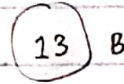


Problem 8.1.

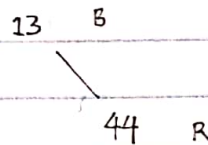
a. [13, 44, 37, 7, 22, 16]

- inserting 13

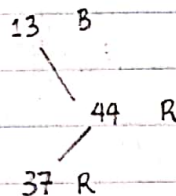


- inserting 44

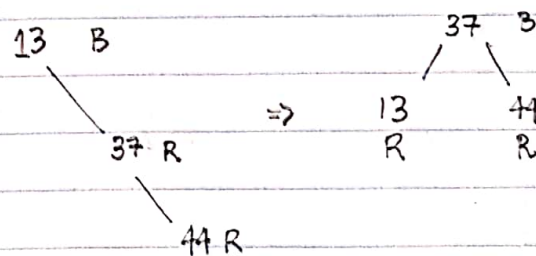
using right subtree property  
since  $44 > 13$ .



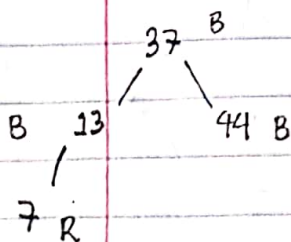
- inserting 37



to resolve violation - rotate 37 with root of tree.

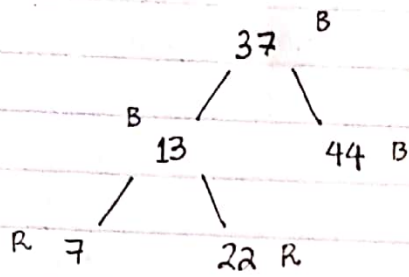


- inserting 7

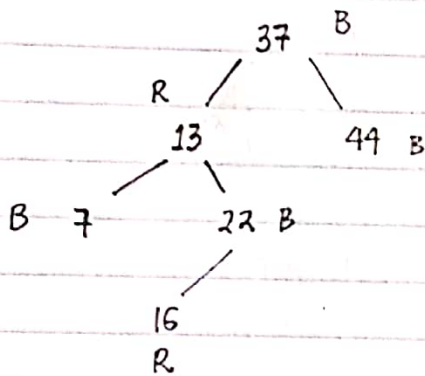


color violations are solved by changing the  
left & right subtree to black.

- inserting 22



- inserting 16

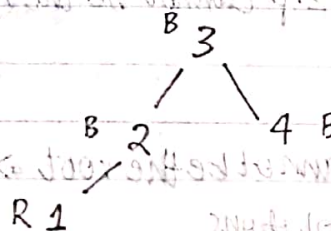
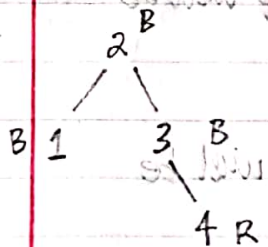
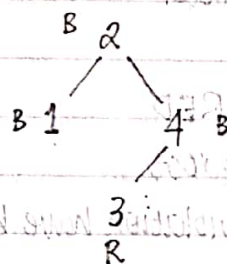
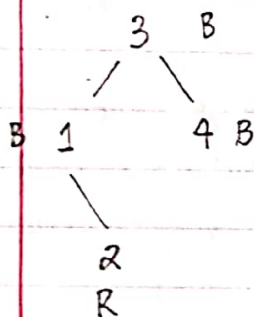


b. All valid RB Trees that store the values [1, 2, 3, 4].

→  $4! = 24$  possible arrangements in which [1, 2, 3, 4]

can be inserted into a red black tree.

⇒ out of this, only 4 trees are valid red black trees.





c. Bonus.

Consider a RB tree formed by inserting  $n$  nodes. Prove that if  $n > 1$ , the tree contains at least one red node.

There are different ways to prove that with  $n > 1$ , a red black tree will have at least 1 red node.

To summarize the cases.

→ if  $z$  and  $z.p.p = \text{RED}$

then  $z$  cannot be the root

thus  $z$  is red after violation have been sorted.

→ if  $z$  and  $z.p = \text{RED}$

and if after rotation  $z.p$  cannot be the root,  $z.p$  will be red after the fix up.

→ if  $z$  is red then  $z$  cannot be the root  $\Rightarrow$  hence  $z$  will be red after fixing violations.

Proof by induction

Base case:  $N = 2$

hypothesis: a tree with  $n$  nodes, with  $1 < n \leq N$ , there is at least 1 red node

Inductive: for a red black tree with  $N+1$  nodes, there is at least 1 red node.

$\Rightarrow$  trivial case = Red node  $N+1$  is the child of a black node

$\Rightarrow$  if  $N+1$  is inserted as the child of a red node - then we can look at different fixing strategies.

1. if  $X$  remains red after recoloring the parent, g. parent & uncle and when  $X$  is moved to Grandparent  $X$  - we shall have at least 1 red node.
2. the parent of  $X$  is red after rotation of  $X$  about the parent & g. parent. the parent remains red after recoloring  $X$  &  $bl$ . we will still have a red node.
3. if  $parent(x)$  is rotated about  $g.parent(x)$  -  $X$  remains red then  $X = Red$  after recoloring  $parent(x)$  &  $g.parent(x)$ . at least 1 red node.

thus, in all cases of fixing the RB tree with nodes  $n > 1$ , there exists at least 1 red node.