

1. Translate the following into predicate logic. You can assume that the domain is people, and so you don't need an additional predicate symbol for " x is a person".
 - (a) There is a person who loves all people who love her. (Use Lxy for " x loves y ".)
 - (b) Every lover loves herself. (A "lover" is somebody who loves at least one person.)
 - (c) There are exactly two people.

2. Could the following sentence be true? Explain your answer.

$$(\neg p \vee q) \wedge ((q \rightarrow (\neg r \wedge \neg p)) \wedge (p \vee r))$$

3. Explain what's wrong with the following attempted proof:

1	(1) Fa	A
	(2) $Fa \rightarrow Fa$	1,1 CP
	(3) $\forall y(Fy \rightarrow Fa)$	2 UI
	(4) $\exists x\forall y(Fy \rightarrow Fx)$	3 EI

4. Prove the following sequent. You can use cut/replacement, but only if you prove the relevant sequents in your exam booklet, and clearly cross-reference them.

$$\vdash \exists x\forall y(Fy \rightarrow Fx)$$

5. Give a rigorous, but informal, proof of the following fact of set theory:

$$C - (A \cap B) \subseteq (C - A) \cup (C - B)$$

Here we use the definition:

$$\forall x((x \in (C - X)) \leftrightarrow (x \in C \wedge x \notin X)).$$

6. Let A be the set defined inductively by:

- $p \in A$
- If $\phi \in A$ and $\psi \in A$ then $(\phi \rightarrow \psi) \in A$.

Show that for every $\phi \in A$, $\neg\phi \vdash \neg p$. [Spring 2019: we didn't do problems like this on a pset, so it's not likely that we would ask a question like this one.]

7. Provide a countermodel to show that the sentence on the left does *not* imply the sentence on the right. In your countermodel, you should explicitly specify a domain, and extensions for all the predicate symbols.

$$\exists x(Fx \rightarrow \exists yGy) \qquad \exists xFx \rightarrow \exists yGy$$