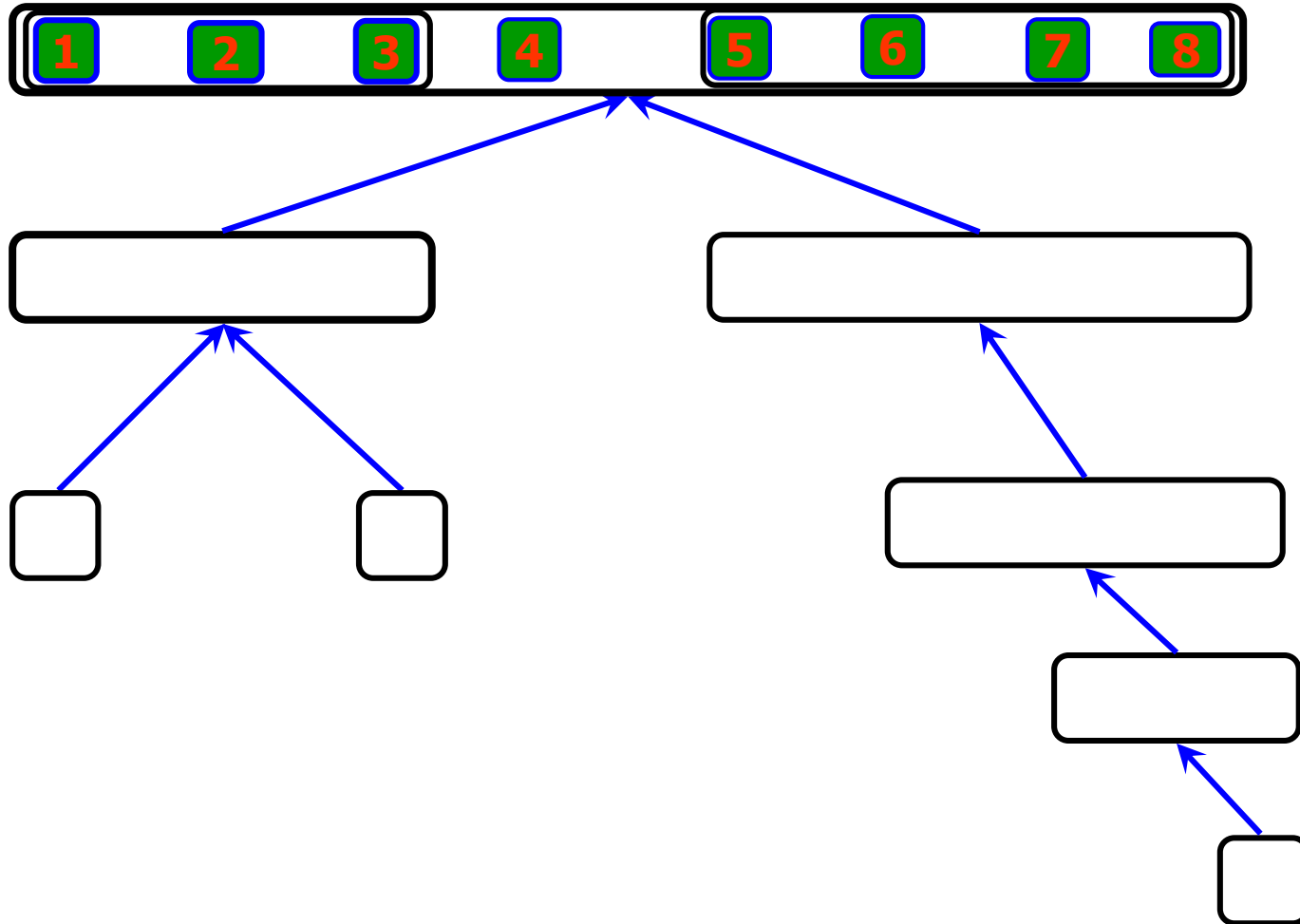
Quick-Sort(A, 0, 7)

Quick-Sort(A, 4, 7) , return



Quick-Sort(A, 0, 7)

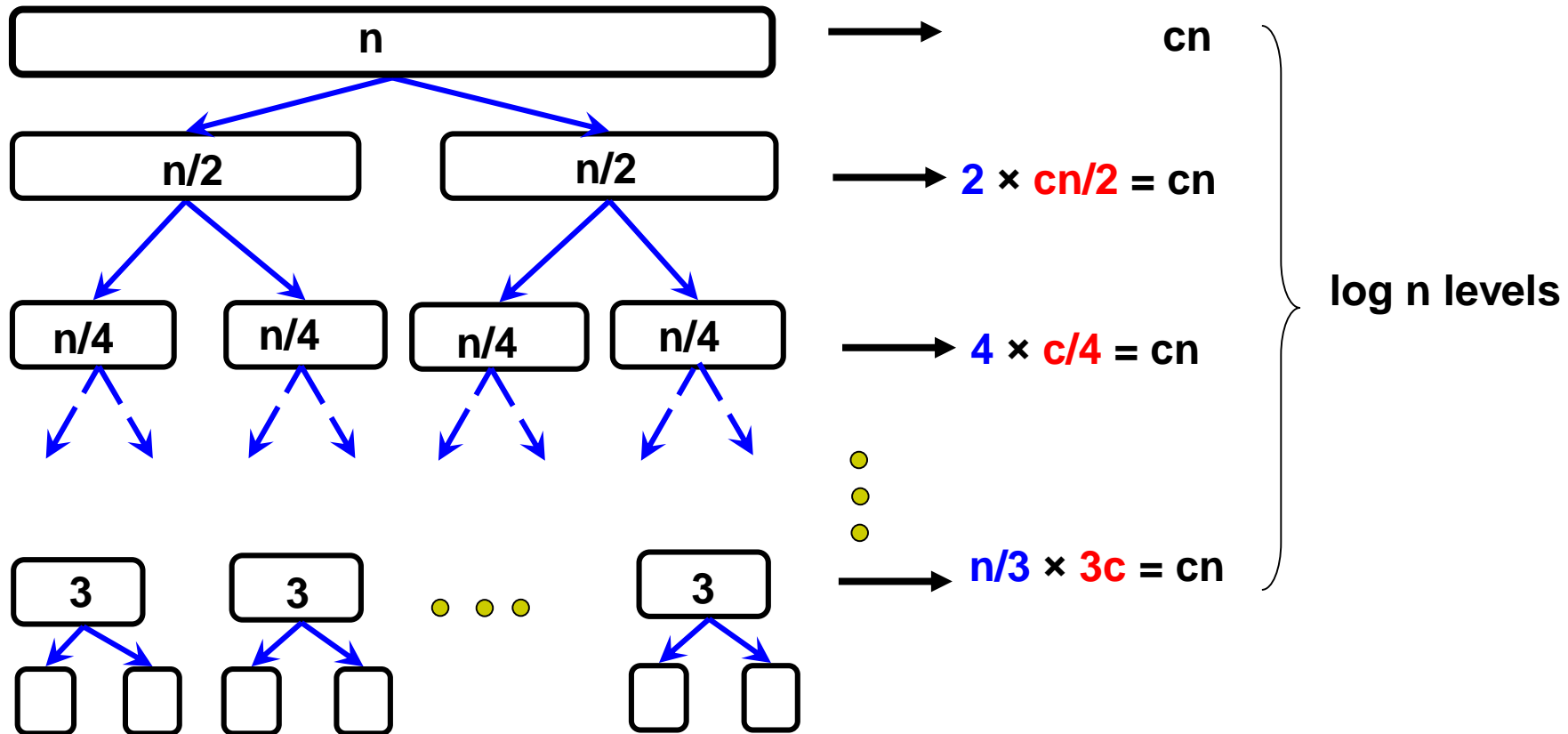
Quick-Sort(A, 0, 7) , **done!**





Quick-Sort: Best Case

- Even Partition

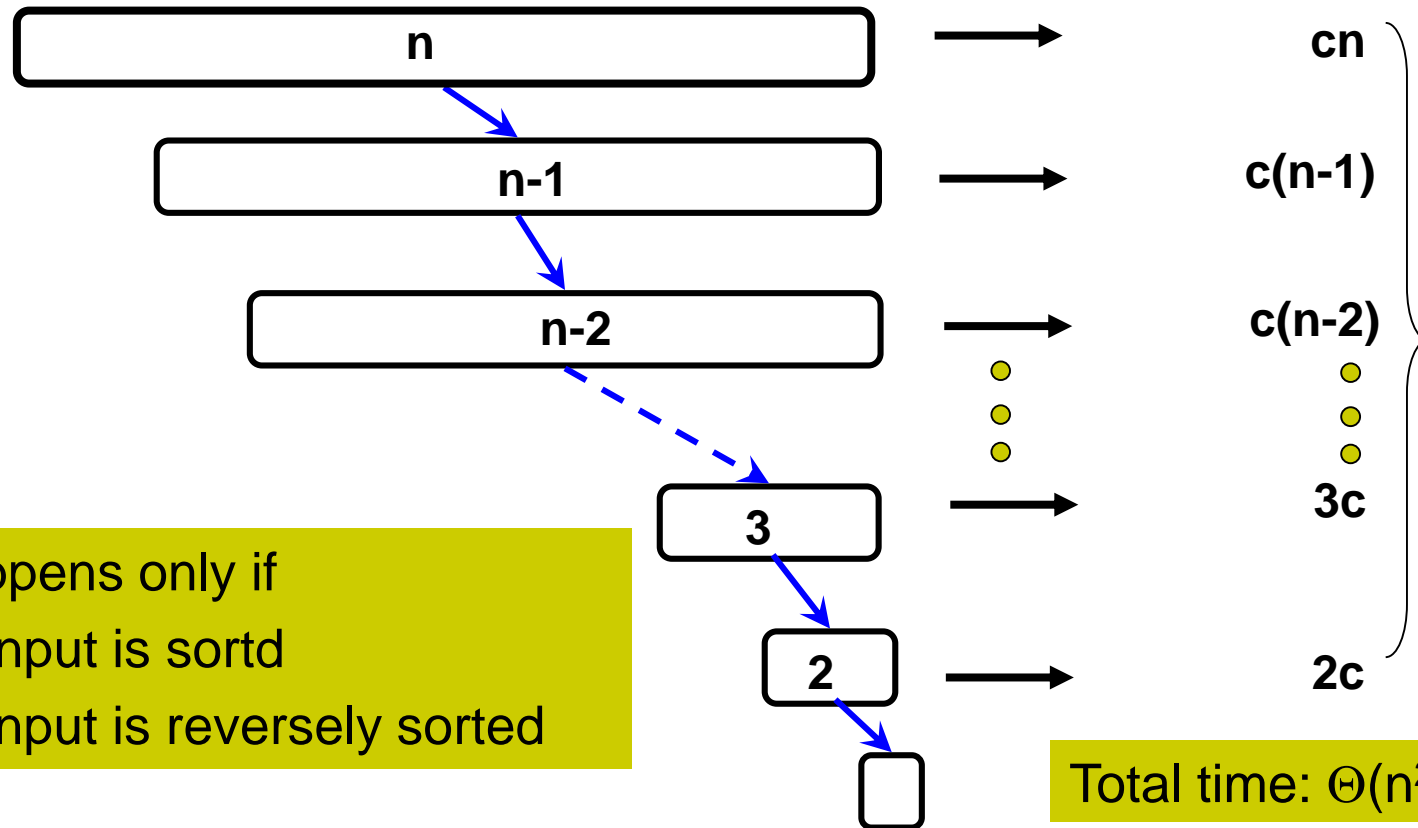


Total time: $\Theta(n \log n)$



Quick-Sort: Worst Case

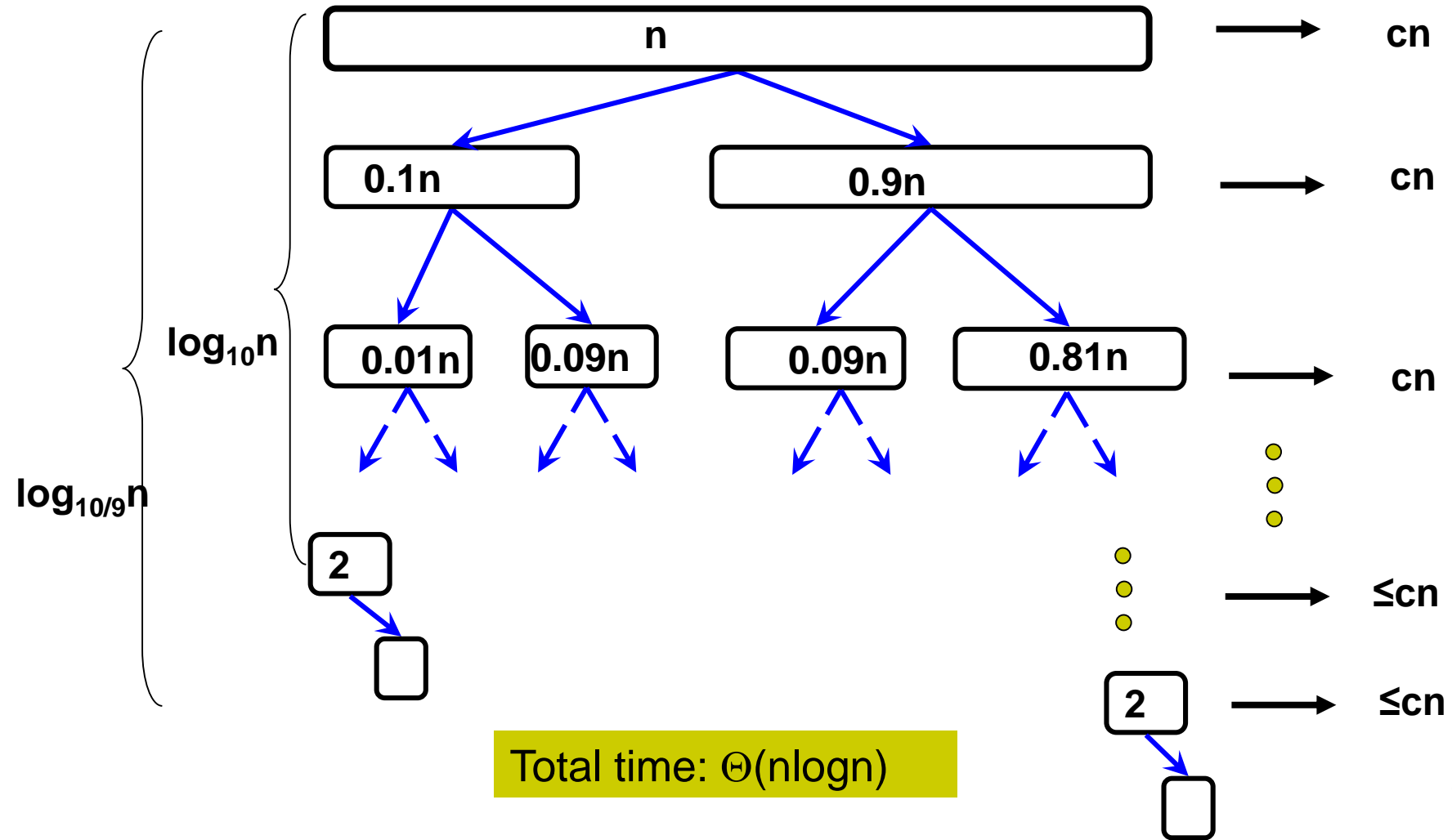
- Unbalanced Partition



Quick-Sort: an Average Case



- Suppose the split is $1/10 : 9/10$



Quick-Sort Summary



- Time

- Most of the work done in partitioning.
- Average case takes $\Theta(n \log(n))$ time.
- Worst case takes $\Theta(n^2)$ time

Space

- Sorts in-place, i.e., does not require additional space

Summary



- Quick-Sort
 - Most of the work done in partitioning
 - Average case takes $\Theta(n \log(n))$ time
 - Worst case takes $\Theta(n^2)$ time
 - $\Theta(1)$ space