

Graph Homomorphisms Through Graph Powers

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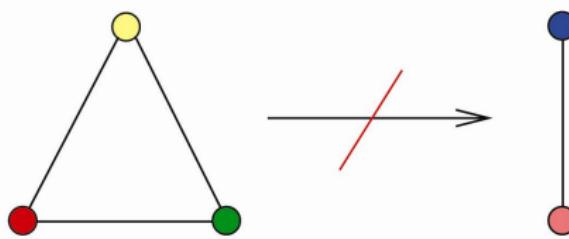
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GRAPHS HOMOMORPHISM

Graph Homomorphism

A **homomorphism** $f : G \rightarrow H$ from a graph G to a graph H is a map $f : V(G) \rightarrow V(H)$ such that if $uv \in E(G)$ then $f(u)f(v) \in E(H)$. Also, the existence of a homomorphism is indicated by the symbols $G \rightarrow H$.

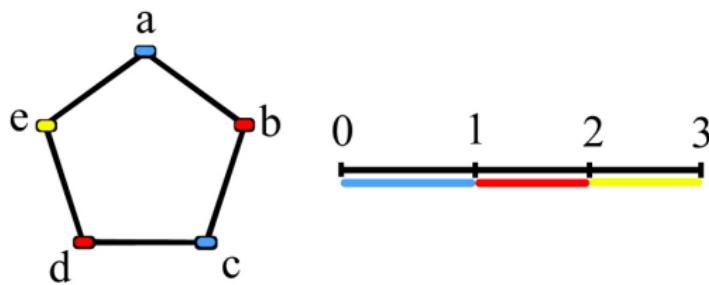
- ① What is \leq ? $G \leq H$ means that $G \rightarrow H$.
- ② What is $=$? $G \longleftrightarrow H$ ($[G] = [H]$) means that $G \rightarrow H$ and $H \rightarrow G$.
- ③ What is $<$? $G < H$ means that $G \rightarrow H$ and there is **no homomorphism** from H to G .



CHROMATIC NUMBER

Chromatic number

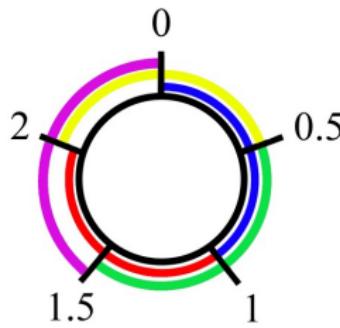
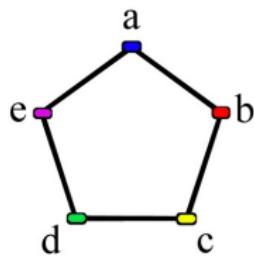
Let I be a line of length r . An r -coloring of a graph G is a mapping c which assigns to each vertex v of G an open unit length interval $c(v)$ of I , such that for every edge $uv \in E(G)$ we have $c(u) \cap c(v) = \emptyset$. The chromatic number of a graph, denoted by $\chi(G)$, is defined as, $\chi(G) = \inf\{r : G \text{ admits an } r - \text{coloring}\}$.



CIRCULAR CHROMATIC NUMBER

Circular Chromatic number

Let C be a circle of (Euclidean) length r . An r -circular coloring of a graph G is a mapping c which assigns to each vertex v of G an open unit length arc $c(v)$ of C , such that for every edge $uv \in E(G)$ we have $c(u) \cap c(v) = \emptyset$. The circular chromatic number of a graph, denoted by $\chi_c(G)$, is defined as, $\chi_c(G) = \inf\{r : G \text{ admits an } r - \text{circular coloring}\}$.



Andrew Vince

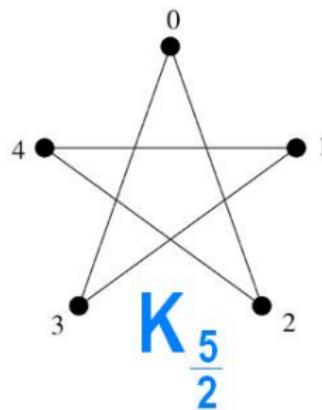


Xuding Zhu

CIRCULAR CHROMATIC NUMBER

Circular Complete Graphs

The circular complete graph $K_{\frac{p}{q}}$ has the vertex set $\{0, 1, \dots, p - 1\}$ and the edge set $\{ij : q \leq |i - j| \leq p - q\}$.



Circular chromatic number: $\chi_c(G) = \inf\left\{\frac{p}{q} \mid G \rightarrow K_{\frac{p}{q}}\right\}$

$$\chi(G) - 1 < \chi_c(G) \leq \chi(G)$$

POWER IN NUMBERS AND GRAPHS

Let G and H be two graphs and a, b, p, q, r and s be positive integers!

Power in Numbers:

- ① Integer Power and Root Power: $a^{\frac{p}{1}}$ and $a^{\frac{1}{q}}$

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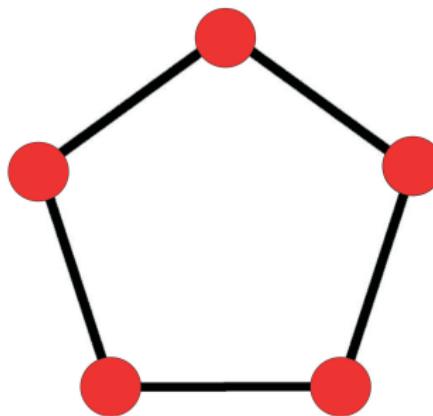
Power in Graphs:

- ① Integer Power and Root Power: $G^{\frac{p}{q}}$ and $G^{\frac{1}{q}}$
- ② Duality: $G \leq H^p$ if and only if $G^{\frac{1}{p}} \leq H$
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GRAPH POWERS

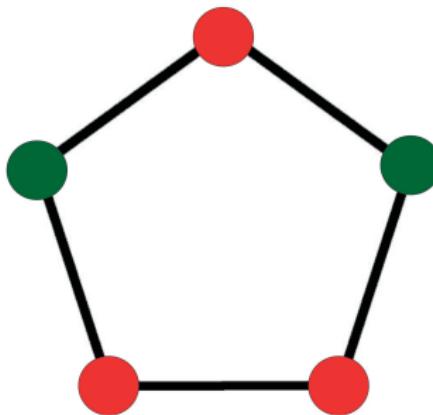
Graph Powers

For a graph G , let G^k be the k th power of G , which is obtained on the vertex set $V(G)$, by connecting any two vertices u and v for which there exists a walk of length k between u and v in G .



Graph Powers

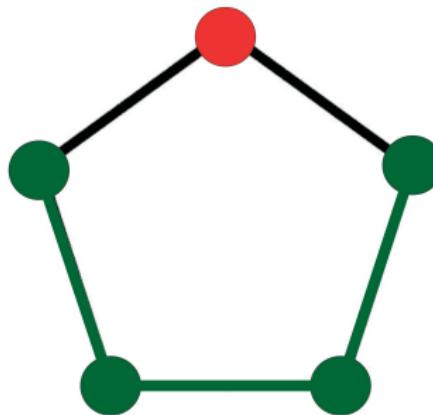
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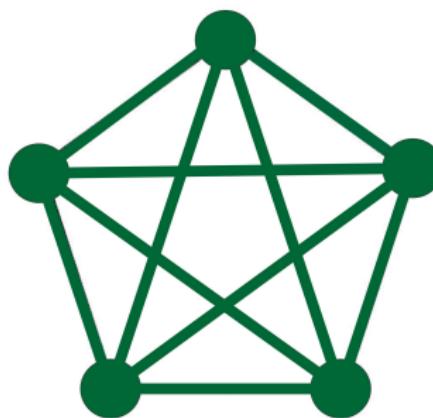
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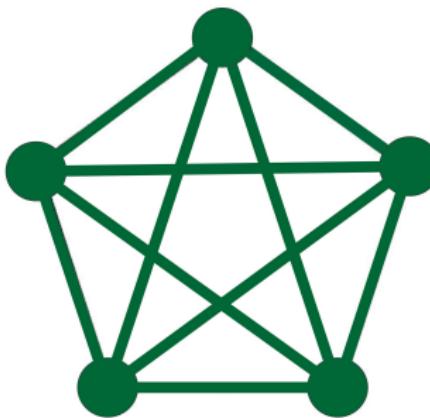
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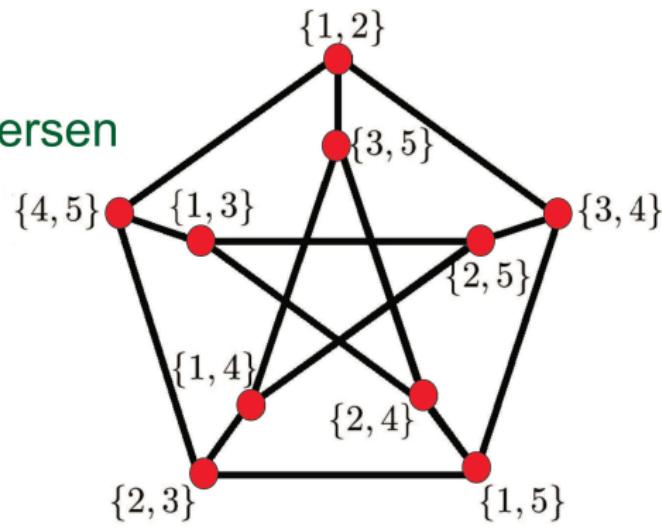
The adjacency matrix of G^k is obtained from the *kth power* of the adjacency matrix of G , by replacing any non-zero entries with one.

GRAPH POWERS

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The Petersen graph

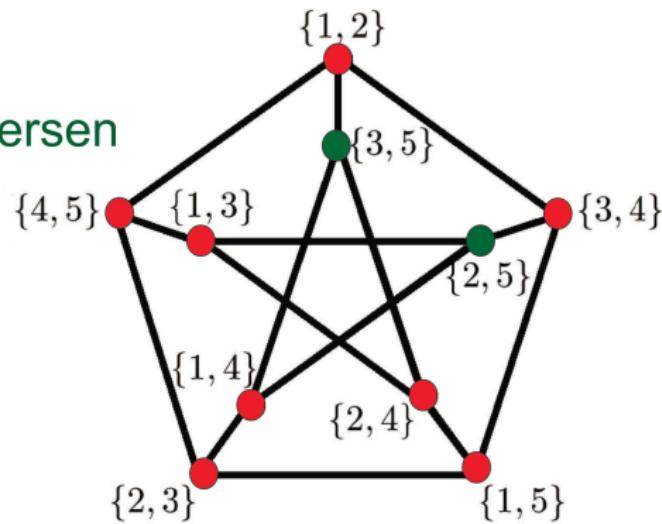


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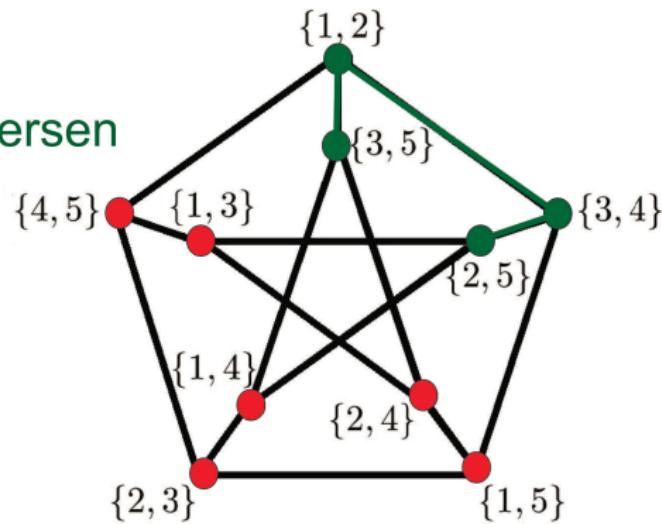


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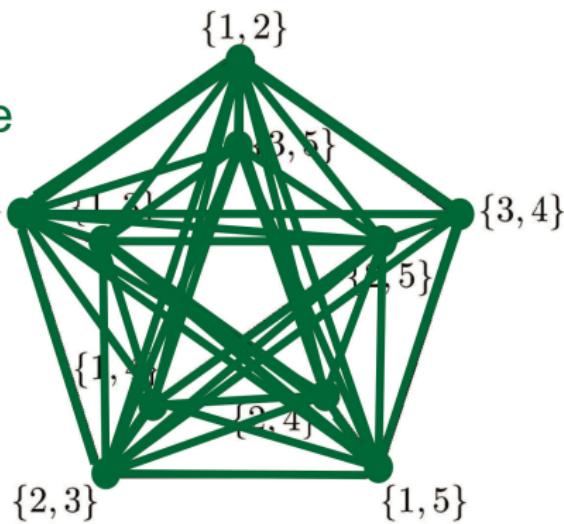


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The Complete
graph K_{10}



GRAPH POWERS

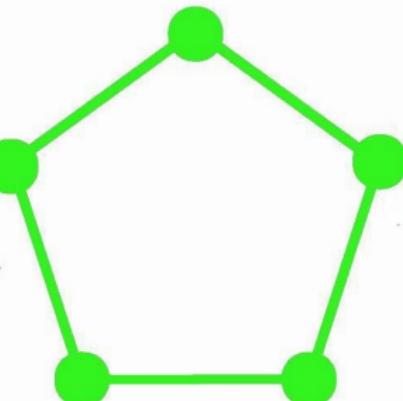
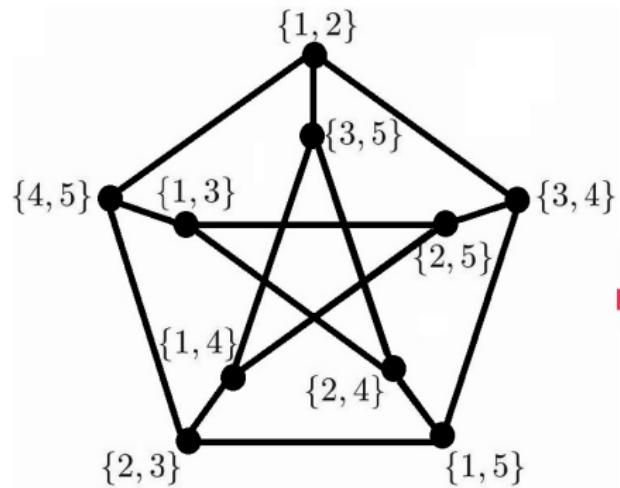
Observation.

Let G and H be two simple graphs such that $\text{Hom}(G, H) \neq \emptyset$. Then, for any positive integer k , $\text{Hom}(G^k, H^k) \neq \emptyset$.

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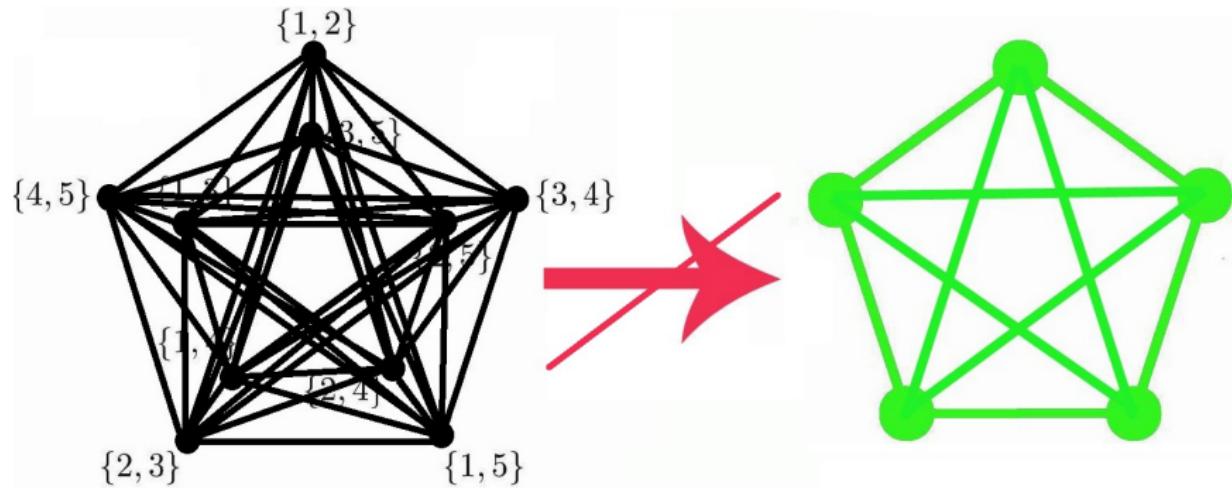
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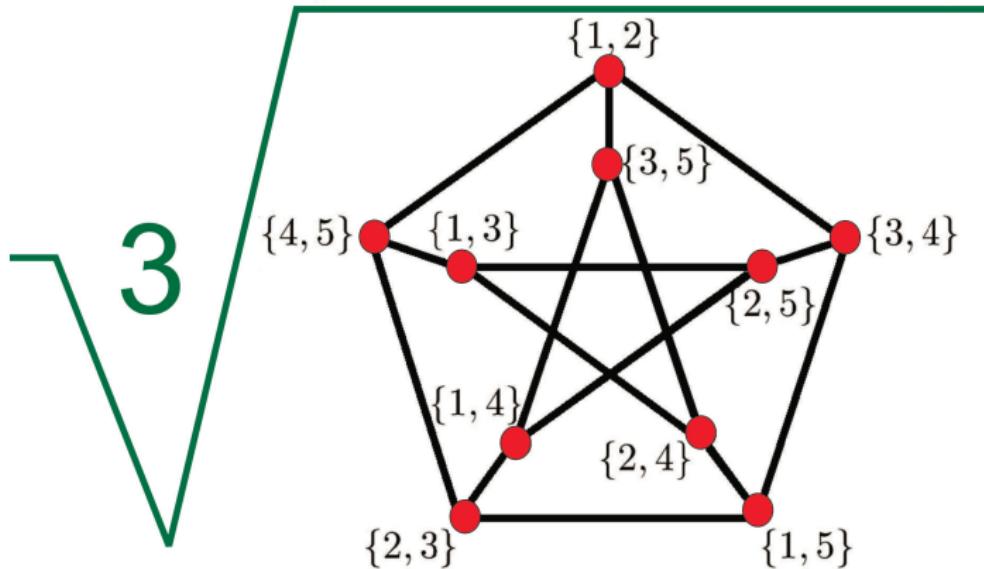
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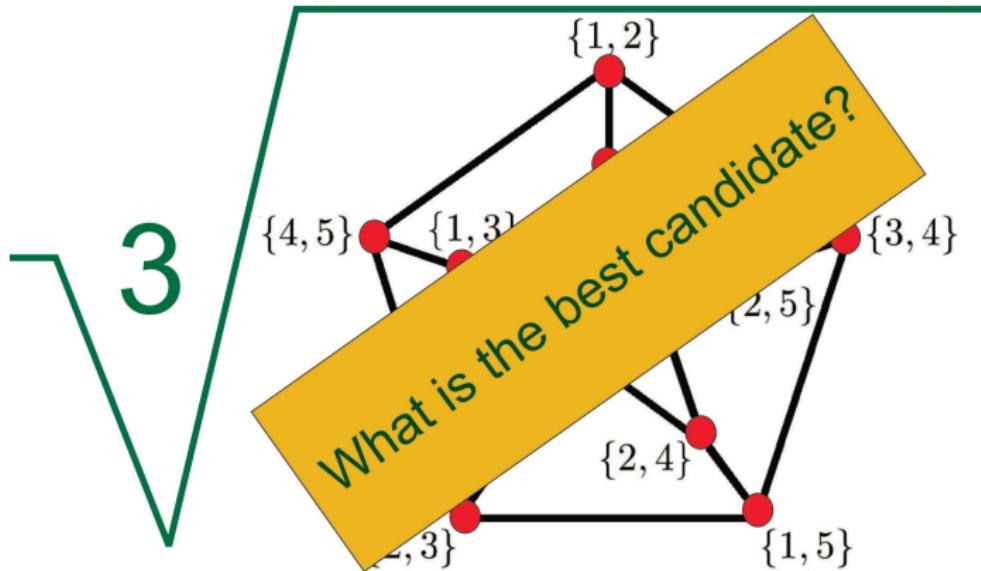
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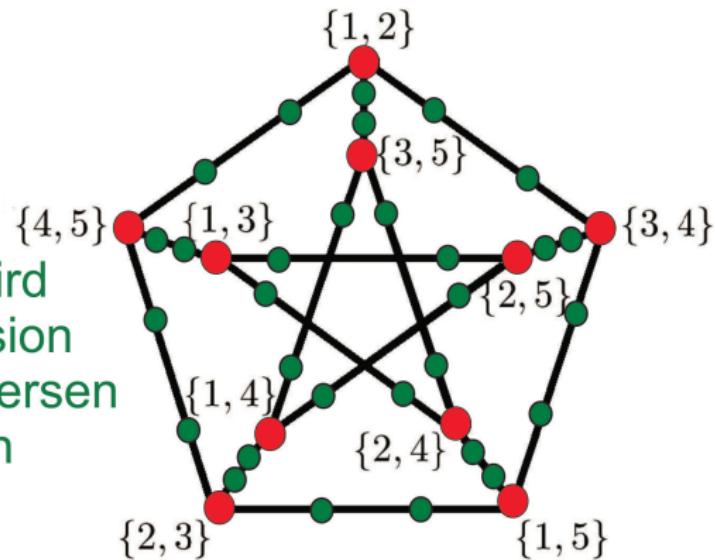


FRACTIONAL POWERS



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The third
subdivision
of the Petersen
graph



FRACTIONAL POWERS

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Lemma.

Let G and H be two graphs. Then, $G^{\frac{1}{2s+1}} \rightarrow H$ if and only if $G \rightarrow H^{2s+1}$ ($G^{\frac{1}{2s+1}} \leq H$ if and only if $G \leq H^{2s+1}$).

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- ③ A rational number r is called positive odd rational number if numerator and denominator are both positive odd integers.

PROPERTIES OF FRACTIONAL POWERS

H.H. and A. Taherkhani, 2010

Lemma. (Deletion Rule)

If a, b and s are positive odd integers, then $G^{\frac{sa}{sb}} \longleftrightarrow G^{\frac{a}{b}}$.

Observation.

Assume that r and s are positive odd rational number. It **does not hold**, in general, that $G^{rs} \longleftrightarrow (G^r)^s$.

Lemma.

If r and s are positive odd rational number, then $(G^r)^s \longrightarrow G^{rs}$.

Theorem. (The Same Base)

Let r and s be positive odd rational numbers and G be a **non-bipartite graph**. If $r < s < \text{og}(G)$, then $G^r < G^s$.

BLOW UP A GRAPH

Observation.

$\omega(G^r)$ stands for the clique number of G^r where r is a positive odd rational number and $r < og(G)$. One can see that if r approaches to $og(G)$ from below, then $\omega(G^r)$ approaches infinity.

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Observation.

Let G and H be non-bipartite graphs. Then there exists a positive odd rational number $r \geq 1$ such that $H \leq G^r$.



BLOW OUT A GRAPH

Observation.

For any positive integer $a \geq 3$, there exists a graph homomorphism from $G^{\frac{1}{a}}$ to C_a .

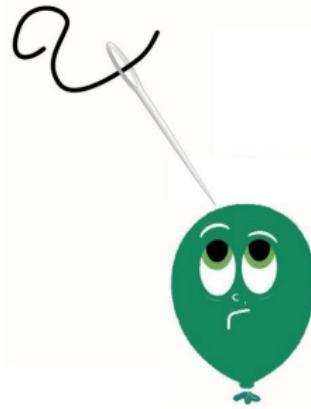
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ODD CYCLES

Theorem. (H.H., 2009)

For any graph G , $\chi(G^{\frac{2k+1}{3}}) \leq 3$ if and only if there exists a graph homomorphism from G to C_{2k+1} .

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Nešetřil's Pentagon Problem, 1999

Let G be a cubic graph of sufficiently large girth, is it true that $\chi(G^{\frac{5}{3}}) \leq 3$?

Jaeger's Conjecture, 1981

Every planar graph with girth at least $4k$ has a homomorphism to C_{2k+1} .

Jaeger's Conjecture, 1981

Let P be a planar graph with girth at least $4k$. Then we have $\chi(P^{\frac{2k+1}{3}}) \leq 3$.

POWER THICKNESS

Observation.

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Question.

If G is a non-bipartite graph, $\chi(G^{\frac{1}{3}}) = 3$ and $\lim_{r \rightarrow og(G)^-} \chi(G^r) = \infty$. Does it hold $\{\chi(G^r) | 0 < r < og(G), r \text{ is an odd } r. \text{ number}\} = \{3, 4, 5, 6, \dots\}$?

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Definition. Power Thickness of Graph

Let G be a **non-bipartite graph**. What is the value of

$$\Theta(G) := \sup\{r \mid r = \frac{2t+1}{2s+1}, \chi(G^r) = \chi(G), r < \text{og}(G)\}?$$

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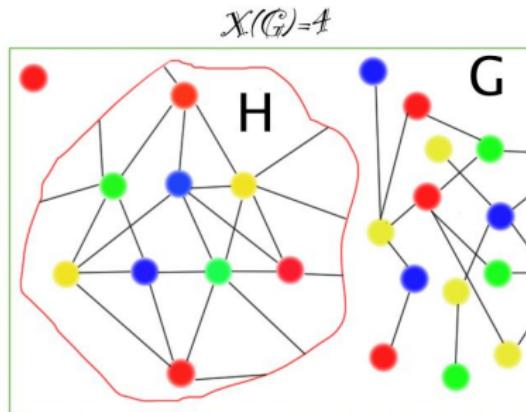
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- ③ For any rational number $r \geq 1$, there exists a graph G with $\theta(G) = r$.

COLORFUL GRAPHS

Definition.

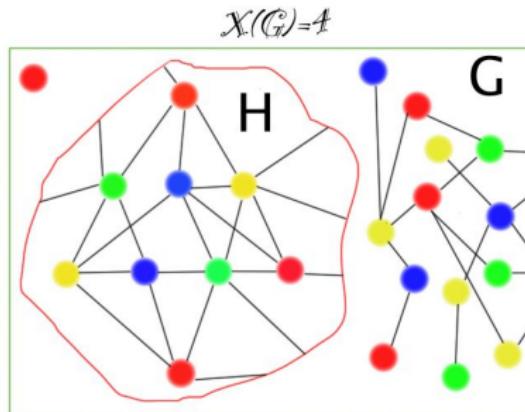
Let G be a graph and $\chi(G) = k$. G is called a **colorful graph** if for any proper k -coloring of G , there exists an induced subgraph H of G such that for any $v \in V(H)$, all colors appear in the closed neighborhood of v .



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Observation.

If G is a **uniquely colorable** graph or $\chi(G) = \omega(G)$, G is a **colourful graph**.
Find another example! Is the **Petersen graph** a colourful graph? Yes!

COLORFUL GRAPHS

Theorem. (H.H. and Ali Taherkhani, 2010)

If G is a non-bipartite colorful graph, then $\theta(G) = 1$. Also, $\chi(G) = \chi_c(G)$.

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Let $m > 2n$. Is the Kneser graph $\text{KG}(m, n)$ a colorful graph?

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Let $m > 2n$. Is the Kneser graph $\text{KG}(m, n)$ a colorful graph?

Theorem. (P.-A. Chen, 2011)

For any $m \geq 2n$, the Kneser graph $\text{KG}(m, n)$ is a colorful graph.

Furthermore, $\chi_c(\text{KG}(m, n)) = \chi(\text{KG}(m, n))$.

Later, G.J. Chang, D.D. Liu, and X. Zhu (2012) have presented a short and interesting proof for this result.

CIRCULAR COLORINGS

Let G be a **non-bipartite graph** and Q_0 be the set of **odd rational numbers**.

Observation.

For any non-bipartite graph G , we have $\chi(G^{\frac{1}{3}}) = 3$.

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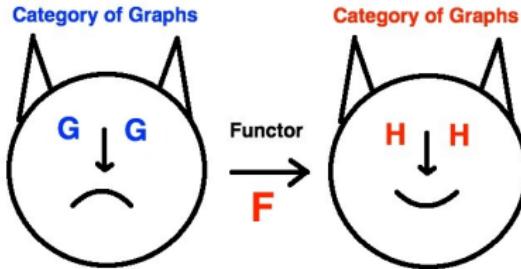
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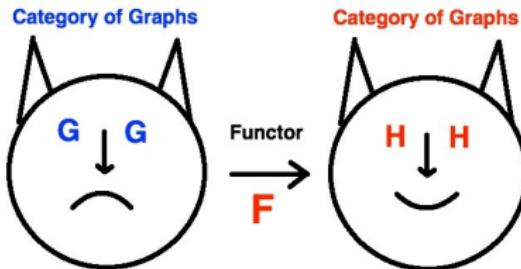
There exists a positive odd rational number $r > \frac{\chi(G)}{3(\chi(G)-2)}$ for which $\chi(G^r) = 3$ if and only if $\chi(G) \neq \chi_c(G)$.

CIRCULAR AND FRACTIONAL COLORINGS



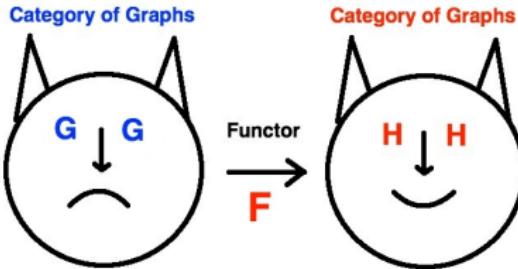
- ① It can be of interest if we have a **functor** or a map F such that one can compute $\chi_c(F(G))$ (resp. $\chi_f(F(G))$) in terms of $\chi_c(G)$ (resp. $\chi_f(G)$).

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- ③ For any graph G , we have

$$\chi_c(G^{\frac{1}{2s+1}}) = \frac{(2s+1)\chi_c(G)}{s\chi_c(G)+1}, \quad \chi_f(G^{\frac{1}{2s+1}}) \leq \frac{(2s+1)\chi_f(G)}{s\chi_f(G)+1}.$$

DUAL POWER

Theorem. (H.H. and A. Taherkhani, 2010)

Let G and H be two graphs. We have $G^{\frac{2a+1}{2b+1}} \rightarrow H$ if and only if
 $G \xrightarrow{\frac{2b+1}{2a+1}} H$.

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Let G and H be two graphs. We have $G^{\frac{2a+1}{2b+1}} \rightarrow H$ if and only if
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The **dual power** inherits (like power) some properties from power in number theory. But, for any graphs G and H , **there is no reason** that there exists a non-negative integer s such that $G^{\frac{1}{2s+1}} \leq H$.

Theorem. (H.H., 2009)

Let m, n , and k be positive integers with $m \geq 2n$. Then

$\chi(\text{KG}(m, n)^{\frac{1}{2k-1}}) = m - 2n + 2$. Moreover, if **m is an even positive integer**,
then $\chi_c(\text{KG}(m, n)^{\frac{1}{2k-1}}) = m - 2n + 2$.

HISTORY OF DUAL POWER

In 2004, A. Gyärfás, T. Jensen, and M. Stiebitz defined the graph $K_n^{\frac{1}{3}}$ to show that there is a n -chromatic graph with an n -coloring where the neighbors of each color-class form an independent set.



A. Gyärfás, T. Jensen, and M. Stiebitz

In 2005, C. Tardif defined the graph $G^{\frac{1}{3}}$ to show that the circular complete graph $K_{\frac{p}{q}}$, for $\frac{p}{q} < 4$, is a multiplicative graph, i.e., $G \times H \rightarrow K_{\frac{p}{q}} \implies G \rightarrow K_{\frac{p}{q}}$ or $H \rightarrow K_{\frac{p}{q}}$.



Claude Tardif

In 2006, G. Simonyi and G. Tardos showed that $K_n^{\frac{1}{2k+1}}$ is an n -critical graph.



Gabor Simonyi and Gabor Tardos

OPEN PROBLEMS

Problem. (Circular Chromatic Number of Planar Graphs)

In view of Four Color Theorem, one can conclude that for every planar graph P , we have $\chi(P^{\frac{3}{5}}) \leq 3$ which is equivalent to $\chi_c(P) \leq 4.5$. Prove it without Four Color Theorem!

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Question. (Power Thickness of Graph)

Let G be a non-bipartite graph. Set

$$\theta(G)_{\chi(G)-1} := \sup\{r \mid r = \frac{2t+1}{2s+1}, \chi(G^r) \leq \chi(G) - 1, r < og(G)\}.$$

Is it true that $\theta(G)_{\chi(G)-1} < 1$?

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Question. (Graph Homomorphism)

Suppose that G and H are two graphs. Also, assume that for any positive integer s , we have $G^{\frac{1}{2s+1}} \leq H^{\frac{1}{2s-1}}$. Is it true that $G \leq H$?

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