My thoughts on supply chain planning

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Here I document some thoughts on supply chain planning. I define the problem and formulate the mixed integer linear program that I used for solving the problem. At the end I briefly comment on its complexity. Before I start with the problem definition I introduce some notations that I use later.

I refer to the power set of set P as 2^P , which is the set of all subsets of P, i.e. $\{P' \mid P' \subseteq P\}$. I use $\dot{\cup}$ for disjoint union.

I start with the definition of supply chain.

Definition 1 (Supply Chain) Let P be a set of products with $o: P \to \mathbb{N}$ as its order function such that o(p) is the number of ordered product $p \in P$. Moreover, C is a set of components with stock function defined by $s: C \to \mathbb{N}_0$. A product or a component is called entity and the set of all entities $\{P \cup C\}$ is referred to as E. Let $C_s \subseteq C$ the component that are directly supplied by a supplier. Each product or component is chained to the network by function $ch: E \to 2^C$. Then $SC = (E = \{P \cup C\}, o, s, ch, C_s)$ is a supply chain.

Chain function ch specifies how an entity is connected to other entities in the chain. Formally speaking, entity e is directly supplied by all entities in ch(e). The type of supply is recognized by partitioning entities into sets \mathcal{A} and \mathcal{O} corresponding and and or delivery, respectively. The set of entities that are supplied by entity e is denoted by $\delta(e)$, i.e. $\{e' \mid e \in ch(e)\}$.

Entities can flow through the chain down from suppliers and inventories as the source and be delivered as the products. An assignment of the flows to the chain that satisfies some constraints is called a plan.

Definition 2 (Plan) A plan is an assignment of flows going out of suppliers $(s_c \text{ for } c \in C_s)$, going out of inventories $(v_c \text{ for } c \in C)$, passing through edges $y_{e,f}$ for all $e \in E$, for all $f \in ch(e)$ and being delivered as an entity $(x_e \text{ for } c \in C)$

 $p \in P$) that satisfy the following constrains

$$0 \le x_p \le o(p), \qquad \forall \ p \in P \tag{1}$$

$$0 \le v_c \le s(c), \qquad \forall \ c \in C \tag{2}$$

$$x_e = y_{e,f}, \quad \forall \ e \in \mathcal{A}, \forall \ f \in ch(e)$$
 (3)

$$x_e \le y_{e,f} + (1 - z_{e,f})\bar{x}_e, \quad \forall e \in \mathcal{O}. \ \forall f \in ch(e)$$
 (4)

$$x_e \ge y_{e,f} - (1 - z_{e,f})\bar{x}_e, \quad \forall e \in \mathcal{O}. \ \forall f \in ch(e)$$
 (5)

$$x_{e} \leq y_{e,f} + (1 - z_{e,f})\bar{x}_{e}, \qquad \forall e \in \mathcal{O}. \ \forall f \in ch(e)$$

$$x_{e} \geq y_{e,f} - (1 - z_{e,f})\bar{x}_{e}, \qquad \forall e \in \mathcal{O}. \ \forall f \in ch(e)$$

$$\sum_{f \in \delta(c)} y_{c,f} = x_{c} + s_{c} + v_{c}, \qquad \forall c \in C_{s}$$

$$\sum_{f \in \delta(c)} y_{c,f} = x_{c} + v_{c}, \qquad \forall c \in C \setminus C_{s}$$

$$\sum_{f \in ch(e)} z_{e,f} = 1, \qquad \forall e \in \mathcal{O}$$

$$z_{e,f} \in \{0,1\}, \qquad \forall e \in \mathcal{O}$$

$$s_{c} \geq 0, \qquad \forall c \in C_{s}$$

$$(9)$$

$$\sum_{f \in \delta(c)} y_{c,f} = x_c + v_c, \qquad \forall \ c \in C \setminus C_s \tag{7}$$

$$\sum_{f \in ch(e)} z_{e,f} = 1, \qquad \forall \ e \in \mathcal{O}$$
(8)

$$z_{e,f} \in \{0,1\}, \quad \forall \ e \in \mathcal{O}$$
 (9)

$$s_c \ge 0, \qquad \forall \ c \in C_s$$
 (10)

In Ineqs. (4) and (5), term \bar{x}_e refers to any upper bound on the flow delivered to entity e. It can be easily computed by a bachward search through the chain.

I define two problems related to supply chain. Supply Chain Planning problem is a decision problem to check whether a supply chain can satisfy its order.

Problem 1 (Supply Chain Planning (SCP)) Given a supply chain SC as defined by Def. 1, SCP decides if there is a plan satisfying $x_p = o(p)$ for all $p \in P$.

The optimization version of Prob. 1 is defined next.

Problem 2 (Maximal Supply Chain Plan (MSCP)) Given a supply chain SC as defined by Def. 1 and an affine function r on flow variables defined in described in Def. 2, MSCP is the following optimization problem:

$$\max_{x} r(x, y, v, s)$$
 subject to
$$\text{constrains (1) to (10)}$$

It is straightforward to prove Prob. 1 is NP-complete. One can easily encode subset sum problem as SCP problem. I also think Prob. 2 is NP-hard, but I have no proof for it.