

Data structures.

Array.

- continuous, fixed index, can do random Access,
- memory allocated
- o static - fixed \rightarrow can't resize

dynamic

- 排序

complexity ★

o insert

o delete

o backup

= algorithm

$O(n)$, $O(n \log n)$...

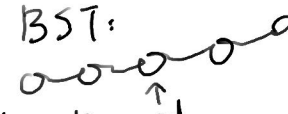
- B = random Access $O(1)$

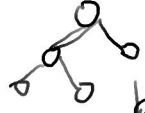
Random Access: $O(1)$

sequential ^{顺序的} search = $O(n)$

\Rightarrow after sorting = $O(\log n)$

ex - binary search (rule)

BST:
sorted: 
(skewed \rightarrow bad search) mid


balanced
(good search)
efficient (adj)

- P = 空间利用率 (在不確定資料量的情況下) 低

\Rightarrow hash table (use array) \Rightarrow if function design badly,
there may be lot of
space don't be used.

linked list

- Node & pointer
- traverse to find data = $O(n)$
- Random Access : not exist.
- sequential search : $O(n)$

Stack / Queue

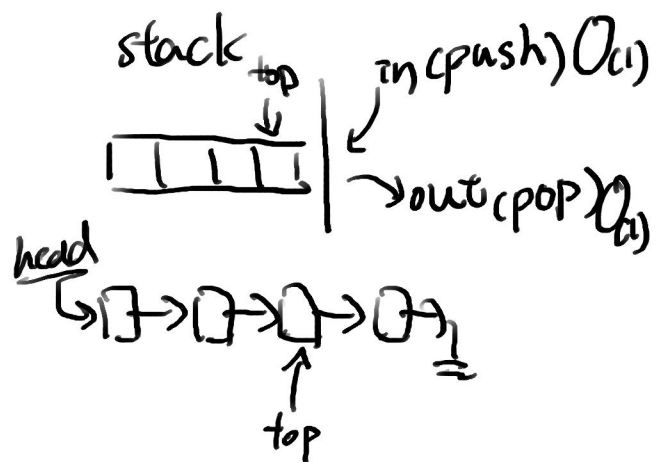
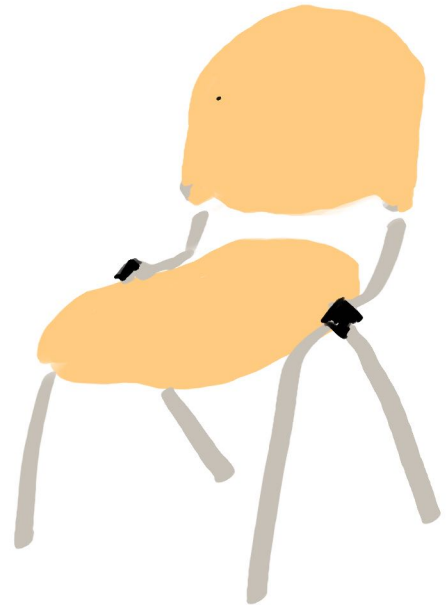
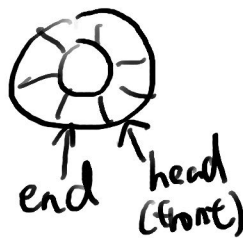
↓
FIFO :

↳ FILO : function call.

implement :

Array

↳ linked list



hash

Array + linked list

bucket



→ travers bad (in linked list)

- collision

- Open Addressing



no \rightarrow (can use random Access)

- linear probing

primary clustering

(某個地方特別擠)

- quadratic probing

secondary clustering

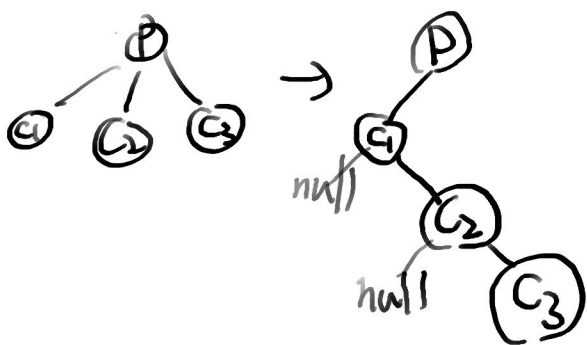
- Double hashing

less clustering

Tree

general \rightarrow binary tree
 \uparrow degree ≤ 2

* left child, right sibling



binary tree \rightarrow BST

\hookrightarrow complete binary tree

\hookrightarrow Heap OAD $\begin{matrix} \nearrow \text{max} \\ \searrow \text{min} \end{matrix}$
 \hookrightarrow priority queue

traversal: $O(\log n)$

Graph

$G(V, E)$

Adjacet. (2)