

## Problems: Week 2

- \*\* 1. Prove by induction on  $M \in \Lambda$  that  $N \in \text{sub}(M)$  implies  $N \sqsubseteq M$ .  
Hint:  $N \in S_1 \cup S_2 \cup S_3$  can be understood as a *disjunction* of three possibilities for sets that contain  $N$ .

- \*\* 2. Prove, by induction on  $M \in \Lambda$ , that if  $x \neq y$  and  $x \notin \text{FV}(Q)$  then:

$$M[P/x][Q/y] = M[Q/y][P[Q/y]/x]$$

Hint: in the variable case, there are three subcases to consider.

- \* 3. Perform the following substitutions:  
a)  $(\lambda x.(\lambda y.xz)z)[(\lambda z.z)/z]$   
b)  $(\lambda x.yx)[yz/x]$   
c)  $(\lambda y.xy)[yx/x]$
- \* 4. Perform one step of reduction for each of the following terms:  
a)  $(\lambda xy.x)(\lambda x.x)(\lambda z.z)$   
b)  $(\lambda xyz.xz(yz))(\lambda xy.x)$   
c)  $(\lambda x.xx)(\lambda x.xx)$   
d)  $(\lambda xy.x)((\lambda z.zz)(\lambda x.x))$
- \* 5. Draw the reduction graph of the term  $(\lambda xy.yy)((\lambda z.zz)(\lambda z.zz))(\lambda x.x)$ .
- \*\* 6. Give an example of a *closed* term  $M$  for each of the following properties:  
a)  $M$  is in  $\beta$ -normal form.  
b)  $M$  has exactly one reduct.  
c)  $M$  has 2 or more reducts.  
d)  $M$  contains fewer redexes than one of its reducts.  
e) A reduct of  $M$  contains a redex that did not occur anywhere in  $M$ .
- \*\* 7. a) State the induction principle associated with the inductively defined set  $\rightarrow_\beta$ , which is a set of pairs of  $\lambda$ -terms.  
b) Prove, by induction on  $M \rightarrow_\beta N$  that:  
if  $M \rightarrow_\beta N$  then  $\text{FV}(N) \subseteq \text{FV}(M)$ .  
You may omit one of the cases (AppL) or (AppR) because it is similar to the other.  
c) Give an example of a step  $M \rightarrow_\beta N$  where  $\text{FV}(M) \neq \text{FV}(N)$ .