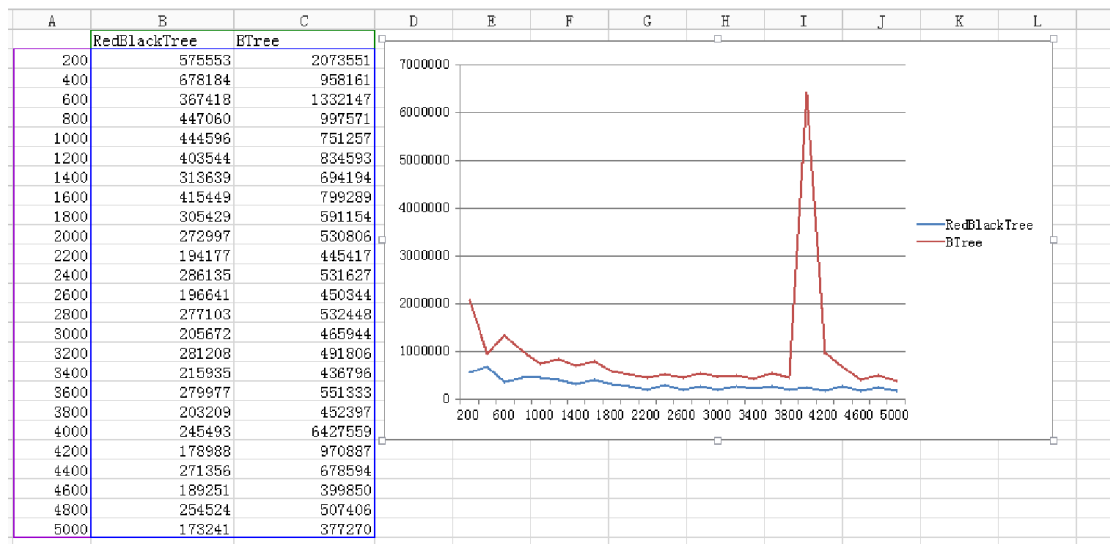


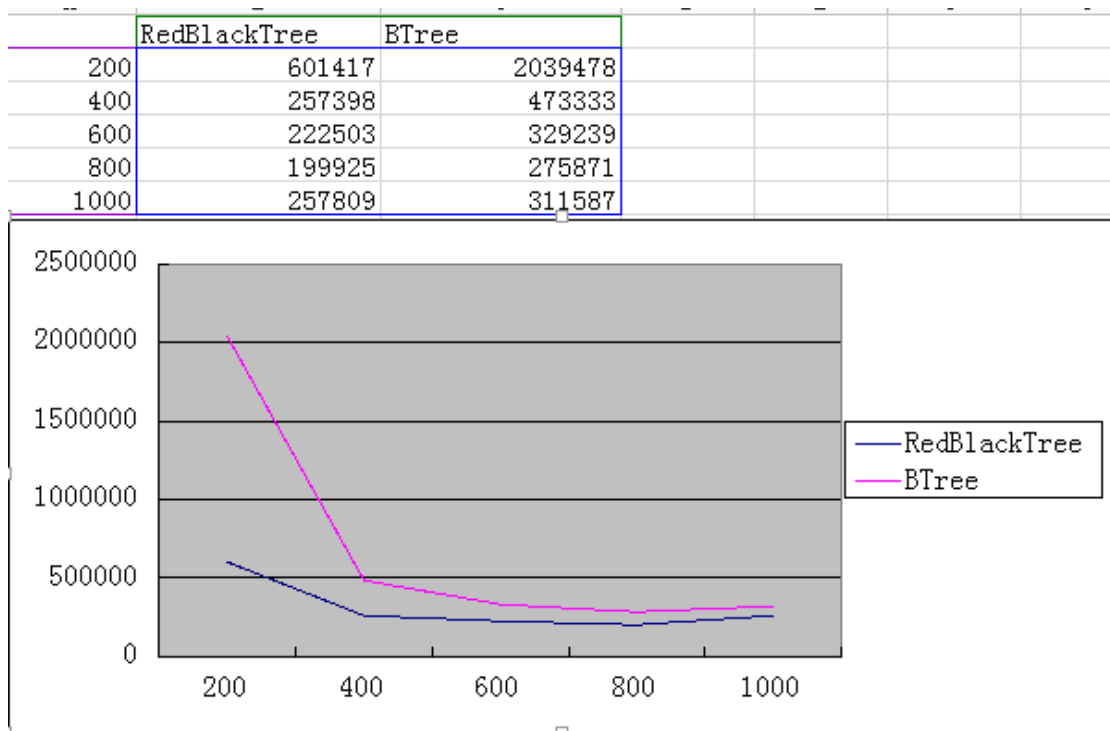
# Analysis work

Following pictures are my test for running time.

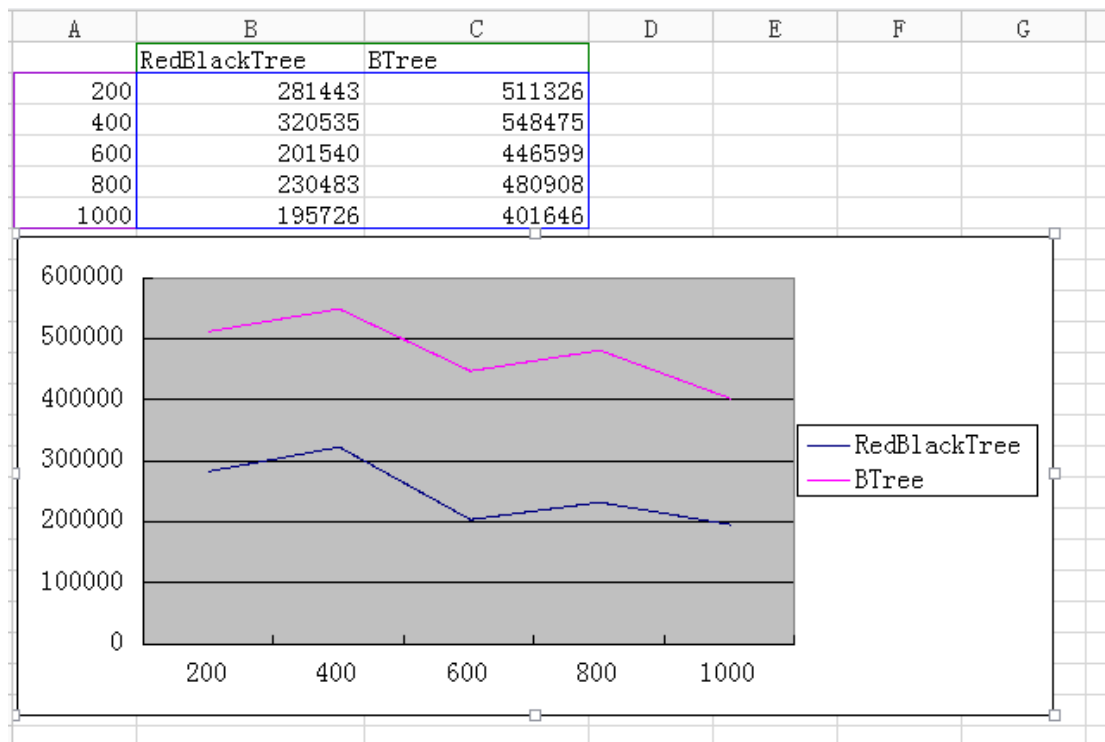
## Initial



## Insert



## Delete



## Analysis

1). At first, we can find that Red-Black-Tree runs faster than B-Tree. Because every node in B-Tree is too large, we costs too much time on copying arrays and search the location of key value in one node. We need at most  $2t$  times for copying and  $2t$  times for searching in one node. However, our operations on Red-Black-Tree, they just need change the pointer. Every change costs a constant time. We need at most  $\lg N$  times changes for every fix up. For every insert, Red-Black-Tree costs at most  $2\lg N$ , B-Tree costs at most  $2t \cdot (\log_{2t}(N))$ . The running time is decided to  $t$  and  $N$ . In this program, Red-Black-Tree runs

faster. And in fact, if we really use B-Tree, there must need disk read or write. So, it will be slower.

2). Besides, we find that the two kinds of tree runs low at start and runs faster and stable at end when initial and insert. Because, we insert a lot of data, the tree is balance and the height of the tree will not change too much. So, we just need search and insert directly. There is nearly no cost for fix up. So, it will be fast and stable.

3). As for delete, at start, the tree is stable. So, we don't need too much cost for fix up. The running time is decided on  $N$ . In the end,  $N$  reduce too much. So, the running time reduce. And there are still some peak. They are just the big fix up.