

# $\infty$ -Harmonic Functions On $SG$

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September 20, 2011

# The Infinity Laplacian $\Delta_\infty$

- The infinity Laplacian is constructed differently on different domains.

## Definitions

- **Continuous Domains:**

$$\Delta_\infty u(x) = \left( \frac{\nabla u(x)}{\|\nabla u(x)\|} \right)^\top D^2 u(x) \left( \frac{\nabla u(x)}{\|\nabla u(x)\|} \right),$$

where  $D^2 f$  is the Hessian matrix of  $f$ .

- **Graph Domains:**

$$\Delta_\infty u(x) = \frac{1}{2} \left( \max_{y \in N(x)} u(y) + \min_{y \in N(x)} u(y) \right) - u(x)$$

- **Fractal Domains: ??**

# Constructing $\Delta_\infty$ on $SG$

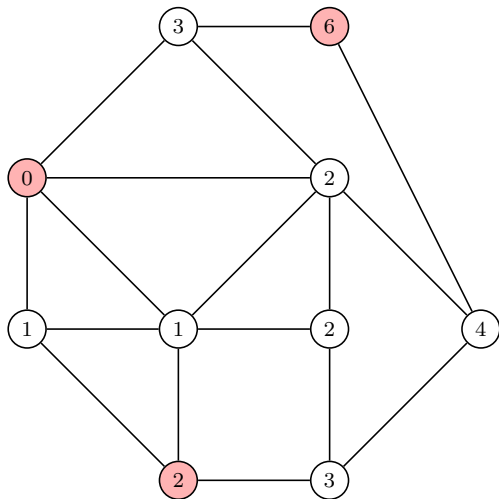
- There is currently no construction for  $\Delta_\infty$  on  $SG$ . We would like to base a construction on a limit of graph approximations.
- This requires a better understanding of the graph infinity Laplacian.
- We can try to first understand  $\infty$ -harmonic functions on a graph:

$$0 = \Delta_\infty u(x)$$

$$0 = \frac{1}{2} \left( \max_{y \in N(x)} u(y) + \min_{y \in N(x)} u(y) \right) - u(x)$$

$$u(x) = \boxed{\frac{1}{2} \left( \max_{y \in N(x)} u(y) + \min_{y \in N(x)} u(y) \right)}$$

- This condition holds everywhere in a graph but at **boundary** nodes, where values are prescribed and fixed.



# $\infty$ -Harmonic Functions on $SG$

- As is standard, let  $\{\Gamma_n\}_{n \in \mathbb{N}}$  be the sequence of graph approximations to  $SG$ .
- $V_i$  is the vertex set of  $\Gamma_i$ .
- For fixed boundary data, let  $u_n : \Gamma_n \rightarrow \mathbb{R}$  be the (discrete)  $\infty$ -harmonic function on  $\Gamma_n$  (with graph  $\infty$ -Laplacian  $\Delta_\infty^{(n)}$ ).
- Define  $u : SG \rightarrow \mathbb{R}$  to be  $\infty$ -harmonic on  $SG$  if

$$u(x) = \lim_{n \rightarrow \infty} u_n(x), \quad \Delta_\infty^{(n)} u_n = 0$$

for all  $n$  such that  $x \in \Gamma_n$ .

# Existence of $\infty$ -Harmonic functions

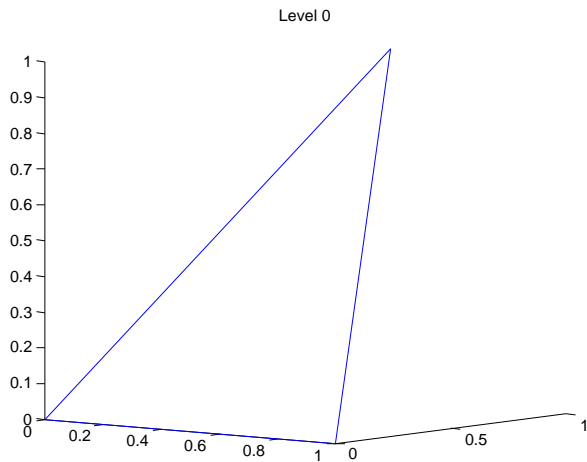
- The goal of this talk is to develop the machinery necessary to prove that the limit defined in the previous slide exists.
- In particular, we will show that, for any boundary data, there exists an  $N$  such that

$$u_n(x)|_{V_1} = u_m(x)|_{V_1}$$

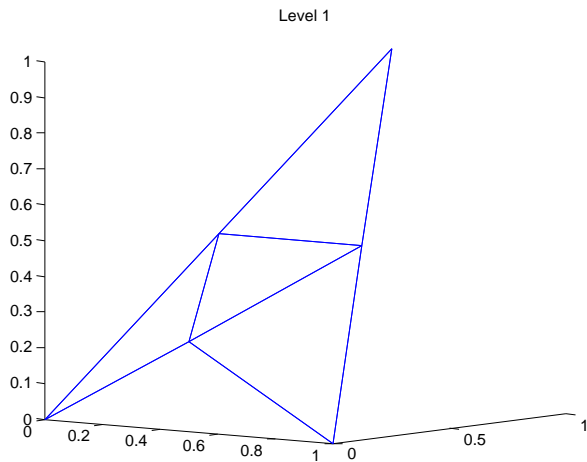
for all  $m, n > N$ .

- This result implies that the same is true for any  $V_j, j \geq 1$ .

# Examples: The Symmetric Case

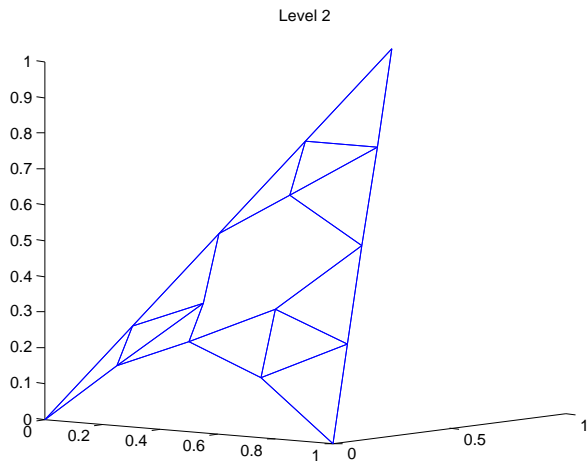


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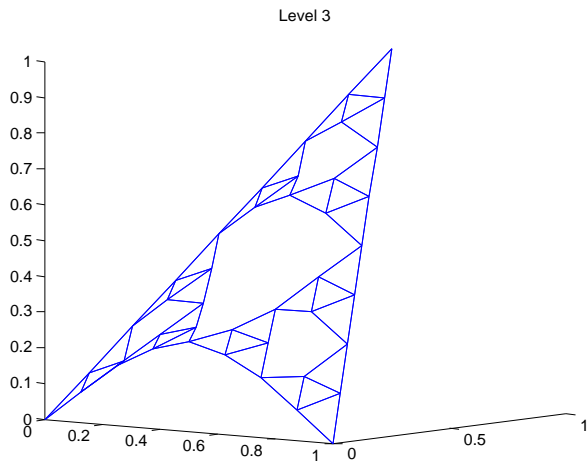




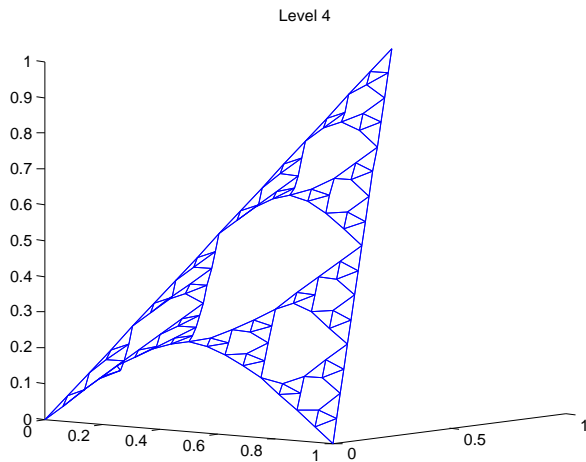
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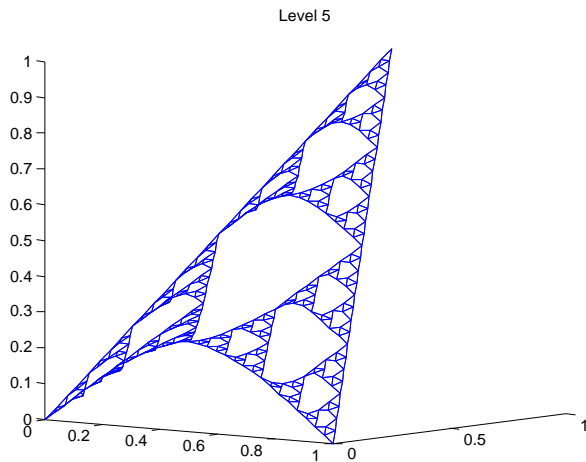
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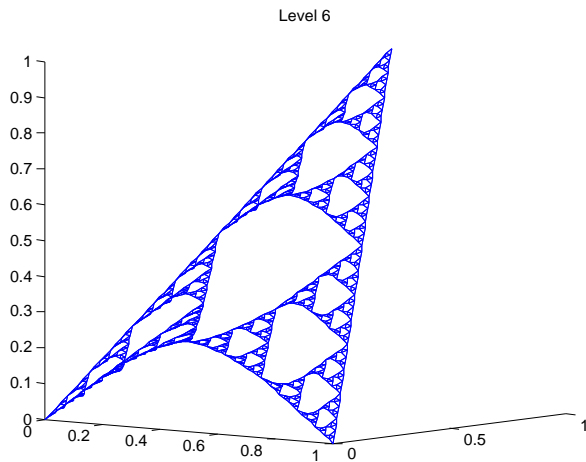
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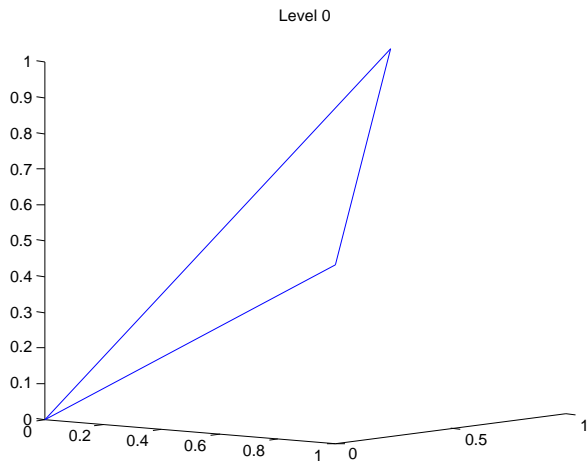
# Examples: The Symmetric Case



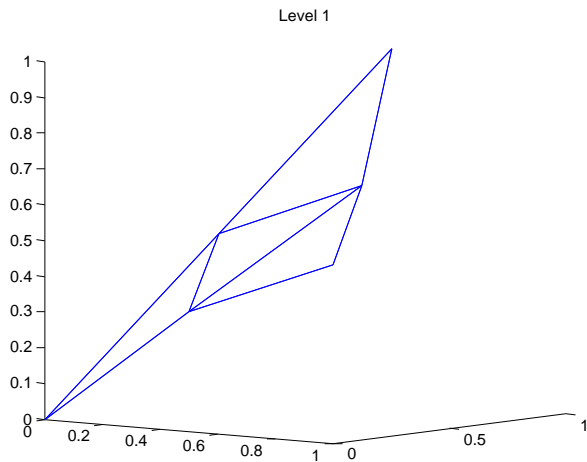
# Examples: The Symmetric Case



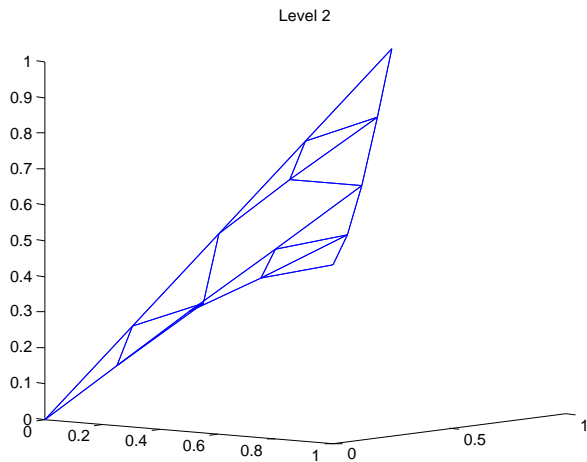
# Examples: The Skew-Symmetric Case



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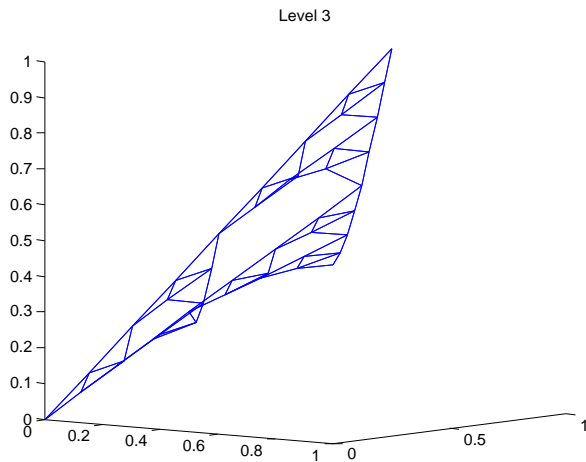


# Examples: The Skew-Symmetric Case

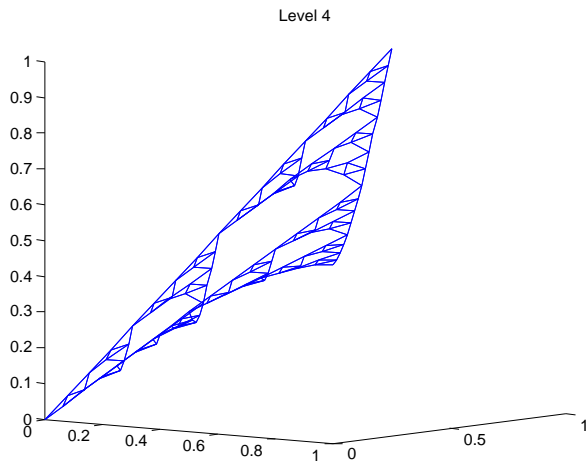




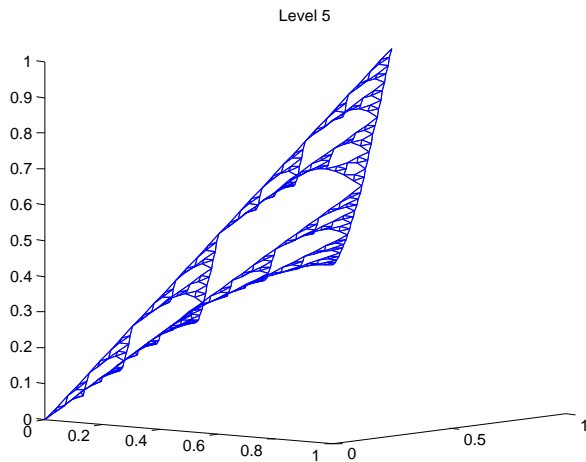
# Examples: The Skew-Symmetric Case



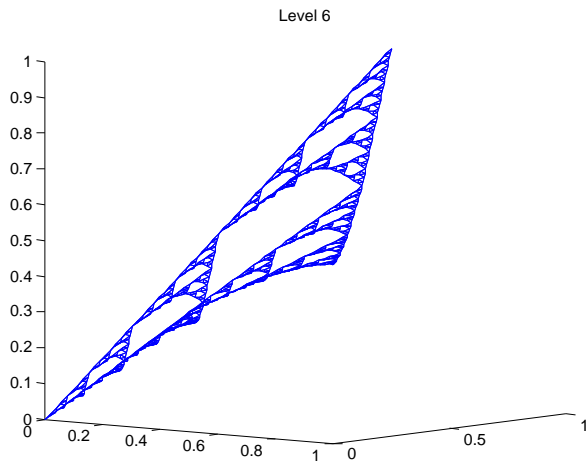
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# The Importance Of Eccentricity

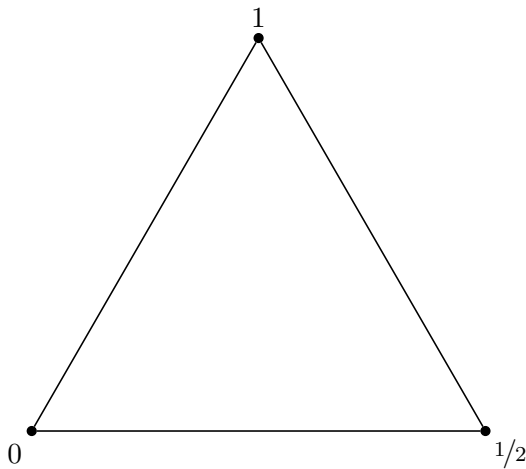
- $\infty$ -harmonic functions (on  $\Gamma_n$  or  $SG$ ) are not additive, but they do have the following properties:

$$\text{If } \Delta_\infty u = 0, \text{ then } \Delta_\infty(cu) = 0 \quad \forall c \in \mathbb{R}.$$

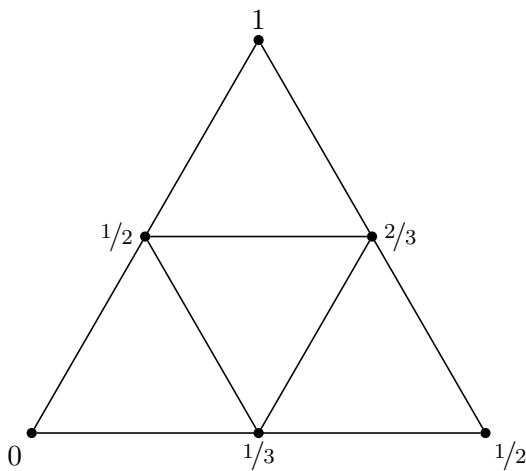
$$\text{If } \Delta_\infty u = 0, \text{ then } \Delta_\infty(u - c) = 0 \quad \forall c \in \mathbb{R}.$$

- These transformations let us reduce the 3-dimensional space of  $\infty$ -harmonic functions on  $SG$  to a 1-dimensional space with boundary data  $(0, e, 1)$  for  $e \in [0, 1/2]$ .
- The eccentricity of  $\Gamma_0$  determines how the  $V_1$  values change on refinement. Numerical evidents suggests the following:
- If  $e \geq 1/3$ ,  $u_n|_{V_1} = u_1|_{V_1}$  for all  $n \geq 1$ .
- As  $e \rightarrow 0$ , the level of refinement at which  $u_n|_{V_1}$  stops changing increases.

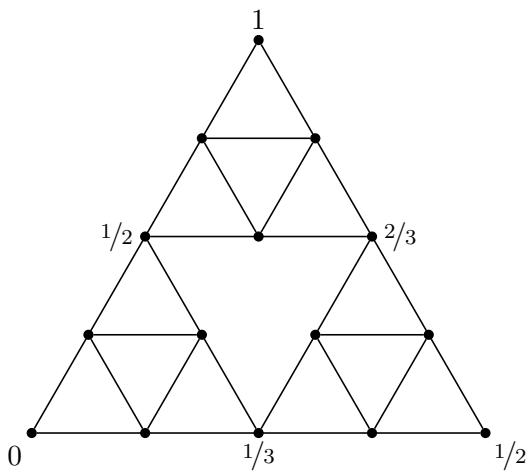
## An Example Which Does Not Change



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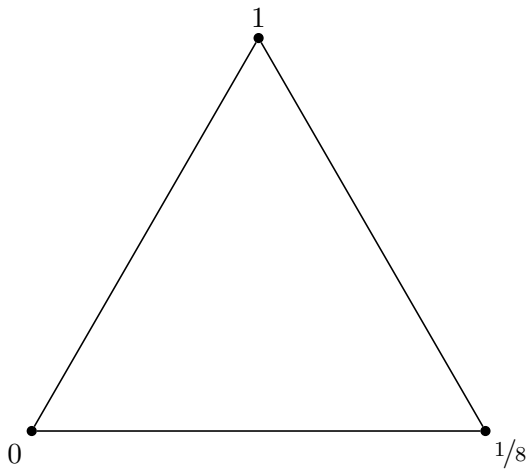


## An Example Which Does Not Change

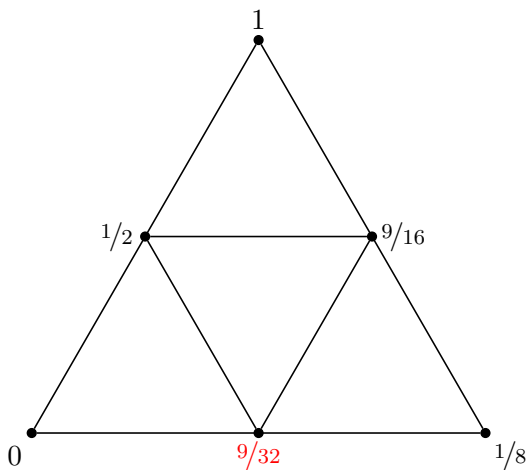




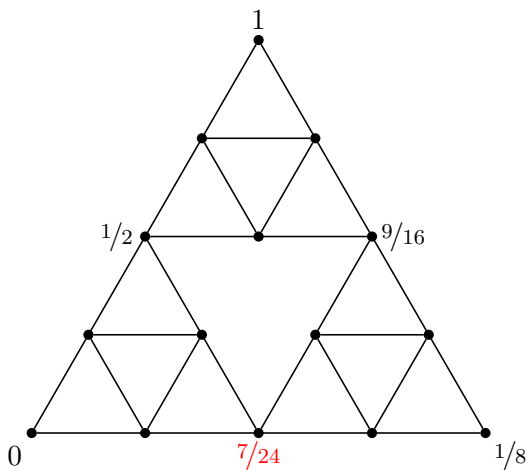
## An Example Which *Does* Change



## An Example Which *Does* Change



# An Example Which Does Not Change



# The Point of this Talk

- We use the concept of eccentricity to state the main theorem to be proved:

**Theorem 1:** For any  $e \in [0, 1/2]$ , there exists a constant  $N$  depending on  $e$  such that  $u_n|_{V_1} = u_m|_{V_1}$  for all  $n, m \geq N$ .

- Proving this requires a better understanding of the relations between successive graph approximations to  $SG$ .
- The behavior of the **Lazarus algorithm** with regards to this graph refinement is of primary importance.

# The Lazarus Algorithm

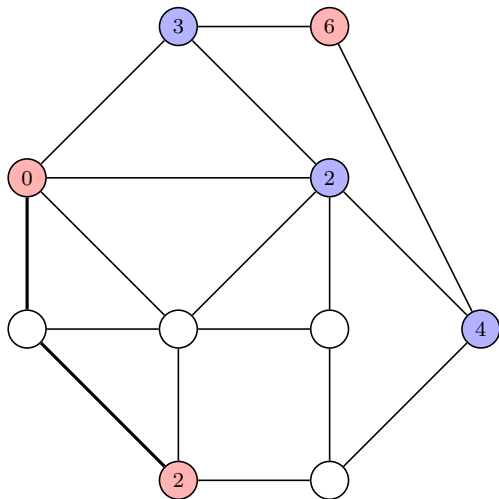
- The **Lazarus algorithm** provides a slow but non-iterative procedure for computing  $\infty$ -harmonic functions on graphs.
- For a graph  $G$  as before and a set of accepted nodes  $A$ , define a path  $x_1 - x_2 - \dots - x_n$  to be **admissible** if  $x_1, x_n \in A$  and  $x_i \notin A$  for all  $i \in [2, n-1]$ .
- For an admissible path with  $n$  nodes and given values  $f(\mathbf{x})$  for  $\mathbf{x} \in A$ , define the **slope** of the path to be  $(f(\mathbf{x}_n) - f(\mathbf{x}_1))/n$

Given this, the Lazarus algorithm proceeds as follows.

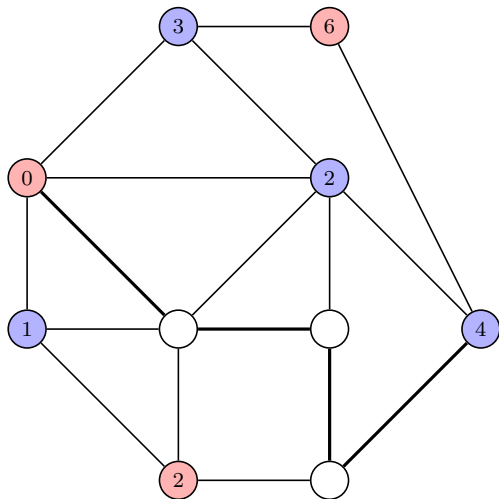
1. In our graph  $G$ , with nodes  $X$  and boundary nodes  $B$ , assume there is some boundary data  $q(\mathbf{x})$  for  $\mathbf{x} \in B$ . Then, initialize  $A := B$ .
2. For each pair of nodes  $\mathbf{x}_i, \mathbf{x}_j \in A$ , compute the slope of the shortest admissible path from  $\mathbf{x}_i$  to  $\mathbf{x}_j$ , if one exists.
3. Along the path with highest slope computed in 2, our  $\infty$ -harmonic function  $u$  is linear. Fill in the new values along this path, and add the unaccepted nodes of the path to  $A$ .
4. If any nodes remain, return to 2.













# Lazarus Paths and Graph Refinements

- Understanding how Lazarus paths change between consecutive graph approximations to  $SG$  is key to proving convergence of  $\infty$ -harmonic functions.
- $G_1, G_2$  graphs with vertex sets  $V_1, V_2$  such that  $V_1 \subset V_2$ . Assume metrics  $d_1, d_2$ .
- For every  $x, y \in V_1$  such that  $y \in N_1(x)$ , the shortest path connecting  $x$  and  $y$  in  $G_2$  **contains no other elements** of  $V_1$ .
- A path  $\rho_2 \subseteq G_2$  is a **refinement** of  $\rho_1 \subseteq G_1$  if  $x \in \rho_1 \Rightarrow x \in \rho_2$ , the terminal nodes of the paths are the same, and the  $\rho_2$  subpath connecting each successive  $x_i, x_{i+1}$  is a  $(d_2)$  geodesic.

# Lazarus Path Refinements on $\{\Gamma_n\}$

- A crucial proposition, due to A. Vladimirovsky, is the following:

**Proposition 1:**  $\ell^k(\Gamma_{m+1})$  is a refinement of  $\ell^k(\Gamma_m)$  for all  $k > 0$  and all  $m \geq k$ .

- With this Proposition, it is easy to show that Theorem 1 is true for  $e \in [1/3, 1/2]$  with  $N = 2$ .

# The $<$ and $\prec$ Graph Relations

- If each path  $\rho_1$  in  $G_1$  has a refinement in  $G_2$ , let us say that

$$G_1 < G_2.$$

- $B_1^0$  is the **boundary set** for the Lazarus algorithm on  $G_1$ ,  $B_2^0$  for  $G_2$ .
- If  $B_1^0 = B_2^0$  and  $d_1(x, y) = d_2(x, y)$  for all  $x, y \in V_1$ , let us say that

$$(G_1, B_1^0) \prec (G_2, B_2^0)$$

$(G_1 \prec G_2$  when there is no danger of ambiguity).

- Then, *the first Lazarus path  $\ell^1(G_2)$  on  $G_2$  is a refinement of the first Lazarus path  $\ell^1(G_1)$  on  $G_1$ .*
- This implies that the  $\infty$ -harmonic function values are equal on the  $V_1$  vertices.

## Problems with $\prec$ !

- The boundary sets for the second Lazarus path are  $B_1^1 = B_1^0 \cup \ell^1(G_1)$  and  $B_2^1 = B_2^0 \cup \ell^1(G_2)$ .
- $B_1^1 \neq B_2^1$ , so

$$(G_1, B_1^1) \not\prec (G_2, B_2^1).$$

- This machinery cannot tell us anything about the equality of any Lazarus paths beyond the first.
- A new graph relation is needed.

## The $\triangleleft$ Graph Relation

- The conditions prescribed by the  $\prec$  relation are stronger than necessary for showing that Lazarus paths are equal.
- For a fixed graph  $G$  and boundary set  $B$ , define

$$\mathcal{B} = \{x \in B \mid \text{for some bdry data, } x \text{ is a terminal node of } \ell^1(G)\}.$$

- In order for  $\mathcal{B}$  to be well-defined, assume a fixed tie-breaking procedure.
- For  $G_1 < G_2$  as before with boundary sets  $B_1, B_2$  and  $\mathcal{B}_1 \subseteq B_1, \mathcal{B}_2 \subseteq B_2$ , say

$$(G_1, \mathcal{B}_1) \triangleleft (G_2, \mathcal{B}_2)$$

if  $\mathcal{B}_1 = \mathcal{B}_2$  and  $d_1(x, y) = d_2(x, y)$  for all  $x, y \in \mathcal{B}_1$ .

- With this relation, it can be shown that  $\ell^1(G_2)$  is a refinement of  $\ell^1(G_1)$ .

# A Reinterpretation of the Lazarus Algorithm

- on a graph  $G^0$  with vertices  $V^0$ , edges  $E^0$ , and boundary set  $B^0$ , the first Lazarus path is some  $\ell^1(G^0) = (x_1, \dots, x_n), x_i \in V^0$ .
- Create a new graph  $G^1$  (the **Lazarus subgraph** of  $G^0$ ) with vertex set  $V^1 = V^0$ , boundary set  $B^1 = B^0 \cup \ell^1(G^0)$ , and edge set  $E^1 = E^0 - \ell^1(G^0)$ .

- Then,

$$\ell^1(G^1) = \ell^2(G^0).$$

- We can create a  $G^2$  analogously from  $G^1$  such that

$$\ell^1(G^2) = \ell^3(G^0),$$

and so on.

- From this perspective, it is enough to talk about the Lazarus path  $\ell(G^i)$  on a graph, since each has only one.



$\triangleleft$  and  $\{\Gamma_n\}$

- Define  $\Gamma_j^0 := \Gamma_j$  for  $j \in \mathbb{N}_0$ , and  $\Gamma_j^i$  to be the **Lazarus subgraph** of  $\Gamma_j^{i-1}$  for  $i \geq 1$ .
- A priori* we know the following to be true:

$$\begin{array}{ccccccccc}
 \Gamma_0^0 & \prec & \Gamma_1^0 & \prec & \Gamma_2^0 & \prec & \Gamma_3^0 & \prec & \dots \\
 \vee & & \vee & & \vee & & \vee & & \vee \\
 \Gamma_0^1 & & \Gamma_1^1 & & \Gamma_2^1 & & \Gamma_3^1 & & \dots \\
 \vee & & \vee & & \vee & & \vee & & \vee \\
 \Gamma_0^2 & & \Gamma_1^2 & & \Gamma_2^2 & & \Gamma_3^2 & & \dots \\
 \vee & & \vee & & \vee & & \vee & & \vee \\
 \Gamma_0^3 & & \Gamma_1^3 & & \Gamma_2^3 & & \Gamma_3^3 & & \dots \\
 & & \vee & & \vee & & \vee & & \vee \\
 & & \vdots & & \vdots & & \vdots & & \ddots
 \end{array}$$

$\triangleleft$  and  $\{\Gamma_n\}$

- To prove **Proposition 1**, it suffices to show that the following relations hold true:

$$\begin{array}{ccccccc}
 \Gamma_0^0 & \prec & \Gamma_1^0 & \prec & \Gamma_2^0 & \prec & \Gamma_3^0 & \prec & \dots \\
 \vee & & \vee & & \vee & & \vee & & \vee \\
 \Gamma_0^1 & & \Gamma_1^1 & \triangleleft & \Gamma_2^1 & \triangleleft & \Gamma_3^1 & \triangleleft & \dots \\
 \vee & & \vee & & \vee & & \vee & & \vee \\
 \Gamma_0^2 & & \Gamma_1^2 & & \Gamma_2^2 & \triangleleft & \Gamma_3^2 & \triangleleft & \dots \\
 \vee & & \vee & & \vee & & \vee & & \vee \\
 \Gamma_0^3 & & \Gamma_1^3 & & \Gamma_2^3 & & \Gamma_3^3 & \triangleleft & \dots \\
 & & \vee & & \vee & & \vee & & \vee \\
 & & \vdots & & \vdots & & \vdots & \triangleleft & \ddots
 \end{array}$$

- This is the motivation for

**Proposition 2:** Let  $\{\Gamma_n\}$  be the standard sequence of graph approximations to  $SG$ . Then,  $\Gamma_m^j \triangleleft \Gamma_{m+1}^j$  for all  $j \geq 0$  and all  $m \geq j$ .

# A Sketch of a Proof for Proposition 2

- The proof relies on each  $\Gamma_m^j$  having the following two properties for  $m \geq j$ :
  - P1.  $\mathcal{B}_m^j = \mathcal{B}_{j-1}^j$
  - P2. If the shortest path between  $x, y \in B_m^j$  in  $\Gamma_{m+1}^j$  is not a refinement of the shortest path in  $\Gamma_m^j$ , then  $x$  or  $y$  is in  $B_m^j - B_{m-1}^j$ .
- These properties can be directly verified for  $\Gamma_m^0$ , and then the proof proceeds inductively.
- The proof that Proposition 2  $\Rightarrow$  Proposition 1 is straightforward.

## A generalization of Proposition 2

- We were seeking a sequence of Lazarus paths on members of  $\{\Gamma_n\}$  which partitioned the graphs into subgraphs isomorphic to members of  $\{\Gamma_n\}$ .
- It is useful to consider Lazarus paths partitioning graphs into other subgraphs, as well.
- A **regular finite fractafold** (rff) is the quotient space

$$\left( \bigsqcup_{i=1}^n SG \right) / \varphi,$$

where  $\varphi$  is a quotient map such that in the  $i$ th copy of  $SG$  (call it  $SG_i$ ),  $\varphi$  identifies each boundary point of  $SG_i$  with at most one boundary point of at most one  $SG_j$ ,  $i \neq j$ , and  $\varphi$  never identifies two different boundary points in  $SG_i$  with boundary points in a single  $SG_j$ .

# Regular Finite Fractafolds

- If  $F = (\bigsqcup_{i=1}^n SG) / \varphi$ , define

$$F_k = \left( \bigsqcup_{i=1}^n \Gamma_k \right) / \varphi.$$

- Using [Proposition 2](#), it can be shown that the analogue of Proposition 2 holds:  $F_m^j \triangleleft F_{m+1}^j$  for  $j > 0$ ,  $m \geq j$ .
- If we start with  $\{\Gamma_n\}$  approximations to  $SG$ , we can directly show that each Lazarus path partitions the approximation graphs into rff's.

# Convergence results on $SG$

- Using the generalization of Proposition 2, we finally get convergence bounds which allow us to prove Theorem 1.

**Result:**  $u_n|_{V_1}$  does not change upon refinement for all eccentricities  $e > \frac{2}{3^{n-2}-3}$ ,  $n > 4$ .

- This proves Theorem 1 for all  $e \in (0, 1/2]$ . The  $e = 0$  case can be taken care of by an argument due to R. Strichartz, completing the proof of Theorem 1!

# Conclusion, Future Work

- This gives us a plausible definition for  $\Delta_\infty u = 0$ , but nothing with nonzero RHS.
- We would like to prove something analogous for functions  $u_n : \Gamma_n \rightarrow \mathbb{R}$  that are solutions to

$$-c^n \Delta_\infty^{(n)} u_n = f_n,$$

where  $c$  is some renormalization constant.