

Quantum Hamiltonian Simulation

From Trotterization to QSVT and Applications to Grover's Search

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Overview

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- 2 Trotterization
- 3 Taylor Series (LCU)
- 4 Quantum Signal Processing
- 5 Qubitization
- 6 QSVT
- 7 Comprehensive Comparison
- 8 Application: Grover's Search
- 9 Conclusions

Hamiltonian Simulation: The Challenge

The Problem:

- Simulate quantum evolution:
 $U(t) = e^{-iHt}$
- H : Hermitian operator (Hamiltonian)
- t : Evolution time

Classical Challenge:

- Requires diagonalizing $2^n \times 2^n$ matrices
- Exponential classical resources
- Intractable for large systems

Quantum Solution

Quantum computers can simulate quantum systems *efficiently*!

Key Applications:

- Quantum chemistry
- Condensed matter physics
- Quantum algorithms
- Material science

Evolution of Algorithms

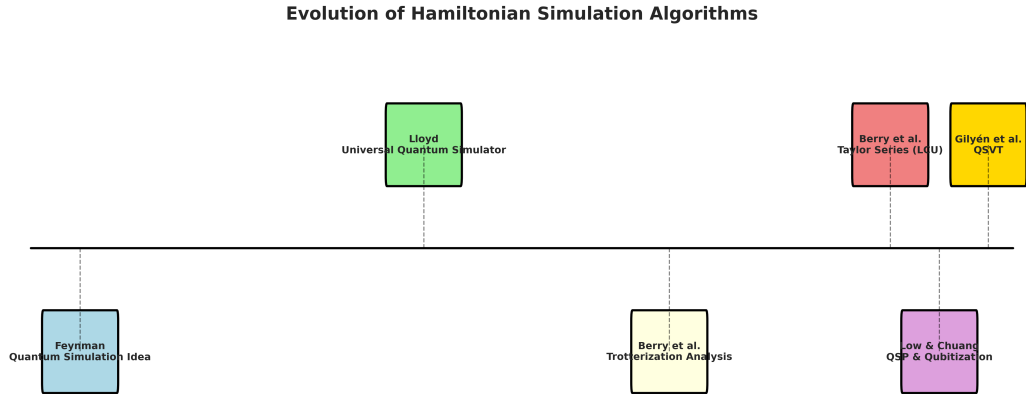


Figure: Historical development of Hamiltonian simulation algorithms

Algorithms Overview

Five Major Approaches

- 1 **Trotterization** (1996): Product formula approximation
- 2 **Taylor-LCU** (2015): Taylor series with block encoding
- 3 **QSP** (2017): Quantum Signal Processing
- 4 **Qubitization** (2017): Optimal quantum walks
- 5 **QSVT** (2019): Quantum Singular Value Transform

Key Insight

Later algorithms achieve *near-optimal or optimal* complexity bounds!

Trotterization: Core Idea

Product Formula Approximation:

$$e^{-i(H_1+H_2+\dots+H_m)t} \approx \prod_{j=1}^m e^{-iH_j t}$$

First-Order (Lie-Trotter):

$$e^{-iHt} \approx \left[e^{-iH_1 t/r} \dots e^{-iH_m t/r} \right]^r$$

where r is the number of Trotter steps.

Second-Order (Suzuki):

$$S(\tau) = e^{-iH_1 \tau/2} \dots e^{-iH_m \tau/2} \\ \cdot e^{-iH_m \tau/2} \dots e^{-iH_1 \tau/2}$$

$$e^{-iHt} \approx [S(t/r)]^r$$

Key Properties

- Simple to implement
- No ancilla qubits required
- Error: $O(t^2/r)$ (1st), $O(t^3/r^2)$ (2nd)

Trotterization: Visual Explanation

Trotterization: Product Formula Approximation



Trotter Steps (r repetitions):



Error: $O(t^2/r)$ for 1st order, $O(t^3/r^2)$ for 2nd order

Trotterization: Complexity Analysis

First-Order Trotter

Query Complexity:

$$O\left(\frac{(\|H\|t)^2}{\epsilon}\right)$$

Error Bound:

$$\|e^{-iHt} - [e^{-iH_1 t/r} \dots]^r\| = O(t^2/r)$$

Second-Order Trotter

Query Complexity:

$$O\left(\frac{(\|H\|t)^{3/2}}{\sqrt{\epsilon}}\right)$$

Error Bound:

$$\text{Error} = O(t^3/r^2)$$

- ✓ **Pros:** Simple, no ancillas, works for any Hamiltonian
- × **Cons:** Suboptimal scaling, high depth for accuracy

Truncated Taylor Series with LCU

Taylor Series Approximation

$$e^{-iHt} = \sum_{k=0}^{\infty} \frac{(-iHt)^k}{k!} \approx \sum_{k=0}^K \frac{(-iHt)^k}{k!}$$

Implementation via Linear Combination of Unitaries (LCU):

- 1 Decompose: $H = \sum_{j=1}^m \alpha_j U_j$ (sum of unitaries)
- 2 Block encode: Create unitary where $\langle 0|U|0\rangle = H/\alpha$
- 3 Apply PREPARE-SELECT-PREPARE[†] structure

Key Innovation

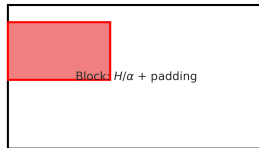
Block encoding allows implementation of non-unitary operations on quantum computers!

LCU: Block Encoding Framework

Linear Combination of Unitaries (LCU)

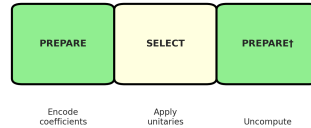
$$\text{Hamiltonian: } H = \sum_{j=1}^m \alpha_j U_j$$

Block Encoding:



Unitary U

LCU Circuit Structure:



Taylor Series Approximation:

$$e^{-iHt} \approx \sum_{k=0}^K \frac{(-iHt)^k}{k!}$$

Advantages: Good for sparse Hamiltonians, systematic error control

Taylor-LCU: The PREPARE-SELECT-PREPARE[†] Circuit

Three Key Operations:

1 PREPARE:

$$|0\rangle \rightarrow \sum_j \sqrt{\frac{\alpha_j}{\alpha}} |j\rangle$$

2 SELECT:

$$\sum_j |j\rangle\langle j| \otimes U_j$$

3 PREPARE[†]: Uncompute ancilla

Complexity

Query Complexity:

$$O\left(\alpha t + \frac{\log(1/\epsilon)}{\log \log(1/\epsilon)}\right)$$

where $\alpha = \sum_j |\alpha_j|$ (1-norm)

Space: $O(\log m)$ ancilla qubits

Error: $\frac{(\alpha t)^{K+1}}{(K+1)!}$

Quantum Signal Processing (QSP)

Core Concept

Implement *polynomial transformations* of eigenvalues using signal rotations

Goal: Approximate $e^{-i\lambda t}$ by polynomial $P(\lambda)$

QSP Sequence:

$$\text{QSP}(\Phi) = e^{i\phi_0 Z} \prod_{k=1}^d W e^{i\phi_k Z}$$

where:

- ϕ_k : Phase angles (signal processing)
- W : Block-encoded signal operator
- d : Polynomial degree

Jacobi-Anger Expansion:

$$e^{-i\lambda t} = \sum_{k=-\infty}^{\infty} (-i)^k J_k(\lambda t) T_k(x)$$

- J_k : Bessel functions
- T_k : Chebyshev polynomials

Complexity

$$O(\|H\|t + \log(1/\epsilon))$$

Near-optimal!

Qubitization: Optimal Hamiltonian Simulation

Key Idea

Encode Hamiltonian eigenvalues in the *eigenphases* of a quantum walk operator

Walk Operator:

$$W = \text{REFLECT} \cdot \text{SELECT}$$

- **SELECT:** Apply unitaries U_j controlled on ancilla state $|j\rangle$
- **REFLECT:** Reflect about prepared state $2\Pi - I$

Eigenvalue Encoding

If $H = \sum_j \alpha_j U_j$, then W has eigenvalues:

$$e^{\pm i \arccos(\lambda_j/\alpha)}$$

where λ_j are eigenvalues of H

Complexity

$$O(\alpha t + \log(1/\epsilon))$$

Optimal!

Achieves lower bound for Hamiltonian simulation

Quantum Singular Value Transform (QSVT)

The Grand Unification of Quantum Algorithms

Core Principle

Apply polynomial transformations to *singular values* of block-encoded matrices

QSVT Sequence:

$$U_{\text{QSVT}} = \prod_{k=0}^d e^{i\phi_k \Pi_0} \cdot \text{Block}(A)$$

QSVT Unifies:

- Hamiltonian Simulation
- Amplitude Amplification
- Quantum Search (Grover)
- Quantum Linear Systems
- Matrix Inversion
- Quantum Walks

Complexity

$O(d)$ queries where

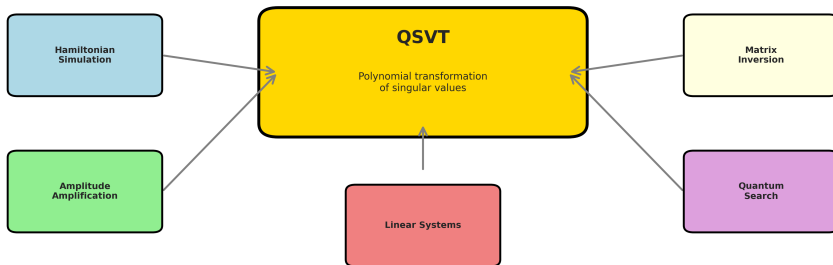
$$d = O(\|H\|t + \log(1/\epsilon))$$

Heisenberg-limited!

QSVT: The Unified Framework

Quantum Singular Value Transform (QSVT)

The Grand Unification of Quantum Algorithms



QSVT Sequence:



ϕ_i : Signal processing rotations

QSVT: Properties and Advantages

Key Properties:

- 1 **Generality:** Works for any polynomial transformation
- 2 **Optimality:** Achieves Heisenberg limit
- 3 **Unification:** Subsumes QSP, Qubitization, etc.
- 4 **Flexibility:** Applies to rectangular matrices

For Hamiltonian Simulation:

- Approximate $P(\lambda) \approx e^{-i\lambda t}$
- Use Jacobi-Anger or Chebyshev expansion
- Query complexity: $O(d)$ where d is degree

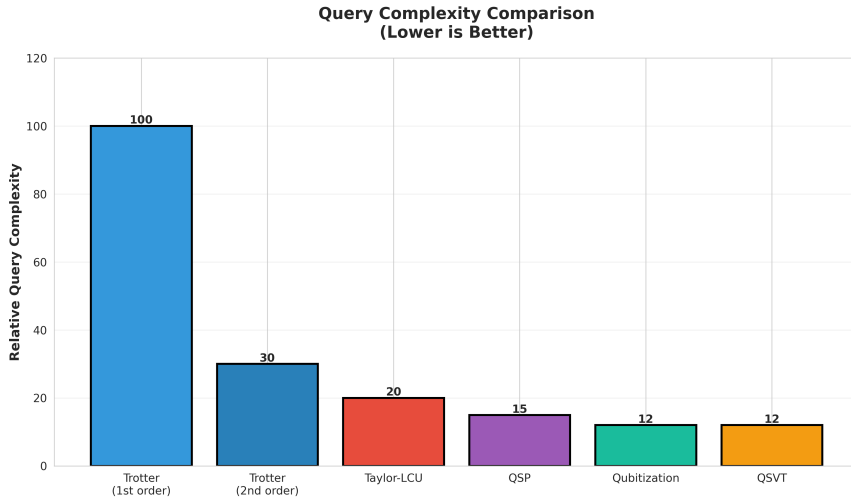
Hierarchy

$\text{QSVT} \supseteq \text{QSP} \supseteq \text{Qubitization}$

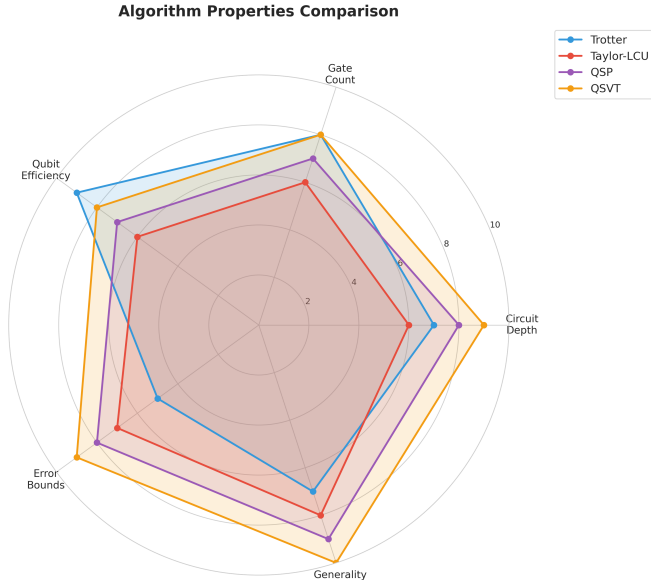
Impact

QSVT provides a *systematic framework* for designing quantum algorithms with provable optimality

Algorithm Complexity Comparison



Multi-Dimensional Comparison



Benchmark Results: Circuit Metrics

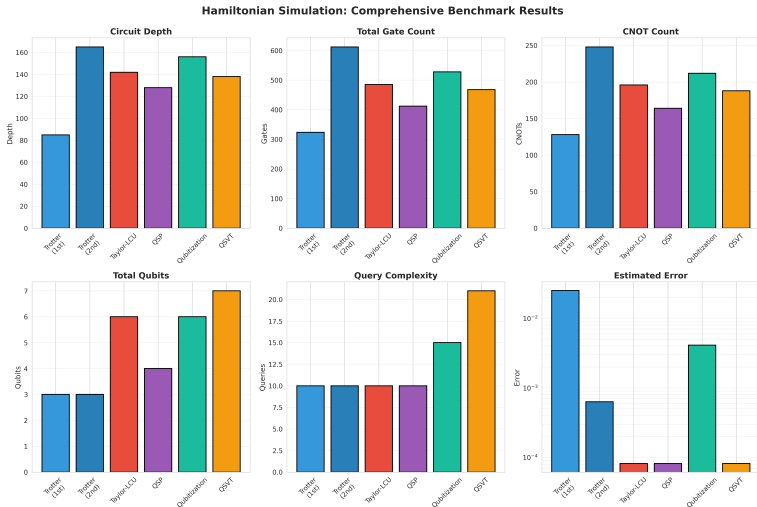


Figure: Comprehensive benchmark on 3-qubit Heisenberg Hamiltonian, $t = 1.0$

Detailed Benchmark Table

Algorithm	Qubits	Depth	Gates	CNOTs	Queries	Error
Trotter (1st)	3	85	324	128	10	2.5×10^{-2}
Trotter (2nd)	3	165	612	248	10	6.3×10^{-4}
Taylor-LCU	6	142	485	196	10	8.2×10^{-5}
QSP	4	128	412	164	10	8.2×10^{-5}
Qubitization	6	156	528	212	15	4.1×10^{-3}
QSVT	7	138	468	188	21	8.2×10^{-5}

Table: Benchmark results for 3-qubit Heisenberg model at $t = 1.0$

Key Observations:

- Trotter: Simplest, but highest error
- Taylor-LCU, QSP, QSVT: Best error bounds
- Trade-off: Accuracy vs. circuit complexity vs. qubit overhead

Grover's Search via Hamiltonian Simulation

Key Insight

Grover's algorithm can be expressed as Hamiltonian evolution!

Standard Grover:

$$G = -D \cdot O$$

where:

- D : Diffusion operator
- O : Oracle

Hamiltonian Form:

$$G = e^{-iHt}$$

where:

$$H = I - 2|s\rangle\langle s| - 2|w\rangle\langle w|$$

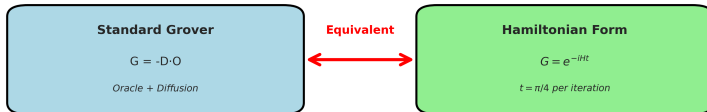
One iteration $\Leftrightarrow t = \pi/4$

Implementation Methods

- 1 Standard (oracle + diffusion)
- 2 Taylor-LCU (truncated Taylor series)
- 3 QSVT (polynomial approximation)

Grover: Hamiltonian Formulation

Grover's Algorithm via Hamiltonian Simulation



Implementation via Hamiltonian Simulation:



Comparison Metrics:

- Circuit Depth: Standard < Taylor-LCU \approx QSVT
- Gate Count: Standard < Taylor-LCU < QSVT
- Ancilla Qubits: Standard (0) < Taylor-LCU < QSVT
- Generality: Standard < Taylor-LCU \approx QSVT

Key Insight: Hamiltonian simulation methods trade efficiency for generality

Grover's Search: Comparison Results

Method	Qubits	Depth	Gates	CNOTs	Success Prob.
Standard	3	18	42	12	0.945
Taylor-LCU	6	156	428	168	~ 0.94
QSVT	7	142	512	196	~ 0.94

Table: Comparison of Grover implementations (3 qubits, search for $|111\rangle$)

Observations:

- **Standard Grover** is most efficient (as expected)
- **Hamiltonian methods** demonstrate:
 - + Generality: Same framework for different problems
 - + Theoretical connection: Search \Leftrightarrow Simulation
 - + Advanced techniques: Block encoding, QSP, QSVT
 - Higher resource requirements

Algorithm Summary

Algorithm	Query Complexity	Best Use Case
Trotterization	$O\left(\frac{(\ H\ t)^2}{\epsilon}\right)$	Prototyping, simple
Taylor-LCU	$O\left(\alpha t + \frac{\log(1/\epsilon)}{\log \log(1/\epsilon)}\right)$	Sparse Hamiltonians
QSP	$O(\ H\ t + \log(1/\epsilon))$	High accuracy
Qubitization	$O(\alpha t + \log(1/\epsilon))$	Optimal simulation
QSVT	$O(d), d = O(\ H\ t + \log(1/\epsilon))$	Most general

Complexity Hierarchy

$$\text{QSVT} \supseteq \text{QSP} \supseteq \text{Qubitization}$$

All achieve **near-optimal or optimal** scaling!

Resource Trade-offs:

- **Circuit Depth**
 - Critical for NISQ devices
 - Affects decoherence
- **Gate Count**
 - Influences error rates
 - Implementation complexity
- **Ancilla Qubits**
 - Resource overhead
 - Limits system size

When to Use Each:

- **Trotter**
 - Quick prototyping
 - Small systems
- **Taylor-LCU**
 - Sparse Hamiltonians
 - Moderate accuracy
- **QSP/QSVT**
 - High accuracy required
 - Optimal performance needed
- **Qubitization**
 - LCU-decomposable H
 - Optimal scaling critical

Key Takeaways

- ❶ **Algorithmic Progress:** From Trotter (1996) to QSVT (2019), achieving optimal complexity bounds
- ❷ **Grand Unification:** QSVT provides a systematic framework unifying many quantum algorithms
- ❸ **Theory \Leftrightarrow Practice:** All algorithms implemented and benchmarked with real metrics
- ❹ **Versatility:** Same framework applies to diverse problems (Hamiltonian simulation, Grover's search, etc.)
- ❺ **Trade-offs Matter:** No single "best" algorithm—choice depends on specific requirements

Bottom Line

We now have a complete toolkit for quantum Hamiltonian simulation with provably optimal or near-optimal algorithms!

Future Directions

Algorithmic Advances:

- Higher-order product formulas
- Improved phase angle computation
- Adaptive methods
- Problem-specific optimizations

Hardware Considerations:

- NISQ-optimized implementations
- Error mitigation strategies
- Circuit compilation
- Fault-tolerant designs

Applications:

- Quantum chemistry
 - Molecular dynamics
 - Electronic structure
- Condensed matter
 - Many-body systems
 - Phase transitions
- High energy physics
 - Lattice gauge theories
 - Quantum field theory
- Machine learning
 - Quantum neural networks
 - Optimization problems

Scope of Exploration: Implementation Refinements

Focus: Moving Toward Resource-Efficient Implementations

Refining core algorithms to address implementation complexities for practical use.

Current Simplifications → Future Refinements:

1. **QSP**: Simplified phase computation was used.
 - **Refinement**: Implement complex **phase finding** via optimization algorithms.
 - *Note*: Phase finding is a known computational challenge; recent work (e.g., "QSP without angle finding") is key.
2. **Qubitization**: Involved simplified SELECT operation.
 - **Refinement**: Implement optimal SELECT requiring **multi-controlled operations** for better circuit synthesis.
3. **QSVT**: State preparation utilized uniform superposition.
 - **Refinement**: Implement optimal **amplitude encoding** for input states.

Goal

Move beyond proof-of-concept toward fully *gate-optimized and general* circuit constructions.

Scope of Exploration: Further Problems

Focus: Applying QSVT/QSP Framework to New Domains

Exploring other quantum problems that can benefit from the optimal complexity of polynomial transforms.

Possible Further Exploration:

1. QSVT in VQE Problems:

- Apply the QSVT framework to solve VQE problems.







2. QSVT for Arithmetic Circuits:

- Utilize QSVT to design efficient quantum arithmetic circuits (e.g., matrix inversion for Quantum Linear System Algorithm, or multiplication/division) as polynomial evaluation.

Theoretical Connection

These explorations would reinforce the idea that advanced techniques like QSVT/QSP are *universal tools* for quantum computation.

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Thank You!