

Johnson's algorithm for All-pairs shortest paths

- Johnson's Algorithm uses both Dijkstra's Algorithm and Bellman-Ford Algorithm.
- Johnson's Algorithm uses the technique of "**reweighting**."
- If all edge weights w in a graph $G = (V, E)$ are nonnegative, we can find the shortest paths between all pairs of vertices by running Dijkstra's Algorithm once from each vertex.
- If G has negative - weight edges, we compute a new - set of non - negative edge weights that allows us to use the same method.

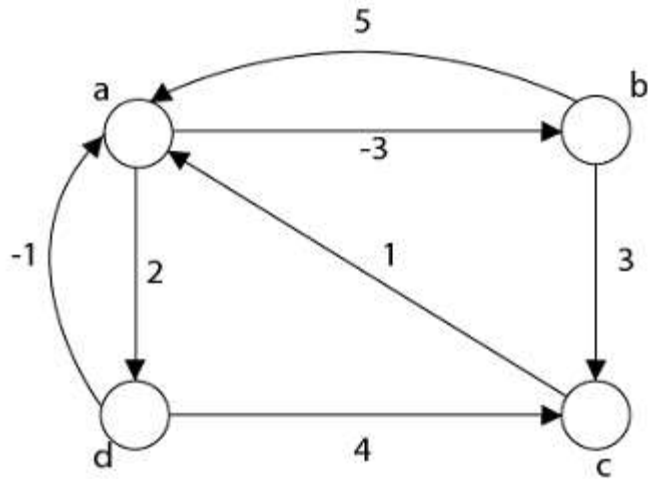
The new set of edges weight \hat{w} must satisfy two essential properties:

1. For all pairs of vertices $u, v \in V$, a path p is a shortest path from u to v using weight function w if and only if p is also a shortest path from u to v using weight function \hat{w} .
2. For all edges (u, v) , the new weight $\hat{w}(u, v)$ is nonnegative.

- For each edge $(u, v) \in E$ define

$$\hat{w}(u, v) = w(u, v) + h(u) - h(v) .$$

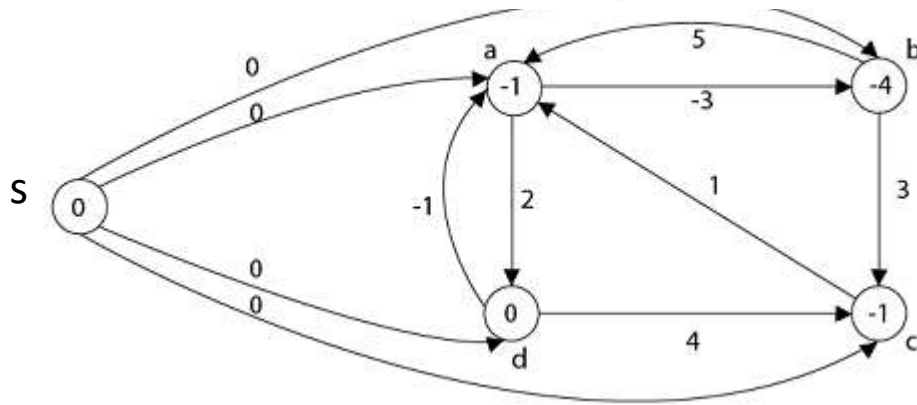
- Where $h(u) = \text{label of } u$
 $h(v) = \text{label of } v$



Step1: Take any source vertex 's' outside the graph and make distance from 's' to every vertex '0'.

Step2: Apply Bellman-Ford Algorithm and calculate minimum weight on each vertex.

$$dist^k[u] = \min \{dist^{k-1}[u], \min_i \{dist^{k-1}[i] + cost[i, u]\}\} \quad k = 2, 3, \dots, n-1.$$



Step3: $w(a, b) = w(a, b) + h(a) - h(b)$
 $= -3 + (-1) - (-4)$
 $= 0$

$$w(b, a) = w(b, a) + h(b) - h(a)$$

$$= 5 + (-4) - (-1)$$

$$= 2$$

$$w(b, c) = w(b, c) + h(b) - h(c)$$

$$w(b, c) = 3 + (-4) - (-1)$$

$$= 0$$

$$w(c, a) = w(c, a) + h(c) - h(a)$$

$$w(c, a) = 1 + (-1) - (-1)$$

$$= 1$$

$$w(d, c) = w(d, c) + h(d) - h(c)$$

$$w(d, c) = 4 + 0 - (-1) \\ = 5$$

$$w(d, a) = w(d, a) + h(d) - h(a)$$

$$w(d, a) = -1 + 0 - (-1) \\ = 0$$

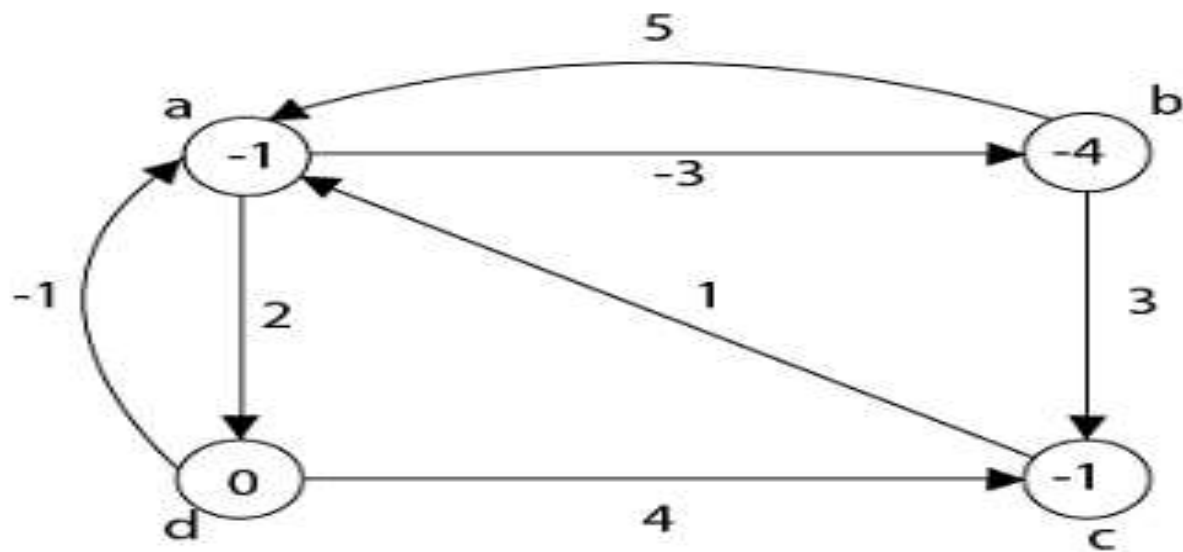
$$w(a, d) = w(a, d) + h(a) - h(d)$$

$$w(a, d) = 2 + (-1) - 0 = 1$$

- **Step 4:** Now all edge weights are positive and now we can apply Dijkstra's Algorithm on each vertex and make a matrix corresponds to each vertex in a graph

	a	b	c	d
a	0	0	0	1
b	1	0	0	2
c	1	1	0	2
d	0	0	0	0

Step5:



$$d_{uv} \leftarrow \delta(u, v) + h(v) - h(u)$$

$$d(a, a) = 0 + (-1) - (-1) = 0$$

$$d(a, b) = 0 + (-4) - (-1) = -3$$

$$d(a, c) = 0 + (-1) - (-1) = 0$$

$$d(a, d) = 1 + (0) - (-1) = 2$$

$$d(b, a) = 1 + (-1) - (-4) = 4$$

$$d(b, b) = 0 + (-4) - (-4) = 0$$

$$d(c, a) = 1 + (-1) - (-1) = 1$$

$$d(c, b) = 1 + (-4) - (-1) = -2$$

$$d(c, c) = 0$$

$$d(c, d) = 2 + (0) - (-1) = 3$$

$$d(d, a) = 0 + (-1) - (0) = -1$$

$$d(d, b) = 0 + (-4) - (0) = -4$$

$$d(d, c) = 0 + (-1) - (0) = -1$$

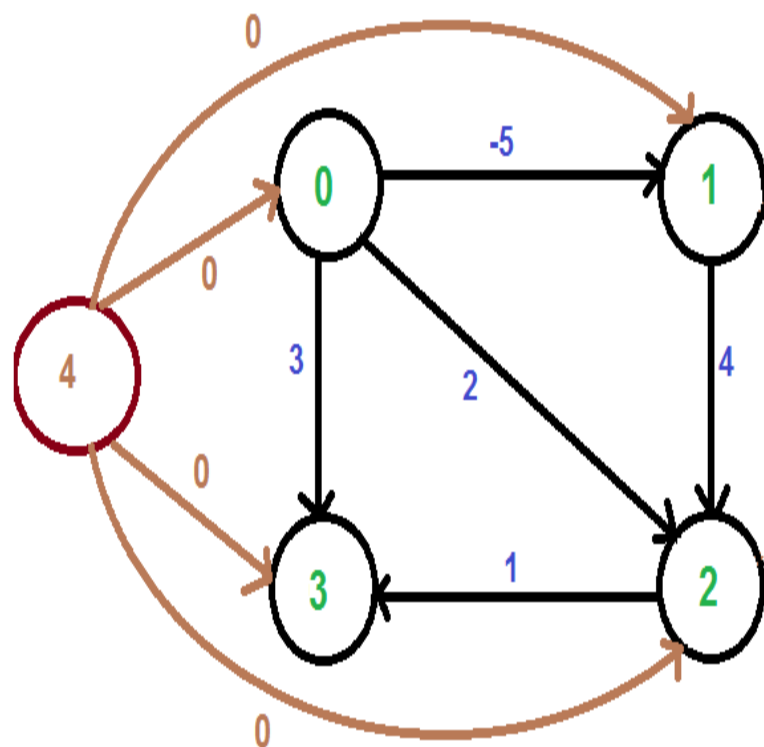
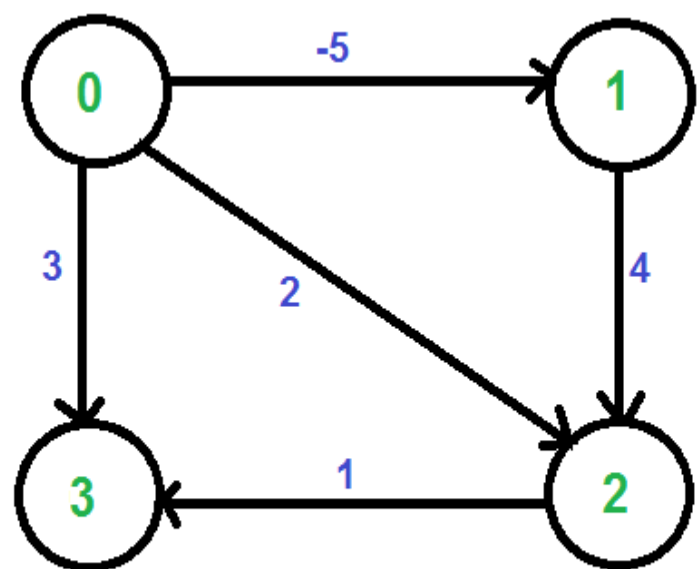
$$d(d, d) = 0$$

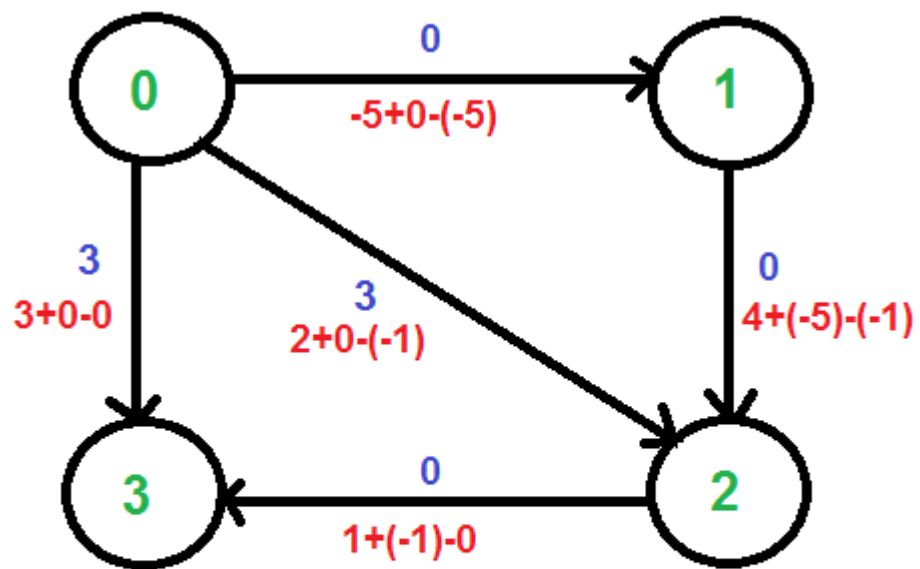
	a	b	c	d
a	0	-3	0	2
b	4	0	3	6
c	1	-2	0	3
d	-1	-4	-1	0

Algorithm

JOHNSON (G)

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1. Compute  $G'$  where  $V[G'] = V[G] \cup \{s\}$  and  
    $E[G'] = E[G] \cup \{(s, v) : v \in V[G]\}$   
  
2. If  $BELLMAN-FORD(G', w, s) = FALSE$   
   then "input graph contains a negative weight cycle"  
   else  
     for each vertex  $v \in V[G']$   
       do  $h(v) \leftarrow \delta(s, v)$   
     Computed by Bellman-Ford algorithm  
     for each edge  $(u, v) \in E[G']$   
       do  $w(u, v) \leftarrow w(u, v) + h(u) - h(v)$   
     for each edge  $u \in V[G]$   
       do run DIJKSTRA( $G, w, u$ ) to compute  
          $\delta(u, v)$  for all  $v \in V[G]$   
       for each vertex  $v \in V[G]$   
         do  $d_{uv} \leftarrow \delta(u, v) + h(v) - h(u)$   
Return D.
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Distances from 4 to 0, 1, 2 and 3 are 0, -5, -1 and 0 respectively.