#### 2. Simulation of Random Variables

# Why Monte-Carlo simulation?

See the section 'Why Monte-Carlo simulation?' in sampling\_ex.pdf.

Often we are interested in the distribution of a random variable X which is complicated, but which can none-the-less be built up from simple components such as independent rv's with known distributions.

Monte-Carlo simulation is an excellent tool for such problem: we seek to generate a random sample from the distribution of X, which we can use to estimate its mean, median, mode, percentiles, etc.

The starting point for any simulation is the generation of r.v.s with known distributions (binomial, poisson, exponential, normal, etc.), which are the building blocks for more complicated distributions. It turns out that all random variables can be generated by manipulating U(0,1) rv's.

# Seeding

If the random number generator is initialised (called the seed) before you start generating pseudo-random numbers, then you can reproduce the whole sequence exactly. This is a very good idea from a scientific point of view; being able to repeat an experiment means that your results are verifiable.

If the random number generator is not initialised, then R initialises it using a value taken from the system clock.

To generate n pseudo-random numbers in R, use runif(n). For a given value of seed (assumed integer), the command set.seed(seed) always puts you at the same point on the cycle of pseudo-random numbers.

## Seeding

```
> set.seed(42)
> runif(2)
[1] 0.9148060 0.9370754
> runif(2)
[1] 0.2861395 0.8304476
> set.seed(42)
> runif(4)
[1] 0.9148060 0.9370754 0.2861395 0.8304476
```

## Simulating discrete random variables

## Simulating discrete random variables

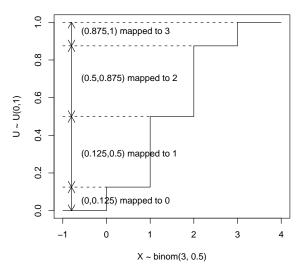
Let X be a discrete random variable taking values in the set  $\{0,1,\ldots\}$  with cdf F and pmf p. The following snippet of code takes a uniform random variable U and returns a discrete random variable X with cdf F.

```
# given U ~ U(0,1)
X <- 0
while (F(X) < U) {
    X <- X + 1
}</pre>
```

When the algorithm terminates we have  $F(X) \geq U$  and F(X-1) < U, that is  $U \in (F(X-1), F(X)]$ . That is,

$$\mathbb{P}(X = x) = \mathbb{P}(U \in (F(x - 1), F(x)]) = F(x) - F(x - 1) = p(x).$$

#### simulating from a binom(3, 0.5) c.d.f.



To simulate binomial, geometric, negative-binomial or Poisson rv's in R, use rbinom, rgeom, rnbinom or rpois.

For simulating other (finite) discrete rv's R provides sample(x, size, replace = FALSE, prob = NULL). The inputs are

x A vector giving the possible values the rv can take;

- size How many rv's to simulate;
- replace Set this to TRUE to generate an iid sample, otherwise the rv's will be conditioned to be different from each other;
  - prob A vector giving the probabilities of the values in x. If omitted then the values in x are assumed to be equally likely.

See the section 'Simulating other (finite) discrete rv using sample function' in sampling\_ex.pdf.

# Simulating continuous random variables

## Simulating continuous random variables

Suppose that we are given  $U \sim U(0,1)$  and want to simulate a continuous rv X with cdf  $F_X$ .

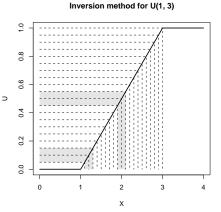
Put  $Y = F_X^{-1}(U)$  then we have

$$F_Y(y) = \mathbb{P}(Y \le y) = \mathbb{P}(F_X^{-1}(U) \le y) = \mathbb{P}(U \le F_X(y)) = F_X(y).$$

That is, Y has the same distribution as X.

Thus, if we can simulate a U(0,1) rv, then we can simulate any continuous rv X for which we know  $F_X^{-1}$ . This is called the *inverse transformation method* or simply the *inversion method*.

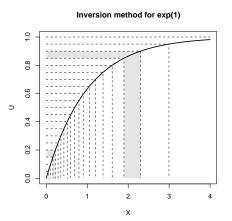
# Simulating $X \sim U(1,3)$



If  $X \sim U(1,3)$  then  $F_X(x) = (x-1)/2$  for  $x \in (1,3)$  and thus  $F_X^{-1}(y) = 2y+1$  for  $y \in (0,1)$ .

Example: see the section 'Simulating X  $\sim$  U(1,3)' in sampling\_ex.pdf.

# Simulating $X \sim \exp(\lambda)$



If 
$$X \sim \exp(\lambda)$$
 then  $F_X(x) = 1 - e^{-\lambda x}$  for  $x \geq 0$  and thus  $F_X^{-1}(y) = -\frac{1}{\lambda}\log{(1-y)}$ .

Example: see the section 'Simulating X  $\sim \exp(1)$ ' in sampling\_ex.pdf.

#### Random variable simulators in R

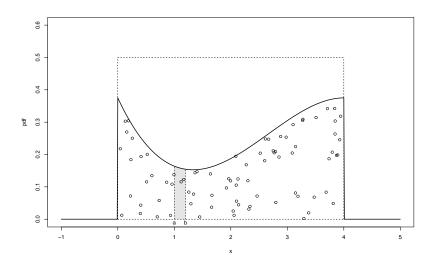
Distribution	R command
binomial	rbinom
Poisson	rpoisson
geometric	rgeom
negative binomial	rnbinom
uniform	runif
exponential	rexp
normal	rnorm
gamma	rgamma
beta	rbeta
student t	rt
F	rf
chi-squared	rchisq
Weibull	rweibull

## The rejection method

## The rejection method

The inversion method works well if we can find  $F^{-1}$  analytically. If not, one method in this situation, which is often faster, is the rejection method.

We start with an example. Suppose that we have a continuous random variable X with pdf  $f_X$  concentrated on the interval (0,4). We imagine 'sprinkling' points  $P_1,P_2,\ldots$ , uniformly at random under the density function, and consider the distribution of  $X_1$ , the x coordinate of  $P_1$ .



Let R be the shaded region under  $f_X$  between a and b, then

$$\begin{split} \mathbb{P}(a < X_1 < b) &= \mathbb{P}(P_1 \text{ hits R}) \\ &= \frac{\text{Area of R}}{\text{Area under density}} \\ &= \frac{\int_a^b f_X(x) dx}{1} \\ &= \int_a^b f_X(x) dx. \end{split}$$

So  $X_1$  has the same distribution as X.

But how do we generate the points  $P_i$  uniformly under  $f_X$ ? The answer is to generate points at random in the rectangle  $[0,4] \times [0,0.5]$ , and then *reject* those that fall above the pdf.

**Rejection method (uniform envelope)** Suppose that  $f_X$  is non-zero only on [a, b], and  $f_X \leq k$ .

- 1. Generate  $X \sim U(a,b)$  and  $Y \sim U(0,k)$  independent of X (so P = (X,Y) is uniformly distributed over the rectangle  $[a,b] \times [0,k]$ ).
- 2. If  $Y < f_X(X)$  then return X, otherwise go back to step 1.

Example: consider the triangular pdf  $f_X$  defined as

$$f_X(x) = \begin{cases} x & \text{if } 0 < x < 1; \\ (2 - x) & \text{if } 1 \le x < 2; \\ 0 & \text{otherwise.} \end{cases}$$

We apply the rejection method as follows: see rejecttriangle.pdf

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# General rejection method

Our rejection method uses a rectangular envelope to cover the target density  $f_X$ . What to do if  $f_X$  is unbounded?

Let X have pdf h and let  $Y \sim U(0, kh(X))$ , then (X, Y) is uniformly distributed under the curve kh:

$$\mathbb{P}((X, Y) \in (x, x + dx) \times (y, y + dy))$$

$$= \mathbb{P}(Y \in (y, y + dy) | X \in (x, x + dx)) \mathbb{P}(X \in (x, x + dx))$$

$$= \frac{dy}{kh(x)} h(x) dx = \frac{1}{k} dx dy.$$

Suppose we wish to simulate from the density  $f_X$ . Let h be a density we can simulate from, and choose k such that

$$k \ge k^* = \sup_x \frac{f_X(x)}{h(x)}.$$

Then kh forms an envelope for  $f_X$ , and we can generate points uniformly within this envelope. By accepting points below the curve  $f_X$ , we get the general rejection method:

#### General rejection method

To simulate from the density  $f_X$ , we assume that we have envelope density h from which you can simulate, and that we have some  $k<\infty$  such that  $\sup_x f_X(x)/h(x) \leq k$ .

- 1. Simulate X from h.
- 2. Generate  $Y \sim U(0, kh(X))$ .
- 3. If  $Y < f_X(X)$  then return X, otherwise go back to step 1.

Suppose that  $f_X^* \propto f_X$ ,  $h^* \propto h$ , and  $f_X^*(x)/h^*(x) \leq 1$ , then the general rejection method is equivalent to:

- 1. Simulate X from h.
- 2. Generate  $Y \sim U(0,1)$ .
- 3. If  $Y < f_X^*(X)/h^*(X)$  then return X, otherwise go back to step 1.

This equivalence is useful when it is hard to get an exact form for  $f_X$  (e.g., it is hard to get a normalisation factor in the posterior distribution).

## Efficiency

The efficiency of the rejection method is measured by the expected number of times you have to generate a candidate point (X, Y).

The area under the curve kh is k and the area under the curve  $f_X$  is 1, so the probability of accepting a candidate is 1/k.

Thus the number of times N we have to generate a candidate point has distribution  $\mathrm{geom}(1/k)$ , with mean

$$\mathbb{E}N = k$$
.

So, the closer h is to  $f_X$ , the smaller we can choose k, and the more efficient the algorithm.

#### Example: gamma

For  $m, \lambda > 0$  the  $\Gamma(\lambda, m)$  density is

$$f(x) = \lambda^m x^{m-1} e^{-\lambda x} / \Gamma(m)$$
, for  $x > 0$ ,

There is no explicit formula for the cdf F or its inverse, so we will use the rejection method to simulate from f.

We will use an exponential envelope  $h(x) = \mu e^{-\mu x}$ , for x>0. Using the inversion method we can easily simulate from h using  $-\log(U)/\mu$ , where  $U\sim U(0,1)$ .

To envelop f we need to find

$$k^* = \sup_{x>0} \frac{f(x)}{h(x)} = \sup_{x>0} \frac{\lambda^m x^{m-1} e^{(\mu-\lambda)x}}{\mu\Gamma(m)}.$$

Clearly  $k^*$  will be infinite if m<1 or  $\lambda\leq\mu$ . For m=1 the gamma is just an exponential. Thus we will assume m>1 and choose  $\mu<\lambda$ .

For  $m \in (0,1)$  the rejection method can still be used, but a different envelope is required.

To find  $k^*$  we take the derivative of the right-hand side above and set it to zero, to find the point where the maximum occurs. You can check that this is at the point  $x=(m-1)/(\lambda-\mu)$ , which gives

$$k^* = \frac{\lambda^m (m-1)^{m-1} e^{-(m-1)}}{\mu(\lambda - \mu)^{m-1} \Gamma(m)}.$$

To improve efficiency we would like to choose our envelope to make  $k^*$  as small as possible. Looking at the formula for  $k^*$  this means choosing  $\mu$  to make  $\mu(\lambda-\mu)^{m-1}$  as large as possible. Setting the derivative with respect to  $\mu$  to zero, we see that the maximum occurs when  $\mu=\lambda/m$ . Plugging this back in we get  $k^*=m^me^{-(m-1)}/\Gamma(m)$ .

We can now code up our rejection algorithm: see gamma\_sim.pdf

# The Gibbs sampler

Notes by Guoqi Qian

The Gibbs sampler is a random vector generation method which does not require the complete information of the target multivariate probability distribution. Instead, it only requires the information of a set of the associated conditional distributions. Gibbs sampler will be computationally feasible if the conditional distributions are ready to be simulated.

- ▶ Suppose we want to generate a vector  $\mathbf{u} = (u_1, \cdots, u_K)$  from the random vector  $\mathbf{U} = (U_1, \cdots, U_K)$  which has a joint pdf  $f(\mathbf{u})$ .
- ▶ Suppose the pdf  $f(\mathbf{u})$  has a very complicated form. But for each  $k=1,\cdots,K$ , the conditional pdf of  $U_k$  given  $\mathbf{U}_{-k}=(U_1,\cdots,U_{k-1},U_{k+1},\cdots,U_K)$  is known and relatively easy to simulate. We use

$$f(u_k|\mathbf{u}_{-k}) \equiv f(u_k|u_1,\dots,u_{k-1},u_{k+1},\dots,u_K)$$

to denote such a conditional pdf.

Then the Gibbs sampler used to generate one observation of U can be described as following:



#### **Algorithm A.1** [Gibbs sampler]

- $1^{\circ}$  Arbitrarily generate/assign an initial vector  $\mathbf{u}^{(0)} = (u_1^{(0)}, \cdots, u_K^{(0)}) \text{ from the support of } f(\mathbf{u}).$
- 2° Generate a value  $u_1^{(1)}$  from  $f(u_1|u_2=u_2^{(0)},\cdots,u_K=u_K^{(0)});$  then generate a value  $u_2^{(1)}$  from  $f(u_2|u_1=u_1^{(1)},u_3=u_3^{(0)},\cdots,u_K=u_K^{(0)});$  continue until generate a value  $u_{K-1}^{(1)}$  from  $f(u_{K-1}|u_1=u_1^{(1)},\cdots,u_{K-2}=u_{K-2}^{(1)},u_K=u_K^{(0)}),$  and generate a value  $u_K^{(1)}$  from  $f(u_K|u_1=u_1^{(1)},\cdots,u_{K-1}=u_{K-1}^{(1)}).$  The generated values are delivered as  $\mathbf{u}^{(1)}=(u_1^{(1)},\cdots,u_K^{(1)}).$

#### **Algorithm A.1** [Gibbs sampler (continued)]

3° For  $j=1,2,\cdots$ , do the following: Generate a value  $u_1^{(j)}$  from  $f(u_1|u_2=u_2^{(j-1)},\cdots,u_K=u_K^{(j-1)})$ ; then generate a value  $u_2^{(j)}$  from  $f(u_2|u_1=u_1^{(j)},u_3=u_3^{(j-1)},\cdots,u_K=u_K^{(j-1)})$ ; continue until generate a value  $u_{K-1}^{(j)}$  from  $f(u_{K-1}|u_1=u_1^{(j)},\cdots,u_{K-2}=u_{K-2}^{(j)},u_K=u_K^{(j-1)})$ , and generate a value  $u_K^{(j)}$  from  $f(u_K|u_1=u_1^{(j)},\cdots,u_{K-1}=u_{K-1}^{(j)})$ .

The generated values are delivered as  $\mathbf{u}^{(j)} = (u_1^{(j)}, \cdots, u_K^{(j)}).$ 

Note the conditioning is always based on the latest values of  $(u_1, \cdots, u_K)$ .

## Remarks on Gibbs sampler

Using theory of Markov chain, it can be shown that

$$\mathbf{u}^{(j)} \stackrel{d}{\to} f(\mathbf{u}) \quad \text{as } j \to \infty$$

under fairly general conditions. Details are not pursued here.

- ▶ This implies that  $\mathbf{u}^{(j)}$  can be roughly regarded as an observation of  $\mathbf{U}$  from pdf  $f(\mathbf{u})$  when j is sufficiently large. Empirical methods are available (details not pursued here) to determine how large j should be.
- In practice, we usually generate a long sequence  $\mathbf{u}^{(1)}, \cdots, \mathbf{u}^{(m)}, \mathbf{u}^{(m+1)}, \cdots, \mathbf{u}^{(m+J)}$  using the Gibbs sampler; then ignore the **burn-in** sequence  $\mathbf{u}^{(1)}, \cdots, \mathbf{u}^{(m)}$  and use only the part  $\mathbf{u}^{(m+1)}, \cdots, \mathbf{u}^{(m+J)}$  as a random sample from  $f(\mathbf{u})$ .
- $\mathbf{u}^{(m+1)}, \cdots, \mathbf{u}^{(m+J)}$  are not independent of each other but constitute a Markov chain.

Eggs of certain species of insects hatch under appropriate conditions. Suppose there are N eggs of this species of insects in a particular area, and X eggs of them will hatch. Also let p be the probability that such an egg will hatch.

- 1. A Bayesian approach may be used to model (X,p,N), in which one can reasonably assume  $N \stackrel{\mathrm{d}}{=} \mathsf{Poisson}(\lambda)$ ,  $p \stackrel{\mathrm{d}}{=} \mathsf{Beta}(a,b)$  and  $(X|p,N) \stackrel{\mathrm{d}}{=} \mathsf{binomial}(N,p)$ . Suppose  $N \stackrel{\mathrm{d}}{=} \mathsf{Poi}(16)$  and  $p \stackrel{\mathrm{d}}{=} \mathsf{Beta}(2,4)$  based on a previous study on these insects. Find the joint pdf of (X,p,N) and the marginal pdf of X.
- 2. Derive a Gibbs sampler for generating (X, p, N) and X.
- 3. Implement the algorithm in R and check its performance.



When  $N \stackrel{\mathrm{d}}{=} \mathsf{Poi}(16)$  and  $p \stackrel{\mathrm{d}}{=} \mathsf{Beta}(2,4)$ , the joint pdf of (X,p,N) is

$$\begin{split} f(x,p,N) &= f(x|p,N) \cdot f(p) \cdot f(N) \\ &= \binom{N}{x} p^x (1-p)^{N-x} \cdot \frac{\Gamma(a+b)}{\Gamma(a)\Gamma(b)} p^{a-1} (1-p)^{b-1} \cdot \frac{\lambda^N}{N!} e^{-\lambda} \\ &= 20 e^{-16} \binom{N}{x} \frac{16^N}{N!} p^{x+1} (1-p)^{N-x+3}; \\ &\quad x = 0,1,\cdots,N; \ 0 \leq p \leq 1; \ N = 0,1,2,\cdots. \end{split}$$

The marginal pdf of X is

$$f(x) = \sum_{N=x}^{\infty} \int_{0}^{1} f(x, p, N) dp = \sum_{N=x}^{\infty} \frac{20e^{-16}(x+1)(N-x+3)(N-x+2)(N-x+1)16^{N}}{(N+5)!}$$

which is difficult to be further simplified.

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But it is not difficult to find the full conditional pdf's:

$$(X|p,N) \stackrel{\mathrm{d}}{=} \mathrm{bin}(N,p), \qquad (p|x,N) \stackrel{\mathrm{d}}{=} \mathrm{Beta}(x+2,N-x+4),$$
 and  $(N-x|x,p) \stackrel{\mathrm{d}}{=} \mathrm{Poi}(16(1-p)).$ 

#### Proof:

$$f(p|x,N) = \frac{f(x,p,N)}{f(x,N)} = \frac{20e^{-16}\binom{N}{x}\frac{16^N}{N!}p^{x+1}(1-p)^{N-x+3}}{\int_0^1 20e^{-16}\binom{N}{x}\frac{16^N}{N!}p^{x+1}(1-p)^{N-x+3}dp}$$
$$= \frac{p^{x+1}(1-p)^{N-x+3}}{\int_0^1 p^{x+1}(1-p)^{N-x+3}dp} = \frac{\Gamma(N+6)}{\Gamma(x+2)\Gamma(N-x+4)}p^{x+1}(1-p)^{N-x+3}$$

which is a pdf of Beta(x+2, N-x+4).

$$f(N|x,p) = \frac{f(x,p,N)}{f(x,p)} = \frac{20e^{-16} \binom{N}{x} \frac{16^N}{N!} p^{x+1} (1-p)^{N-x+3}}{\sum_{N=x}^{\infty} 20e^{-16} \binom{N}{x} \frac{16^N}{N!} p^{x+1} (1-p)^{N-x+3}}$$

$$= \frac{\binom{N}{x} [16(1-p)]^N \frac{1}{N!}}{\sum_{\tilde{N}=x}^{\infty} \binom{\tilde{N}}{x} \frac{16^{\tilde{N}}}{\tilde{N}!} (1-p)^{\tilde{N}}} = \frac{\frac{1}{(N-x)!} [16(1-p)]^N}{\sum_{\tilde{N}=x}^{\infty} \frac{1}{(\tilde{N}-x)!} [16(1-p)]^{\tilde{N}}}$$

$$= \frac{[16(1-p)]^{N-x}}{(N-x)!} e^{-16(1-p)}; \quad N = x, x+1, x+2, \cdots.$$

This implies that  $(N - x | x, p) \stackrel{\text{d}}{=} \mathsf{Poisson}(16(1 - p)).$ 

Now a sample of (X,p,N) can be generated using the following Gibbs sampler algorithm

#### **Algorithm A.2** [Gibbs sampler algorithm for (X, p, N)]

- 1° Arbitrarily choose an initial vector, e.g.  $(x^{(0)}, p^{(0)}, N^{(0)}) = (8, 0.5, 16).$
- $2^{\rm o}~(x^{(1)},p^{(1)},N^{(1)})$  is obtained by: generating  $x^{(1)}$  from  ${\rm Bin}(N^{(0)},p^{(0)});$  generating  $p^{(1)}$  from  ${\rm Beta}(x^{(1)}+2,N^{(0)}-x^{(1)}+4);$  and generating an  $N^{(1)}$  from  ${\rm Poi}(16(1-p^{(1)}))+x^{(1)}.$
- 3°  $(x^{(j)}, p^{(j)}, N^{(j)})$ ,  $j = 1, 2, \cdots$ , is obtained by: generating  $x^{(j)}$  from  $\text{Bin}(N^{(j-1)}, p^{(j-1)})$ ; generating  $p^{(j)}$  from  $\text{Beta}(x^{(j)} + 2, N^{(j-1)} x^{(j)} + 4)$ ; and generating an  $N^{(j)}$  from  $\text{Poi}(16(1 p^{(j)})) + x^{(j)}$ .

Once a sample of (X, p, N) is obtained, a sample of X can be read off from it.

The above algorithm was implemented in R. Then a sample of 1000 (X,p,N) samples are generated, and the simulated marginal pdf's of X, p and N are obtained. The results are shown in simulation\_Gibbs.pdf.