Maths for Project Euler

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Abstract

Project Euler is a website that provides hundreds of problems that require programming to solve, and some of them are so difficult that there is no way to independently come up with the maths used for the problems. That's why I like to cheat a bit and read what maths are used in other people's solutions.

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1 Number theory

1.1 Pythagorean triple

(Related problem: Q9, Q39, Q75, Q86)

Definition 1.1.1. A **Pythagorean triple** consists of three positive integers a, b, c such that

$$a^2 + b^2 = c^2$$

For example, $3^2 + 4^2 = 5^2$, so (3, 4, 5) is a Pythagorean triple.

Definition 1.1.2. A **primitive Pythagorean triple** is a Pythagorean triple in which a, b, c are coprime. (In other words, gcd(a, b, c) = 1.)

For example, (3,4,5) is a primitive Pythagorean triple, but (6,8,10) is not.

1.1.1 Generating a triple

Theorem 1.1.1. (Euclid's formula) Given a pair of integers m , n with m>n>0 , the formula

$$a = m^2 - n^2$$
, $b = 2mn$, $c = m^2 + n^2$

form a Pythagorean triple. The triples generated is primitive if and only if m and n are coprime and one of them is even.

Moreoever, every primitive Pythagorean triple can be expressed by Euclid's formula . The n, m are unque for a given triple.

Proof. [1] First, note that if $a = m^2 - n^2$, b = 2mn, $c = m^2 + n^2$, then

$$a^{2} + b^{2} = (m^{2} - n^{2}) + (2mn)^{2}$$

$$= m^{4} - 2m^{2}n^{2} + n^{4} + 4m^{2}n^{2}$$

$$= m^{4} + 2m^{2}n^{2} + n^{4}$$

$$= (m^{2} + n^{2})^{2}$$

$$= c^{2}$$

Thus a, b, c form a Pythagorean triple.

Now we prove that every primitive Pythagorean triple can be expressed by Euclid's formula.

All such primitive triples can be written as (a, b, c) where $a^2 + b^2 = c^2$ and a, b, c are coprime. Thus, a, b, c are pairwise coprime (which means gcd(a, b) = gcd(a, c) = gcd(b, c) = 1).

To see why they are pairwise coprime, suppose two of the numbers are not coprime, say a,b, which have a common factor k>1. Let a=kr and b=ks. Then

$$a^{2} + b^{2} = c^{2}$$
$$(kr)^{2} + (ks)^{2} = c^{2}$$
$$k^{2}(r^{2} + s^{2}) = c^{2}$$
$$k\sqrt{r^{2} + s^{2}} = c$$

This means k is also a factor of c, so a, b, c is not a primitive triple, and we have a contradiction. Thus a, b, c must be pairwise coprime.

Now back to business. As a and b are coprime, at least one of them is odd, so we may suppose a is odd (if not, just let a and b switch place). Then b must be even, and c must be odd.

To see why, suppose b is odd. Then c is even, and we can let a=2k+1, b=2l+1 for some non-negative integer k,l. Then

$$a^{2} + b^{2} = (2k+1)^{2} + (2l+1)^{2}$$
$$= 4k^{2} + 4k + 1 + 4l^{2} + 4l + 1$$
$$= 4(k^{2} + k + l^{2} + l) + 2$$

We see that a^2+b^2 is congruent to 2 modulo 4, but c^2 must be a multiple of 4 (since c is even), thus it is impossible that $a^2+b^2=c^2$, and there is a contradiction.

This implies that b must be even, so c is odd. (odd + even = odd)

From
$$a^2 + b^2 = c^2$$
 we obtain $c^2 - a^2 = b^2$ and hence $(c - a)(c + a) = b^2$. Then $\frac{(c + a)}{b} = \frac{b}{(c - a)}$

Since $\frac{(c+a)}{b}$ is rational, we set it equal to $\frac{m}{n}$ in lowest terms. (In other words, $\frac{m}{n}$ is a reduced fraction and $\gcd(m,n)=1$. There must uniquely exist such m and n since every fraction has a reduced form.) Thus $\frac{(c-a)}{b}=\frac{n}{m}$, being the reciprocal of $\frac{(c+a)}{b}$. Then solving

$$\frac{c}{b} + \frac{a}{b} = \frac{m}{n}, \qquad \frac{c}{b} - \frac{a}{b} = \frac{n}{m}$$

for $\frac{c}{b}$ and $\frac{a}{b}$ gives

$$\frac{c}{b} = \frac{1}{2}(\frac{m}{n} + \frac{n}{m}) = \frac{m^2 + n^2}{2mn}, \qquad \frac{a}{b} = \frac{1}{2}(\frac{m}{n} - \frac{n}{m}) = \frac{m^2 - n^2}{2mn}$$

As $\frac{m}{n}$ is fully reduced, m and n are coprime, and they cannot be both even. Note that they also cannot be both odd.

To see why, suppose m and n are both odd. From $\frac{a}{b} = \frac{m^2 - n^2}{2mn}$ we obtain

$$a(2mn) = b(m^2 - n^2)$$

Note that $m^2 - n^2$ is a multiple of 4 (because an odd square is congruent to 1 modulo 4), so LHS is also a multiple of 4. Note that mn is odd since both m and n is odd. This means a must be even to make LHS a multiple of 4. But this contradicts the earlier assumption that a is odd.

Thus one of m and n must be odd and the other is even, and the numerators of the two fractions $\frac{m^2+n^2}{2mn}$ and $\frac{m^2-n^2}{2mn}$ are odd (odd² \pm even² = odd \pm even = odd).

Then we claim that these fractions are fully reduced.

To see why, suppose they are not fully reduced. Then both the numerator and denominator are divisible by some prime p. Note that $p \neq 2$ since numerator is odd.

Since p divides the denominator 2mn, it divides m or n (by Euclid's lemma), but it cannot divide both since m and n are coprime. Thus p can only divide either m or n. Thus, p does not divide the numerator $m^2 \pm n^2$, and there is a contradiction.

divide the numerator $m^2\pm n^2$, and there is a contradiction. To see why p cannot divide $m^2\pm n^2$, suppose (WLOG) p divides m and p divides $m^2\pm n^2$. we can let m=rp and $m^2\pm n^2=sp$.

So $(rp)^2 \pm n^2 = sp$ which means $n^2 = \pm p(s - r^2p)$. Since p divides n^2 , p also divides n (by Euclid's lemma).

Thus it is impossible that p divides m and p divides $m^2 \pm n^2$ but p does not divide n .

Thus, when p can only divide either m or n, p doe not divide $m^2 \pm n^2$.

Since there is a contradiction, no such p exists, and the fractions $\frac{m^2 \pm n^2}{2mn}$ are reduced fractions. Recall that we have

$$\frac{c}{b} = \frac{m^2 + n^2}{2mn}, \qquad \frac{a}{b} = \frac{m^2 - n^2}{2mn}$$

And since $\frac{a}{b}$ and $\frac{c}{b}$ are also reduced fractions (as a, b, c are pairwise coprime as shown in the beginning), we can equate numerators with numerators and denominators with denominators, giving Euclid's formula

$$a = m^2 - n^2$$
, $b = 2mn$, $c = m^2 + n^2$

So we have shown that every primitive triple has some m, n with m > n > 0 satisfying Euclid's formula.

Now we prove that for a given primitive triple, the m, n satisfying Euclid's formula are unique.

Suppose there are m, n and m', n' satisfying Euclid's formula, namely,

$$a = m^2 - n^2 = m'^2 - n'^2$$
, $b = 2mn = 2m'n'$, $c = m^2 + n^2 = m'^2 + n'^2$

Then
$$\frac{c+a}{b} = \frac{(m^2+n^2)+(m^2-n^2)}{2mn} = \frac{m}{n}$$
. Similarly, $\frac{c+a}{b} = \frac{m'}{n'}$. Thus, $\frac{m}{n} = \frac{m'}{n'}$. Suppose $m \neq m'$ and $n \neq n'$.

Since $\frac{m}{n}$ is a reduced fraction, $\frac{m'}{n'}$ must be a non-reduced fraction with the same value, which means m' = km and n' = kn for some k > 1.

To see why, note that we have mn' = nm', so m divides nm'. Since m and n are coprime, we have m divides m' by Euclid's lemma (or called 'property of divisibility inheritence' in number thoery subsection of Toddler Probability). This means m'=km for some k . Since $m\neq m'$, we have k > 1. Similarly we have n' = kn for some k > 1.

But then
$$a = m'^2 - n'^2 = (km)^2 - (kn)^2 = k^2(m^2 - n^2)$$
, and $b = 2m'n' = 2(km)(kn) = k^2(2mn)$.

We see that a and b share a common divisor $k^2 > 1$, but this is impossible since a, b, c are primitive triple, so a, b should be coprime. So there is a contradiction.

Thus, it can only be the case that m = m' and n = n'.

 (\Rightarrow) Now we show that given some m,n with m>n>0, if the triple generated by m,n is primitive, then m and n are coprime and one of them is even.

If the triple is primitive, then exactly one of a and b is odd, and c is even, as shown above.

Note that b = 2mn, so b must be even. That means $a = m^2 - n^2$ must be odd.

If m and n are both even or both odd, then a will be even (odd - odd = even, even - even =even), which cannot be true. Thus it can only be the case that exactly one of m and n is even.

To show that m and n are coprime, suppose they are not coprime. Then they have a common divisor k > 1 . Let m = kr and n = ks . Then

$$a = m^{2} - n^{2} = (kr)^{2} - (ks)^{2} = k^{2}(r^{2} - s^{2})$$

$$b = 2mn = 2(kr)(ks) = k^{2}(2rs)$$

$$c = m^{2} + n^{2} = (kr)^{2} + (ks)^{2} = k^{2}(r^{2} + s^{2})$$

This means a, b and c have a common divisor $k^2 > 1$, so the triple (a, b, c) is not primitive, which is a contradiction.

 (\Leftarrow) Now we show that given some m, n with m > n > 0, if m and n are coprime and one of them is even, then the triple generated by m, n is primitive.

If m and n are coprime and (exactly) one of them is even, then first, note that $\frac{c}{b} = \frac{m^2 + n^2}{2mn}$.

We repeat the argument as before to show that $\frac{m^2 + n^2}{2mn}$ is a reduced fraction, so $\frac{c}{b}$ is also reduced fraction (because $c = m^2 + n^2$ and b = 2mn), which means c, b are coprime, so a, b, c must be primitive triple. (because if a, b, c is not primitive, then c, b cannot be coprime.)

Continued fractions 1.2

(Related: Q64, Q65)

(Most of this section and the next two subsection are copied from [2].)

Definition 1.2.1. A simple continued fraction is of the form

$$a_{0} + \frac{1}{a_{1} + \frac{1}{a_{2} + \frac{1}{a_{3} + \frac{1}{\ddots}}}}$$

$$(1)$$

where a_n are integers, and with the exception of a_0 , they are all positive.

For convenience, we write $[a_0; a_1, a_2, a_3, \ldots]$ to represent the continued fraction.

(In this article, when we say continued fraction, we refer to simple continued fraction, where the numerators are all 1.)

The a_n are referred to as quotients of the simple continued fraction.

If the sequence a_n is finite, the continued fraction is called finite continued fraction:

nite, the continued fraction is called finite continued
$$[a_0;a_1,a_2,\dots,a_n]=a_0+\cfrac{1}{a_1+\cfrac{1}{a_2+\cfrac{1}{\ddots}+\cfrac{1}{a_n}}}$$

If the sequence a_n is infinite, the continued fraction is called infinite continued fraction, and it can be defined as

$$[a_0; a_1, a_2, a_3, \ldots] = \lim_{n \to \infty} [a_0; a_1, a_2, \ldots, a_n]$$

1.2.1 Floor and fractional part of a number

The **floor** of a number x returns the greatest integer that is no greater than the number itself, and

The fractional part of a number x is defined as the number minus its floor, and is denoted $\{x\}$ That is

$$\{x\} = x - |x| \tag{2}$$

Note that

$$0 \le \{x\} < 1 \tag{3}$$

, and (for positive numbers) $\{x\}$ is just the part of x after the decimal point.

For example,
$$\{5\} = 0$$
, $\{2.48\} = 0.48$, $\{-8.3\} = 0.7$.

Note that dividing a number up into its floor and its fractional part is the only way that we can divide it into two parts with one part an integer and the other part equal or greater than 0 and still less than 1.

Lemma 1.2.1. If $\alpha = a + x = b + y$ where $a, b \in \mathbb{Z}$ and $0 \le x, y < 1$ then a = b and x = y.

Proof. If

$$a + x = b + y$$

then

$$a - b = y - x$$

Note that $a-b \in \mathbb{Z}$ and $y-x \in (-1,1)$. Because 0 is the only integer in the interval (-1,1) it follows that a = b and x = y.

Representations of number

To represent a number x using continued fraction, first get the floor of x and let it be a_0 . Then express the fractional part $\{x\}$ as the reciprocal of the reciprocal of the fractional part: $\frac{1}{\frac{1}{\{x\}}}$. Note

that $\frac{1}{\{x\}} > 1$ (since $\{x\} < 1$), so we can get its floor to be a_1 and the new fractional part is treated the same as the last. If x is rational, this process is repeated until there is no fractional parts left. If x is irrational, this process is repeated forever. (We'll prove this in the following theorems.)

Let's use $\frac{54}{19}$ as an example:

$$\frac{54}{19} = 2 + \frac{1}{\frac{19}{16}} = 2 + \frac{1}{1 + \frac{1}{\frac{16}{3}}} = 2 + \frac{1}{1 + \frac{1}{5 + \frac{1}{3}}} = [2; 1, 5, 3]$$

Theorem 1.2.2. Every rational number has exactly two finite simple continued fraction expansions, and every finite simple continued fraction expansion represents a rational number.

Proof. Let $\alpha \in \mathbb{Q}$ and divide it into its floor and fractional part.

If the fractional part of α is 0, α is an integer and one simple continued fraction expansion is $[\alpha] = \alpha$ and a second is $[\alpha - 1; 1] = (\alpha - 1) + \frac{1}{1}$.

These are the only two expansions. There are no more with one quotient since α only has one value. There are no more with two quotients because if the second quotient is not 1 then the expression consists of a fraction that is less than 1 and α would have a fractional part. Finally, for the same reason there are no expansions with three or more quotients. Thus there are only two ways to express an integer as a simple continued fraction:

$$[\alpha] = \alpha \qquad \text{or} \qquad [\alpha - 1; 1] = (\alpha - 1) + \frac{1}{1} \tag{4}$$

If the fractional part of α is not 0 then define the 1st residue of alpha, and in general, of any continued fraction, as the reciprocal of its fractional part.

$$\alpha = a_0 + \frac{1}{r_1}$$

where $a_0 \in \mathbb{Z}$ and is the floor of α , and r_1 is the 1st residue of α . Note that $1 < r_1$ because of equation (3).

Now if the k th residue is not an integer it has a fractional part and we can define the (k+1) th residue recursively by the relationship

$$r_k = a_k + \frac{1}{r_{k+1}} \tag{5}$$

where $a_k = \lfloor r_k \rfloor$ and $r_{k+1} = \frac{1}{\{r_k\}}$. Now if $1 < r_k$ then $a_k \in \mathbb{Z}^+$ and of course $1 < r_{k+1}$. Because $1 < r_1$ it follows by induction that $1 < r_n$ for each r_n which is defined.

It can be easily seen that if α is rational and not an integer then r_1 is rational, and also that if r_n is rational and not an integer then r_{n+1} is rational. Now let $\frac{b}{c} = r_k$ where r_k is not an integer and $b, c \in \mathbb{Z}^+$ with their only common factor being 1. That is, $\frac{b}{c}$ is in its lowest terms. Because r_k is not an integer 1 < c. Then from equations (2), (3) and (5)

$$0 < r_k - a_k = \frac{b}{c} - a_k = \frac{b - a_k c}{c} = \frac{d}{c} = \frac{1}{r_{k+1}} < 1$$

where $0 < b - a_k c = d \in \mathbb{Z}$. Then

$$1 < r_{k+1} = \frac{c}{d}$$

Now because r_k has a denominator of c and r_{k+1} has a denominator of d each subsequent rational residue has a smaller integer denominator when reduced to it lowest terms. So then the denominators of the residues form a decreasing sequence of integers that are all greater than zero. Eventually one of the denominators must be 1 and then the residue is an integer.

Let the first residue that is an integer be r_n . Then the recurrence relationship given in equation (5) will not hold as there is no fractional part of r_n . Instead let

$$r_n = a_n$$

Notice that again we have $a_k = \lfloor r_k \rfloor$.

At this point we have n different equations that look like

$$\alpha = a_0 + \frac{1}{r_1}$$

$$\vdots$$

$$r_{n-1} = a_{n-1} + \frac{1}{r_n}$$

$$r_n = a_r$$

Combining these into one equation generates a simple continued fraction:

$$\alpha = a_0 + \frac{1}{a_1 + \frac{1}{\ddots + \frac{1}{a_n}}} = [a_0; a_1, \dots, a_n]$$

To establish uniqueness, assume that

$$\alpha = [a_0; a_1, \dots, a_n] = [b_0; b_1, \dots, b_m]$$

wher each $b_i \in \mathbb{N}$. Also assume without the loss of generality that neither a_n or b_m are 1. Now $0 < [0; b_m] < 1$. Also, because $1 < b_{m-1} + [0; b_m]$ it follows that $0 < [0; b_{m-1}, b_m] < 1$. One can prove by induction that

$$[0; b_i, b_{i+1}, \dots, b_m] \in (0, 1)$$

for $j = 1 \dots m$.

And because

$$a_0 + [0; a_1, \dots, a_n] = b_0 + [0; b_1, \dots, b_m]$$

, by Lemma 1.2.1, $a_0 = b_0$ and

$$a_1 + [0; a_2, \dots, a_n] = b_1 + [0; b_2, \dots, b_m]$$

. Continuing this way one finds that $a_0=b_0$, $a_1=b_1$, ..., $a_n=b_m$ and also that n=m. So if the last partial quotient is not 1 then the simple continued fraction expansion is unique.

However, as stated earlier, an integer can be represented in two ways, given by equations (4). So α can be represented one of two ways as a simple continued fraction;

$$a = [a_0; a_1, \dots, a_n] = [a_0; a_1, \dots, a_n - 1, 1]$$

Proving the second part of the theorem is trivial; it is evident that any finite simple continued fraction represents a rational number as the simple continued fraction can just be simplified from the lower right-hand corner upwards to generate the rational number it represents.

Theorem 1.2.3. Every irrational number has an infinite simple continued fraction expansion.

¹If $a_n=1$ then $r_{n-1}=a_{n-1}+1\in\mathbb{N}$ and so we can replace $[a_0;a_1,\ldots,a_n]=[a_0;a_1,\ldots,a_{n-1},1]$ with $[a_0;a_1,\ldots,a_{n-1}+1]$. Likewise for b_m .

Proof. Let $\alpha \in \mathbb{R}$; $\alpha \notin \mathbb{Q}$. Even though α is irrational, the residues are defined the same as in Theorem 1.2.2 In fact, because α is irrational its fractional part is not 0 and so there must be a 1st residue. Now it is obvious that r_1 is irrational, else α equals the sum of two rational numbers, and would not be irrational. Furthermore, for every irrational r_n, r_{n+1} is irrational, because of equation (5):

$$r_k = a_k + \frac{1}{r_{k+1}} \tag{5}$$

So all the residues are irrational. Thus there will never be a residue such that $r_k = a_k$ and so the irrational number has an infinite simple continued fraction expansion.

Now from the recursion formula of equation (5) an infinite number of equations are produced that look like

$$\alpha = a_0 + \frac{1}{r_1}$$

$$r_1 = a_1 + \frac{1}{r_2}$$

$$r_2 = a_2 + \frac{1}{r_3}$$

$$\vdots$$

We can combine these into one equation and generate an infinite simple continued fraction:

$$a_0 + \cfrac{1}{a_1 + \cfrac{1}{a_2 + \cfrac{1}{a_3 + \cfrac{1}{\ddots}}}} = [a_0; a_1, a_2, a_3, \ldots]$$

1.2.3 Convergents

Definition 1.2.2. If we crop the (finite or infinite) expression $\alpha = [a_0; a_1, \ldots, a_n, a_{n+1}, \ldots]$ to the n th quotient we get the rational number $[a_0; a_1, \ldots, a_n]$. This is called the n th **convergent** of a number α .

Notice for α with a finite number of convergents, the last convergent is equal to α .

Let p_n and q_n denote respectively the numerator and denominator of the n th convergent in its lowest terms. We have

$$p_0 = a_0$$
$$q_0 = 1$$

The 1st convergent is

$$[a_0; a_1] = a_0 + \frac{1}{a_1} = \frac{a_1 a_0 + 1}{a_1}$$

and so

$$p_1 = a_1 a_0 + 1$$
$$q_1 = a_1$$

Continuing in this fashion, the 2nd convergent is

$$[a_0; a_1, a_2] = a_0 + \frac{1}{a_1 + \frac{1}{a_2}} = a_0 + \frac{a_2}{a_2 a_1 + 1} = \frac{a_2 a_1 a_0 + a_2 + a_0}{a_2 a_1 + 1}$$

and so

$$p_2 = a_2 a_1 a_0 + a_2 + a_0 = a_2 p_1 + p_0 (6)$$

$$q_2 = a_2 a_1 + 1 = a_2 q_1 + q_0 (7)$$

This can be generalized in the following theorem.

Theorem 1.2.4. The numerator and denominator of the n th convergent of a real number $\alpha = [a_0; a_1, \ldots]$ where $0 < a_i \in \mathbb{R}$ is given by

$$p_n = a_n p_{n-1} + p_{n-2}$$

$$q_n = a_n q_{n-1} + q_{n-2}$$
(8)

and so the n th convergent is

$$\frac{p_n}{q_n} = [a_0; a_1, \dots, a_n] \tag{9}$$

(This is not necessarily a simple continued fraction so each a_i is not necessarily an integer.)

Proof. From equations (6) and (7) we can see that the theorem holds for n=2, for all $0 < a_0, a_1 \in \mathbb{R}$. Let us now assume that the theorem holds for $n=1\ldots k$ and then we will prove that it holds for n=k+1 and so the result follows by strong induction. Specifically, assume

$$p_k = a_k p_{k-1} + p_{k-2}$$
$$q_k = a_k q_{k-1} + q_{k-2}$$

and

$$\frac{p_k}{q_k} = [a_0; a_1, \dots, a_k] \tag{10}$$

Now note that

$$[a_0; a_1, \dots, a_k, a_{k+1}] = a_0 + \frac{1}{a_1 + \frac{1}{a_1 + \frac{1}{a_{k+1}}}}$$

$$= \left[a_0; a_1, \dots, a_k + \frac{1}{a_{k+1}}\right]$$

$$= (11)$$

Because $0 < a_k + \frac{1}{a_{k+1}} \in \mathbb{R}$, the only difference between the (k+1) th convergent and the k th convergent is that $a_k + \frac{1}{a_{k+1}}$ is in place of a_k . So if we replace a_k in equation (10) we get an expression for (k+1) th convergent. This is valid because it is evident from the recursion relationships of the numerators and denominators of the convergents, given from the induction assumption, that changing the value of a_k to $a_k + \frac{1}{a_{k+1}}$ will have no effect on the value of p_{n-1}, p_{n-2}, \ldots or q_{n-1}, q_{n-2}, \ldots . Fromthe equations (10) and (11):

$$[a_0; a_1, \dots, a_k, a_{k+1}] = \frac{\left(a_k + \frac{1}{a_{k+1}}\right) p_{k-1} + p_{k-2}}{\left(a_k + \frac{1}{a_{k+1}}\right) q_{k-1} + q_{k-2}}$$

$$= \frac{a_{k+1} a_k p_{k-1} + p_{k-1} + a_{k+1} p_{k-2}}{a_{k+1} a_k q_{k-1} + q_{k-1} + a_{k+1} q_{k-2}}$$

$$= \frac{a_{k+1} (a_k p_{k-1} + p_{k-2}) + p_{k-1}}{a_{k+1} (a_k q_{k-1} + q_{k-2}) + q_{k-1}}$$

$$= \frac{a_{k+1} p_k + p_{k-1}}{a_{k+1} q_k + q_{k-1}}$$

$$= \frac{p_{k+1}}{q_{k+1}}$$

and the theorem is proved by the induction principle.

It is common practice to define $p_{-2}=0$, $p_{-1}=1$, $q_{-2}=1$ and $q_{-1}=0$. Then the previous theorem can apply to n = 0 and n = 1, as

$$p_0 = a_0(1) + 0 = a_0$$

$$q_0 = a_0(0) + 1 = 1$$

$$p_1 = a_1(a_0) + 1 = a_1a_0 + 1$$

$$q_1 = a_1(1) + 0 = a_1$$

Because in this section there are no requirements that a_n are integers, we can use the idea of residues to prove this corollary to the previous theorem.

Corollary 1.2.5. If r_n is the n th residue of α and the numerators and denominators of the convergents of α are defined as in Theorem 1.2.4 then

$$\alpha = \frac{r_n p_{n-1} + p_{n-2}}{r_n q_{n-1} + q_{n-2}} \tag{12}$$

Proof. When assembling a simple continued fraction from the recursion relationships given by equation (5), if we stop at the n th equation we will end up with the equation

$$\alpha = a_0 + \frac{1}{a_1 + \frac{1}{\ddots + \frac{1}{a_{n-1} + \frac{1}{r_n}}}}$$

$$= [a_0; a_1, \dots, a_{n-1}, r_n]$$

While this may not be a simple continued fraction because r_n may not be an integer, all the terms of the continued fraction are real, so Theorem 1.4 still applies. So in equation (9) we can replace a_n with r_n to make a modified form of the last convergent, and thus

$$\alpha = \frac{r_n p_{n-1} + p_{n-2}}{r_n q_{n-1} + q_{n-2}}$$

The formula derived in Theorem 1.2.4 lead us to a very important relation that is central in solving Pell's equation.

Theorem 1.2.6. If the numerators and denominators of the convergents of a continued fraction are defined as in Theorem 1.2.4 then

$$p_{n+1}q_n - p_nq_{n+1} = (-1)^n (13)$$

Proof. Using the extended definitions of p_{-2}, p_{-1}, q_{-2} and q_{-1} we can establish the relation true for n=-2.

$$p_{-1}q_{-2} - p_{-2}q_{-1} = 1 \cdot 1 - 0 \cdot 0 = 1 = (-1)^{-2}$$

Now assume the relation true for n = k . That is

$$p_{k+1}q_k - p_k q_{k+1} = (-1)^k (14)$$

Then consider when n = k + 1.

$$p_{k+2}q_{k+1} - p_{k+1}q_{k+2} = (a_{k+2}p_{k+1} + p_k)q_{k+1} - p_{k+1}(a_{k+2}q_{k+1} + q_k)$$

from equations (8)

$$= p_k q_{k+1} - p_{k+1} q_k$$

= $(-1)(p_{k+1}q_k - p_k q_{k+1})$
= $(-1)^{k+1}$

by the induction assumption of equation (14). So the theorem is proved by induction.

In a simple continued fraction the a_n are all integers and it is easy to see from the recurrence relations given in equations (8) that all the p_n and q_n will be integers. They also have an important property described by a corollary of Theorem 1.2.6.

Corollary 1.2.7. When defined by equations (8), the numerators of each convergent of a simple continued fraction share no common factors with their corresponding denominator other than 1, and so the convergents are in their lowest terms.

Proof. Any common factor of p_n and q_n could be factored out of the left-hand side of equation (13) and so must also be a factor of the right-hand side. However, the only factors of the right-hand side are -1 and 1. Thus they share no common factors other than 1, and the convergents are in their lowest terms.

1.2.4 (Extra) Convergence of infinite continued fractions

[3] Let

 $c_n = [a_0; \dots, a_n] = \frac{p_n}{q_n}$

denote the n th convergent.

Theorem 1.2.8. (Extra of Theorem 1.2.6) If the numerators and denominators of the convergents of a continued fraction are defined as in Theorem 1.2.4 then

$$p_n q_{n-1} - q_n p_{n-1} = (-1)^{n-1}$$

$$p_n q_{n-2} - q_n p_{n-2} = (-1)^n a_n$$

Equivalently,

$$\frac{p_n}{q_n} - \frac{p_{n-1}}{q_{n-1}} = (-1)^{n-1} \cdot \frac{1}{q_n q_{n-1}}$$

$$p_n \quad p_{n-2} \quad (-1)^n \quad a_n$$

$$\frac{p_n}{q_n} - \frac{p_{n-2}}{q_{n-2}} = (-1)^n \cdot \frac{a_n}{q_n q_{n-2}}$$

Proof. The first equation is just Theorem 1.2.6 with shifted indices.

For the second equation, we have

$$p_n q_{n-2} - q_n p_{n-2} = (a_n p_{n-1} + p_{n-2}) q_{n-2} - p_{n-2} (a_n q_{n-1} + q_{n-2})$$
$$= a_n (p_{n-1} q_{n-2} - p_{n-2} q_{n-1})$$
$$= (-1)^n a_n$$

Theorem 1.2.9. The even indexed convergents c_{2n} increase strictly with n, and the odd indexed convergents decrease strictly with n. Also, the odd index convergents c_{2n+1} are greater than all of the even indexed convergents c_{2m} .

Proof. The a_n are positive for $n \geq 1$, so the q_n are positive. By Theorem 1.2.8, for $n \geq 2$,

$$c_n - c_{n-2} = (-1)^n \cdot \frac{a_n}{q_n q_{n-2}}$$

which proves the first claim (because RHS will be positive if n is even, and negative if n is odd).

Suppose for the sake of contradiction that there exists integers r,m such that $c_{2m+1} < c_{2r}$. Theorem 1.2.8 implies that for $n \ge 1$,

$$c_n - c_{n-1} = (-1)^{n-1} \cdot \frac{1}{q_n q_{n-1}}$$

has sign $(-1)^{n-1}$, so for all $s \ge 0$ we have $c_{2s+1} > c_{2s}$. Thus it is impossible that r = m. If r < m, then by what we proved in the first paragraph, $c_{2m+1} < c_{2r} < c_{2m}$, a contradiction (with s = m). If r > m, then $c_{2r+1} < c_{2m+1} < c_{2r}$, which is also a contradiction (with s = r).

Theorem 1.2.10. Let a_0, a_1, \ldots be a sequence of integers such that $a_n > 0$ for all $n \ge 1$, and for each $n \ge 0$, set $c_n = [a_0; a_1, \ldots, a_n]$. Then $\lim_{n \to \infty} c_n$ exists.

Proof. For any $m \geq n$, the number c_n is a partial convergent of $[a_0; \ldots, a_m]$. By Theorem 1.2.9 the even convergents c_{2n} form a strictly increasing sequence and the odd convergents c_{2n+1} form a strictly decreasing sequence. Moreoever, the even convergents are all $\leq c_1$ and the odd convergents are all $\geq c_0$. Hence (let) $\alpha_0 = \lim_{n \to \infty} c_{2n}$ and $\alpha_1 = \lim_{n \to \infty} c_{2n+1}$ both exist (because bounded monotonic sequence converges), and $\alpha_0 \leq \alpha_1$. Finally, by Theorem 1.2.8,

$$|c_{2n} - c_{2n-1}| = \frac{1}{q_{2n}q_{2n-1}}m \le \frac{1}{2n(2n-1)} \to 0$$

so $\alpha_0 = \alpha_1$.

Theorem 1.2.11. Let $\alpha \in \mathbb{R}$ be a real number. Then α is the value of the (possibly infinite) simple continued fraction $[a_0, a_1, a_2, \ldots]$ produced by the continued fraction procedure.

Proof. The theorem is obvious if the sequence is finite. Suppose the sequence is infinite. Let r_{n+1} be the (n+1) th residue. Then

$$\alpha = [a_0, a_1, \dots, a_n, r_{n+1}]$$

By Corollary 1.2.5, we have

$$\alpha = \frac{r_{n+1}p_n + p_{n-1}}{r_{n+1}q_n + q_{n-1}}$$

Then if $c_n = [a_0, a_1, \dots, a_n]$, then

$$\alpha - c_n = \frac{r_{n+1}p_n + p_{n-1}}{r_{n+1}q_n + q_{n-1}} - \frac{p_n}{q_n}$$

$$= \frac{(r_{n+1}p_n + p_{n-1})q_n - (r_{n+1}q_n + q_{n-1})p_n}{q_n(r_{n+1}q_n + q_{n-1})}$$

$$= \frac{p_{n-1}q_n - p_nq_{n-1}}{q_n(r_{n+1}q_n + q_{n-1})}$$

$$= \frac{(-1)^n}{q_n(r_{n+1}q_n + q_{n-1})}$$
 (by Theorem 1.2.8)

Thus

 $\begin{aligned} |\alpha - c_n| &= \frac{1}{q_n(r_{n+1}q_n + q_{n-1})} \\ &< \frac{1}{q_n(a_{n+1}q_n + q_{n-1})} \qquad \text{(since } r_{n+1} > a_{n+1}) \\ &= \frac{1}{q_n \cdot q_{n+1}} \qquad \text{(by Theorem 1.2.4)} \\ &\leq \frac{1}{n(n+1)} \to 0 \end{aligned}$

1.3 Quadratic irrationals

Definition 1.3.1. A quadratic irrational is an irrational number in the form $F + G\sqrt{M}$ where $F, G \in \mathbb{Q}$ and $M \in \mathbb{Z}^+$ and is not a perfect square.

We'll prove that every number in this form is indeed irrational.

Theorem 1.3.1. \sqrt{M} is irrational if $M \in \mathbb{Z}^+$ and is not a perfect square.

Proof. (This is madness.) If $M \in \mathbb{Z}^+$ and is not a perfect square then there exists $n \in \mathbb{Z}^+$ such that $n < \sqrt{M} < n+1$ and so

$$n^2 < M < (n+1)^2 \tag{15}$$

If \sqrt{M} is rational then it can be expressed as $\sqrt{M} = \frac{c}{d}$ where $c, d \in \mathbb{Z}^+$ and $\frac{c}{d}$ is in its lowest terms.

$$M = \frac{c^2}{d^2}$$

$$c^2 = Md^2 \tag{16}$$

Now from equations (15) and (16) we get

$$n^{2}d^{2} < Md^{2} < (n+1)^{2}d^{2}$$

 $n^{2}d^{2} < c^{2} < (n+1)^{2}d^{2}$
 $nd < c < (n+1)d$
 $0 < c - nd < d$
 $0 < f < d$ (17)

where $c - nd = f \in \mathbb{Z}^+$.

Again from equations (15) and (16) we get

$$n^{2}c^{2} < Mc^{2}$$
 $< (n+1)^{2}c^{2}$
 $n^{2}c^{2} < M^{2}d^{2}$ $< (n+1)^{2}c^{2}$
 $nc < Md$ $< (n+1)c$
 $0 < Md - nc < c$
 $0 < q$ $< c$ (18)

where $Md - nc = g \in \mathbb{Z}^+$.

Then

$$g^{2} - Mf^{2} = M^{2}d^{2} - 2Mncd + n^{2}c^{2} - M(c^{2} - 2ncd + n^{2}d^{2})$$

$$= (M^{2}d^{2} - Mc^{2}) - (2Mncd - 2Mncd) + (n^{2}c^{2} - Mn^{2}d^{2})$$

$$= 0$$

So

$$M = \frac{g^2}{f^2} \tag{19}$$

But because of equations (17) and (18) in conjunction with equation (19) above, $\frac{c^2}{d^2}$ was not in its lowest terms, contrary to assumption. Hence there is a contradiction and \sqrt{M} must be irrational. \square

It is obvious that adding a rational number to an irrational number will produce an irrational number; if the result were rational then an irrational number could be produced by subtracting a rational from a rational, which is clearly impossible. Furthermore, multiplying an irrational by a rational produces an irrational. Else au irrational could be produced by dividing a rational by a rational, which we know only produces rational numbers. Thus, numbers of the form $F + G\sqrt{M}$ as described previously are also irrational.

Theorem 1.3.2. If $a+b\sqrt{M}=c+d\sqrt{M}$ where $a,b,c,d\in\mathbb{Q}$, $M\in\mathbb{Z}^+$ and is not a perfect square, then a=c and b=d.

Proof. If the above conditions are true and $d-b\neq 0$ then

$$\sqrt{M} = \frac{a-c}{d-b}$$

and is rational, contrary to Theorem 1.3.1. So d - b = 0 and hence a = c also.

Lemma 1.3.3. All quadratic irrationals $\alpha = F + G\sqrt{M}$ where $F, G \in \mathbb{Q}$ and $M \in \mathbb{Z}^+$ and is not a perfect square, solve a quadratic equation with integer coefficients and also have a conjugate of the form $\alpha' = F - G\sqrt{M}$ which satisfies the same quadratic equation.

Proof. The quadratic equation

$$0 = (x - (F + G\sqrt{M}))(x - (F - G\sqrt{M}))$$

= $x^2 - 2Fx + F^2 - G^2M$

has roots α and α' , and has rational coefficients. Now let H be the common denominator of 2F and F^2-G^2M . That is, let H be the least number that makes both 2FH and F^2H-G^2MH integers. Then α and α' are roots of the quadratic equation

$$Hx^2 - 2FHx + F^2H - G^2MH = 0$$

which has integer coefficients.

From now on, let the symbol ' be the symbol for conjugate. So α' means the conjugate of α and $(\alpha + \beta)'$ means the conjugate of the sum of α and β .

Lemma 1.3.4. When applying one of the following operations;

- 1. addition
- 2. subtraction
- 3. multiplication
- 4. division

between two quadratic irrationals involving the same integer as the subject of the square root, conjugating the quadratic irrationals before the operation is equivalent to conjugating the result after the operation.

Proof. Let $\alpha = F + G\sqrt{M}$ be a quadratic irrational and $\beta = H + I\sqrt{M}$ be another. So α and β have the same integer under the square root. Now consider each operation:

1)

$$\alpha' + \beta' = (F + G\sqrt{M})' + (H + I\sqrt{M})'$$

$$= F - G\sqrt{M} + H - I\sqrt{M}$$

$$= (F + H) - (G + I)\sqrt{M}$$

$$= (F + G\sqrt{M} + H + I\sqrt{M})'$$

$$= (\alpha + \beta)'$$

2)

$$\alpha' - \beta' = (F + G\sqrt{M})' - (H + I\sqrt{M})'$$

$$= F - G\sqrt{M} - H + I\sqrt{M}$$

$$= (F - H) - (G - I)\sqrt{M}$$

$$= ((F - H) + (G - I)\sqrt{M})'$$

$$= (F + G\sqrt{M} - H - I\sqrt{M})'$$

$$= (\alpha - \beta)'$$

3)

$$\begin{split} \alpha'\beta' &= (F - G\sqrt{M})(H - I\sqrt{M}) \\ &= (FH + GIM) - (FI + HG)\sqrt{M} \\ &= (FH + GIM + FI\sqrt{M} + HG\sqrt{M})' \\ &= ((F + G\sqrt{M})(H + I\sqrt{M}))' \\ &= (\alpha\beta)' \end{split}$$

4)

$$\begin{split} \frac{\alpha'}{\beta'} &= \frac{(F - G\sqrt{M})}{(H - I\sqrt{M})} \times \frac{(H + I\sqrt{M})}{(H + I\sqrt{M})} \\ &= \frac{(FH - GIM) - (HG - FI)\sqrt{M}}{H^2 - I^2M} \\ &= \left(\frac{(FH - GIM) + (HG - FI)\sqrt{M}}{H^2 - I^2M}\right)' \\ &= \left(\frac{(F + G\sqrt{M})}{(H + I\sqrt{M})} \times \frac{(H - I\sqrt{M})}{(H - I\sqrt{M})}\right)' \\ &= \left(\frac{\alpha}{\beta}\right)' \end{split}$$

Lemma 1.3.5. Every expression in the form $F \pm G\sqrt{M}$ where $F, G \in \mathbb{Q}$; $M \in \mathbb{Z}^+$ and is not a perfect square and G > 0, has an equivalent expression in the form

$$\frac{A\pm\sqrt{D}}{B}$$

where $A \in \mathbb{Z}$; $B, D \in \mathbb{Z}^+$ and D is not a perfect square.

Proof. If $F,G\in\mathbb{Q}$ then they can be expressed $F=\frac{f}{h}$; $G=\frac{g}{h}$ where h is the lowest common denominator of F and G, and $f\in\mathbb{Z}$; $g,h\in\mathbb{Z}^+$. Then

$$F \pm G\sqrt{M} = \frac{f \pm g\sqrt{M}}{h}$$
$$= \frac{f \pm \sqrt{g^2 M}}{h}$$
$$= \frac{A \pm \sqrt{D}}{B}$$

where A, B and D satisfy the previously stated conditions and the two forms are equivalent. \Box

1.3.1 Reduced quadratic irrationals

Definition 1.3.2. A quadratic irrational α is **reduced** if $\alpha > 1$ and $-1 < \alpha' < 0$.

If α is reduced and in regular form $\frac{A+\sqrt{D}}{B}$ then $\alpha'=\frac{A-\sqrt{D}}{B}$ and we can deduce

$$1 < \frac{A + \sqrt{D}}{B}$$

$$B < A + \sqrt{D}$$
(20)

and

$$-1 < \frac{A - \sqrt{D}}{B}$$

$$-B < A - \sqrt{D}$$

$$\sqrt{D} - A < B \tag{21}$$

Also because B > 0 we can deduce

$$\frac{A - \sqrt{D}}{B} < 0$$

$$A - \sqrt{D} < 0$$

$$A < \sqrt{D}$$
(22)

and

$$-1 + 1 < \alpha' + \alpha$$

$$0 < \frac{2A}{B}$$

$$0 < A$$
(23)

From equations (22) and (23) we find

$$0 < A < \sqrt{D} \tag{24}$$

and from equations (20), (21), (24) we conclude

$$0 < \sqrt{D} - A < B < \sqrt{D} + A < 2\sqrt{D} \tag{25}$$

This set of inequalities will allow us to state the next theorem.

Theorem 1.3.6. There are only a finite number of reduced quadratic irrationals associated with any given D.

Proof. We know from equations (24) and (25) that for α to be a reduced quadratic irrational it is necessary that A must be between 0 and \sqrt{D} and B must be between 0 and $2\sqrt{D}$. Furthermore, both A and B must be integers. Thus, if we fix D then there are a finite number of pairs of integers which meet this criteria. Thus there are only a finite number of potential candidates of A and B that make α reduced when D is fixed.

Theorem 1.3.7. If α_n is reduced and $\alpha_n = a_n + \frac{1}{\alpha_{n+1}}$ where a_n is the floor of α_n then α_{n+1} is reduced and has the same integer as the subject of the square root.

Proof. First we show α_{n+1} is reduced.

$$0 < \alpha_n - a_n < 1$$
$$0 < \frac{1}{\alpha_{n+1}} < 1$$

 \mathbf{SO}

$$\alpha_{n+1} > 1 \tag{26}$$

Furthermore, from Theorem 1.3.4

$$(\alpha_n - a_n)' = \left(\frac{1}{\alpha_{n+1}}\right)'$$
$$\alpha'_n - a_n = \frac{1}{\alpha'_{n+1}}$$

then because $1 < a_n$ and $-1 < \alpha'_n < 0$

$$1 < a_n - \alpha_n' = \frac{-1}{\alpha_{n+1}'}$$

$$-1 < \alpha'_{n+1} < 0 \tag{27}$$

Equations (26) and (27) fulfill the requirements for α_{n+1} being reduced.

Now we show the subject of the square root is the same for α_{n+1} as it is for α_n . Let

$$\alpha = \frac{-b + \sqrt{b^2 - 4ac}}{2a} = \frac{A_n + \sqrt{D}}{B_n}$$

So

$$0 = a\alpha_n^2 + b\alpha_n + c$$

$$= a\left(a_n + \frac{1}{\alpha_{n+1}}\right)^2 + b\left(a_n + \frac{1}{\alpha_{n+1}}\right) + c$$

$$= aa_n^2 + \frac{2aa_n}{\alpha_{n+1}} + \frac{a}{\alpha_{n+1}^2} + ba_n + \frac{b}{\alpha_{n+1}} + c$$

$$= (aa_n^2 + ba_n + c)\alpha_{n+1}^2 + (2aa_n + b)\alpha_{n+1} + a$$
(28)

Note that the coefficients of this last quadratic equation are integers. Thus when we solve it for the positive root $\alpha_{n+1} = \frac{A_{n+1} + \sqrt{D_{n+1}}}{B_{n+1}}$ we will get A_{n+1} , B_{n+1} , D_{n+1} all integers. Now solve for D_{n+1} :

$$D_{n+1} = (2aa_n + b)^2 - 4(aa_n^2 + ba_n + c)a$$

$$= 4a^2a_n^2 + 4aa_nb + b^2 - 4a^2a_n^2 - 4aba_n - 4ac$$

$$= b^2 - 4ac$$

$$= D$$

Thus the theorem is proved.

Finally we come to the focal point of this section.

Theorem 1.3.8. If α is a reduced quadratic irrational then its simple continued fraction expansion is purely periodic.

Proof. Theorem 1.2.3 proves that the simple continued fraction of α is infinite. However, Theorem 1.3.7 implies that every residue of α is reduced. As a consequence of these two seemingly opposite statements, at some point there occurs some residue, r_k , that is a repetition of a previous residue, r_i .

Consider the simple continued fraction expansion of α , because a_j and a_k are the largest integers less than r_j and r_k respectively and $r_j = r_k$, it follows that $a_j = a_k$ and

$$r_{j} = r_{k}$$

$$a_{j} + \frac{1}{r_{j+1}} = a_{k} + \frac{1}{r_{k+1}}$$

$$r_{j+1} = r_{k+1}$$

Furthermore, the same reasoning can be applied to show that $r_{j+2}=r_{k+2}$, $r_{j+3}=r_{k+3}$ and so on. Now because $r_{n-1}=a_{n-1}+\frac{1}{r_n}$ we can manipulate two expressions about r_{j-1} and r_{k-1} side by side to eventually show that they are equal.

$$r_{j-1} = a_{j-1} + \frac{1}{r_j}$$

$$r_{k-1} = a_{k-1} + \frac{1}{r_k}$$

$$r'_{j-1} = a_{j-1} + \frac{1}{r'_j}$$

$$r'_{k-1} = a_{k-1} + \frac{1}{r'_k}$$

from Theorem 1.3.4. Because $r_j = r_k$ it follows that $r'_i = r'_k$, and

$$r'_{j-1} - a_{j-1} = r'_{k-1} - a_{k-1}$$

$$a_{j-1} - r'_{j-1} = a_{k-1} - r'_{k-1}$$

Because r_{j-1} and r_{k-1} are reduced, we have $-1 < r_{j-1}, r_{k-1} < 0$ so $0 < |r_{j-1}|, |r_{k-1}| < 1$, and it follows from Lemma 1.2.1 that $a_{j-1} = a_{k-1}$ and $r'_{j-1} = r'_{k-1}$ and thus, $r_{j-1} = r_{k-1}$. As before, the same method shows that $r_{j-2} = r_{k-2}$, $r_{j-3} = r_{k-3}$ and so on, up to $r_{j-(j-1)} = r_{k-(j-1)}$ which means $r_1 = r_{k-j+1}$, and $\alpha = r_{k-j}$ (α can be seen as r_0).

Let m be the value where r_m is the first residue where the value equals α . Then $r_i = r_{m+i}$ for all $i \in \mathbb{Z}^+$. Furthermore, taking the unique integer a_n for α and each of r_n we get $a_0 = a_m$ and $a_i = a_{m+i}$ for all $i \in \mathbb{Z}^+$.

$$a_i = a_{m+i}$$
 for all $i \in \mathbb{Z}^+$.
Thus $\alpha = [\overline{a_0; a_1, \dots, a_{m-1}}]$.

The variable m is known as the length of the period and will be quite important in solving Fell's equation.

1.4 Continued fractions of square roots

Square roots of natural numbers are not reduced quadratic irrationals because their conjugates are never between -1 and 0. Thus they are never purely periodic. However, they do have a special form

Theorem 1.4.1. Simple continued fractions of square roots take the form

$$\sqrt{D} = [a_0; \overline{a_1, a_2, \dots, a_{m-1}, 2a_0}]$$

when $D \in \mathbb{Z}^+$ and is not a perfect square.

Proof. If $D \in \mathbb{Z}^+$ and is not a perfect square then $1 < \sqrt{D}$. This means its conjugate $-\sqrt{D} < -1$ and so \sqrt{D} is not reduced. However, if a_0 is the greatest integer less than \sqrt{D} then

$$1 < a_0 + \sqrt{D}$$

and its conjugate

$$-1 < a_0 - \sqrt{D} < 0$$

So $a_0 + \sqrt{D}$ is reduced and by Theorem 1.3.8 its simple continued fraction representation is purely periodic.

So

$$a_{0} + \sqrt{D} = 2a_{0} + \cfrac{1}{a_{1} + \cfrac{1}{\ddots + \cfrac{1}{a_{m-1} + \cfrac{1}{2a_{0} + \cfrac{1}{\ddots}}}}}$$

and

$$\sqrt{D} = [a_0; \overline{a_1, a_2, \dots, a_{m-1}, 2a_0}]$$

Before we go on to solving Pell's equation we will find the simple continued fraction representation of $\sqrt{7}$ and show that it takes the form above.

First note that
$$a_0 = \lfloor \sqrt{7} \rfloor = 2$$
. So $\sqrt{7} = 2 + \frac{1}{r_1}$ and

$$r_1 = \frac{1}{\sqrt{7} - 2} \cdot \frac{\sqrt{7} + 2}{\sqrt{7} + 2}$$
$$= \frac{\sqrt{7} + 2}{3}$$

²The overhead bar is the standard mathematical notation used when the content below it repeats forever.

Now
$$a_1 = \lfloor r_1 \rfloor = \lfloor \frac{\sqrt{7} + 2}{3} \rfloor = 1$$
. So $r_1 = \frac{\sqrt{7} + 2}{3} = 1 + \frac{1}{r_2}$ and
$$r_2 = \frac{1}{\left(\frac{\sqrt{7} + 2}{3} - 1\right)}$$
$$= \frac{3}{\sqrt{7} - 1} \cdot \frac{\sqrt{7} + 1}{\sqrt{7} + 1}$$
$$= \frac{\sqrt{7} + 1}{2}$$

Now
$$a_2=\lfloor r_2\rfloor=\lfloor \frac{\sqrt{7}+1}{2}\rfloor=1$$
 . So $r_2=\frac{\sqrt{7}+1}{2}=1+\frac{1}{r_3}$ and

$$r_3 = \frac{1}{\left(\frac{\sqrt{7}+1}{2}-1\right)}$$
$$= \frac{2}{\sqrt{7}-1} \cdot \frac{\sqrt{7}+1}{\sqrt{7}+1}$$
$$= \frac{\sqrt{7}+1}{3}$$

Now
$$a_3=\lfloor r_3\rfloor=\lfloor \frac{\sqrt{7}+1}{3}\rfloor=1$$
 . So $r_3=\frac{\sqrt{7}+1}{3}=1+\frac{1}{r_4}$ and

$$r_4 = \frac{1}{\left(\frac{\sqrt{7}+1}{3}-1\right)}$$
$$= \frac{3}{\sqrt{7}-2} \cdot \frac{\sqrt{7}+2}{\sqrt{7}+2}$$
$$= \sqrt{7}+2$$

Now
$$a_4=\lfloor r_4\rfloor=\lfloor \sqrt{7}+2\rfloor=4$$
 . So $r_4=\sqrt{7}+2=4+\frac{1}{r_5}$ and

$$r_5 = \frac{1}{\sqrt{7} - 2}$$
$$= r_1$$

As a consequence, $a_5=a_1$, $a_6=a_2$, Thus the simple continued fraction representation is

$$\sqrt{7} = [2; \overline{1, 1, 1, 4}] \tag{29}$$

which is consistent with Theorem 1.4.1.

1.5 Pell's equation

(Related: Q66, Q94)

Definition 1.5.1. Pell's equation is an equation of the form

$$x^2 - Dy^2 = 1$$

where D is a given non-square positive integer, and integer solutions are sought for x and y.

The trivial solution is x = 1, y = 0 or x = -1, y = 0, but that is not very interesting. So we want x, y to be positive integers.

Note that D must be non-square for x, y to have non-trivial solutions, because if D is perfect square, then we can let $D = c^2$ for some positive integer c and we have

$$x^{2} - c^{2}y^{2} = 1$$
$$(x + cy)(x - cy) = 1$$

Because the only possible factors of 1 are -1 and 1, this leaves us with two options.

The first is that

$$x + cy = x - cy = 1$$

then x = 1 and cy = 0, so y = 0. This is the trivial solution.

The second is that

$$x + cy = x - cy = -1$$

then x = -1, and cy = 0 so y = 0 which is another trivial solution.

1.5.1 Solving Pell's equation with continued fractions

Theorem 1.5.1. Let $D \in \mathbb{Z}^+$ and not be a perfect square, so $\sqrt{D} = [a_0; \overline{a_1, a_2, \dots, a_{m-1}, 2a_0}]$. Also let p_n and q_n be defined as in Theorem 1.2.4.

If the length of the period, m, is even then $(x,y)=(p_{m-1},q_{m-1})$ solves the Pell equation $x^2-Dy^2=1$ for integers. If the length of the period, m, is odd then $(x,y)=(p_{m-1},q_{m-1})$ solves the negative Pell equation $x^2-Dy^2=-1$ for integers and $(x,y)=(p_{2m-1},q_{2m-1})$ solves the Pell equation $x^2-Dy^2=1$ for integers.

Proof. Because $\sqrt{D} = [a_0; \overline{a_1, a_2, \dots, a_{m-1}, 2a_0}]$ it follows that

$$\sqrt{D} = a_0 + \frac{1}{a_1 + \frac{1}{\cdots + \frac{1}{a_{m-1} + \frac{1}{a_0 + \sqrt{D}}}}}$$
(30)

Now from Theorem 1.2.5 we get

$$\sqrt{D} = \frac{(a_0 + \sqrt{D})p_{m-1} + p_{m-2}}{(a_0 + \sqrt{D})q_{m-1} + q_{m-2}}$$
$$(a_0 + \sqrt{D})q_{m-1}\sqrt{D} + q_{m-2}\sqrt{D} = (a_0 + \sqrt{D})p_{m-1} + p_{m-2}$$
$$q_{m-1}D + (a_0q_{m-1} + q_{m-2})\sqrt{D} = a_0p_{m-1} + p_{m-2} + p_{m-1}\sqrt{D}$$

and from theorem 1.3.2 it follows that

$$q_{m-1}D = a_0p_{m-1} + p_{m-2}$$
 and $p_{m-1} = a_0q_{m-1} + q_{m-2}$

so

$$p_{m-2} = q_{m-1}D - a_0p_{m-1}$$
 and $q_{m-2} = p_{m-1} - a_0q_{m-1}$ (31)

We can adjust the formula given in Theorem 1.2.6 by letting n = m - 2 to show

$$(-1)^{m-2} = p_{m-1}q_{m-2} - p_{m-2}q_{m-1}$$

= $p_{m-1}(p_{m-1} - a_0q_{m-1}) - (q_{m-1}D - a_0p_{m-1})q_{m-1}$

from equations (31). So

$$p_{m-1}^2 - Dq_{m-1}^2 = (-1)^m (32)$$

So when m is even $(x,y) = (p_{m-1},q_{m-1})$ solve the Pell equation, and when m is odd it solves the negative Pell equation.

Note that when setting up this proof in equation (30) we did not need to stop at the end of the first period. Instead, we could have stopped at the end of any period. If we stopped at the end of the second period, equation (30) would look like

$$\sqrt{D} = a_0 + \frac{1}{a_1 + \frac{1}{\cdots + \frac{1}{a_{2m-1} + \frac{1}{a_0 + \sqrt{D}}}}}$$

From this equation the previous logic can be carried out the same with the only difference being that m-1 is replaced by 2m-1 and m-2 is replaced by 2m-2. The new version of equation (32) will then look like

$$p_{2m-1}^2 - Dq_{2m-1}^2 = (-1)^m (33)$$

Thus when m is odd and it is not sufficient to stay in the first period to solve the Pell equation, $(x,y) = (p_{2m-1}, q_{2m-1})$ will give a solution in integers.

The general form of equation (33) for the k th period is

$$p_{km-1}^2 - Dq_{km-1}^2 = (-1)^{km}$$

for $k \in \mathbb{Z}^+$. This shows when m is even, $(x,y)=(p_{km-1},q_{km-1})$ will solve Pell's equation for all k, and when m is odd, $(x,y)=(p_{km-1},q_{km-1})$ will solve Pell's equation for all even k, and will solve the negative Pell equation for all odd k.

A direct consequence of this is that for any given non-square positive integer D, there are an infinite number of solutions for (x,y).

Example

Let's solve the Pell equation $x^2 - 7y^2 = 1$.

We know from equation (29) that $\sqrt{7} = [2; \overline{1,1,1,4}]$. We can see that the length of the period is 4 so Theorem 1.4.2 tells us that the numerator and the denominator of the 3rd convergent will solve the Pell equation for D=7. The numerator and denominators of the first 7 convergents are calculated from Theorem 1.2.4:

$$p_n = a_n p_{n-1} + p_{n-2}$$
$$q_n = a_n q_{n-1} + q_{n-2}$$

which are as follows:

	-2									
a_n			2	1	1	1	4	1	1	1
p_n	0 1	1	2	3	5	8	37	45	82	127
q_n	1	0	1	1	2	3	14	17	31	48

We can see that $p_3=8$ and $q_3=3$. This indeed is a solution to Pell's equation with D=7 as $8^2-7\times 3^2=1$. Furthermore, calculating into the second period shows that $p_7=127$ and $q_7=48$. A quick check with a calculator shows that it is true that $127^2-7\times 48^2=1$. If more solutions are required this is easily extended into the 3rd period or further.

1.5.2 Generating more solutions given a solution

But there is a faster method for finding other solutions once we find the first solution. Notice that we can 'factorize' the Pell's equation:

$$x^{2} - Dy^{2} = 1$$
$$(x + y\sqrt{D})(x - y\sqrt{D}) = 1$$

Then the LHS is the product of a pair of quadratic irrationals which are conjugate of each other. (Now we copy from another source [4]. So there is a slight change of notation.)

Definition 1.5.2. Let $\mathbb{Z}(\sqrt{D})$ denote the set of all numbers in the form $x+y\sqrt{D}$, $x,y\in\mathbb{Z}$. Symbolically,

 $\mathbb{Z}(\sqrt{D}) = \{x + y\sqrt{D} \mid x, y \in \mathbb{Z}\}\$

Definition 1.5.3. The **norm** of a number $u = x + y\sqrt{D}$ is defined as

$$N(u) = uu' = (x + y\sqrt{D})(x - y\sqrt{D}) = x^2 - Dy^2$$

Theorem 1.5.2. Let $u, v \in \mathbb{Z}(\sqrt{D})$, then u + v, $uv \in \mathbb{Z}(\sqrt{D})$. If $u \in \mathbb{Z}(\sqrt{D})$ and N(u) = 1, then $u^{-1} \in \mathbb{Z}(\sqrt{D})$.

Proof. Let $u=p+q\sqrt{D}$ and $v=r+s\sqrt{D}$, where $p,q,r,s\in\mathbb{Z}$. Then

$$u + v = (p + q\sqrt{D}) + (r + s\sqrt{D}) = (p + r) + (q + s)\sqrt{D} \quad \in \mathbb{Z}(\sqrt{D})$$

$$uv = (p + q\sqrt{D})(r + s\sqrt{D}) = (pr + qsD) + (ps + qr)\sqrt{D} \in \mathbb{Z}(\sqrt{D})$$

For the if-then statement, first note that

$$(p + q\sqrt{D})^{-1} = \frac{1}{(p + q\sqrt{D})(p - q\sqrt{D})}(p - q\sqrt{D})$$
 (2)

If $u \in \mathbb{Z}(\sqrt{D})$ and N(u) = 1, then

$$u^{-1} = \frac{1}{N(u)}(p - q\sqrt{D}) = \frac{1}{1}(p - q\sqrt{D}) = u' \in \mathbb{Z}(\sqrt{D})$$

Theorem 1.5.3. The norm and the conjugate are multiplicative in u:

$$N(uv) = N(u)N(v)$$
 and $(uv)' = u'v'$

Proof.

$$u'v' = (p - q\sqrt{D})(r - s\sqrt{D}) = (pr + qsD) - (ps + qr)\sqrt{D}$$
$$= (uv)'$$

Then

$$N(uv) = (uv)(uv)' = uvu'v' = (uu')(vv') = N(u)N(v)$$

Note that for $u=x+y\sqrt{D}$, a pair of integers (x,y) is a solution to Pell's equation if and only if N(u)=1 .

Lemma 1.5.4. Let (x,y) be an integer solution to Pell's equation and $u=x+y\sqrt{D}$.

- 1. If x > 0 and y > 0, then u > 1;
- 2. If x > 0 and y < 0, then 0 < u < 1;
- 3. If x < 0 and y > 0, then -1 < u < 0;

4. If x < 0 and y < 0, then u < -1;

Proof. Suppose x>0 and y>0. Since $(x-y\sqrt{D})(x+y\sqrt{D})=1$, we have $x-y\sqrt{D}>0$ and $x+y\sqrt{D}>x-y\sqrt{D}$. Hence u>1 and u'<1. This proves the first two statements. The third and the fourth statements follow from the first two.

Definition 1.5.4. Let (a,b) be a non-trivial solution to Pell's equation with positive integer components $a>0,\ b>0$. We say that this solution is **fundamental** if the number $u=a+b\sqrt{D}$ takes the minimal possible value.

Note that the number u is uniquely determined by Theorem 1.3.2 . Let's also note that u > 1 by Lemma 1.5.4.

Theorem 1.5.5. Let (x_1, y_1) be a fundamental solution to Pell's equation and $u = x_1 + y_1 \sqrt{D}$. Let

$$u^n = x_n + y_n \sqrt{D}$$
, $n = 0, 1, 2, \dots$

Then $(\pm x_n, \pm y_n)$, $n = 0, 1, 2, \ldots$, is the complete set of solutions to Pell's equation.

In other words, every positive solution is a pair of coefficients of the expanded form of u^n for $n = 0, 1, 2, \dots$

Proof. The trivial solution (1,0) is in this set and we get it for n=0. Let (x,y) be an arbitrary non-trivial solution to Pell's equation, and $v=x+y\sqrt{D}$. We may assume x>0 and y>0. All we need to show is that $v=x+y\sqrt{D}$ can be represented as u^n for some positive integer n.

Let us assume the contrary. As x > 0 and y > 0, we know that v > 1. Since u > 1 the terms of the sequence $1, u, u^2, \ldots, u^n, \ldots$ get arbitrary large, thus there exists n such that $u^n < v < u^{n+1}$. Let us multiply this inequality by $(u^n)^{-1}$. We get

$$1 < v(u^n)^{-1} < u \tag{3}$$

Let's make a numbe of observations.

Firstly, note that $(u^n)^{-1} = (u^{-1})^n = (x_1 - y_1\sqrt{D})^n$. By Theorem 1.5.2, this means $(u^n)^{-1} \in \mathbb{Z}(D)$ and hence $v(u^n)^{-1} \in \mathbb{Z}(D)$. We can let $v(u^n)^{-1} = \bar{x} + \bar{y}\sqrt{D}$ for some $\bar{x}, \bar{y} \in \mathbb{Z}$.

Secondly, by Theorem 1.5.3, $N(v(u^n)^{-1}) = N(v)N((u^{-1})^n) = N(v)(N(u^{-1}))^n = N(v)(N(u'))^n = (1)(1)^n = 1$ by so (\bar{x}, \bar{y}) is a solution to Pell's equation.

Thirdly, by equation (3), $1 < v(u^n)^{-1}$, so by Lemma 1.5.4 we get $\bar{x} > 0, \bar{y} > 0$ (the if-then condition also works in reverse since any other possibilities of \bar{x}, \bar{y} will make $v(u^n)^{-1} < 1$).

Finally, this contradicts to equation (3), namely to $v(u^n)^{-1} < u$, since u was fundamental. To elaborate, we have a value $v(u^n)^{-1} = \bar{x} + \bar{y}\sqrt{D}$ that is smaller than value of the fundamental solution $u = x_1 + y_1\sqrt{D}$. This is impossible since u is supposed to be the minimum value.

Thus the theorem is proved.

Example

Let's use this theorem in an example. Let's say for the Pell's equation

$$x^2 - 2y^2 = 1$$

we find the fundamental solution (3,2). (check: $3^2 - 2(2)^2 = 1$)

Then every positive solution is the coefficient of the expanded form of $u^n = (3 + 2\sqrt{2})^n$. To find the next solution, let n = 2 and expand:

$$(3+2\sqrt{2})^2 = 9 + (2)(3)(2\sqrt{2}) + 8$$
$$= 17 + 12\sqrt{2}$$

Then (17, 12) is the next solution to the equation (check: $17^2 - 2(12^2) = 289 - 288 = 1$).

To find the third solution, let n = 3 and expand:

$$(3+2\sqrt{2})^3 = (17+12\sqrt{2})(3+2\sqrt{2})$$
$$= 99+70\sqrt{2}$$

Then (99,70) is the third solution to the equation (check: $99^2 - 2(70^2) = 9801 - 9800 = 1$). We can find every solution by repeating this process.

And to find the fundamental solution, we can find use the continued fraction method in section 1.5.1. The first solution is guaranteed to be the fundamental solution.

Theorem 1.5.6. Let $D \in \mathbb{Z}^+$ and not be a perfect square, so $\sqrt{D} = [a_0; \overline{a_1, a_2, \dots, a_{m-1}, 2a_0}]$. Also let p_n and q_n be defined as in Theorem 1.2.4.

If the length of the period, m, is even then $(x,y) = (p_{m-1},q_{m-1})$ is the fundamental solution to Pell's equation. If the length of the period, m, is odd then $(x,y) = (p_{2m-1},q_{2m-1})$ is the fundamental solution.

Proof. From Theorem 1.2.4,

$$p_n = a_n p_{n-1} + p_{n-2}$$
$$q_n = a_n q_{n-1} + q_{n-2}$$

We want to show that p_n is a strictly increasing sequence for $n \geq 0$.

Note that $p_{-1} = 1$ and $p_0 = a_0 > 0$, and $p_1 = a_1 p_0 + 1 > p_0$.

Assume that p_0, \ldots, p_k is a strictly increasing sequence, where $k \geq 1$. Then $p_k > p_0 > 0$ and $p_{k-1} \geq p_0 > 0$, and

$$p_{k+1} - p_k = a_{k+1}p_k + p_{k-1} - p_k$$
$$= p_k(a_{k+1} - 1) + p_{k-1}$$

Since $p_k(a_{k+1}-1) \ge 0$ and $p_{k-1}>0$, we have $p_{k+1}-p_k>0$, which means p_0,\ldots,p_{k+1} is also strictly increasing. By induction, p_n is a strictly increasing sequence for $n\ge 0$.

Since $q_0 = 1$, $q_1 = a_1 > 0$ and $q_2 = a_2q_1 + 1 > q_1$, we can use similar argument to show that q_n is a strictly increasing sequence for $n \ge 1$.

Now, since n = m - 1 (if m is even) or n = 2m - 1 (if m is odd) is the minimum n such that (p_n, q_n) solves Pell's equation, and p_n is strictly increasing, we have

$$p_{m-1} < p_k$$
 (or $p_{2m-1} < p_k$) and $q_{m-1} < q_k$ (or $q_{2m-1} < q_k$)

for any k > m-1 (or k > 2m-1) such that (p_k, q_k) solve Pell's equation. This means

$$p_{m-1} + q_{m-1}\sqrt{D} < p_k + q_k\sqrt{D}$$
(or)
$$p_{2m-1} + q_{2m-1}\sqrt{D} < p_k + q_k\sqrt{D}$$

so (p_{m-1}, q_{m-1}) (or (p_{2m-1}, q_{2m-1})) is the fundamental solution.

Now we give the 'general' formula for x and y respectively.

Theorem 1.5.7. Let (x_1, y_1) be a fundamental solution to Pell's equation and $u = x_1 + y_1 \sqrt{D}$, $u' = x_1 - y_1 \sqrt{D}$. Let

$$u^n = x_n + y_n \sqrt{D}, \qquad n = 0, 1, 2, \dots$$

Then

$$x_n = \frac{u^n + (u')^n}{2}$$
, $y_n = \frac{u^n - (u')^n}{2\sqrt{D}}$

Proof. For n = 0 we have $(x_0, y_0) = (1, 0)$, $u^0 = 1 + 0\sqrt{D} = 1$. So

$$\frac{1^0 + 1^0}{2} = 1 = x_0$$
, $\frac{1^n - 1^n}{2\sqrt{D}} = 0 = y_0$

For n=1,

$$\frac{(x_1 + y_1\sqrt{D}) + (x_1 - y_1\sqrt{D})}{2} = x_1 , \quad \frac{(x_1 + y_1\sqrt{D}) - (x_1 - y_1\sqrt{D})}{2\sqrt{D}} = y_1$$

Assume that

$$x_k = \frac{u^k + (u')^k}{2}$$
, $y_k = \frac{u^k - (u')^k}{2\sqrt{D}}$

When n = k + 1,

$$u^{k+1} = (x_k + y_k \sqrt{D})(x_1 + y_1 \sqrt{D})$$

= $x_k x_1 + y_k y_1 D + (x_k y_1 + y_k x_1) \sqrt{D}$

So for LHS, $x_{n+1} = x_k x_1 + y_k y_1 D$ and $y_{n+1} = x_k y_1 + y_k x_1$. By induction hypothesis,

$$x_{n+1} = \left(\frac{u^k + (u')^k}{2}\right) x_1 + \left(\frac{u^k - (u')^k}{2\sqrt{D}}\right) y_1 D$$

$$= \frac{(u^k + (u')^k) x_1 + (u^k - (u')^k) y_1 \sqrt{D}}{2}$$

$$= \frac{u^k (x_1 + y_1 \sqrt{D}) + (u')^k (x_1 - y_1 \sqrt{D})}{2}$$

$$= \frac{u^{k+1} + (u')^{k+1}}{2}$$

$$= \text{RHS}$$

And

$$y_{n+1} = \left(\frac{u^k + (u')^k}{2}\right) y_1 + \left(\frac{u^k - (u')^k}{2\sqrt{D}}\right) x_1$$

$$= \frac{(u^k + (u')^k) y_1 \sqrt{D} + (u^k - (u')^k) x_1}{2\sqrt{D}}$$

$$= \frac{u^k (x_1 + y_1 \sqrt{D}) - (u')^k (x_1 - y_1 \sqrt{D})}{2\sqrt{D}}$$

$$= \frac{u^{k+1} - (u')^{k+1}}{2\sqrt{D}}$$

$$= \text{RHS}$$

Thus the formula holds for all non-negative integers n .

1.6 Divisibility and modular arithmetic

1.6.1 Some properties of gcd and congruence

(It is assumed that the modulus n is always a positive integer.)

Lemma 1.6.1. If a, n are coprime and $ab \equiv ac \pmod{n}$, then $b \equiv c \pmod{n}$. (cancellation property)

Proof. $ab \equiv ac \pmod{n} \Rightarrow ab - ac \equiv 0 \pmod{n} \Rightarrow a(b-c) \equiv 0 \pmod{n}$

- $\Rightarrow n \mid a(b-c)$
- \Rightarrow Since n, a are coprime, $n \nmid a \Rightarrow n \mid (b-c)$ by Euclid's lemma
- $\Rightarrow b \equiv c \pmod{n}$ (difference property)

Lemma 1.6.2. Given integers a, b, c, gcd(a, bc) = 1 if and only if gcd(a, b) = 1 and gcd(a, c) = 1. (property of coprime sharing)

Proof. This is fairly straightforward if we allow the use of the Fundamental Theorem of Arithmetic.

- (\Leftarrow) If gcd(a,b) = gcd(a,c) = 1, then a does not share a prime factor with b and c. Since any prime factor of bc comes from either b or c, it cannot be found in a.
- (\Rightarrow) If gcd(a,bc) = 1, then a does not share any prime factor with bc, which means b and c cannot have a prime factor found in a. (Otherwise the product bc will share a prime factor with a, a contradiction.)

Lemma 1.6.3. Given integers a, b, gcd(a, b) = gcd(a + bk, b) for all integers k. (GCD reduction property)

Proof. First we show that gcd(a, b) = gcd(a - b, b) as follows:

Let $d = \gcd(a, b)$ and $d' = \gcd(a - b, b)$.

By definition, $d \mid a$ and $d \mid b \Rightarrow d \mid (a - b)$ (property of linear combination).

Since $d \mid b$ and $d \mid (a - b)$, we have $d \mid \gcd(a - b, b)$ (GCD universal property) $\Rightarrow d \mid d'$

Also, by definition $d' \mid b$ and $d' \mid (a-b) \implies d' \mid ((a-b)+b) = a$ (property of linear combination)

 $\Rightarrow d' \mid \gcd(a, b) \text{ (GCD universal property)} \Rightarrow d' \mid d$

Thus, d = d', as desired.

By flipping the + and - signs in the proof above, we also get gcd(a,b) = gcd(a+b,b). Thus it follows inductively that gcd(a,b) = gcd(a+bk,b) for all integers k.

Corollary 1.6.4. $gcd(a, b) = gcd(a \mod b, b)$ (modulo reduction property)

Proof. Let $a \mod b = r$. Then by definition of modulo a = bk + r for some integer k.

 $\Rightarrow r = a - bk \Rightarrow \gcd(a \mod b, b) = \gcd(r, b) = \gcd(a - bk, b) = \gcd(a, b)$ by previous lemma.

1.6.2 Euler's totient function

(Related: Q69, 70, 72, 73, 243)

Definition 1.6.1. Euler's **totient function**, $\phi(n)$, is the number of positive integers not exceeding n which are coprime to n.

For example, $\phi(9) = 6$ since 1, 2, 4, 5, 7, 8 are all less than or equal to nine and coprime to nine.

Also, $\phi(1) = 1$ since gcd(1, 1) = 1 and 1 is the only positive integer not exceeding itself.

Also, $\phi(p) = p - 1$ for prime p since p is coprime to $1, \dots, p - 1$.

Theorem 1.6.5. Totient function is **multiplicative** for two coprime numbers. i.e. If m, n are coprime, then $\phi(mn) = \phi(m) \cdot \phi(n)$.

Proof. The proof requires Chinese remainder theorem.

Let m, n be coprimes. Consider the set $\mathbb{Z}_{mn} = \{0, 1, 2, \dots, mn-1\}$ and the sets $\mathbb{Z}_m = \{0, 1, 2, \dots, m-1\}$ and $\mathbb{Z}_n = \{0, 1, 2, \dots, n-1\}$. We want to show that $|\mathbb{Z}_{mn}| = |\mathbb{Z}_m \times \mathbb{Z}_n|$, that is, \mathbb{Z}_{mn} has the same number of elements as $\mathbb{Z}_m \times \mathbb{Z}_n$ (even though it is obvious).

Define a function $f: \mathbb{Z}_{mn} \to \mathbb{Z}_m \times \mathbb{Z}_n$, i.e., f maps \mathbb{Z}_{mn} to the set $\mathbb{Z}_m \times \mathbb{Z}_n$. For each a in $\mathbb{Z}_{mn} = \{0, 1, 2, \dots, mn - 1\}$, define f(a) = (c, d) where $c \in \mathbb{Z}_m$ with $c \equiv a \pmod{m}$ and $d \in \mathbb{Z}_n$ with $d \equiv a \pmod{n}$, and (c, d) is an ordered pair.

Note that c,d are necessarily unique for the same a (so the function is valid) because we have $0 \le c < m$, so $c = c \ /. \ m = a \ /. \ m$, and similarly, $0 \le d < n$, so $d = d \ /. \ n = a \ /. \ n$.

To show that f is a bijection, let $(c,d) \in \mathbb{Z}_m \times \mathbb{Z}_n$. Consider the equations $x \equiv c \pmod m$ and $x \equiv d \pmod n$. By Chinese remainder theorem, these two equations have a solution a that is unique modulo mn, i.e. there is a unique solution a where $0 \le a < mn$, which means there is exactly one $a \in \mathbb{Z}_{mn}$ such that f(a) = (c,d) for every $(c,d) \in \mathbb{Z}_m \times \mathbb{Z}_n$. Thus f is a bijection.

This means that $|\mathbb{Z}_{mn}| = |\mathbb{Z}_m \times \mathbb{Z}_n|$.

Now, let m, n still be coprimes. Let \mathbb{Z}_{mn}^* be the set of all elements in \mathbb{Z}_{mn} that is coprime to mn. Mathamatically, $\mathbb{Z}_{mn}^* = \{\gcd(a,mn) = 1 : a \in \mathbb{Z}_{mn}\}$. Let \mathbb{Z}_m^* and \mathbb{Z}_n^* be defined similarly. Note that $|\mathbb{Z}_m^*| = \phi(m)$, $|\mathbb{Z}_n^*| = \phi(n)$ and $|\mathbb{Z}_{mn}^*| = \phi(mn)$.

We want to show that $|\mathbb{Z}_{mn}^*| = |\mathbb{Z}_m^* \times \mathbb{Z}_n^*| = |\mathbb{Z}_m^*| \times |\mathbb{Z}_n^*|$, which means $\phi(mn) = \phi(m)\phi(n)$.

The same mapping f defined above is used. We show that when the domain of f is restricted to the set \mathbb{Z}_{mn}^* , it is a bijection from \mathbb{Z}_{mn}^* onto the set $\mathbb{Z}_m^* \times \mathbb{Z}_n^*$. First, we show that for any $a \in \mathbb{Z}_{mn}^*$, f(a) is indeed in $\mathbb{Z}_m^* \times \mathbb{Z}_n^*$. To see this, note that a and mn are coprime. So it must be that a and m are coprime too and that a and n are coprime (property of coprime sharing).

The function f is one-to-one (an injection) as shown above. The remaining piece is that for each $(c,d) \in \mathbb{Z}_m^* \times \mathbb{Z}_n^*$, there is some $a \in \mathbb{Z}_{mn}^*$ such that f(a) = (c,d). As in the proof of $|\mathbb{Z}_{mn}| = |\mathbb{Z}_m \times \mathbb{Z}_n|$ above, there is some $a \in \mathbb{Z}_{mn}$ such that f(a) = (c,d) for every (c,d). Note that c,m are coprime and d,n are coprime by definition. Since by construction $c \equiv a \pmod{m}$, a = c + km for some k, so $\gcd(c + km, m) = 1$ (property of coprime steps), which means that a and m are coprime. By similar reasoning, a and n are coprime. This means that a and mn are coprime (property of coprime sharing), which means $a \in \mathbb{Z}_{mn}^*$. Thus the function f is a one-to-one from \mathbb{Z}_{mn}^* onto $\mathbb{Z}_m^* \times \mathbb{Z}_n^*$ (a bijection), so we have $|\mathbb{Z}_{mn}^*| = |\mathbb{Z}_m^* \times \mathbb{Z}_n^*| = |\mathbb{Z}_m^*| \times |\mathbb{Z}_n^*|$, and

$$\phi(mn) = \phi(m)\phi(n)$$

Theorem 1.6.6 (Value of ϕ for a prime power). If p is a prime and $k \geq 1$, then

$$\phi(p^k) = p^k - p^{k-1} = p^{k-1}(p-1) = p^k(1 - \frac{1}{p})$$

Proof. Since p is a prime number, , the only possible values of $\gcd(p^k,m)$ for an arbitrary m are $1,p,p^2,\ldots,p^k$. And the only way to have $\gcd(p^k,m)>1$ is if m is a multiple of p, that is, $m\in\{p,2p,3p,\ldots,p^{k-1}p=p^k\}$, and there are p^{k-1} such multiples not greater than p^k . Therefore, the other p^k-p^{k-1} numbers in $\{1,2,\ldots,p^k\}$ are all relatively prime to p^k .

Theorem 1.6.7 (Euler's product formula).

$$\phi(n) = n \prod_{p \mid n} (1 - \frac{1}{p})$$

where the product is over all distinct prime factors of n.

Proof. The fundamental theorem of arithmetic states that if n > 1 there is a unique expression $n = p_1^{k_1} p_2^{k_2} \cdots p_t^{k_t}$, where $p_1 < p_2 < \ldots < p_t$ are prime numbers and each $k_i \ge 1$. (The case n = 1

corresponds to the empty product.) Repeatedly using the multiplicative property gives:

$$\begin{split} \phi(n) &= \phi(p_1^{k_1})\phi(p_2^{k_2})\cdots\phi(p_t^{k_t}) \\ &= p_1^{k_1-1}(p_1-1)\,p_2^{k_2-1}(p_2-1)\cdots p_t^{k_t-1}(p_t-1) \\ &= p_1^{k_1}(1-\frac{1}{p_1})p_2^{k_2}(1-\frac{1}{p_2})\cdots p_t^{k_t}(1-\frac{1}{p_t}) \\ &= p_1^{k_1}p_2^{k_2}\cdots p_t^{k_t}(1-\frac{1}{p_1})(1-\frac{1}{p_2})\dots(1-\frac{1}{p_t}) \\ &= n(1-\frac{1}{p_1})(1-\frac{1}{p_2})\dots(1-\frac{1}{p_t}) \\ &= n\prod_{p\mid n}(1-\frac{1}{p}) \end{split}$$

Thus, we see that $\phi(n)$ is n times the product of 1 minus the reciprocal of each distinct prime factors of n. Note that the big pi notation \prod is the product version of the sigma notation \sum .

Example: To find $\phi(12)$, note that the distinct prime factors of 12 is 2 and 3, so

$$\phi(12) = 12(1 - \frac{1}{2})(1 - \frac{1}{3}) = 4$$

We can verify that there are only 4 coprime positive integers smaller than 12, which are 1, 5, 7, 11.

1.6.3 Fermat's little theorem

Theorem 1.6.8 (Fermat's little theorem). Let p be a prime and let a be an integer such that $p \nmid a$. Then

$$a^{p-1} \equiv 1 \pmod{p}$$

Proof. [5] (P.88) We begin by considering the first p-1 positive multiples of a; that is, the integers

$$a, 2a, 3a, \ldots, (p-1)a$$

None of these numbers is congruent modulo p to any other, nor is any congruent to zero (as none of them has a factor of p). Indeed, if it happened that

$$ra \equiv sa \pmod{p}$$
 $1 \le r < s \le p - 1$

then a could be cancelled to give $r \equiv s \pmod{p}$, which is impossible. (Note: The cancellation property works when a, p are coprime.) Therefore, the previous set of integers must be congruent modulo p to $1, 2, 3, \ldots, p-1$, taken in some order. Multiplying all these congruence together, we find that

$$a \cdot 2a \cdot 3a \cdot \dots (p-1)a \equiv 1 \cdot 2 \cdot 3 \dots (p-1) \pmod{p}$$

whence

$$a^{p-1}(p-1)! \equiv (p-1)! \pmod{p}$$

Once (p-1)! is cancelled from both sides of the congruence (this is possible since p and (p-1)! are coprime³), our line of reasoning culminates in the statement that $a^{p-1} \equiv 1 \pmod{p}$, which is Fermat's little theorem.

This result can be stated in a slightly more general way in which the requirement that $p \not\mid a$ is dropped.

Corollary 1.6.9. Let p is a prime and a be any integer. Then

$$a^p \equiv a \pmod{p}$$

Proof. [5] (P.88) When $p \mid a$, the statement obviously holds, since we also have $p \mid a^p$ and thus $a^p \equiv a \equiv 0$. If $p \not\mid a$, then according to Fermat's little theorem, we have $a^{p-1} \equiv 1 \pmod{p}$. Multiply both sides by a to get $a^p \equiv a \pmod{p}$.

³Any prime factors in (p-1)! is smaller than p, and p has only p as its prime factor. Thus they have no common prime factors.

1.6.4 Euler's totient theorem

Euler's totient theorem generalizes Fermat's little theorem from the case of prime p to any positive integer n.

Theorem 1.6.10 (Euler's totient theorem). If $n \ge 1$ and gcd(a, n) = 1, then

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

where $\phi(n)$ is Euler's totient funtion.

Proof. [5] (P.137) The theorem is trivially true for n = 1 (check: $a^1 \equiv 1 \pmod{1}$ which is true because every integer are congruent to 0 under modulo 1).

Thus we assume n > 1. First we prove a lemma:

Lemma. Let n > 1 and gcd(a, n) = 1. If $x_1, x_2, \ldots, x_{\phi(n)}$ are all the positive integers less than n and coprime to n, then

$$ax_1, ax_2, \ldots, ax_{\phi(n)}$$

are congruent modulo n to $x_1, x_2, \ldots, x_{\phi(n)}$ in some order.

Proof of lemma:

Observe that no two of the integers $ax_1, ax_2, \ldots, ax_{\phi(n)}$ are congruent modulo n to each other. For if $ax_i \equiv ax_j \pmod{n}$, with $1 \le i < j \le \phi(n)$, then the cancellation property yields $x_i \equiv x_j \pmod{n}$ and thus $x_i = x_j$, a contradiction (because when $i \ne j$ then we should have $x_i \ne x_j$). Furthermore, because $\gcd(x_i, n) = 1$ for all i and $\gcd(a, n) = 1$, Lemma 1.6.2 (property of coprime sharing) guarantees that each of the ax_i is coprime to n.

Fixing on a particular ax_i , there exists a unique integer x_i' , where $0 \le x_i' < n$, for which $ax_i \equiv x_i' \pmod{n}$ (because the remainder of modulo operation is unique). By modulo reduction property (Corollary 1.6.4),

$$\gcd(x_i', n) = \gcd(ax_i, n) = 1$$

Thus, x_i' being coprime to n, must be one of the integers $x_1, x_2, \dots, x_{\phi(n)}$. And note that the corresponding x_i' for each ax_i is different, which is given by the first observation. All told, this proves that the numbers $ax_1, ax_2, \dots, ax_{\phi(n)}$ and the numbers $a_1, a_2, \dots, a_{\phi(n)}$ are identical (modulo n) in a certain order.

Now, let $x_1, x_2, \ldots, x_{\phi(n)}$ be defined the same as above. We have

$$ax_1 \equiv x_1' \pmod{n}$$

$$ax_2 \equiv x_2' \pmod{n}$$

$$\vdots \qquad \vdots$$

$$ax_{\phi(n)} \equiv x_{\phi(n)}' \pmod{n}$$

where $x_1', x_2', \dots, x_{\phi(n)}'$ are the integers $x_1, x_2, \dots, x_{\phi(n)}$ in some order. On taking the product of these $\phi(n)$ congruences, we get

$$(ax_1)(ax_2)\dots(ax_{\phi(n)}) \equiv x_1'x_2'\dots x_{\phi(n)}' \pmod{n}$$

 $\equiv x_1x_2\dots x_{\phi(n)} \pmod{n}$ (commutative property)

and so

$$a^{\phi(n)}(x_1x_2\dots x_{\phi(n)}) \equiv x_1x_2\dots x_{\phi(n)} \pmod{n}$$

Because $gcd(x_i, n) = 1$ for each i, the property of coprime sharing implies that $gcd(x_1x_2...x_{\phi(n)}, n) = 1$. Therefore, we may divide both sides of the foregoing congruence by the common factor $x_1x_2...x_{\phi(n)}$, leaving us with

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

1.6.5 Order of an integer modulo n

Lemma 1.6.11. The linear congruence

$$ax \equiv b \pmod{n}$$

has a solution if and only if $gcd(a, n) \mid b$.

Proof. (\Rightarrow) Suppose x = t is a solution. Then $at \equiv b \pmod{n}$.

By difference property, at - b = kn for some k. $\Rightarrow at - kn = b$.

Since gcd(a, b) divides both a and n, it divides at - kn = b (property of linear combination).

 (\Leftarrow) Suppose $gcd(a, n) \mid b$. Then $b = k \cdot gcd(a, n)$ for some k.

Also, by Bezout's identity, there exists integers s, t such that $as + nt = \gcd(a, n)$.

- $\Rightarrow \quad k(as+nt) = k \cdot \gcd(a,n) \quad \Rightarrow \quad a(ks) + n(kt) = b \quad \Rightarrow \quad a(ks) b = n(-kt)$
- $\Rightarrow a(ks) \equiv b \pmod{n}$ (difference property)
- $\Rightarrow x = ks$ is a solution to the linear congruence $ax \equiv b \pmod{n}$.

Theorem 1.6.12 (Condition for existence of multiplicative order). Given a and n, there exists k such that

$$a^k \equiv 1 \pmod{n}$$

if and only if a and n are coprime.

Proof. (\Rightarrow) Suppose $a^k \equiv 1 \pmod{n}$ for some k. Write $aa^{k-1} \equiv 1 \pmod{n}$. By previous lemma, since this linear congruence has a solution (in this case, a^{k-1}), then $\gcd(a,n)$ divides 1, which means $\gcd(a,n)=1$, i.e. a and n are coprime.

 (\Leftarrow) If a and n are coprime, then by Euler's totient theorem (Theorem 1.6.10), we have

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

, so $k = \phi(n)$ is always a solution to the congruence.

Definition 1.6.2. Let n > 1 and gcd(a, n) = 1. The (multiplicative) **order** of a modulo n is the smallest positive integer k such that $a^k \equiv 1 \pmod{n}$.

If the modulus n is clear, we may simply say the order of a.

The previous theorem ensures that such k always exists for coprime a, n, and we can leave the order undefined for gcd(a, n) > 1.

Theorem 1.6.13. Let the integer a have order k modulo n. Then $a^h \equiv 1 \pmod{n}$ if and only if $k \mid h$.

This theorem tells us that when we multiply by a repeatedly, only the positive powers of a^k will be congruent to 1 (mod n).

Proof. (\Rightarrow) Suppose that we begin with $k \mid h$, so that h = jk for some integer j. Because $a^k \equiv 1 \pmod{n}$, power property yields $(a^k)^j \equiv 1^j \pmod{n}$ or $a^h \equiv 1 \pmod{n}$.

(\Leftarrow) Conversely, let h be any positive integer satisfying $a^h \equiv 1 \pmod{n}$. By the division algorithm, there exists q and r such that h = qk + r, where $0 \le r < k$. Consequently,

$$a^h = a^{qk+r} = (a^k)^q a^r$$

By hypothesis, both $a^h \equiv 1 \pmod{n}$ and $a^k \equiv 1 \pmod{n}$, the implication of which is that $a^r \equiv 1 \pmod{n}$. Because $0 \le r < k$, we end up with r = 0; otherwise, the choice of k as the smallest positive integer such that $a^k \equiv 1 \pmod{n}$ is contradicted. Hence, h = qk, and $k \mid h$.

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Corollary 1.6.14. Let the integer a have order k modulo n. Then $k \mid \phi(n)$.

Proof. $a^{\phi(n)} \equiv 1 \pmod{n}$ by Euler's totient theorem, so the order $k \mid \phi(n)$ by previous theorem. \square So when attempting to find the order of a, we only need to consider the divisors of $\phi(n)$.

Theorem 1.6.15. If the integer a has order k modulo n, then $a^i \equiv a^j \pmod n$ if and only if $i \equiv j \pmod k$.

Proof. First, suppose that $a^i \equiv a^j \pmod{n}$, where $i \geq j$. Because a is coprime to n, we may cancel a power of a to obtain $a^{i-j} \equiv 1 \pmod{n}$. By previous theorem, this congruence holds only if $k \mid i-j$, which is just another way of saying $i \equiv j \pmod{k}$.

Conversely, let $i \equiv j \pmod{k}$. Then we have i = j + qk for some integer q. By the definition of k, $a^k \equiv 1 \pmod{n}$, so that

$$a^i \equiv a^{j+qk} \equiv a^j (a^k)^q \equiv a^j \pmod{n}$$

which is the desired conclusion.

1.7 Repeating decimal

1.7.1 Basics

Definition 1.7.1. A decimal representation of a non-negative real number α is the concatenation of a sequence of digits $b_i \in \{0, \dots, 9\}$ and $a_i \in \{0, \dots, 9\}$, separated by the decimal point.

$$\alpha = b_k b_{k-1} \dots b_0 \cdot a_1 a_2 \dots$$

for which the following infinite (or finite) sum holds:

$$\alpha = \sum_{i=0}^{k} b_i 10^i + \sum_{i=1}^{\infty} \frac{a_i}{10^i}$$

The sequence before the decimal point $[b_i]$ is called the **integer part**, and its length is finite. The sequence after decimal point $[a_i]$ is called the **fractional part**, and its length may be finite or infinite.

Example: $\sqrt{2} = 1.41421356..., 1 = 0.999...$

Note: The presence of the decimal point . implies that $b_k b_{k-1} \dots b_0 \cdot a_1 a_2 \dots$ is read as the concatenation of the variables instead of their product.

Definition 1.7.2. A **repeating decimal** is a decimal representation of a number whose digits are eventually periodic ^a, and the repeating digits are neither only '0' nor only '9'. ^b

If the repeating digits consists of only '0' (or equivalently there are a finite amount of digits after decimal point), the decimal is a **terminating decimal**.

Definition 1.7.3. Notation: The fractional part of a repeating decimal can be written as:

$$0.a_1...a_m\overline{a_{m+1}...a_{m+p}}$$

The minimum p for which the repeating decimal can be written in this form is called the (decimal) **period** of the repeating decimal (denoted T), and the repeating digits $\overline{a_{m+1}...a_{m+T}}$ is called the **reptend**.

For terminating decimal, we define the period to be zero.

Example of repeating decimal:

$$\frac{1}{7} = 0.\overline{142857} = 0.142857142857\dots; \text{ period} = 6, \quad \frac{1}{6} = 0.1\overline{6} = 0.16666\dots; \text{ period} = 1.$$

Example of terminating decimal: $\frac{1}{2} = 0.5$, $\frac{1}{8} = 0.125$, both period = 0.

Theorem 1.7.1. Every rational number can be expressed as either a terminating or repeating decimal, and this representation is unique.

Proof. Let the rational number be n/d. We can use long division to get its decimal representation:

By division algorithm, given n and d we can uniquely write $n = bd + r_0$ where $0 \le r_0 < d$ and $b \ge 0$ are integers.

Then $\frac{n}{d} = \frac{bd + r_0}{d} = b + \frac{r_0}{d}$, which means b is the integer part of the decimal representation.

^aThat is, after some place, the same sequence of digits is repeated forever.

 $[^]b\mathrm{I}$ exclude 'trivial repeating decimals' like 1.4999... for simplicity, since its value equals a terminating decimal.

Now consider the fractional part r_0/d of the rational number. Write

$$\frac{r_0}{d} = \frac{1}{10} \left(\frac{10r_0}{d} \right)
= \frac{1}{10} \left(\frac{a_1 d + r_1}{d} \right) \qquad (10r_0 = a_1 d + r_1 \text{ by division algorithm, } 0 \le r_1 < d \text{ and } a_1 \ge 0)
= \frac{a_1}{10} + \frac{1}{10} \left(\frac{r_1}{d} \right)$$

Note that $a_1d \leq 10r_0 < 10d \implies a_1d < 10d \implies a_1 < 10$, which means a_1 is a valid digit.

Now repeat the process for r_1/d :

$$\frac{r_1}{d} = \frac{a_2}{10} + \frac{1}{10} \left(\frac{r_2}{d}\right)$$

giving

$$\frac{r_0}{d} = \frac{a_1}{10} + \frac{1}{10} \left(\frac{a_2}{10} + \frac{1}{10} \left(\frac{r_2}{d} \right) \right)$$
$$= \frac{a_1}{10} + \frac{a_2}{10^2} + \frac{1}{10^2} \left(\frac{r_2}{d} \right)$$

Continue the process for successive remainders:

$$\frac{r_0}{d} = \frac{a_1}{10} + \frac{a_2}{10^2} + \frac{a_3}{10^3} + \dots = \sum_{k=1}^{\infty} \frac{a_i}{10^i}$$

If $r_k = 0$ for some k, then it's easy to see that $a_i = 0$ and $r_i = 0$ for all i > k, so we can terminate the process once we see the remainder is 0, and we get a terminating decimal $0.a_1a_2...a_k$.

Otherwise, since all the remainders r_i are less than d and larger than zero, only finitely many different remainders can occur. So after some point, we must have for the first time $r_k = r_j$ for $0 \le j < k$. Then $a_{k+1} = a_{j+1}$ and $r_{k+1} = r_{j+1}$ since the division algorithm gives the same quotient and remainder when given the same dividend (r_j) and divisor (d).

Inductively it follows that $a_{k+i}=a_{j+i}$ for all i>0, so the digits repeat with a period of k-j (since $a_{k+1+(k-j)}=a_{j+1+(k-j)}=a_{k+1}$). We get the repeating decimal $0 \cdot a_1 \dots a_j \overline{a_{j+1} \dots a_k}$.

As for why the above representation is unique, assume that there are two different decimal representations $0.a_1a_2...$ and $0.b_1b_2...$ for the same proper fraction. Subtract the two representations to get

$$0 = \frac{a_1 - b_1}{10} + \frac{a_2 - b_2}{10^2} + \frac{a_3 - b_3}{10^3} + \dots$$

Since the representations are different, there is a least k such that $a_k - b_k \neq 0$. Without loss of generality assume $a_k - b_k > 0$. Then for the equality to hold, we must have

$$0 = \frac{a_k - b_k}{10^k} + \sum_{i=k+1}^{\infty} \frac{a_i - b_i}{10^i}$$

$$\sum_{i=k+1}^{\infty} \frac{b_i - a_i}{10^i} = \frac{a_k - b_k}{10^k}$$

For LHS to sum to RHS which is at least $\frac{1}{10^k}$, we can only make each $b_i - a_i$ in LHS the largest possible, namely 9, since

$$\sum_{i=k+1}^{\infty} \frac{9}{10^i} = \frac{9}{10^{k+1}} + \frac{9}{10^{k+2}} + \frac{9}{10^{k+3}} + \dots$$

$$= \frac{\frac{9}{10^{k+1}}}{1 - \frac{1}{10}} \qquad \text{(sum of geometric sequence formula)}$$

$$= \frac{1}{10^k}$$

And the RHS can only be the smallest possible, $\frac{1}{10^k}$. This means $a_k = b_k + 1$ and $b_i = 9$, $a_i = 0$ for i > k, giving the only two possible representations:

$$0.0...0(b_k + 1)00000...$$
 and $0.0...0b_k 99999...$

This means if a proper fraction has two different decimal representations, then one must be a terminating decimal and the other must be a 'trivial repeating decimal'. But the 'trivial repeating decimal' $0.0...0b_k99999...$ is excluded from my definition of repeating decimals. So the only valid representation is the terminating decimal.

And since it is impossible for repeating decimals to have two different representations, the representation of repeating decimals is also unique.

Corollary 1.7.2. If a rational number with denominator d has a repeating decimal representation, then its period is at most d-1.

Proof. From the proof of above theorem, we know that the period is k-j where $r_k = r_j$ is the first instance of repetition of remainders (with j < k). Since there are only d-1 distinct remainders from 1 to d-1, the index of r_k and r_j can differ by at most d-1.

Lemma 1.7.3. Multiplying a number by 10^n shifts the decimal point n digits to the right, and multiplying a number by 10^{-n} shifts the decimal point n digits to the left. (The number is padded with '0' if it runs out of digits.)

Example: $10^4 \times 1.23 = 12300$., $10^{-3} \times 45.67 = 0.04567$, $10^5 \times 1.\overline{142857} = 114285.7\overline{142857}$.

Proof. Let $\alpha = a_k a_{k-1} \dots a_0 \cdot a_{-1} a_{-2} \dots = a_k 10^k + a_{k-1} 10^{k-1} + \dots$. Then

$$10^{n}\alpha = 10^{n}(a_{k}10^{k} + a_{k-1}10^{k-1} + \dots + a_{-n}10^{-n} + \dots)$$
$$= a_{k}10^{n+k} + a_{k-1}10^{n+k-1} + \dots + a_{-n} + \dots$$

So a_{-n} , which is originally n digit to the right of the unit digit a_0 , now becomes the new unit digit. This is equivalent to putting the decimal point n digit to the right.

And since multiplying by 10^{-n} undoes the operation, it puts the decimal point n unit to the left.

Theorem 1.7.4. Every repeating or terminating decimal is a rational number.

Proof. If the terminating decimal is $x = b \cdot a_1 a_2 ... a_n$ (where b may be any non-negative integer), then we can multiply x by 10^n to shift the decimal point all the way to the right and get:

$$10^{n} x = ba_{1}a_{2}...a_{n}.$$

$$x = \frac{ba_{1}a_{2}...a_{n}.}{10^{n}}$$

which is the ratio of two integers.

If the repeating decimal is $x = b \cdot a_1 \dots a_n \overline{d_1 \dots d_j}$, first let $c = 0 \cdot \overline{d_1 \dots d_j}$ and find its value. Multiply c by 10^j to shift the decimal point j digits to the right:

$$10^{j}c = d_{1} \dots d_{j} \cdot \overline{d_{1} \dots d_{j}}$$

$$10^{j}c = d_{1} \dots d_{j} \cdot + c$$

$$(10^{j} - 1)c = d_{1} \dots d_{j} \cdot$$

$$c = \frac{d_{1} \dots d_{j}}{10^{j} - 1}$$

Note that c is a rational number. Now we can multiply x by 10^n to separate the repeating and terminating groups:

$$10^{n} x = ba_{1} \dots a_{n} \cdot \overline{d_{1} \dots d_{j}}$$

$$10^{n} x = ba_{1} \dots a_{n} \cdot + c$$

$$x = \frac{ba_{1} \dots a_{n} \cdot + c}{10^{n}}$$

Since all the components in the fraction are rational, x is also a rational number.

This implies by contraposition that if x is irrational, then it doesn't have a terminating or repeating decimal representation. For example, since $\sqrt{2}$ is irrational, its digits must go on forever without repeating.

Theorem 1.7.5 (Condition for terminating decimal). A reduced fraction has a terminating decimal representation if and only if its denominator is of the form $2^n 5^m$ for non-negative integers n, m.

Proof. If x is terminating decimal, then we can write $x = \frac{a}{10^n}$ where the denominator is of the form $10^n = 2^n 5^n$.

Conversely, if the denominator of x is of the form $2^n 5^m$, then $x = \frac{a}{2^n 5^m}$ for some a and we can write

$$x = \frac{2^m 5^n a}{2^{n+m} 5^{n+m}} = \frac{2^m 5^n a}{10^{n+m}}$$

Since $2^m 5^n a$ is an integer, it has a finite amount of digits, and dividing it by 10^{n+m} only shifts the decimal point to the left by n+m which is also a finite amount. So $x=\frac{2^m 5^n a}{10^{n+m}}$ has a terminating decimal representation.

Lemma 1.7.6. Let a number r have a decimal period of T. Then r can be written in the form

$$b.\overline{a_1a_2\dots a_p}$$

if and only if $T \mid p$ and r can be written as $b \cdot \overline{a_1 \dots a_T}$.

Proof. (\Rightarrow) Suppose $r = b \cdot \overline{a_1 a_2 \dots a_p}$ has a period of T. First note that $T \leq p$ since p is a valid cycle length. Since the decimals have a cycle of length p, we have $a_i = a_j$ whenever $i \equiv j \pmod{p}$. Also, since the period is T, after some point we also have $a_i = a_j$ whenever $i \equiv j \pmod{p}$. In other words, after some point, for any i we have $a_i = a_{i+kT}$ for $k = 0, 1, 2, \ldots$

Combining the two, if $i + kT \equiv j \pmod{p}$ then $a_{i+kT} = a_j = a_i$. Taking i to be a multiple of p (say mp), the condition simplifies to

$$Tk \equiv j \pmod{p}$$

. Recall Lemma 1.6.11, which says that the linear congruence $ax \equiv b \pmod{n}$ has a solution if and only if $gcd(a, n) \mid b$. In other words, $ax \mod n$ can only be a multiple of gcd(a, n) and it could be any multiple of gcd(a, n) in the range from 0 to n - 1.

In this case (k is the variable of the congruence), $Tk \mod p$ will take on all multiples of $\gcd(T,p)$ from 0 to p-1. This means $a_{mp+l\cdot\gcd(T,p)}=a_{l\cdot\gcd(T,p)}=a_{mp}$ for $l=0,1,2,\ldots$ Simplifying, we get $a_{l\cdot\gcd(T,p)}=a_p$ for $l=1,2,\ldots$ But since the digits cycle every p steps and every T steps, we can shift the index to the right by 1 (take a_{mp+1} to be the new a_{mp}) and repeat the argument to conclude $a_{l\cdot\gcd(T,p)+1}=a_{mp+1}=a_1$ for $l=1,2,\ldots$, and $a_{l\cdot\gcd(T,p)+2}=a_{mp+2}=a_2$ for $l=1,2,\ldots$ and so on. This means the digits actually repeat every $\gcd(T,p)$ steps, giving $T\leq\gcd(T,p)$ by definition of period. But $\gcd(T,p)\leq T$ by definition of GCD, so we must have $T=\gcd(T,p)$. The only way this is true is by T being a divisor of p.

From above we also know that reptend starts right after the decimal point (because $a_1 = a_{\gcd(T,p)+1} = a_{2\gcd(T,p)+1} = \ldots$, $a_2 = a_{\gcd(T,p)+2} = a_{2\gcd(T,p)+2} = \ldots$), and so on.

 (\Leftarrow) Suppose r can be written as $b \cdot \overline{a_1 \dots a_T}$ where T is period, and p is a multiple of T. Then we can write p = mT for some m, and clearly

$$b.\overline{a_1 \dots a_T} = b.\underline{a_1 \dots a_T a_1 \dots a_T \dots a_1 \dots a_T} = b.\overline{a_1 \dots a_p}$$

$$\overline{a_1 \dots a_T} \text{ repeated } m \text{ times}$$

1.7.2 Decimal period of 1/d

(We assume d > 1 in the following subsubsection.)

Theorem 1.7.7 (Period of 1/d). Given an integer d > 1 that is coprime to 10, the decimal period of $\frac{1}{d}$ is the least positive integer n for which

$$10^n \equiv 1 \pmod{d}$$

Moreover, $\frac{1}{d}$ has the decimal form $0.\overline{a_1 \dots a_T}$ where T is the period.

In other words, the period of 1/d is the multiplicative order of 10 (mod d), and the reptend starts right after the decimal point.

Note that if d contains 2 or 5 as a factor, then by Theorem 1.6.12 (condition for multiplicative order), the congruence $10^n \equiv 1 \pmod{d}$ will never hold for any n.

Proof. By Theorem 1.6.12, if d is coprime to 10 there always exists n such that $10^n \equiv 1 \pmod{d}$. So suppose n is the least positive integer satisfying this congruence. Then we can write $10^n - 1 = ad$ for some integer a. Then we have

$$\frac{1}{d} = \frac{a}{10^n - 1}$$

$$= \frac{\frac{a}{10^n}}{1 - \frac{1}{10^n}}$$

$$= \frac{a}{10^n} + \frac{a}{10^{2n}} + \frac{a}{10^{3n}} + \dots$$
 (sum of geometric sequence in reverse)

Now prepend '0's to the left of a until it has n digits. (If a already has n digits then do nothing. a cannot have more than n digits since $a < 10^n - 1$.) Call this new string a^* . Then $\frac{a}{10^n} + \frac{a}{10^{2n}} + \frac{a}{10^{3n}} + \dots$ is exactly the repeating decimal $0.\overline{a^*}$.

For example, if a = 76923 and n = 6, then the string a^* is '076923', and we have:

$$\frac{a}{10^n} + \frac{a}{10^{2n}} + \frac{a}{10^{3n}} + \dots = \frac{76923}{10^6} + \frac{76923}{10^{12}} + \frac{76923}{10^{18}} + \dots$$

$$= 0.076923$$

$$+ 0.000000 \ 076923$$

$$+ 0.000000 \ 000000 \ 076923 + \dots$$

$$= 0.076923 \ 076923 \ 076923 \dots$$

$$= 0.\overline{076923}$$

$$= 0.\overline{a^*}$$

Thus $\frac{1}{d}$ has the form $0.\overline{a^*}$ and the period of the repeating decimal representation is n.

The period cannot be any shorter, because if the period were m which is less than n, then we could use the above arguments in reverse to show that $10^m \equiv 1 \pmod{d}$, which contradicts the assumption that n is the least positive integer satisfying this congruence.

Corollary 1.7.8 (Period divides totient). If gcd(d, 10) = 1, then the period of $\frac{1}{d}$ divides $\phi(d)$.

Proof. The period T is the multiplicative order of 10 modulo d, and by Corollary 1.6.14, $T \mid \phi(d)$. \square

Corollary 1.7.9 (Cyclic numbers are primes). If the period of $\frac{1}{d}$ is d-1, then d is a prime number.

These primes whose reciprocal's period are one less that itself are called **cyclic numbers**.

Proof. Note that $\phi(d) = d - 1$ if and only if d is a prime, because if d were composite, then d has a factor 1 so <math>d and p are not coprime, which means $\phi(d)$ is less than d - 1.

By previous corollary, the period T of $\frac{1}{d}$ divides $\phi(d)$, which means $T \leq \phi(d)$. If d is composite then $\phi(d) < d-1$ which means T < d-1. The corollary follows from the contrapositive of this statement (If T = d-1 then d is not composite, i.e. prime).

Theorem 1.7.10 (Period of multiplies of 1/d). Given an integer d > 1 that is coprime to 10 and $1 \le m < d$, the decimal period of $\frac{m}{d}$ does not exceed the period of $\frac{1}{d}$. If m and d are coprime, then the period of $\frac{m}{d}$ is equal to the period of $\frac{1}{d}$. (Otherwise it may be less than or equal.)

Proof. Let 1/d have period n and $1 \le m < d$. Then by previous theorem, n is the least integer satisfying the congruence $10^n \equiv 1 \pmod{n}$, and we can write $10^n - 1 = ad$ for some positive integer a. Then we have

$$\frac{1}{d} = \frac{a}{10^n - 1}$$

$$\frac{m}{d} = \frac{ma}{10^n - 1}$$

$$= \frac{ma}{10^n} + \frac{ma}{10^{2n}} + \frac{ma}{10^{3n}} + \dots$$

Since $1 \le m < d$, we have $ma < ad = 10^n - 1 < 10^n$, which means ma cannot exceed n digits. So it is valid to prepend '0' (if any) to the left of ma to form the n-digit string a_m^* like before, and obtain the repeating decimal:

$$\frac{m}{d} = 0.\overline{a_m^*}$$

whose period is at most n.

[6] (P.200) For the second statement of the theorem, we impose an extra condition that gcd(m, d) = 1. Suppose to the contrary that m/d has period k < n. Then it must have the form $0.\overline{a_1 \dots a_k}$ (by Lemma 1.7.6). Using the geometric series argument above in reverse, we have

$$\frac{m}{d} = \frac{a'}{10^k - 1}$$

for some integer $1 \le a' < 10^k - 1$. Cross multiplying gives $m(10^k - 1) = da'$. In particular, d divides $m(10^k - 1)$. But because m and d are coprime, d divides $10^k - 1$ by Euclid's lemma. But this is a contradiction, since n is assumed to be the least integer for which n divides $10^n - 1$. The conclusion is therefore that the period of m/d must equal n.

(I thought about this for two days and it turns out Euclid's lemma is the key to the proof.)

Theorem 1.7.11. If integer d is of the form $q \, 2^n 5^m$ where $2^n 5^m > 1$ and q > 1 is coprime to both 2 and 5, then the decimal period of $\frac{1}{d}$ is equal to the period of $\frac{1}{q}$, and there are non-repeating digits after the decimal point (i.e. digits not part of the reptend).

In other words, the factor 2 and 5 in the denominator have no effect on the period.

Proof. We can always separate the fraction into two components to isolate the denominator $2^n 5^m$. First, since q and $2^n 5^m$ are coprime, by Bezout's identity there exists integers x and y such that $2^n 5^m x + qy = 1$. So we can write:

$$\frac{1}{q \, 2^n 5^m} = \frac{2^n 5^m x + qy}{q 2^n 5^m}$$
$$\frac{1}{q \, 2^n 5^m} = \frac{x}{q} + \frac{y}{2^n 5^m}$$

Note that $\frac{x}{q}$ is a non-integer since gcd(x,q) = 1 (otherwise we can factor out their common prime factor p in $2^n 5^m x + qy = 1$ which implies p divides 1, a contradiction). Similarly $\frac{y}{2^n 5^m}$ is non-integer.

By Theorem 1.7.5 (condition for terminating decimal), $\frac{y}{2^n 5^m}$ is a terminating decimal and $\frac{x}{q}$ is a repeating decimal. So adding $\frac{y}{2^n 5^m}$ to $\frac{x}{q}$ affects only a finite number of digits after decimal point, but it must affect at least one digit or else we would be adding a whole number which is a contradiction $(\frac{y}{2^n 5^m})$ must be non-integer). So there must be a non-repeating part before the reptend.

Theorem 1.7.12 (Period of product of denominators). Suppose $gcd(d_1, 10) = gcd(d_2, 10) = gcd(d_1, d_2) = 1$. If $\frac{1}{d_1}$ have period T_1 and $\frac{1}{d_2}$ have period T_2 , then the period of $\frac{1}{d_1d_2}$ is $lcm(T_1, T_2)$.

Example:
$$\frac{1}{7} = 0.\overline{142857}$$
; period = 6, $\frac{1}{41} = 0.\overline{02439}$; period = 5.
$$\frac{1}{7 \times 41} = \frac{1}{287} = 0.\overline{003484320557491289198606271777}$$
; period = 30 = lcm(6, 5).

Proof. [6](P.202-203) Let T_3 be the period of $\frac{1}{d_1d_2}$. We will show that $T_3 = \text{lcm}(T_1, T_2)$ by showing LHS and RHS divide each other.

By definition, $T_1 \mid \text{lcm}(T_1, T_2)$ and $T_2 \mid \text{lcm}(T_1, T_2)$. So we can write $\text{lcm}(T_1, T_2) = kT_1 = mT_2$ for some k, m.

By Theorem 1.7.7 (period of 1/d), T_1, T_2 are the least positive integers such that

$$10^{T_1} \equiv 1 \pmod{d_1}$$
 and $10^{T_2} \equiv 1 \pmod{d_2}$

Multiply the congruences by k and m respectively, we get:

$$(10^{T_1})^k \equiv 1^k \pmod{d_1}$$
 and $(10^{T_2})^m \equiv 1^m \pmod{d_2}$
 $10^{\text{lcm}(T_1, T_2)} \equiv 1 \pmod{d_1}$ and $10^{\text{lcm}(T_1, T_2)} \equiv 1 \pmod{d_2}$

This means d_1 and d_2 both divide the same number $10^{\text{lcm}(T_1,T_2)} - 1$ (difference property). Since d_1 and d_2 are coprime, their product d_1d_2 also divides $10^{\text{lcm}(T_1,T_2)} - 1$ (property of divisor product), which is equivalent to:

$$10^{\operatorname{lcm}(T_1, T_2)} \equiv 1 \pmod{d_1 d_2}$$

Since T_3 is the least positive integer s.t. $10^{T_3} \equiv 1 \pmod{d_1 d_2}$ (Theorem 1.7.7), we have $T_3 \mid \text{lcm}(T_1, T_2)$ (by Theorem 1.6.13).

Now we show that $\operatorname{lcm}(T_1, T_2)$ divides T_3 . From $10^{T_3} \equiv 1 \pmod{d_1 d_2}$ we know $d_1 d_2 \mid 10^{T_3} - 1$, so this of course implies $d_1 \mid 10^{T_3} - 1$. Since $10^{T_3} \equiv 1 \pmod{1}$ (difference property) and T_1 is the least positive integer for $10^{T_1} \equiv 1 \pmod{1}$, by Theorem 1.6.13 again we have $T_1 \mid T_3$. By the same argument, $T_2 \mid T_3$ as well. Since both T_1 and T_2 divide T_3 , it follows that $\operatorname{lcm}(T_1, T_2) \mid T_3$ as well (LCM universal property).

So this gives our desired equality $T_3 = \text{lcm}(T_1, T_2)$.

2 Combinatorics

2.1 Integer partition

2.1.1 Basics

Definition 2.1.1. A partition of a positive integer n is a way of writing n as a sum of positive integers (where order does not matter). The total number of partitions of n is denoted p_n .

For example, all the partitions of 5 is

 $5 \\ 4+1 \\ 3+2 \\ 3+1+1 \\ 2+2+1 \\ 2+1+1+1 \\ 1+1+1+1+1$

We have $p_5 = 7$.

Definition 2.1.2. Let p(n, m) be the number of ways to partition n objects using parts up to m objects (called m-parts).

For example, p(5,2) = 3, as

$$5 = 2 + 2 + 1$$

= $2 + 1 + 1 + 1$
= $1 + 1 + 1 + 1 + 1$

We define p(0, m) = 1 for any m since there is only one way to partition zero objects: doing nothing.

We also define p(n,0) = 0 for $n \neq 0$ since there is no way to partition n using up to 0-parts.

Also, if n < 0, then define p(n, m) = 0. (Helpful for the recursive formula below.)

Also note that for $m \ge n$, $p(n, m) = p_n$. And $p_0 = 1$ using this definition.

Theorem 2.1.1 (Recursive formula for p(n, m)).

$$p(n,m) = \begin{cases} 1 & \text{if } n = 0 \\ 0 & \text{if } m = 0 \text{ or } n < 0 \\ p(n-m,m) + p(n,m-1) & \text{otherwise} \end{cases}$$

Proof. If n, m > 0,

All the ways to partition n using up to m-parts must contain all the ways to partition n using up to (m-1)-parts, because 'up to m-parts' includes 1-parts, 2-parts, ..., (m-1)-parts, and m-parts. Thus, p(n,m) includes p(n,m-1).

The rest of p(n, m) are partitions that use m-parts at least once. To count them, we can exclude one m-part from the partition and look at the remaining n-m objects. We can partition these n-m objects in p(n-m, m) ways. (If n < m, then it doesn't make sense to partition negative objects so set p(n-m, m) = 0 instead).

2.1.2 Formal power series

We can represent all the partition number p_n using generating function, which is a kind of formal power series.

[7] Loosely speaking, a formal power series is an infinite sum

$$\sum_{n=0}^{\infty} a_n x^n = a_0 + a_1 x + a_2 x^2 + \dots$$

where the x^n 's does not stand for anything and only serves to 'store' its coefficients.

In this article we only consider formal power series where $a_0 = 1$.

Product

Define a product of $1 + a_1x + a_2x^2 + \dots$ and $1 + b_1x + b_2x^2 + \dots$ by writing

$$(1 + a_1x + a_2x^2 + \dots)(1 + b_1x + b_2x^2 + \dots) = 1 + c_1x + c_2x^2 + \dots$$

where
$$c_k = a_k + a_{k-1}b_1 + a_{k-2}b_2 + \ldots + b_k$$
.

Note that this multiplication is well-defined because we can compute any particular element of the product in a finite amount of time.

Also, this multiplication is associative since to find out the coefficient of each x^n of the product ABC, we can consider only the terms up to x^n for each A, B, C and discard the rest, and they will multiply out like regular polynomials, which are associative.

Inverse

The inverse of a series $1 + a_1x + a_2x^2 + \dots$ is another series $1 + b_1x + b_2x^2 + \dots$ such that

$$(1 + a_1x + a_2x^2 + \ldots)(1 + b_1x + b_2x^2 + \ldots) = 1$$

Computing this product, we want to find b_i such that

$$1 + (a_1 + b_1)x + (a_2 + a_1b_1 + b_2)x^2 + (a_3 + a_2b_1 + a_1b_2 + b_3)x^3 + \dots = 1$$

Setting each coefficient of x_i to zero, we get a series of equations for the b_i , which have a unique inductive solution:

$$b_1 = a_1$$

$$b_2 = -a_1b_1 - a_2$$

$$b_3 = -a_1b_2 - a_2b_1 - a_3$$

$$\vdots$$

Since the inverse of $1 + a_1x + a_2x^2 + \dots$ is well-defined and unique, we can denote the inverse of $1 + a_1x + a_2x^2 + \dots$ by $\frac{1}{1 + a_1x + a_2x^2 + \dots}$.

Example:
$$\frac{1}{1-x} = 1 + x + x^2 + \dots$$
 since
$$(1-x)(1+x+x^2+\dots) = (1-x) + (x-x^2) + (x^2-x^3) + (x^3-x^4) + \dots = 1$$

2.1.3 Generating function representation

We can represent all the partition number p_n using the generating function below:

Theorem 2.1.2 (Partition generating function).

$$(1+x+x^2+\ldots)(1+x^2+x^4+\ldots)(1+x^3+x^6+\ldots)(\ldots) = \sum p_n x^n$$

Proof. Intuitively, this makes sense. Every term x^n is created by taking a combination of x^{k_1} from the first sum, x^{2k_2} from the second sum, x^{3k_3} from the third sum, and so on. This multiplies to create $x^n = x^{k_1+2k_2+3k_3+...+sk_s}$.

However, we can get more x^n for all the other ways we can choose $k_1 + 2k_2 + 3k + ... + sk_s = n$. Since a partition is uniquely identified by the number of 1-parts it contains (namely k_1) and number of 2-parts it contains (k_2) and so on, each way to choose the k_i s corresponds to a unique and different partition of n. Thus, the coefficient of x^n is the number of partitions of n.

Theorem 2.1.3 (Inverse form of partition generating function).

$$\frac{1}{(1-x)(1-x^2)(1-x^3)\dots} = \sum p_n x^n$$

Proof. This follows immediately from the previous formula by inverses of the series on the left side. \Box

The expanded form of the denominator is given by Euler's pentagonal number theorem.

2.1.4 Euler's pentagonal number theorem

Definition 2.1.3. The *n* th pentagonal number is given by
$$g_n = \frac{n(3n-1)}{2}$$
.

The sequence produced by $n = 1, -1, 2, -2, 3, \dots$ is $1, 2, 5, 7, 12, \dots$ and it is called generalized pentagonal number, which is the one that appears in the theorem.

Theorem 2.1.4 (Euler's pentagonal number theorem).

$$\prod_{n=1}^{\infty} (1 - x^n) = \sum_{k=-\infty}^{\infty} (-1)^k x^{k(3k-1)/2} = 1 + \sum_{k=1}^{\infty} (-1)^k (x^{k(3k-1)/2} + x^{k(3k+1)/2})$$

In other words,

$$(1-x)(1-x^2)(1-x^3)\dots = 1-x-x^2+x^5+x^7-x^{12}-x^{15}+x^{22}+x^{26}-\dots$$

Here is the combinatoric proof given by the American mathematician Franklin in 1881. It surprisingly involves little algebra, unlike Euler's original proof. And it is really mindblowing to me.

Proof. [7] The basic idea is that the series $(1-x)(1-x^2)(1-x^3)\dots$ can be interpreted as a sophisticated count of a certain restricted type of partitions. Let us begin with the following expression, which I have multiplied out a few terms:

$$(1+x)(1+x^2)(1+x^3)\dots = 1+x+x^2+2x^3+2x^4+3x^5+4x^6+5x^7+\dots$$

Looking back at theorem 2.1.2, we see that the product on the left is equal to

$$\sum_{k_1, k_2} x^{k_1 + 2k_2 + 3k_3 + \dots + sk_s}$$

but this time each k_i is either zero or one. This means the coefficient of x^n counts the number of partitions of n with distinct parts, since we can only choose to include zero or one of each k-part in the partition.

For example, the coefficient of x^7 is 5 because there are only five partitions of 7 with distinct parts, namely 7, 6+1, 5+2, 4+3, and 4+2+1.

We aren't quite interested in this series, but instead in

$$(1-x)(1-x^2)(1-x^3)\dots = \sum_{k_1,k_2,\dots} (-1)^{k_1+k_2+k_3+\dots+k_s} x^{k_1+2k_2+3k_3+\dots+sk_s}$$

So this time when we count partitions with distinct parts, we count partitions with an even number of terms positively, but partitions with an odd number of terms negatively. The coefficient of x^n is thus "the number of distinct partitions of n with an even number of terms, minus the number of distinct partitions of n with an odd number of terms."

According to the pentagonal number theorem, this product is

$$1 - x - x^2 + x^5 + x^7 - x^{12} - x^{15} + x^{22} + x^{26} - x^{35} - x^{40} + \dots$$

Notice in particular that the coefficient of x^n is usually zero. Franklin concentrated on that fact, and tried to understand why the number of distinct partitions of n with an even number of terms is usually exactly the same as the number of distinct partitions of n with an odd number of terms. And the explanation he gave is that you can pair up each even partition with a corresponding odd

partition. For example, there are 12 partitions of 11 into distinct parts, and we will see that the appropriate pairing is

$$10+1 \leftrightarrow 11$$

$$9+2 \leftrightarrow 8+2+1$$

$$8+3 \leftrightarrow 7+3+1$$

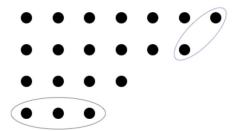
$$7+4 \leftrightarrow 6+4+1$$

$$6+5 \leftrightarrow 5+4+2$$

$$5+3+2+1 \leftrightarrow 6+3+2$$

How is this pairing defined?

Franklin's trick is exposed in the next picture. You could finish the proof without reading further by thinking carefully about this picture. (I can't.) Draw a distinct partition as a pattern of rows of dots (called Ferrers diagram); for instance the picture below corresponds to 20 = 7 + 6 + 4 + 3. Concentrate on the last row, and on the largest diagonal that can be drawn from the top right. The idea is to move the bottom row up to form a new diagonal, or move the diagonal down to form a new row. In the picture below, the bottom row cannot be moved up because that would leave a hanging dot, but the diagonal can be moved down. Notice that moving converts a partition with an even number of terms into a partition with an odd number of terms. In the case illustrated below, it converts 7 + 6 + 4 + 3 into 6 + 5 + 4 + 3 + 2.

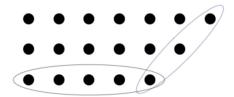


What is the rule for moving? Say the bottom row has a dots and the diagonal on the right has b dots. If we want to move the bottom row up without getting a hanging dot, we need $a \le b$. If we want to move the diagonal down and get a shorter final row, we need a > b.

So when $a \leq b$, we can move up, but not down, and when a > b we can move down but not up.

The one legal move for a given diagram produces a new diagram. This diagram also has a unique legal move. But certainly one thing we can do is to reverse the original move and return to the original diagram, so that must be the unique legal move. It follows that our diagrams are paired: the legal move for each leads to the other. Since the number of rows increases or decreases by one, even partitions are paired with odd partitions, as promised.

The only problem with this argument is that it seems to show that *all* of the coefficients of $(1-x)(1-x^2)(1-x^3)...$ are zero. The truth is that there are "edge cases" where the analysis just given doesn't quite work. These edge cases occur when the row at the bottom and the diagonal strip on the right side share a common corner, as below.



Let us again say that the bottom row has a dots and the diagonal on the right has b dots. If we want to move the bottom row up, then the diagonal length will decrease by one, and so to avoid a hanging dot we need $a \leq (b-1)$. If we want to move the diagonal down and get a shorter final row, then the existing final row length before the motion will decrease by one and we need (a-1) > b. So there are two troublesome cases where neither motion is legal: a = b and a = b + 1. Below are two samples, the first with a = b and the second with a = b + 1. Notice that neither motion is legal for

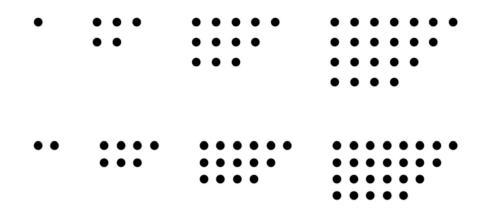
these diagrams. The left diagram corresponds to an odd partition of 12 and the right corresponds to an odd partition of 15. Note that the pentagonal number theorem expansion contains $-x^{12}-x^{15}$.



The first list below contains diagrams with a=b. The second list contains diagrams with a=b+1. Each list extends infinitely to the right. The number of dots in diagrams on the first list is $1,5,12,22,\ldots$ and the number of dots in diagrams on the second list is $2,7,15,26,\ldots$ and these are exactly the exponents which occur in the pentagonal number expansion!

We can check that the numbers of dots from the first list equal $\frac{3k(3k-1)}{2}$ by breaking each diagram into a rectangle of size $k \times (k-1)$ and a triangle of size $k \times k$. The total number of dots is therefore $k(k-1) + \frac{k(k+1)}{2} = \frac{2k(k-1)}{2} + \frac{k(k+1)}{2} = \frac{k(3k-1)}{2}$, as desired.

And the numbers from the second list equal $\frac{k(3k+1)}{2}$ since it is just k more than the corresponding number in the first list, and $\frac{k(3k-1)}{2} + k = \frac{k(3k-1)}{2} + \frac{2k}{2} = \frac{k(3k+1)}{2}$.



2.1.5 A more powerful recurrence relation

The pentagonal number theorem leads to a rapid method of computing the partition numbers using recurrence relations, which allows us to compute even p_{50000} in a reasonable amount of time (a few seconds).

Theorem 2.1.5 (Partition recurrence relation).

$$p_n = p_{n-1} + p_{n-2} - p_{n-5} - p_{n-7} + p_{n-12} + p_{n-15} - \dots$$

In other words,

$$p_n = \sum_{k \neq 0} (-1)^{k-1} p_{n-g_k}$$

where the summation is over all nonzero integers k (positive and negative) and g_k is the k th generalized pentagonal number.

Note that since $p_n = 0$ for all n < 0, the apparently infinite series on the right has only finitely many nonzero terms, enabling an efficient calculation of p_n .

Proof. Recall theorem 2.1.3:

$$\frac{1}{(1-x)(1-x^2)(1-x^3)\dots} = \sum p_n x^n$$

Rewrite the denominator using pentagonal formula and move it to other side of equation. We get:

$$(1-x-x^2+x^5+x^7-x^{12}-x^{15}+\ldots)(1+p_1x+p_2x^2+p_3x^3+\ldots)=1$$

Consequently the coefficient of x^n in the product is zero, and so

$$p_n - p_{n-1} - p_{n-2} + p_{n-5} + p_{n-7} - p_{n-12} - p_{n-15} + \dots = 0$$

Finally, move everything except p_n to right hand side, and we get the desired recurrence relation:

$$p_n = p_{n-1} + p_{n-2} - p_{n-5} - p_{n-7} + p_{n-12} + p_{n-15} - \dots$$

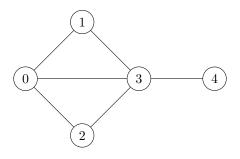
Each of these expressions is a finite sum because $p_0 = 1$ and $p_k = 0$ for negative k by definition. These formulas then allow us to compute the p_n inductively starting with the value $p_0 = 1$. Thus $p_1 = p_0 = 1$. Then $p_2 = p_1 + p_0 = 2$, etc.

3 Graph theory

(Related: Q107, Q83)

Definition 3.0.1. A graph G = (V, E) is a set of vertices and edges E where each edge (u, v) is a connection between vertices, where $u, v \in V$.

For example, in the following graph,



$$V = \{0, 1, 2, 3, 4\}$$

$$E = \{(0,1), (0,2), (0,3), (1,3), (2,3), (3,4)\}$$

3.1 Prim's algorithm

(Q107)[8]

Definition 3.1.1. G(V, E, w) is an **edge weighted graph** if there exists a weight function $w: E \to \mathbb{R}$ that assigns a weight to every edge $e \in E$.

Definition 3.1.2. G(V, E) is **connected** if there exists a path between any two vertices. G(V, E) is a **tree** if it is connected and has no cycles.

Definition 3.1.3. Let H(V', E') be a subgraph of G(V, E). We say that H is a **spanning tree** of G if H is a tree and V' = V.

In other words, a spanning tree of G is a tree that contains every vertex of G.

Definition 3.1.4. The weight of a graph is the sum of the weights of all edges.

Weight
$$=\sum_{e\in E}w(e)$$

Definition 3.1.5. Let G(V, E, w) be an edge-weighted graph where $w: E \to \mathbb{N}$. H(V, E') is a **minimum spanning tree** (MST) of G if it is a spanning tree with weights less than or equal to the weight of any other spanning tree of G, i.e., $\sum_{e \in E'} w(e) \leq \sum_{e \in E''} w(e)$ for all other spanning trees H'(V, E'') of G.

3.1.1 The procedure

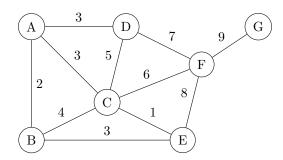
The objective of Prim's algorithm is to find the minimum spanning tree of a weighted undirected graph.

The algorithm may informally be described as performing the following steps:

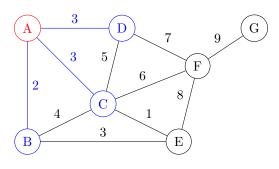
- 1. Initialize a tree with a single vertex, chosen arbitrarily from the graph.
- 2. Grow the tree by one edge: Of the edges that connect the tree to vertices not yet in the tree, find the minimum-weight edge, and transfer it to the tree.

3. Repeat step 2 (until all vertices are in the tree).

Example [9]

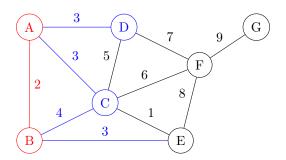


Let's start with vertex A. First add A to the visited list and examine all vertices reachable from A.



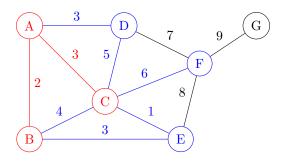
 ${\rm Visited} = \{A\}$

Choose the smallest edge that connects to an unvisited vertex, which is B. Add B to the visited list. Now look at all the vertices reachable from A and B.



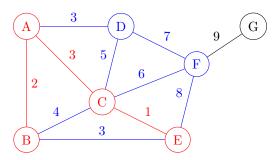
$$Visited = \{A, B\}$$

The three edges all have a weight of 3. Picking any one of these will work. Let's pick AC .

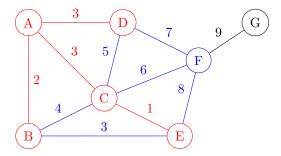


 $Visited = \{A, B, C\}$

Continue in this manner, each time picking the smallest edge that connects to an unvisited vertex.

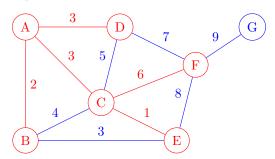


 $Visited = \{A, B, C, E\}$



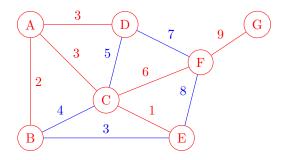
 $Visited = \{A, B, C, D, E\}$

Notice at this point the edge BE with weight of 3 is the smallest edge but both vertices are already in MST so we do not pick it. Instead we will choose to add F to the MST.



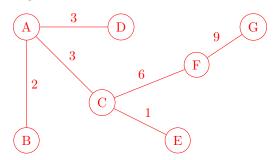
$$Visited = \{A, B, C, D, E, F\}$$

The only unvisited vertex remaining is G so we add this to the minimum spanning tree.



 ${\bf Visited}=\{A,B,C,D,E,F\}$

All the vertices are now connected in a tree. If we drop the blue edges, we get the minimum spanning tree with a total weight of 2+3+3+1+6+9=24.



3.1.2 Proof of optimality

- [8] Let's recall two facts about spanning trees.
 - 1. Let T be a spanning tree of G. If you add any edge $e \notin T$ to T then $T' = T \cup \{e\}$ contains a cycle. This is easy to see. Let e = (u, v) be an edge not in T. Since T is a spanning tree, there is already a path between any two vertices of G; in particular there is a path between u and v. Adding e to T will thus close the cycle from u to v.
 - 2. Consider $T' = T \cup \{e\}$ as defined above, and let \mathcal{C} be the cycle created by adding the edge e. If you remove any edge $e' \in \mathcal{C}$ from the cycle, you will get a new spanning tree of G; since removing an edge from a cycle will not disconnect the graph.

Now let's get into the proof.

Proof. Let T be the spanning tree return by the algorithm, and suppose there doesn't exist any MST of G consistent with T (which means T is not minimum spanning tree). Consider an optimal MST O of G:

```
Algorithm 1 Prim's Algorithm
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Input: A weighted graph G(V, E, w)
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Output: A spanning tree that minimizes $\sum_{e \in E'} w(e)$

1: $U \leftarrow V$

 $\triangleright U$ is the set of unvisited vertices

- 2: Pick an arbitrary start vertex $s \in V$
- 3: $T = \{s\}$
- $4:\ U \leftarrow U \backslash \{s\}$
- 5: while $U \neq \emptyset$ do
- 6: Choose $u \in U$ adjacent to a $v \in T$ such that w(u, v) is the smallest out of such vertices
- 7: $T \leftarrow T \cup \{u\}$
- 8: $U \leftarrow U \setminus \{u\}$
- 9: end while

return T

Let e = (x, y) be the first edge chosen by the algorithm that is inconsistent with any MST of G (which means e is not contained in any MST), and let T' be the subtree of T created by the algorithm just before the edge e was chosen. Let P_{xy} be the path between x and y in O. Let C be the cycle created in O from adding e to P_{xy} (Fact 1).

 P_{xy} must contain an edge e'=(a,b) that has an end point in T' and end point outside of T'. Why? Since otherwise, P_{xy} would be fully contained in T', and choosing e(x,y) next would create a cycle in $T' \cup e$, a contradiction to step 6 of the Algorithm.

Therefore both e(x,y) and e'(a,b) have an end point in T' and an end point outside of T'. Since the algorithm chose e instead of e', it follows that $w(e) \leq w(e')$. By Fact (2), we can break $\mathcal C$ by removing e'(a,b) and obtaining a new MST of G, call it O'. Notice that the total weight of O' is w(O') = w(O) + w(e) - w(e'). Since $w(e) \leq w'(e)$, it follows that $w(O') \leq w(O)$. Thus there exists an optimal MST of G, O', that contains the edge e(x,y) (which is a contradiction).

Therefore, in order to show that Prim's Algorithm does indeed produce an optimal MST for G, it suffices to repeat this argument for every new edge \hat{e} chosen by the algorithm, such that \hat{e} doesn't appear in any optimal solution.

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