

logic precept 2

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Review pset1

- Minor problems with translation: never write a non-atomic sentence as if it were atomic
- Minor problems with proofs: each rule applies to a sentence type, and can only be applied to entire sentences of that type

E.g. DN applied to $P \rightarrow \neg\neg Q$ does not give $P \rightarrow Q$

E.g. MP cannot be applied to $P \vee (Q \rightarrow R)$

- The harder problems in part C

- (1) $P \vee (Q \wedge R)$ A
- (2) $P \vee Q$ 1 $\wedge E$

Line (2) is not a valid application of $\wedge E$. To apply $\wedge E$ to a line, the formula on that line needs to be a conjunction. But the formula on line (1) is a disjunction.

This new rule would be bad. Consider, for example, the following "proof" that would be permitted by the new rule.

- (1) $(P \wedge Q) \rightarrow R$ A
- (2) $P \rightarrow R$ new rule

But we do not consider such inferences to be valid. For example, let P = "I want a new bicycle" and Q = "I am able to buy a new bicycle", and R = "I buy a new bicycle". Then the premise could be true while the conclusion is false -- so the argument is not valid.

Warmup: Old proof rules in new form

Exercise 2.5 ($\wedge E$, $\wedge I$, $\vee I$, MP, MT, DN)

$$1. P \wedge (Q \wedge R) \dashv\vdash (P \wedge Q) \wedge R$$

1	(1)	$p \wedge (q \wedge r)$	A
1	(2)	p	1 &E
1	(3)	$q \wedge r$	1 &E
1	(4)	q	3 &E
1	(5)	r	3 &E
1	(6)	$p \wedge q$	2,4 &I
1	(7)	$(p \wedge q) \wedge r$	6,5 &I

Here we see, among other things, that and introduction tells us to combine dependencies, but that we don't need to include redundancies. In particular, we write just 1 as dependency in lines 6 and 7 instead of 1,1 or 1,1,1.

$$2. P \dashv\vdash P \wedge P$$

$$3. P \rightarrow \neg Q, Q \vdash \neg P$$

$$4. \neg\neg P \vdash \neg\neg P \wedge (P \vee Q)$$

$$5. \neg(P \wedge Q) \rightarrow R, \neg R \vdash P$$

$$6. P \rightarrow (Q \wedge R), A \rightarrow \neg R, P \vdash \neg A$$

$$7. \neg P \rightarrow \neg Q, Q \vdash P$$

$$8. P \vdash \neg\neg(P \vee Q)$$

New proof rules

We will work in blocks. First block ends in problems A1 and A2 on the pset. The focus is on “pure” applications of CP.

$$P \rightarrow R \vdash (P \wedge Q) \rightarrow R$$

$$P \rightarrow Q \vdash (R \rightarrow P) \rightarrow (R \rightarrow Q)$$

$$P \rightarrow Q \vdash (Q \rightarrow R) \rightarrow (P \rightarrow R)$$

$$(A1) P \rightarrow Q \vdash P \rightarrow (Q \vee R)$$

$$(A2) P \rightarrow (Q \rightarrow R) \vdash Q \rightarrow (P \rightarrow R)$$

Second block ends in problems A3 and A4. The focus is on the “contrapositive maneuver”.

$$P \rightarrow Q \vdash \neg Q \rightarrow \neg P$$

$$\vdash (P \wedge Q) \rightarrow P$$

$$(A3) \neg P \vdash \neg(P \wedge Q)$$

$\vdash P \rightarrow (P \vee Q)$

(A4) $\neg(P \vee Q) \vdash \neg P$

Third block ends in problems A5 and A6. We come back to this if we have time.

$P \vdash (P \rightarrow Q) \rightarrow Q$

(A5) $P \vdash (P \rightarrow \neg P) \rightarrow \neg P$

The fourth block ends in problems B1 and B2.

$(P \wedge Q) \vee (P \wedge R) \vdash P$

$P \vee Q \vdash Q \vee P$

(B1) $P \vee (Q \wedge R) \vdash P \vee Q$

The fifth block ends in problems B3 and B4.

$\neg P \vdash P \rightarrow Q$

$P, \neg P \vdash Q$

(B3) $P \vee Q, \neg P \vdash Q$

(B4) $(P \rightarrow R) \wedge (Q \rightarrow R) \vdash (P \vee Q) \rightarrow R$

Evaluating proofs

Exercise Which of the following proofs with CP is correct? If a proof is not correct, explain what is wrong with it, and say whether there is a simple fix, or whether it is fatally flawed. (The following “proofs” use a slightly different notation – one that is easier to input via keyboard. Hopefully the relation between the two notations will become clear from the context.)

```
1  (1) p&q      A
1  (2) p        1 &E
1  (3) q        2 &E
    (4) p>q     2,3 CP
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1  (1) q      A
2  (2) p      A
1  (3) p>q    2,1 CP
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Exercise Which of the following proofs with or-elimination is correct? If a proof is not correct, explain what is wrong with it, and say whether there is a simple fix, or whether it is fatally flawed. (The following “proofs” use a slightly different notation – one that is easier to input via keyboard. Hopefully the relation between the two notations will become clear from the context.)

```
1  (1) p|p      A
```

2	(2) p	A
1	(3) p	1,2,2,2,2 E
1	(1) p q	A
2	(2) p	A
3	(3) q	A
2,3	(4) p&q	2,3 &I
2,3	(5) p	4 &E
1	(6) p	1,2,2,3,5 E

Additional practice problems

$$P \vdash Q \rightarrow (P \wedge Q)$$

$$(P \rightarrow Q) \wedge (P \rightarrow R) \vdash P \rightarrow (Q \wedge R)$$

$$P \rightarrow (P \rightarrow Q) \vdash P \rightarrow Q$$

$$(P \vee Q) \rightarrow R \vdash P \rightarrow R$$

$$P \rightarrow (Q \rightarrow R), P \rightarrow Q \vdash P \rightarrow R$$

$$P \rightarrow (Q \rightarrow R) \vdash (P \rightarrow Q) \rightarrow (P \rightarrow R)$$

$$P \rightarrow Q \vdash \neg Q \rightarrow \neg P$$

$$(P \wedge Q) \rightarrow R \vdash P \rightarrow (Q \rightarrow R)$$

$$P \rightarrow (Q \rightarrow R) \vdash P \rightarrow (\neg R \rightarrow \neg Q)$$

$(P \rightarrow Q) \rightarrow P \vdash (P \rightarrow Q) \rightarrow Q$ (Hint: Not as difficult as it looks. Assume $(P \rightarrow Q) \rightarrow P$ and $P \rightarrow Q$. The latter can be used both as the antecedent of a conditional, and as a conditional itself.)

$$\vdash (P \wedge Q) \rightarrow (Q \wedge P)$$

$$\vdash Q \rightarrow (P \rightarrow Q)$$

$$\vdash Q \rightarrow (P \rightarrow P)$$

$$(P \rightarrow Q) \rightarrow P, Q \vdash P$$

$$(P \rightarrow Q) \rightarrow P \vdash \neg P \rightarrow P$$

$$(P \rightarrow R) \wedge (Q \rightarrow R) \vdash (P \vee Q) \rightarrow R$$

$$P \vee (Q \vee R) \dashv\vdash (P \vee Q) \vee R$$

$$P \wedge (Q \vee R) \dashv\vdash (P \wedge Q) \vee (P \wedge R)$$

$$P \vee (Q \wedge R) \dashv\vdash (P \vee Q) \wedge (P \vee R)$$

$$\neg P \vee Q \dashv\vdash P \rightarrow Q$$

$$\neg P \vee \neg Q \vdash \neg(P \wedge Q)$$

$$P \rightarrow Q \vdash \neg(P \wedge \neg Q)$$

$$\neg(P \wedge Q) \vdash \neg P \vee \neg Q$$

$$\neg(P \rightarrow Q) \vdash P \wedge \neg Q$$

$$\vdash (P \rightarrow Q) \vee (Q \rightarrow P)$$

$$P \rightarrow (Q \vee R) \vdash (P \rightarrow Q) \vee (P \rightarrow R)$$

$$(P \wedge Q) \rightarrow \neg Q \vdash P \rightarrow \neg Q$$

$$P \wedge \neg Q \vdash \neg(P \rightarrow Q)$$

$\vdash ((P \rightarrow Q) \rightarrow P) \rightarrow P$ (Hint: One possibility is to first prove $\vdash P \vee \neg P$, and then argue by cases. The first case is easy if you remember “positive paradox”. For the second case, remember “negative paradox”, i.e. that $\neg P$ implies $P \rightarrow Q$.)

$$P \rightarrow (Q \vee R) \vdash \neg R \rightarrow (\neg Q \rightarrow \neg P)$$

$$P \rightarrow \neg Q \vdash (P \wedge Q) \rightarrow R$$

$$P \rightarrow \neg Q \vdash P \rightarrow (Q \rightarrow R)$$

$$\neg(P \rightarrow Q) \vdash P \rightarrow \neg Q$$

$$P \dashv\vdash (P \wedge Q) \vee (P \wedge \neg Q)$$

$$P \rightarrow (Q \rightarrow R) \dashv\vdash (P \rightarrow Q) \rightarrow (P \rightarrow R)$$

$$P \rightarrow \neg P \dashv\vdash \neg P$$

$$P \rightarrow (Q \rightarrow \neg Q) \dashv\vdash P \rightarrow \neg Q$$

$$(P \rightarrow Q) \rightarrow (P \rightarrow \neg Q) \dashv\vdash P \rightarrow \neg Q$$

$$P \dashv\vdash P \wedge (Q \vee \neg Q)$$

$$P \dashv\vdash P \vee (Q \wedge \neg Q)$$

$$(P \rightarrow Q) \rightarrow Q \vdash (Q \rightarrow P) \rightarrow P$$

$$(P \rightarrow Q) \rightarrow R \vdash (P \rightarrow R) \rightarrow R$$

$(P \rightarrow R) \rightarrow R \dashv\vdash P \vee R$ (Hint: it’s easy to derive $\neg P \rightarrow R$ from the sentence on the left.)

$(P \rightarrow Q) \rightarrow P \dashv\vdash P$ (Hint: assume $\neg P$ and derive $P \rightarrow Q$.)