

Graph Algorithms I

Breadth First Search

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Graphs

Graph $G = (V, E)$

V : Set of Nodes (Vertices)

E : Set of Edges



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V : Set of Nodes (Vertices)

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Let $|V| = n$

Let $|E| = m$



Undirected graph:



Undirected graph:

Each edge is an **unordered** pair (u, v)

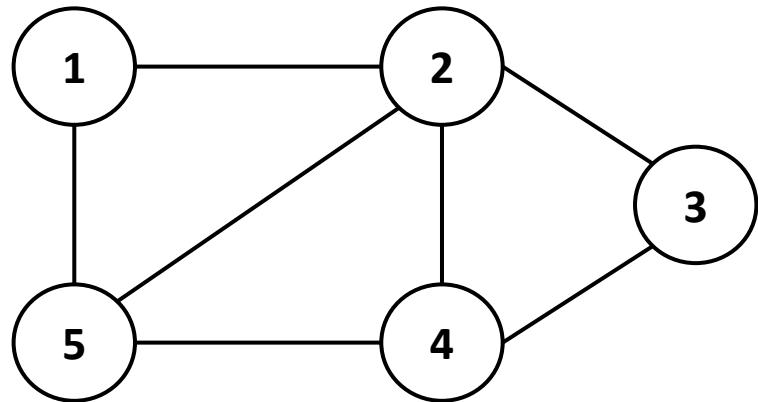
i.e. $(u, v) = (v, u)$



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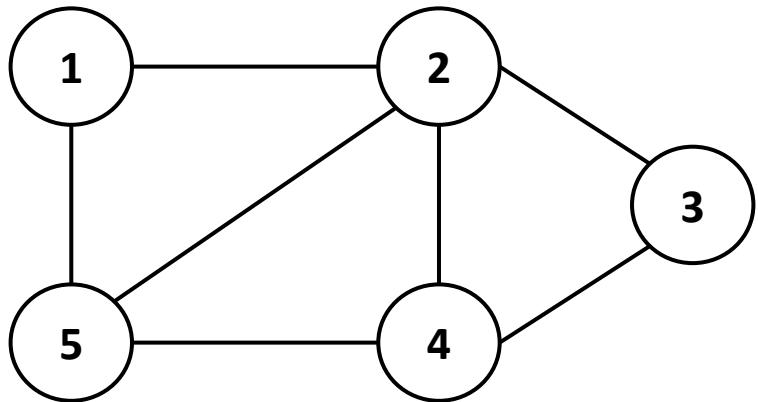
Example:



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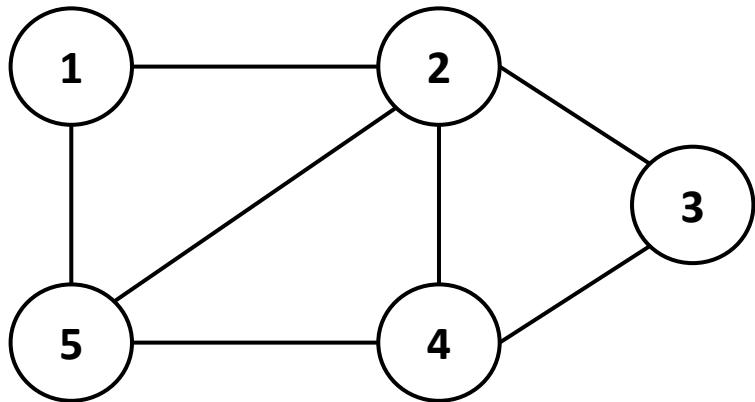
$$V = \{1, 2, 3, 4, 5\}$$



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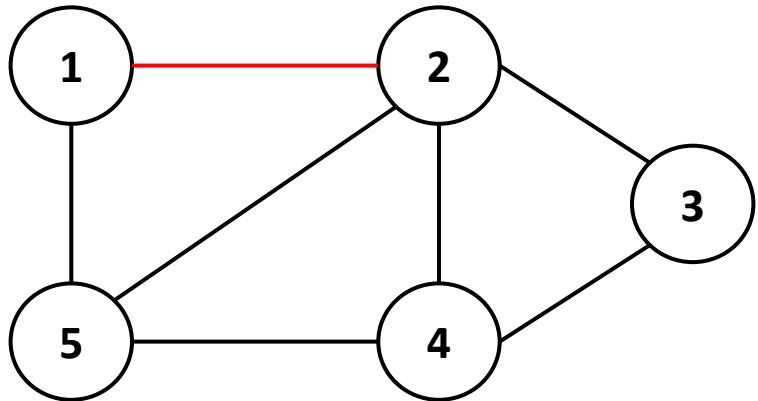
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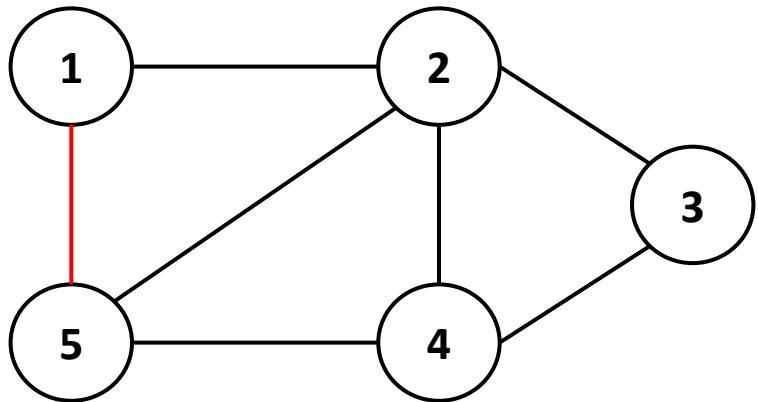
$$E = \{ (1, 2), \\ \} \quad \quad \quad$$



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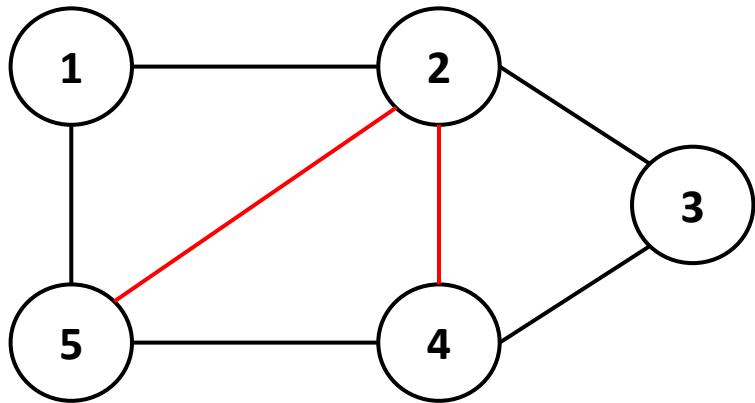
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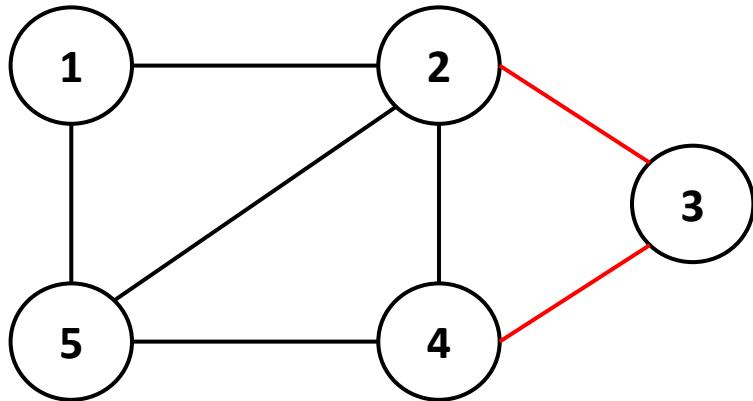
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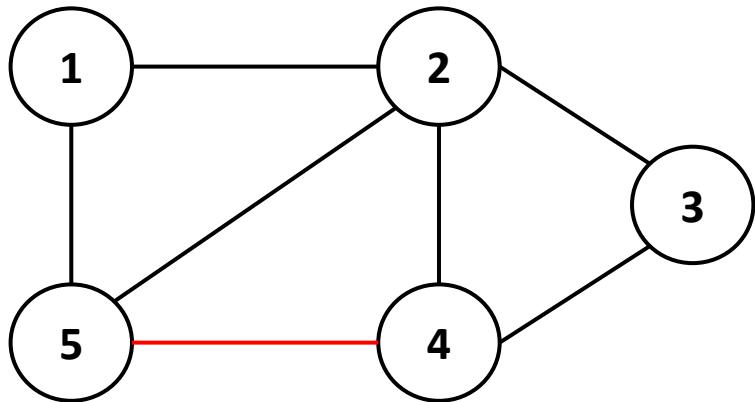
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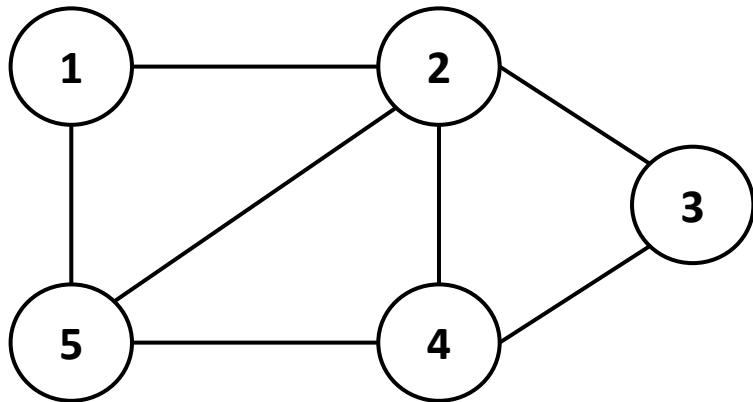
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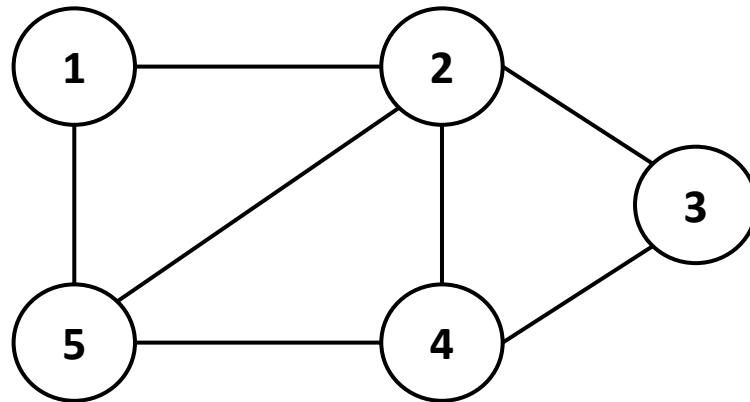
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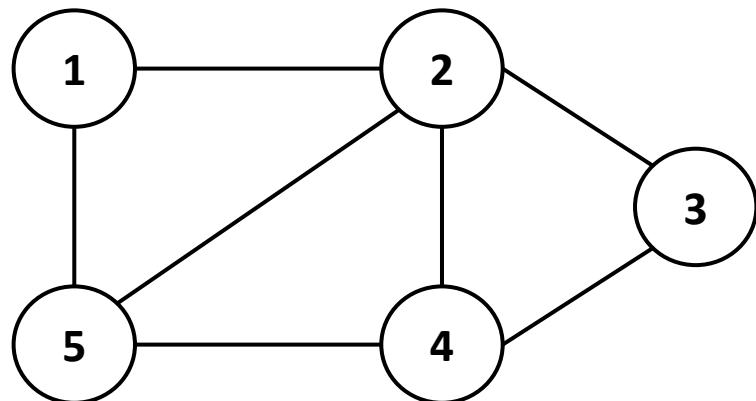
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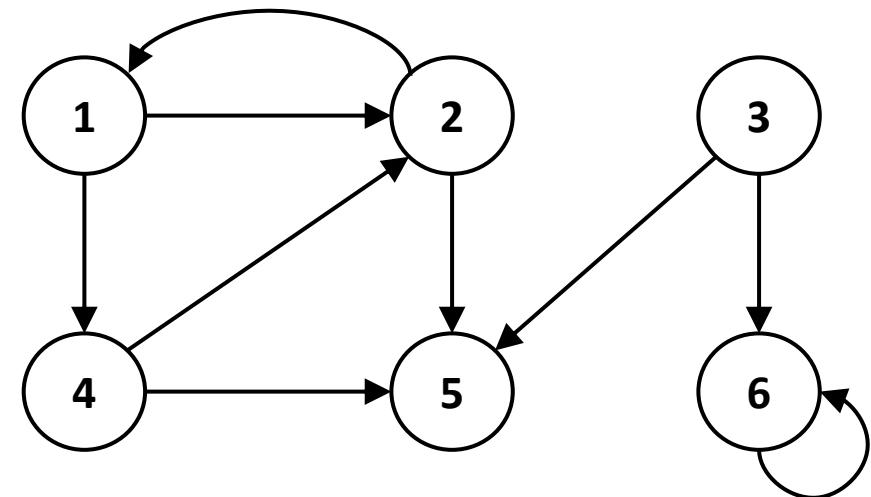
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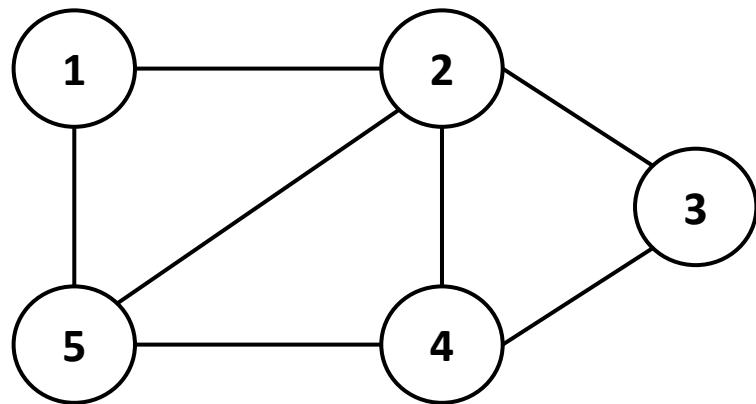
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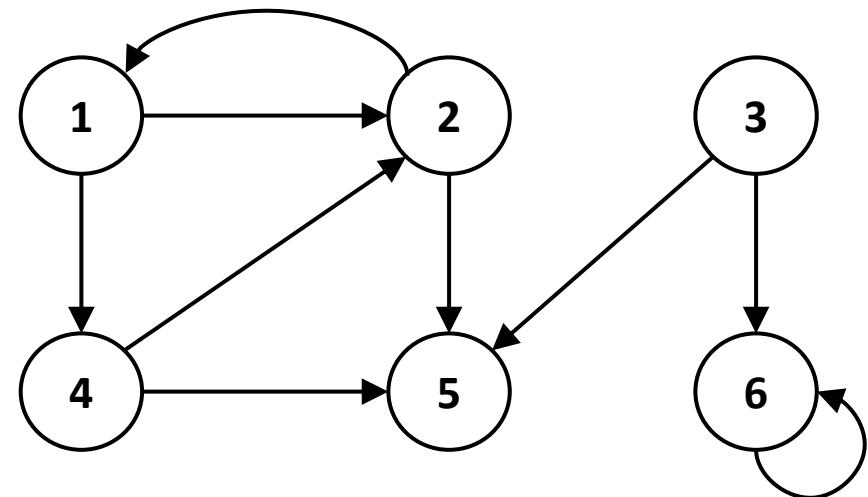
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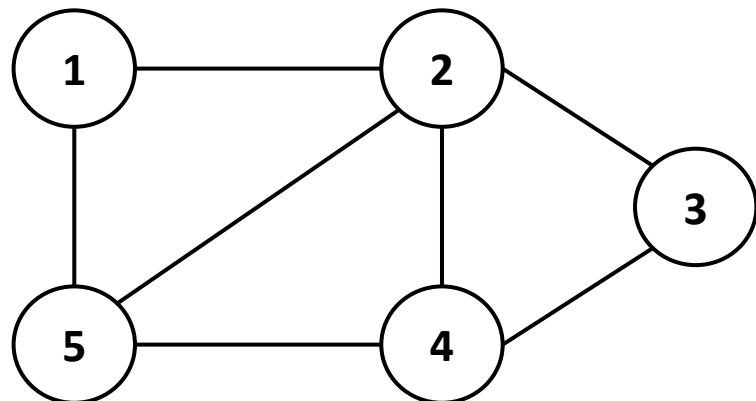
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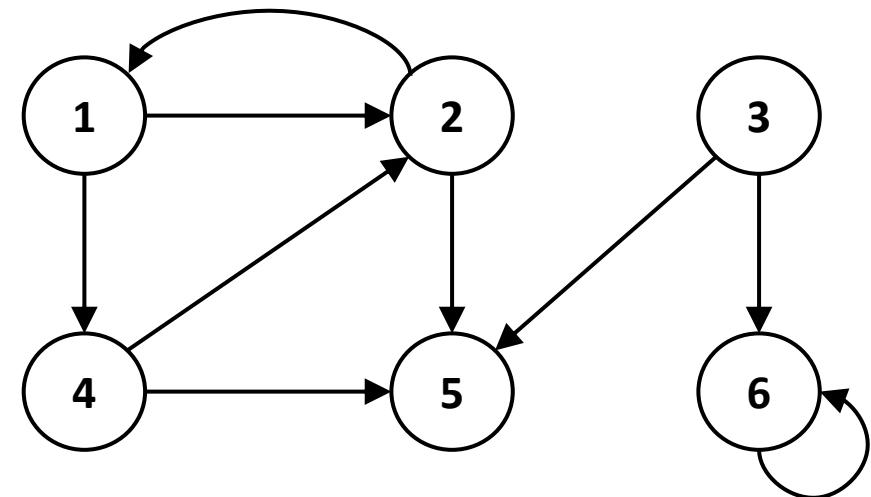
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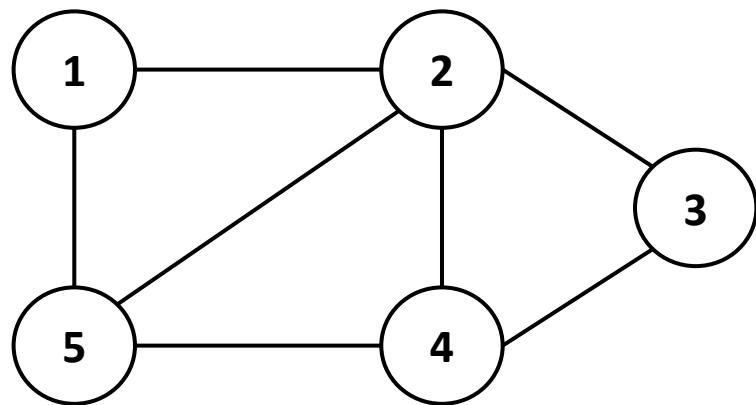
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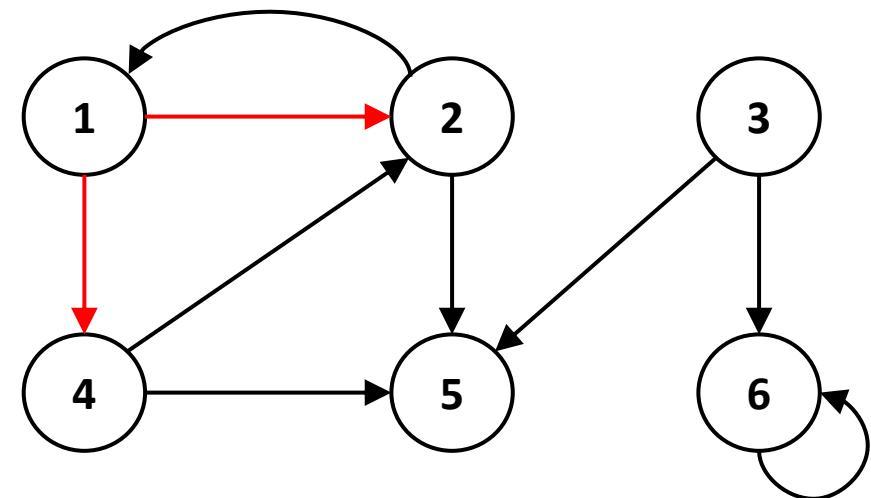
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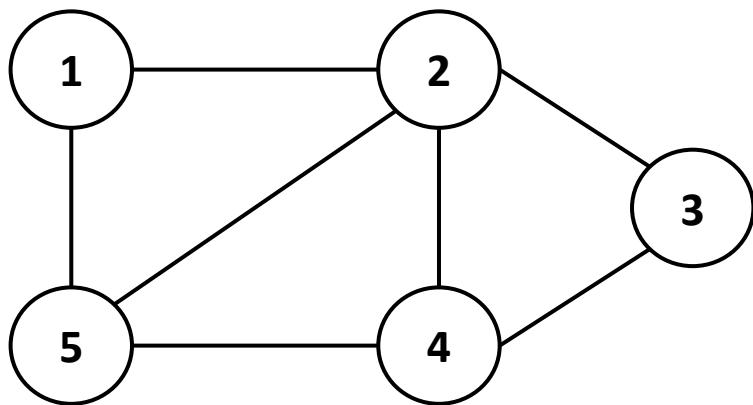
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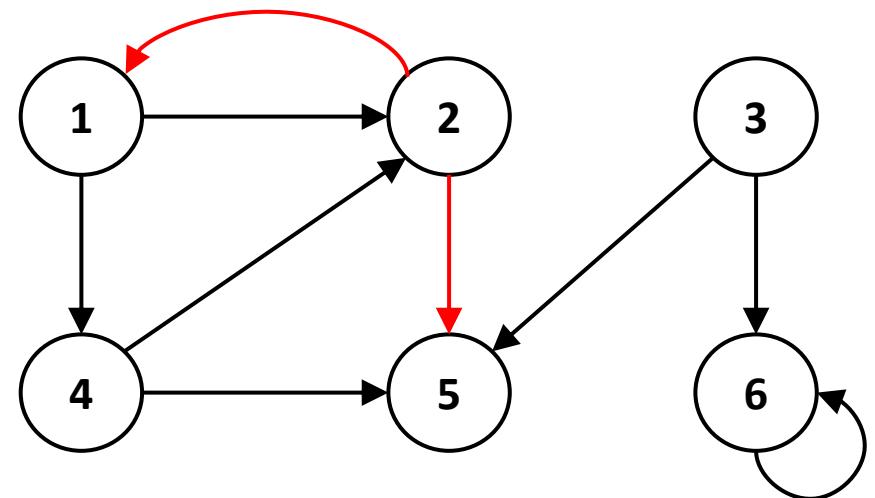
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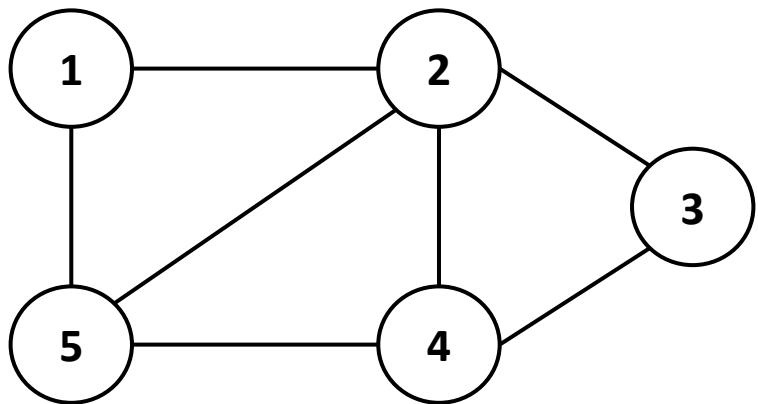
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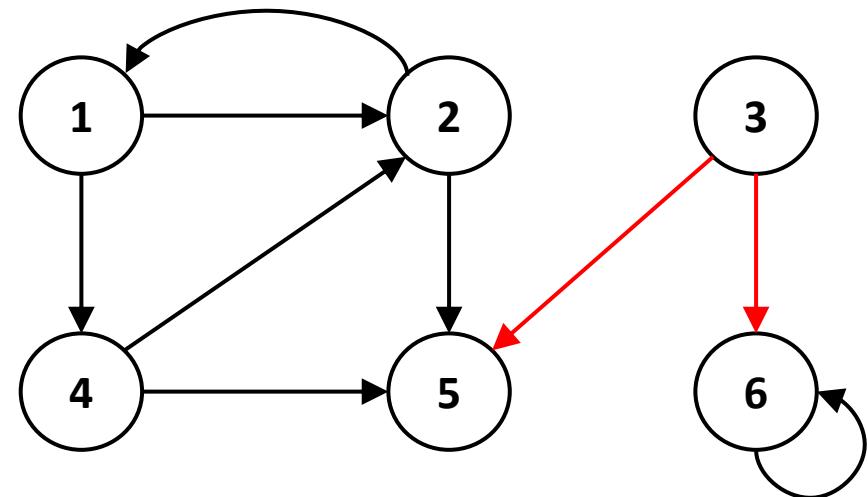
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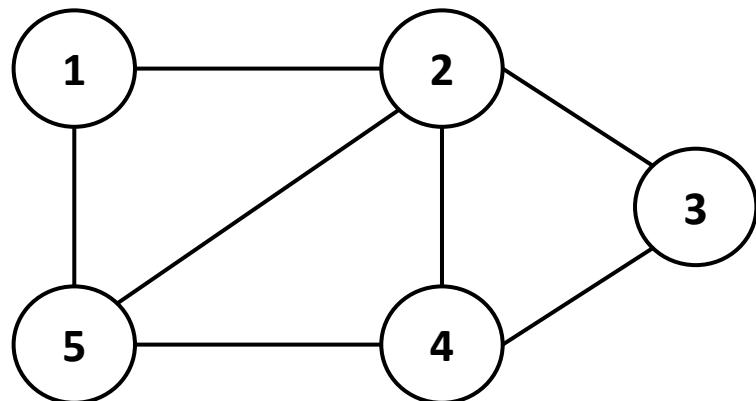
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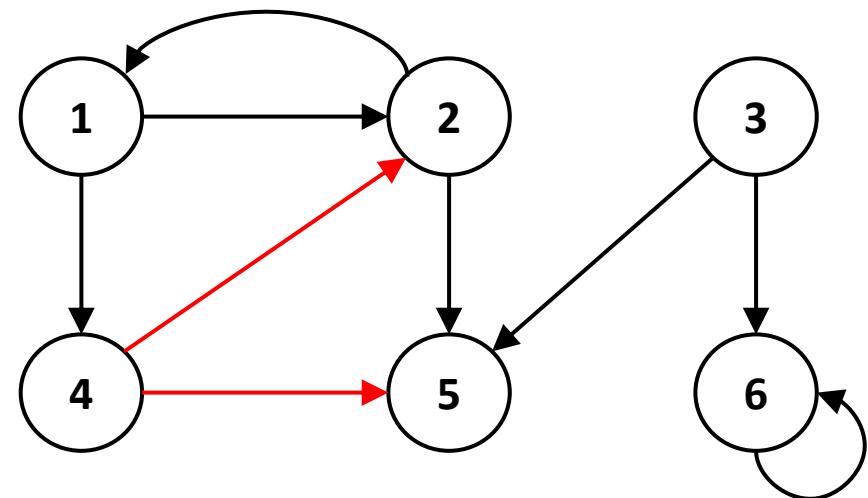
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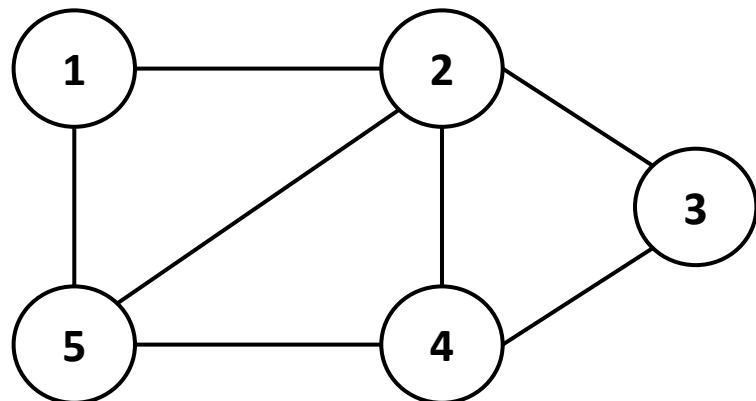
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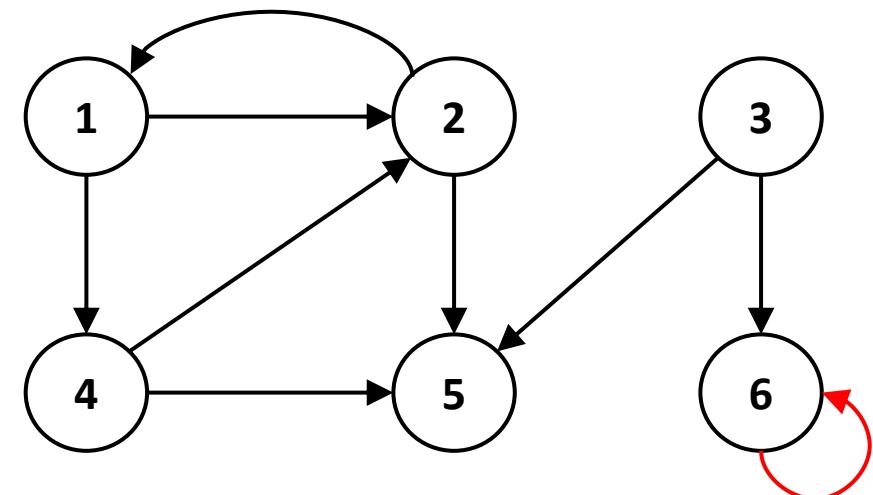
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Representing graphs

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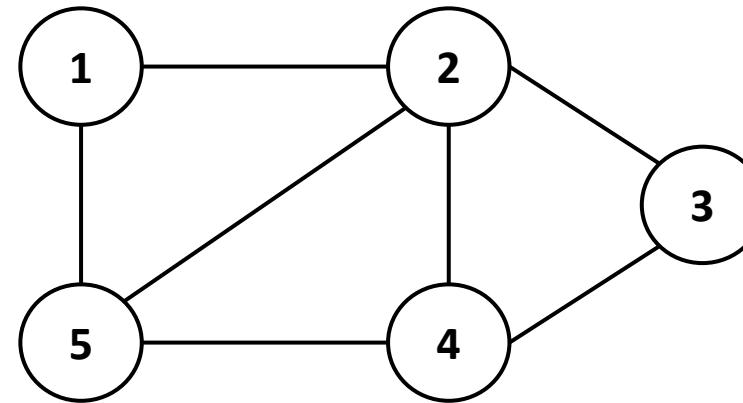
Representing graphs

Adjacency List of $G = (V, E)$:

Array **Adj[1...n]** of n lists, one for each node $u \in V$,

Adj[u] : list of all nodes v such that $(u,v) \in E$

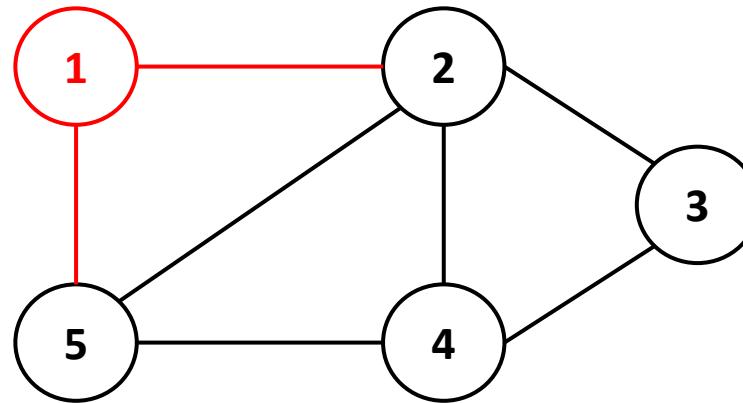




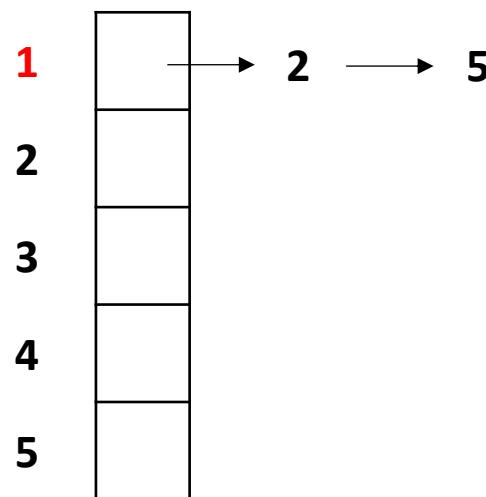
Adjacency list

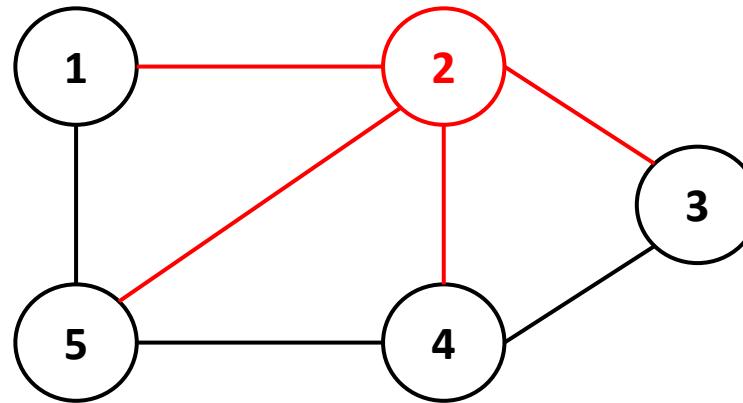
| | |
|---|--|
| 1 | |
| 2 | |
| 3 | |
| 4 | |
| 5 | |



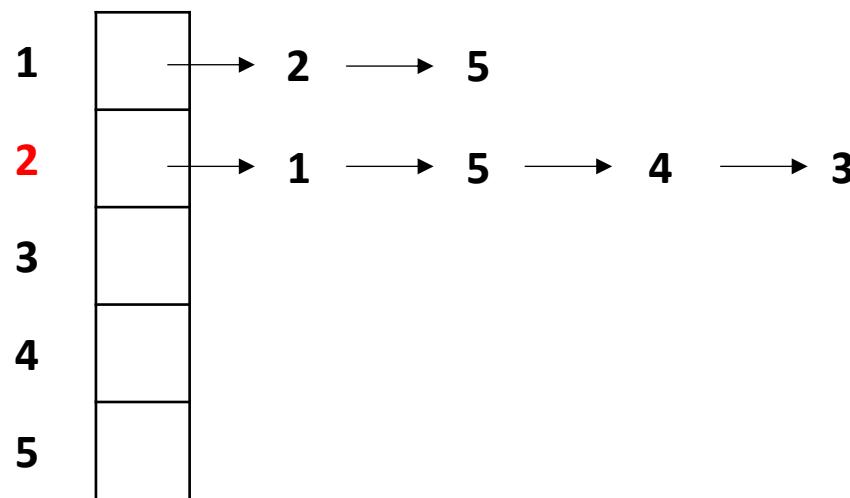


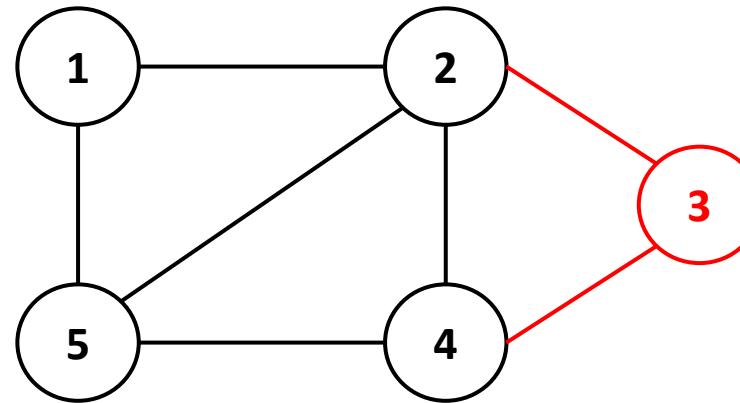
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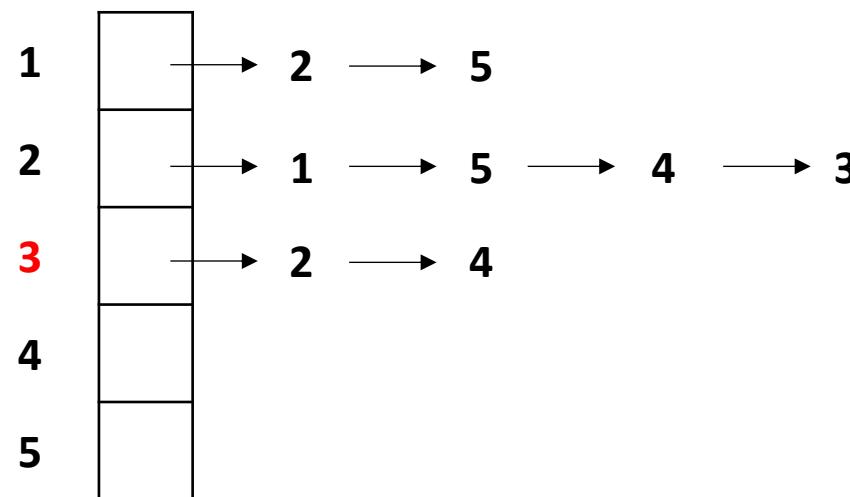


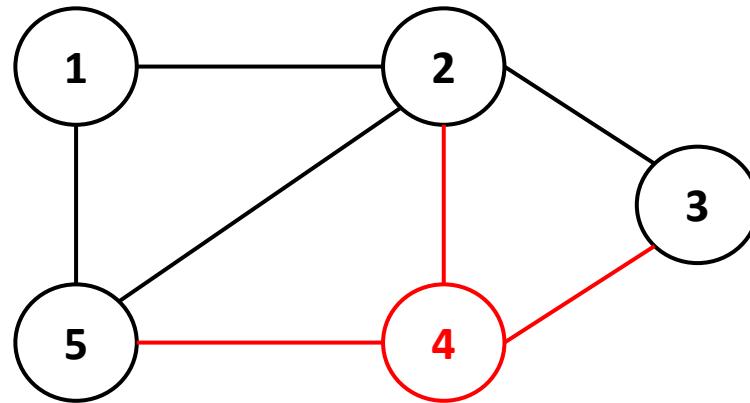
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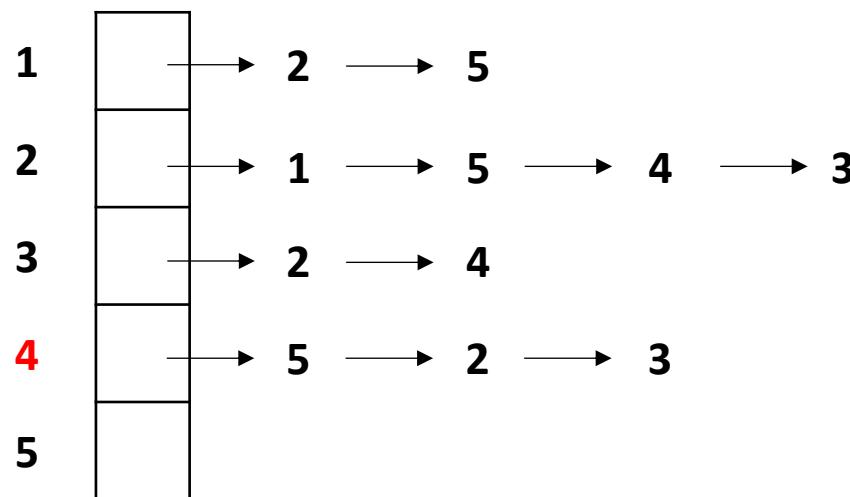


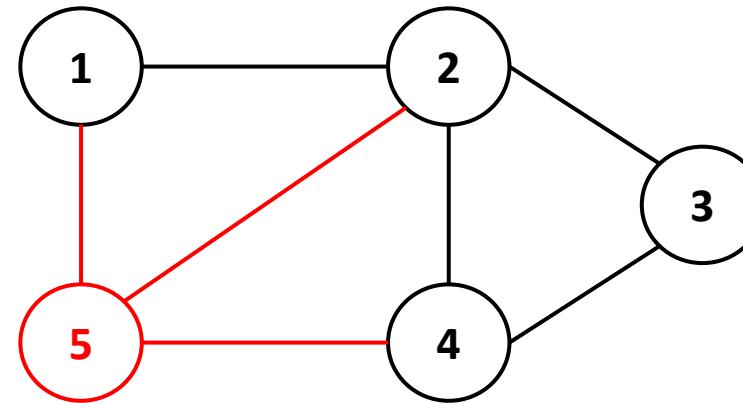
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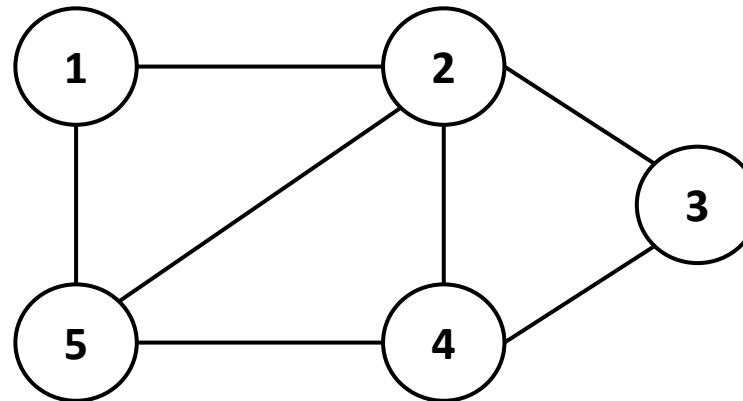




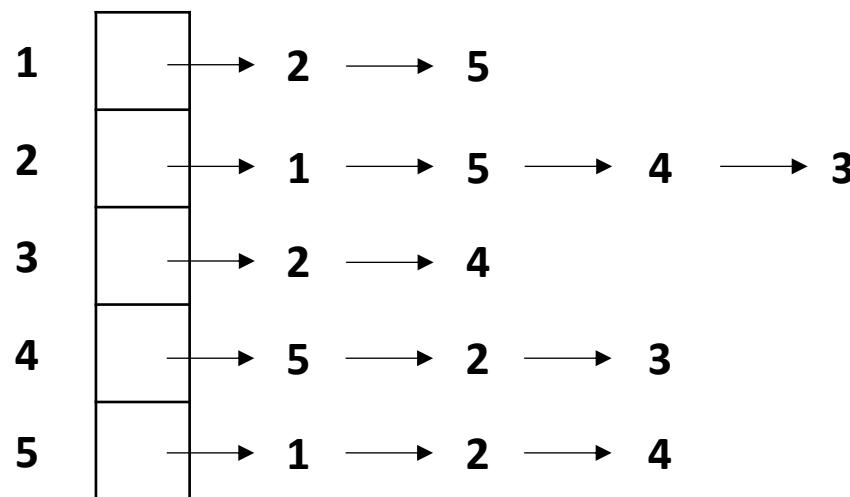
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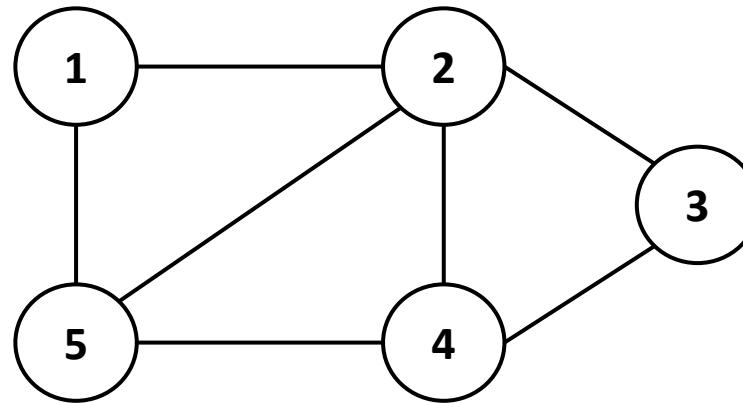
| | |
|---|-----------------|
| 1 | → 2 → 5 |
| 2 | → 1 → 5 → 4 → 3 |
| 3 | → 2 → 4 |
| 4 | → 5 → 2 → 3 |
| 5 | → 1 → 2 → 4 |



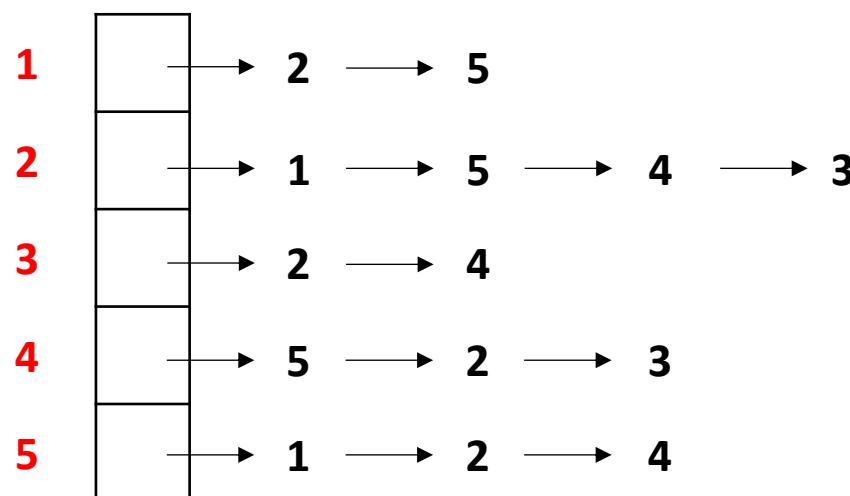


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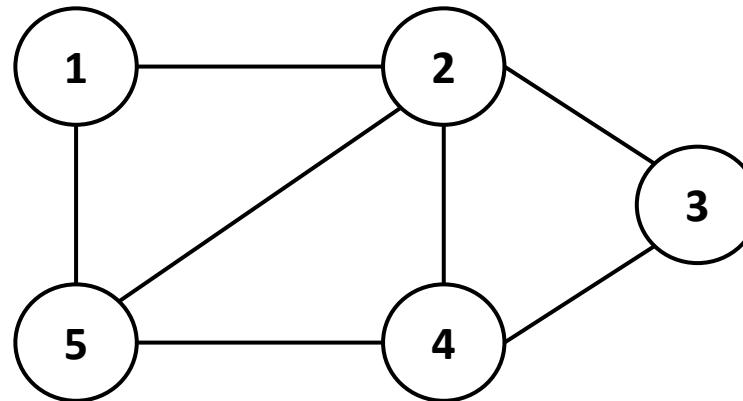


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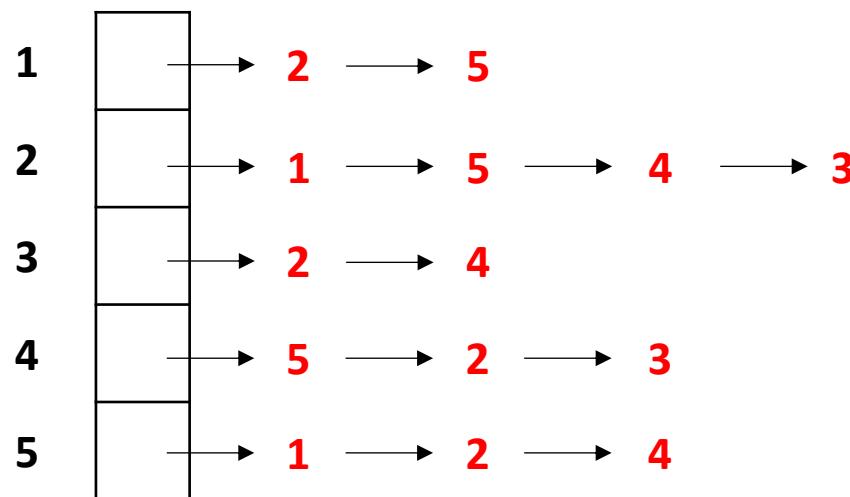


number of nodes

Size: $\Theta(n)$

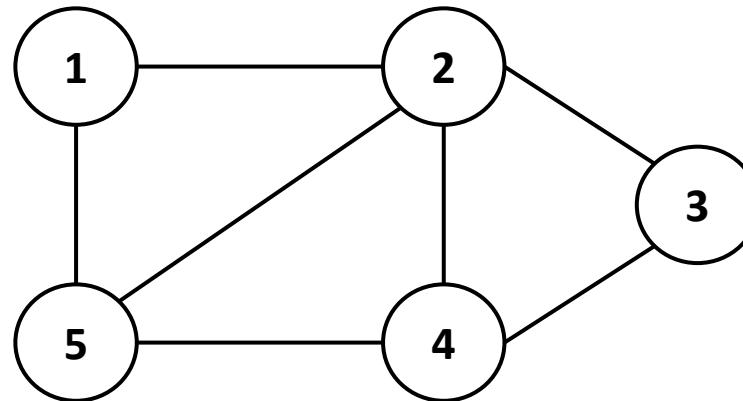


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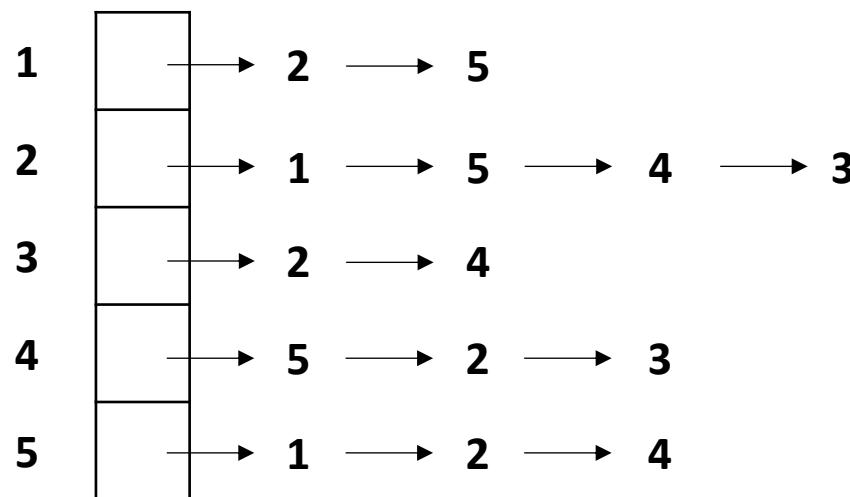


number of nodes number of edges
 Size: $\Theta(n + m)$

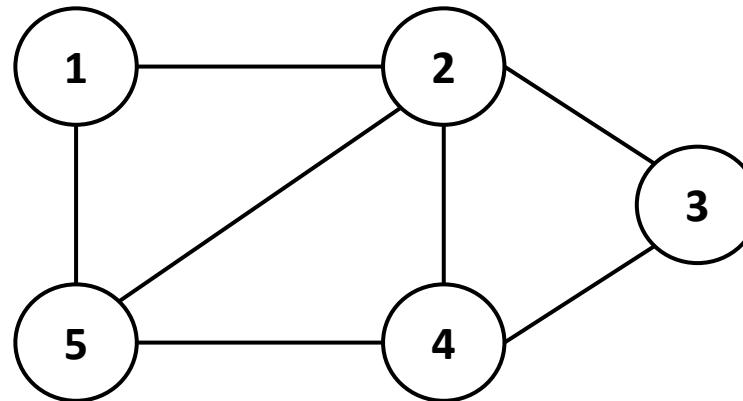




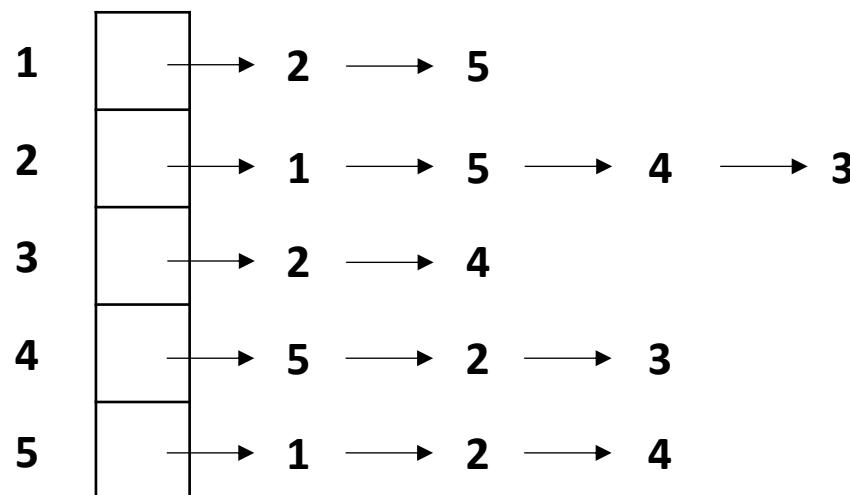
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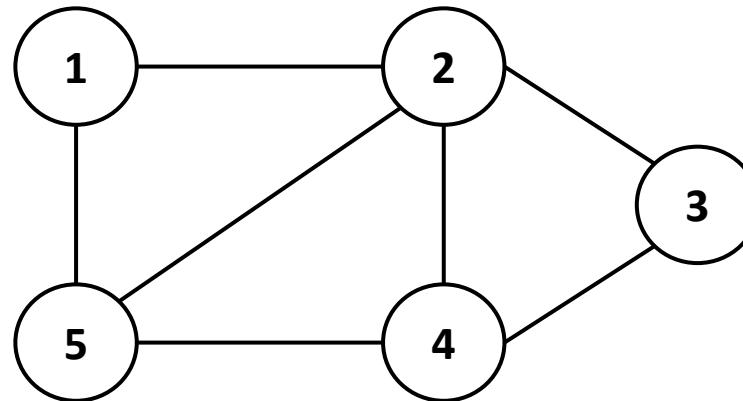


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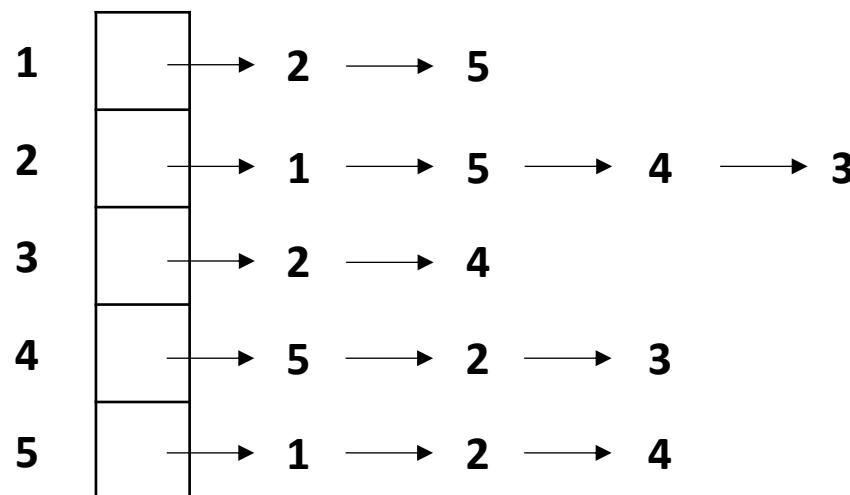
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Note that $m \leq$





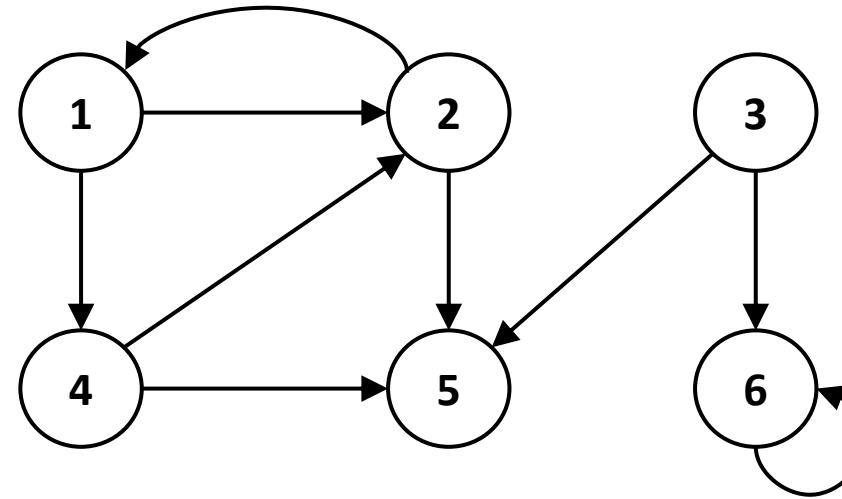
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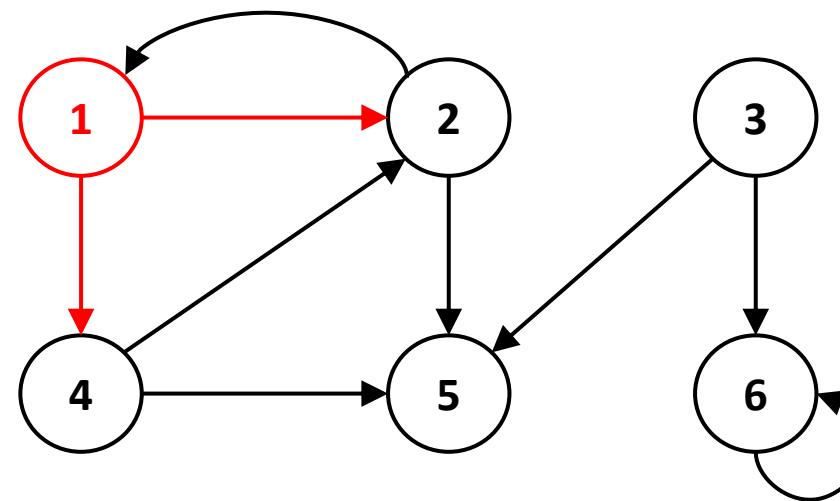
Note that $m \leq n^2$



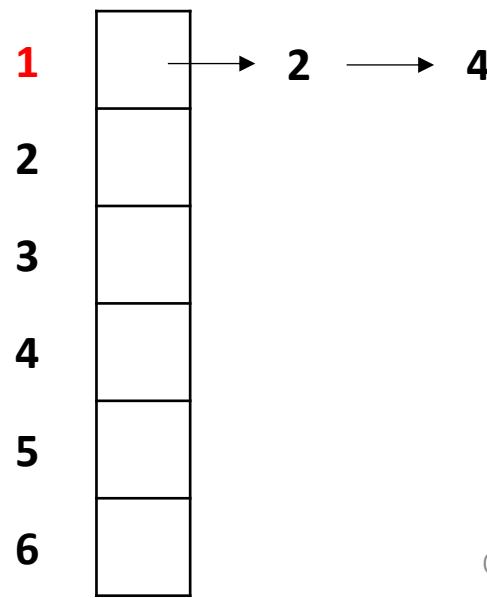
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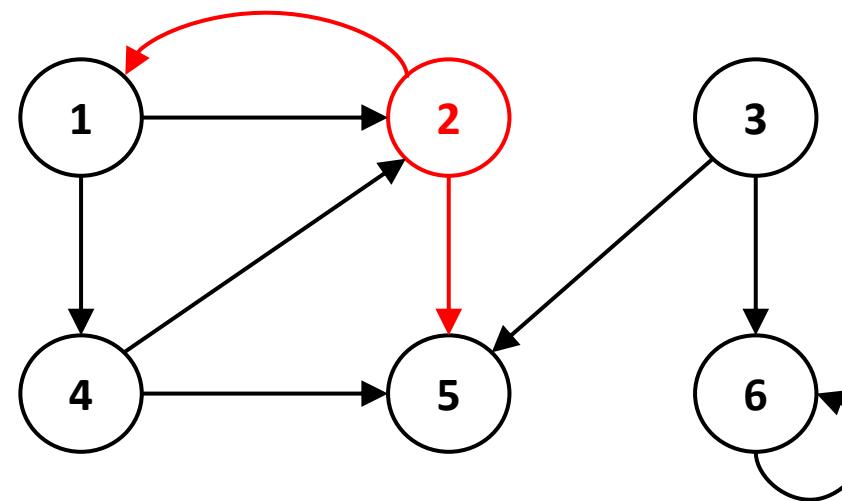
| | |
|---|--|
| 1 | |
| 2 | |
| 3 | |
| 4 | |
| 5 | |
| 6 | |





Adjacency list

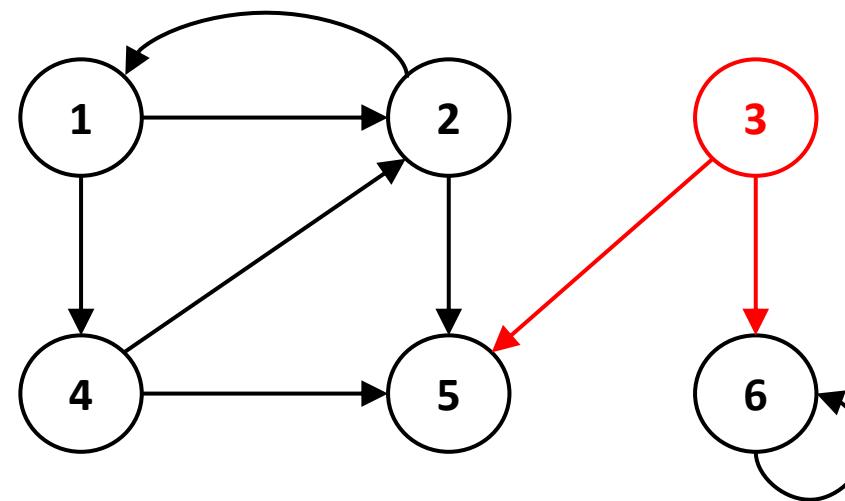




Adjacency list

| | | | | |
|---|---|---|---|---|
| 1 | → | 2 | → | 4 |
| 2 | → | 5 | → | 1 |
| 3 | | | | |
| 4 | | | | |
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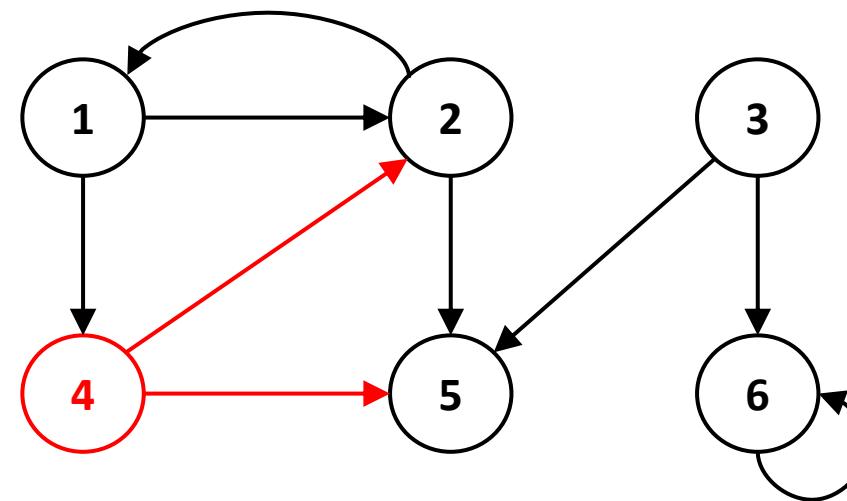




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| 6 | | | | |

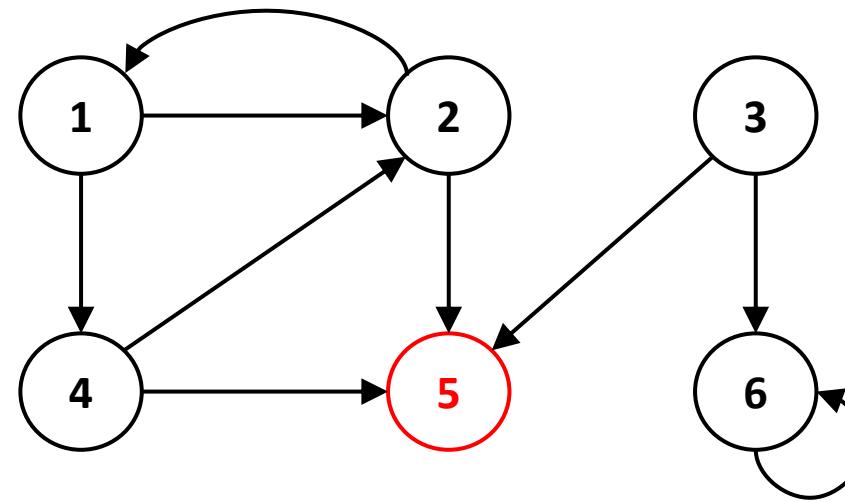




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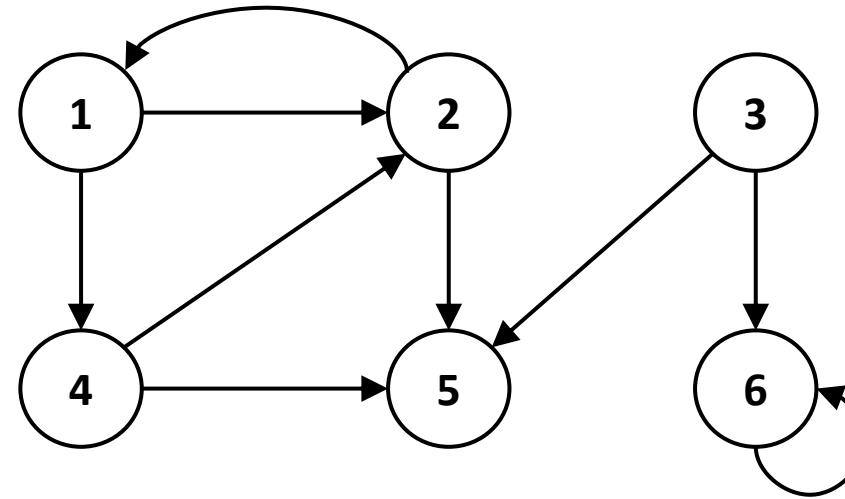




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| 4 | → 2 → 5 |
| 5 | |
| 6 | |

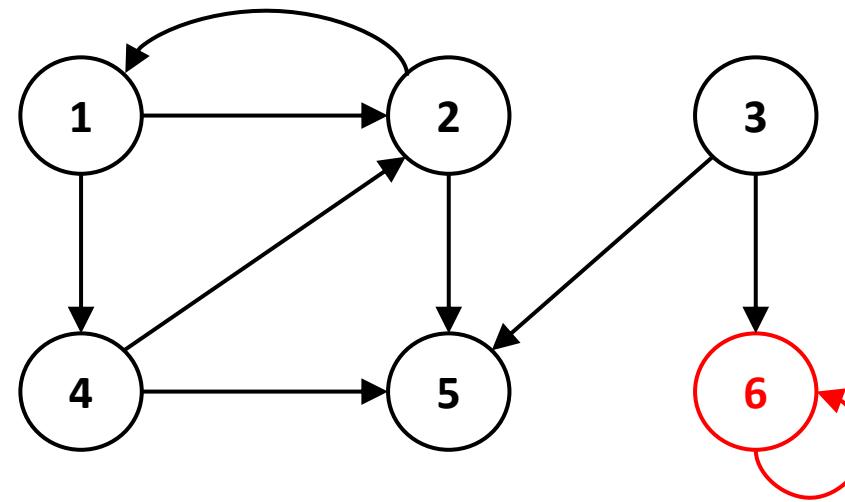




Adjacency list

| | | | | |
|---|---|---|---|---|
| 1 | → | 2 | → | 4 |
| 2 | → | 5 | → | 1 |
| 3 | → | 5 | → | 6 |
| 4 | → | 2 | → | 5 |
| 5 | | | | |
| 6 | → | 6 | | |

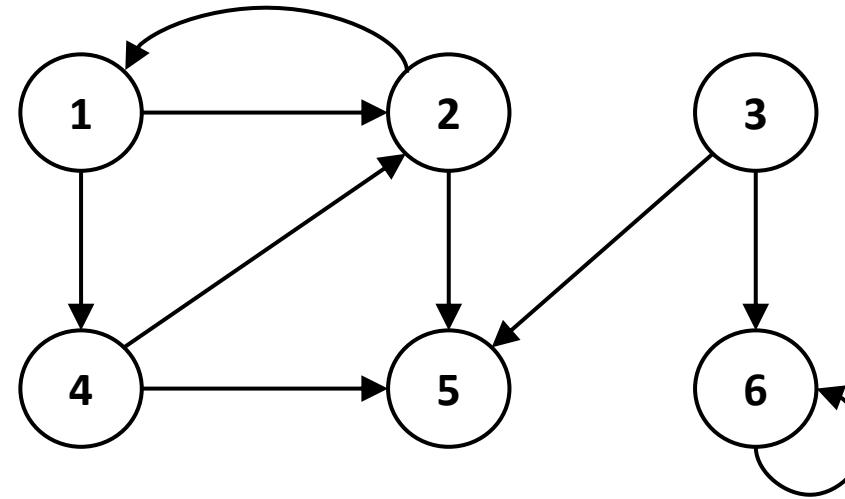




Adjacency list

| | |
|---|---------|
| 1 | → 2 → 4 |
| 2 | → 5 → 1 |
| 3 | → 5 → 6 |
| 4 | → 2 → 5 |
| 5 | |
| 6 | → 6 |





Adjacency list

| | | | | |
|---|---|---|---|---|
| 1 | → | 2 | → | 4 |
| 2 | → | 5 | → | 1 |
| 3 | → | 5 | → | 6 |
| 4 | → | 2 | → | 5 |
| 5 | → | | | |
| 6 | → | 6 | | |

Size: $\Theta(n + m)$



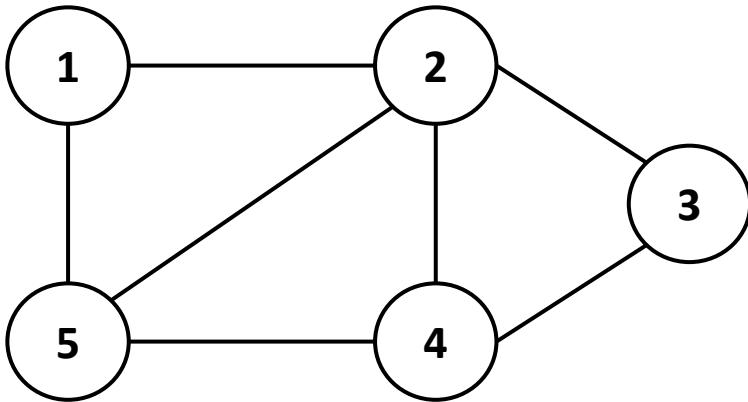
Representing graphs

Adjacency Matrix of $G = (V, E)$:

$n \times n$ matrix A such that for $1 \leq i, j \leq n$:

$$A[i, j] = \begin{cases} 1 & \text{if } (i, j) \in E \\ 0 & \text{otherwise} \end{cases}$$

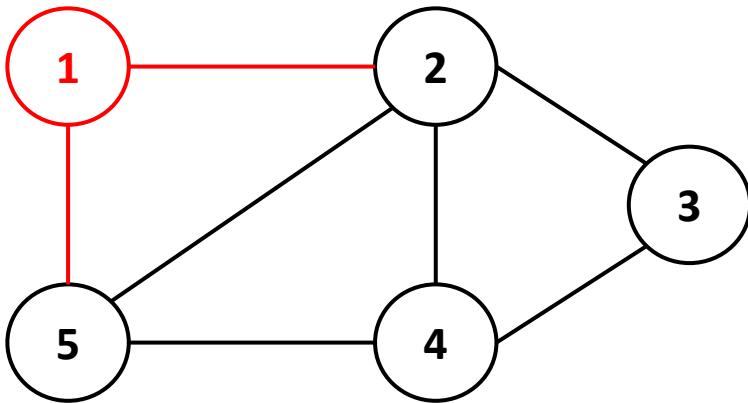




Adjacency matrix

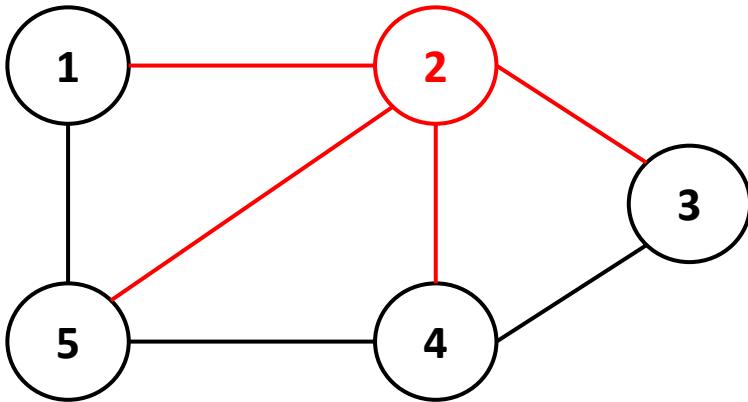
| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | | | | | |





Adjacency matrix

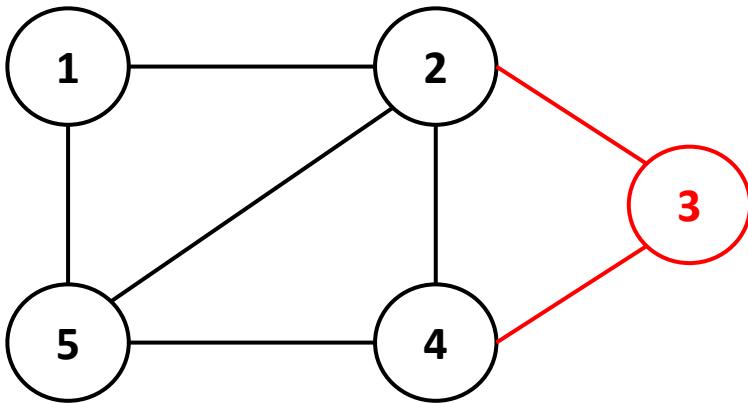
| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 0 | 1 |
| 2 | | | | | |
| 3 | | | | | |
| 4 | | | | | |
| 5 | | | | | |



Adjacency matrix

| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 0 | 1 |
| 2 | 1 | 0 | 1 | 1 | 1 |
| 3 | | | | | |
| 4 | | | | | |
| 5 | | | | | |

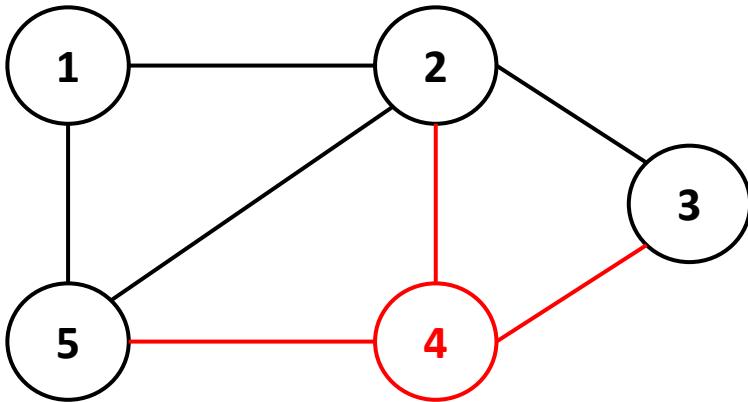




Adjacency matrix

| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 0 | 1 |
| 2 | 1 | 0 | 1 | 1 | 1 |
| 3 | 0 | 1 | 0 | 1 | 0 |
| 4 | | | | | |
| 5 | | | | | |

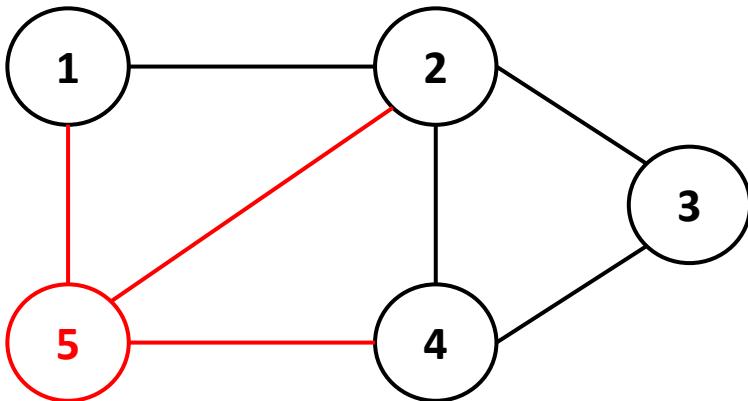




Adjacency matrix

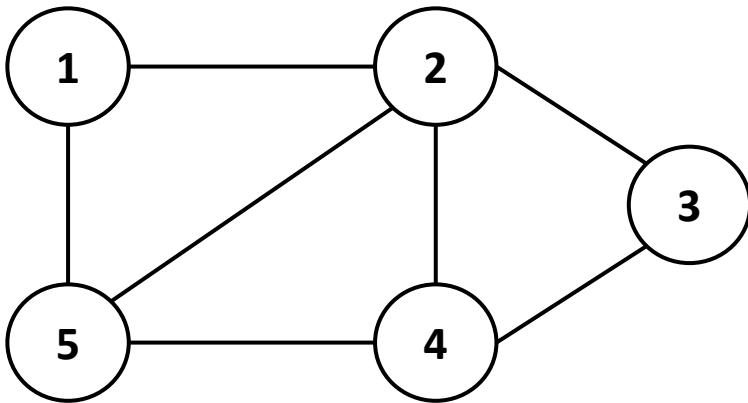
| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 0 | 1 |
| 2 | 1 | 0 | 1 | 1 | 1 |
| 3 | 0 | 1 | 0 | 1 | 0 |
| 4 | 0 | 1 | 1 | 0 | 1 |
| 5 | | | | | |





Adjacency matrix

| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 0 | 1 |
| 2 | 1 | 0 | 1 | 1 | 1 |
| 3 | 0 | 1 | 0 | 1 | 0 |
| 4 | 0 | 1 | 1 | 0 | 1 |
| 5 | 1 | 1 | 0 | 1 | 0 |

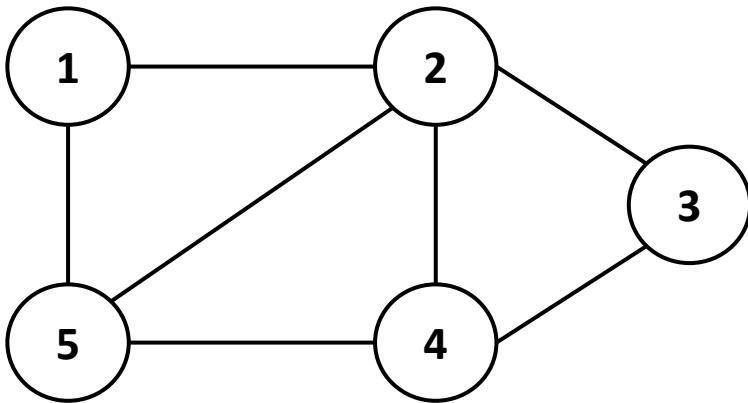


Adjacency matrix

Adjacency matrix of an undirected graph is **symmetric** about its diagonal

| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 0 | 1 |
| 2 | 1 | 0 | 1 | 1 | 1 |
| 3 | 0 | 1 | 0 | 1 | 0 |
| 4 | 0 | 1 | 1 | 0 | 1 |
| 5 | 1 | 1 | 0 | 1 | 0 |

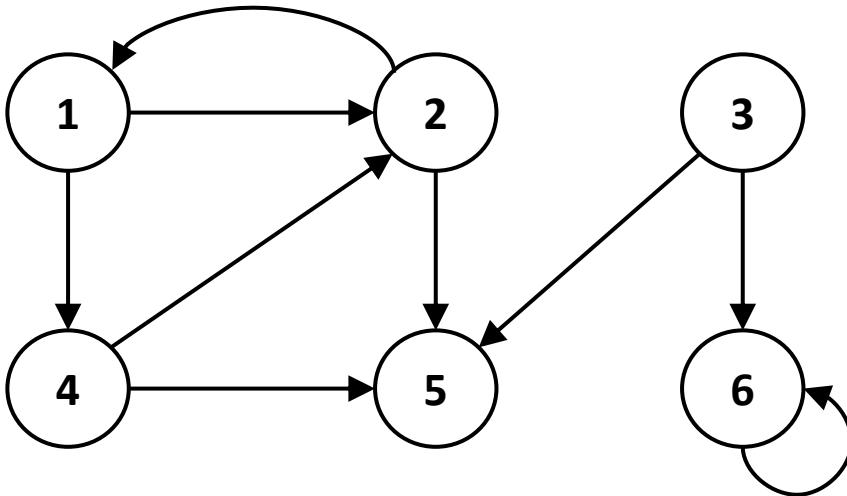




Adjacency matrix

| | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 0 | 1 |
| 2 | 1 | 0 | 1 | 1 | 1 |
| 3 | 0 | 1 | 0 | 1 | 0 |
| 4 | 0 | 1 | 1 | 0 | 1 |
| 5 | 1 | 1 | 0 | 1 | 0 |



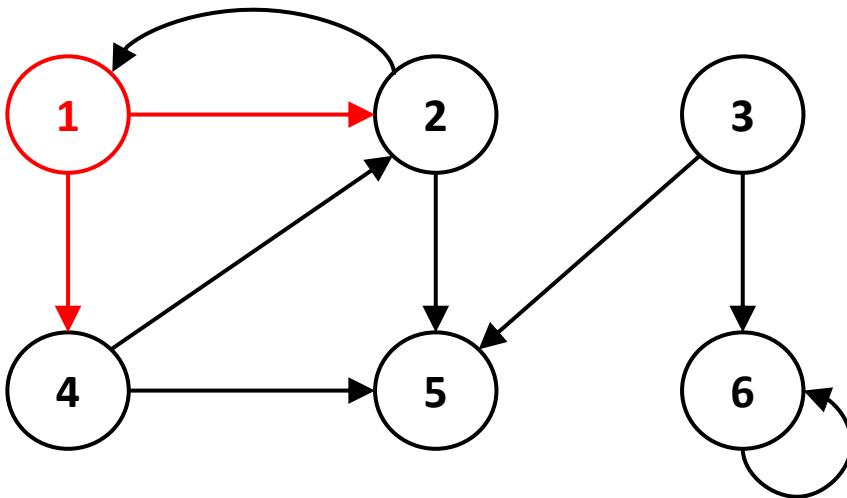


Adjacency matrix

1 2 3 4 5 6

| | | | | | |
|---|--|--|--|--|--|
| 1 | | | | | |
| 2 | | | | | |
| 3 | | | | | |
| 4 | | | | | |
| 5 | | | | | |
| 6 | | | | | |

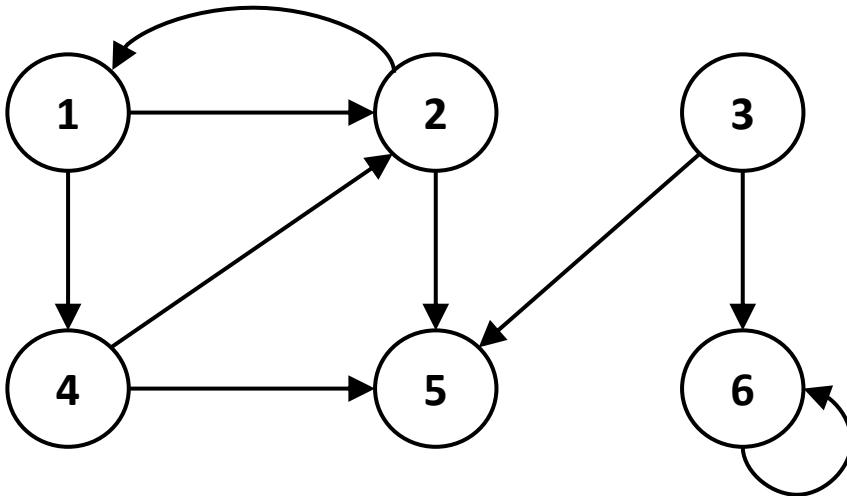




Adjacency matrix

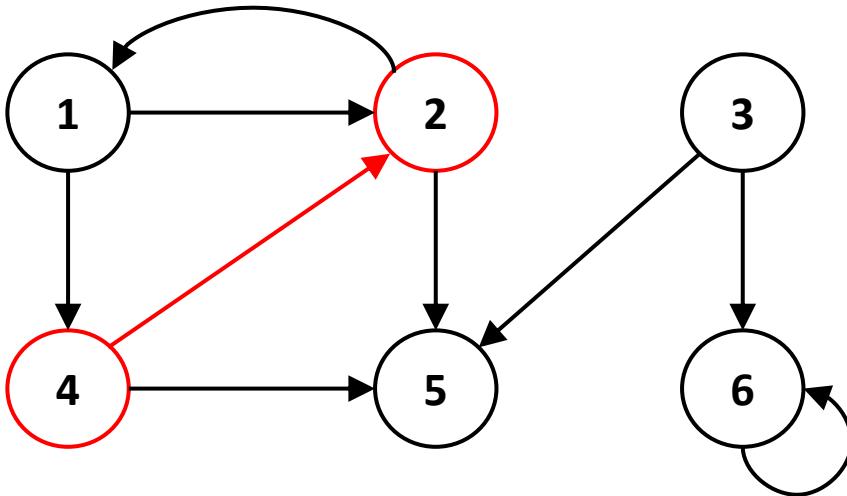
| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 2 | | | | | | |
| 3 | | | | | | |
| 4 | | | | | | |
| 5 | | | | | | |
| 6 | | | | | | |





Adjacency matrix

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 |
| 4 | 0 | 1 | 0 | 0 | 1 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 |
| 6 | 0 | 0 | 0 | 0 | 0 | 1 |

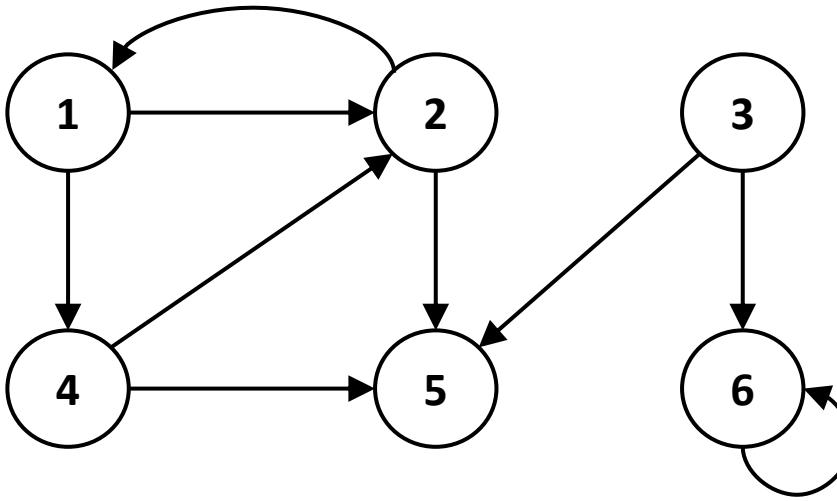


Adjacency matrix

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 |
| 4 | 0 | 1 | 0 | 0 | 1 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 |
| 6 | 0 | 0 | 0 | 0 | 0 | 1 |

$$A[2, 4] \neq A[4, 2]$$

Adjacency matrix of a directed graph is not necessarily symmetric about its diagonal.



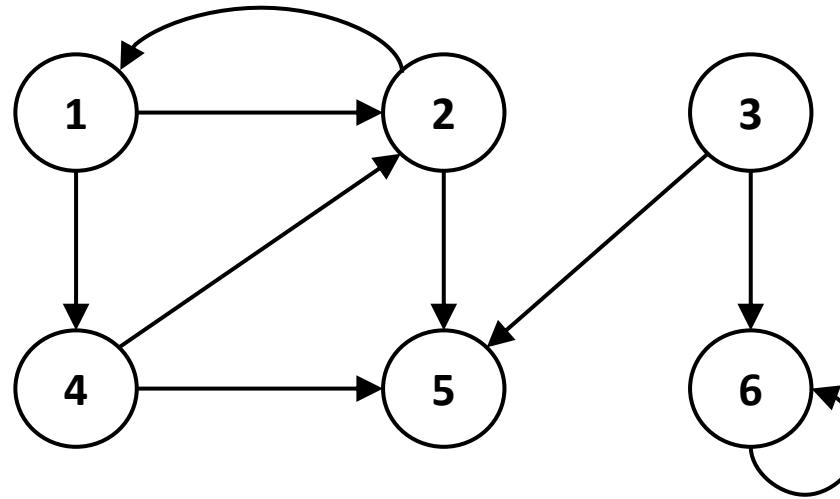
Adjacency matrix

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 |
| 4 | 0 | 1 | 0 | 0 | 1 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 |
| 6 | 0 | 0 | 0 | 0 | 0 | 1 |

Size: $\Theta(n^2)$

When to use one representation over the other?



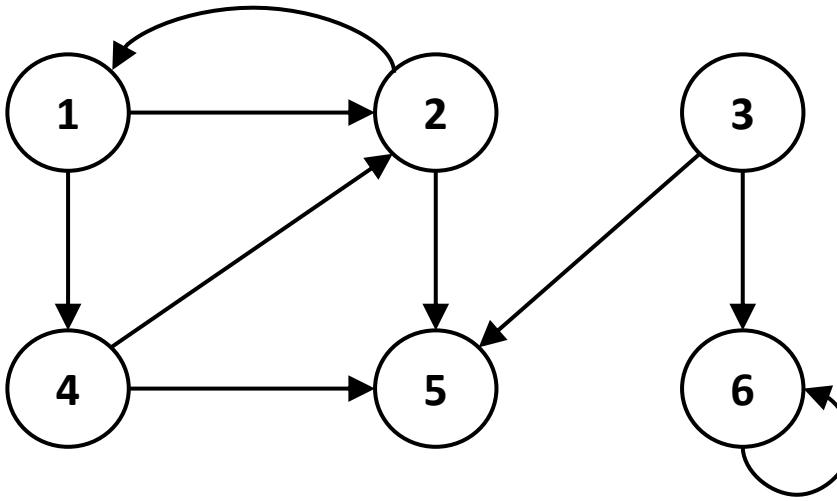


Adjacency list Size: $\Theta(n + m)$

| | | | | |
|---|---------------|---|-------------------|---|
| 1 | \rightarrow | 2 | \longrightarrow | 4 |
| 2 | \rightarrow | 5 | \longrightarrow | 1 |
| 3 | \rightarrow | 5 | \longrightarrow | 6 |
| 4 | \rightarrow | 2 | \longrightarrow | 5 |
| 5 | \rightarrow | 6 | | |
| 6 | \rightarrow | 6 | | |

Adjacency matrix Size: $\Theta(n^2)$

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 |
| 4 | 0 | 1 | 0 | 0 | 1 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 |
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Adjacency list Size: $\Theta(n + m)$

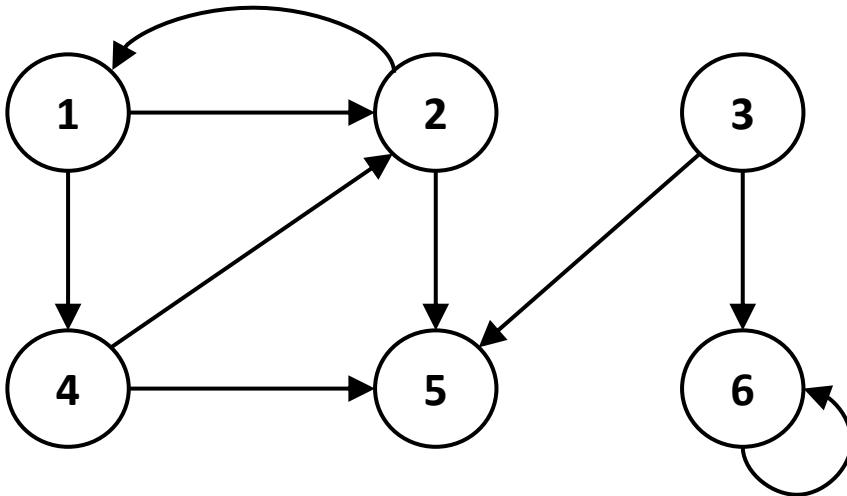
| | | | | |
|---|---------------|---|-------------------|---|
| 1 | \rightarrow | 2 | \longrightarrow | 4 |
| 2 | \rightarrow | 5 | \longrightarrow | 1 |
| 3 | \rightarrow | 5 | \longrightarrow | 6 |
| 4 | \rightarrow | 2 | \longrightarrow | 5 |
| 5 | \rightarrow | | | |
| 6 | \rightarrow | 6 | | |

edge (3,6) is in G?

Adjacency matrix Size: $\Theta(n^2)$

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 |
| 4 | 0 | 1 | 0 | 0 | 1 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 |
| 6 | 0 | 0 | 0 | 0 | 0 | 1 |





Adjacency list Size: $\Theta(n + m)$

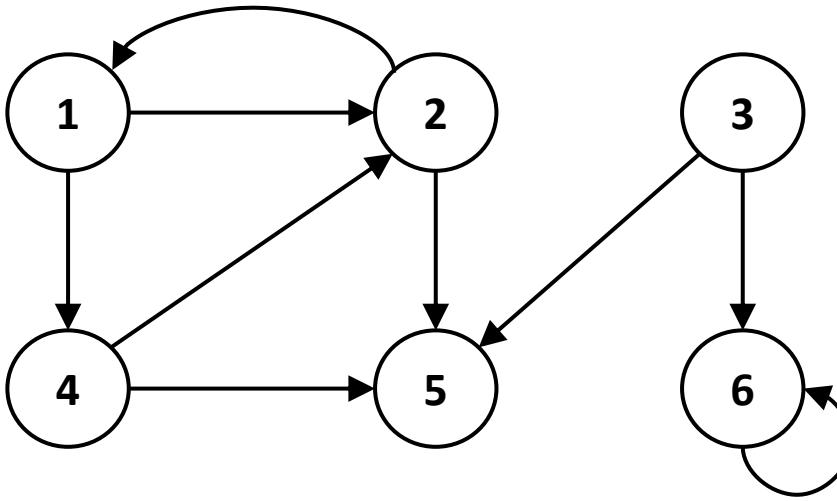
| | | | | |
|---|---------------|---|-------------------|---|
| 1 | \rightarrow | 2 | \longrightarrow | 4 |
| 2 | \rightarrow | 5 | \longrightarrow | 1 |
| 3 | \rightarrow | 5 | \longrightarrow | 6 |
| 4 | \rightarrow | 2 | \longrightarrow | 5 |
| 5 | | | | |
| 6 | \rightarrow | 6 | | |

edge (3,6) is in G?

Adjacency matrix Size: $\Theta(n^2)$

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 |
| 4 | 0 | 1 | 0 | 0 | 1 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 |
| 6 | 0 | 0 | 0 | 0 | 0 | 1 |





Adjacency list Size: $\Theta(n + m)$

| | | | | |
|---|---------------|---|-------------------|---|
| 1 | \rightarrow | 2 | \longrightarrow | 4 |
| 2 | \rightarrow | 5 | \longrightarrow | 1 |
| 3 | \rightarrow | 5 | \longrightarrow | 6 |
| 4 | \rightarrow | 2 | \longrightarrow | 5 |
| 5 | \rightarrow | 6 | | |
| 6 | \rightarrow | 6 | | |

edge (3,6) is in G?

$A[3,6] = 1$?

Adjacency matrix Size: $\Theta(n^2)$

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 |
| 4 | 0 | 1 | 0 | 0 | 1 | 0 |
| 5 | 0 | 0 | 0 | 0 | 0 | 0 |
| 6 | 0 | 0 | 0 | 0 | 0 | 1 |



When to use one representation over the other?

- For sparse graphs ($m \ll n^2$), Adjacency List is more space-efficient.



When to use one representation over the other?

- For sparse graphs ($m \ll n^2$), Adjacency List is more space-efficient.
- Finding whether a certain $(u, v) \in E$:
 - Is $O(1)$ with Adjacency Matrix
 - Is $O(n)$ with Adjacency List



Graph Search

Systematic exploration of graph:



Graph Search

Systematic exploration of graph:

- Start at some node and explore it



Graph Search

Systematic exploration of graph:

- Start at some node and explore it
 - i.e., follow its edges to discover new nodes



Graph Search

Systematic exploration of graph:

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 - i.e., follow its edges to discover new nodes
- Explore newly discovered nodes, i.e., follow edges from those nodes etc...



Graph Search

Systematic exploration of graph:

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Graph searches reveal structural properties of G:



Graph Search

Systematic exploration of graph:

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Graph searches reveal structural properties of G:

Is G connected?



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Graph searches reveal structural properties of G:

Is G connected?

Does G have a cycle?



Graph Search

Systematic exploration of graph:

- Start at some node and explore it
 - i.e., follow its edges to discover new nodes
- Explore newly discovered nodes, i.e., follow edges from those nodes etc...

Graph searches reveal structural properties of G:

Is G connected?

Does G have a cycle?

Shortest path info,

etc...



Graph Search

Two basic types of graph searches:



Graph Search

Two basic types of graph searches:

- Breadth First Search (BFS)
- Depth First Search (DFS)



Breadth First Search



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Breadth First Search

BFS started at a node s



Breadth First Search

BFS started at a node s

- (1) Explore s (i.e. follow all the edges out of s to discover new nodes)



Breadth First Search

BFS started at a node s

- (1) Explore s (i.e. follow all the edges out of s to discover new nodes)
- (2) Explore discovered nodes in the **order of their discovery**



Breadth First Search

BFS started at a node s

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[“First Discovered-First Explored” policy]



Breadth First Search

BFS started at a node s

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[“First Discovered-First Explored” policy]

To do so, put them in a Queue and explore them in FIFO order



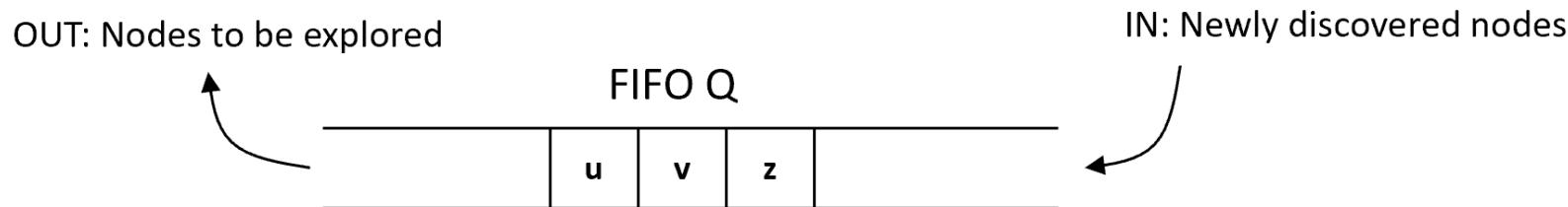
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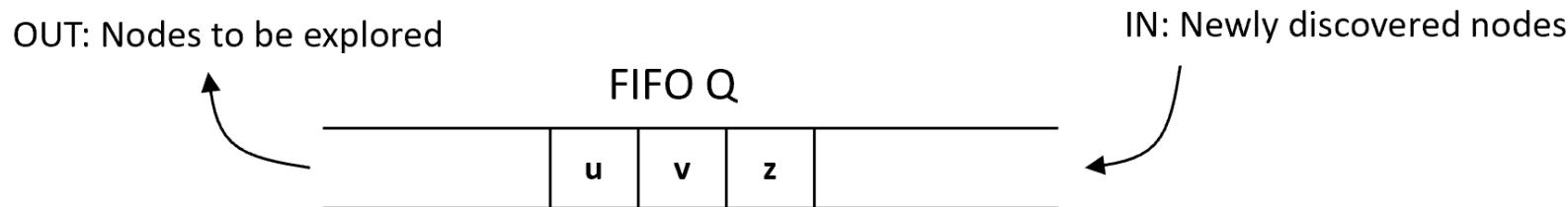
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[“First Discovered-First Explored” policy]

To do so, put them in a Queue and explore them in FIFO order



Do (2) until all discovered nodes are explored



BFS algorithm maintains the following information for each node v :



BFS algorithm maintains the following information for each node v :

1. $\text{color}[v] =$



BFS algorithm maintains the following information for each node v :

1. $\text{color}[v] = \begin{cases} \text{white} & : v \text{ is undiscovered} \end{cases}$



BFS algorithm maintains the following information for each node v :

1. $\text{color}[v] = \begin{cases} \text{white} & : v \text{ is undiscovered} \\ \text{grey} & : v \text{ was discovered but not yet explored} \end{cases}$



BFS algorithm maintains the following information for each node v :

1. $\text{color}[v] = \begin{cases} \text{white} & : v \text{ is undiscovered} \\ \text{grey} & : v \text{ was discovered but not yet explored} \\ \text{black} & : v \text{ was discovered and explored} \end{cases}$



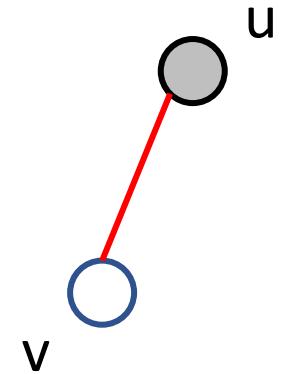
BFS algorithm maintains the following information for each node v :

2. $p[v]$



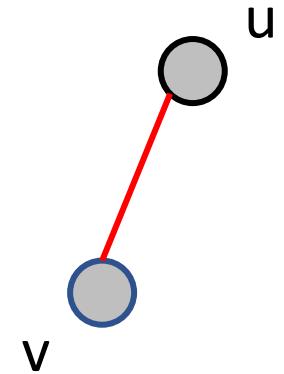
BFS algorithm maintains the following information for each node v :

2. $p[v] = u$: node v was discovered while exploring u



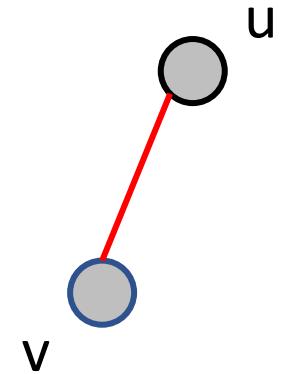
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s



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$$s \rightarrow u_1$$



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: i.e. u discovered v

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$$s \rightarrow u_1 \rightarrow u_2$$



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2. $p[v] = u$: node v was discovered while exploring u
: i.e. u discovered v

3. $d[v]$: Length of **discovery path** from s

$s \rightarrow u_1 \rightarrow u_2 \rightarrow \dots$



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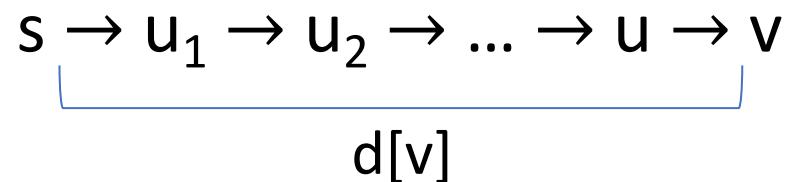
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$s \rightarrow u_1 \rightarrow u_2 \rightarrow \dots \rightarrow u \rightarrow v$



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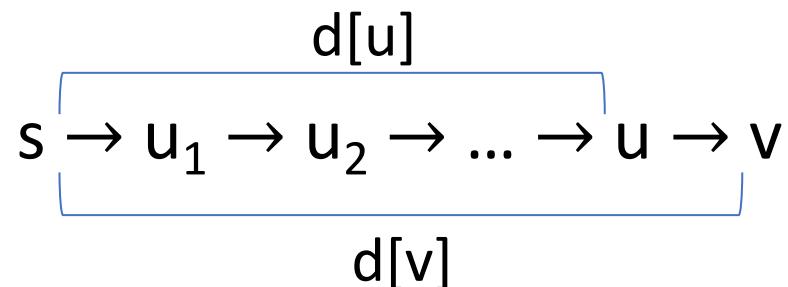
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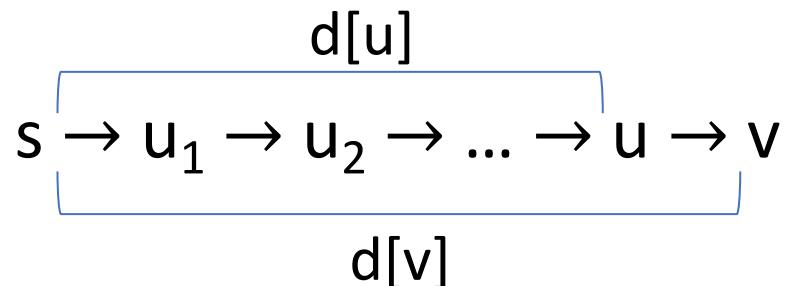
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BFS algorithm maintains the following information for each node v :

2. $p[v] = u$: node v was discovered while exploring u
: i.e. u discovered v

3. $d[v]$: Length of discovery path from s



Clearly, $d[v] = d[u] + 1$



BFS(G, s)

/ G = (V, E) and s ∈ V */*



BFS(G, s)

/* G = (V, E) and s ∈ V */

color[s] ← grey ; d[s] ← 0 ; p[s] ← NIL



BFS(G, s)

/* G = (V, E) and s ∈ V */

color[s] ← grey ; d[s] ← 0 ; p[s] ← NIL

For each v ∈ V – {s} **do**



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/* G = (V, E) and s ∈ V */

color[s] ← grey ; d[s] ← 0 ; p[s] ← NIL

For each v ∈ V – {s} **do**

color[v] ← white



BFS(G, s)

/* G = (V, E) and s ∈ V */

color[s] ← grey ; d[s] ← 0 ; p[s] ← NIL

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For each v ∈ V – {s} **do**

 color[v] ← white

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/* G = (V, E) and s ∈ V */

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For each v ∈ V – {s} **do**

 color[v] ← white

 d[v] ← ∞

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BFS(G, s)

/* G = (V, E) and s ∈ V */

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 color[v] ← white

 d[v] ← ∞

 p[v] ← NIL

Q ← empty ;

/* Q: nodes that are discovered but not yet explored */



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

color[v] \leftarrow white

$d[v] \leftarrow \infty$

$p[v] \leftarrow \text{NIL}$

$Q \leftarrow \text{empty} ; \text{ENQ}(Q, s)$

/* Q : nodes that are discovered but not yet explored */



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

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$d[v] \leftarrow \infty$

$p[v] \leftarrow \text{NIL}$

$Q \leftarrow \text{empty} ; \text{ENQ}(Q, s)$

/* Q : nodes that are discovered but not yet explored */

While Q is not empty **do**



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

color[v] \leftarrow white

$d[v] \leftarrow \infty$

$p[v] \leftarrow \text{NIL}$

$Q \leftarrow \text{empty} ; \text{ENQ}(Q, s)$

/* Q : nodes that are discovered but not yet explored */

While Q is not empty **do**

$u \leftarrow \text{DEQ}(Q)$

/* Explore u */



BFS(G, s)

/ G = (V, E) and $s \in V$ */*

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

 color[v] \leftarrow white

$d[v] \leftarrow \infty$

$p[v] \leftarrow \text{NIL}$

$Q \leftarrow \text{empty} ; \text{ENQ}(Q, s)$

/ Q: nodes that are discovered but not yet explored */*

While Q is not empty **do**

$u \leftarrow \text{DEQ}(Q)$

/ Explore u */*

For each $(u, v) \in E$ **do**

/ Explore edge (u,v) */*



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

 color[v] \leftarrow white

$d[v] \leftarrow \infty$

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While Q is not empty **do**

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/* Explore u */

For each $(u, v) \in E$ **do**

/* Explore edge (u, v) */

If color[v] = **then do**

/* If v is first discovered */



BFS(G, s)

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color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

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 /* Explore edge (u, v) */

If color[v] = white **then do**

 /* If v is first discovered */

 color[v] \leftarrow



BFS(G, s)

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color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

 color[v] \leftarrow white

$d[v] \leftarrow \infty$

$p[v] \leftarrow \text{NIL}$

$Q \leftarrow \text{empty}$; ENQ(Q, s)

/* Q : nodes that are discovered but not yet explored */

While Q is not empty **do**

$u \leftarrow \text{DEQ}(Q)$

 /* Explore u */

For each $(u, v) \in E$ **do**

 /* Explore edge (u,v) */

If color[v] = white **then do**

 /* If v is first discovered */

 color[v] \leftarrow grey



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

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 /* Explore u */

For each $(u, v) \in E$ **do**

 /* Explore edge (u, v) */

If color[v] = white **then do**

 /* If v is first discovered */

 color[v] \leftarrow grey

$d[v] \leftarrow$



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

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While Q is not empty **do**

$u \leftarrow \text{DEQ}(Q)$

 /* Explore u */

For each $(u, v) \in E$ **do**

 /* Explore edge (u, v) */

If color[v] = white **then do**

 /* If v is first discovered */

 color[v] \leftarrow grey

$d[v] \leftarrow d[u] + 1$



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

 color[v] \leftarrow white

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$d[v] \leftarrow d[u] + 1$

$p[v] \leftarrow$



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$p[v] \leftarrow u$



BFS(G, s)

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$\text{color}[v] \leftarrow \text{white}$

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$Q \leftarrow \text{empty}$; $\text{ENQ}(Q, s)$

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While Q is not empty **do**

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/* Explore u */

For each $(u, v) \in E$ **do**

/* Explore edge (u, v) */

If $\text{color}[v] = \text{white}$ **then do**

/* If v is first discovered */

$\text{color}[v] \leftarrow \text{grey}$

$d[v] \leftarrow d[u] + 1$

$p[v] \leftarrow u$

$\text{ENQ}(Q, v)$



BFS(G, s)

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color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

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If color[v] = white **then do**

 /* If v is first discovered */

 color[v] \leftarrow grey

$d[v] \leftarrow d[u] + 1$

$p[v] \leftarrow u$

 ENQ(Q, v)

End If



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

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 /* If v is first discovered */

 color[v] \leftarrow grey

$d[v] \leftarrow d[u] + 1$

$p[v] \leftarrow u$

 ENQ(Q, v)

End If

End For



BFS(G, s)

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For each $(u, v) \in E$ **do**

/* Explore edge (u, v) */

If color[v] = white **then do**

/* If v is first discovered */

 color[v] \leftarrow grey

$d[v] \leftarrow d[u] + 1$

$p[v] \leftarrow u$

 ENQ(Q, v)

End If

End For

 color[u] \leftarrow

/* Done exploring u */



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

$\text{color}[s] \leftarrow \text{grey}$; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

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$\text{color}[v] \leftarrow \text{white}$
 $d[v] \leftarrow \infty$
 $p[v] \leftarrow \text{NIL}$

$Q \leftarrow \text{empty}$; $\text{ENQ}(Q, s)$

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$u \leftarrow \text{DEQ}(Q)$

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For each $(u, v) \in E$ **do**

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If $\text{color}[v] = \text{white}$ **then do**

/* If v is first discovered */

$\text{color}[v] \leftarrow \text{grey}$

$d[v] \leftarrow d[u] + 1$

$p[v] \leftarrow u$

$\text{ENQ}(Q, v)$

End If

End For

$\text{color}[u] \leftarrow \text{black}$

/* Done exploring u */



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

color[s] \leftarrow grey ; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

 color[v] \leftarrow white

$d[v] \leftarrow \infty$

$p[v] \leftarrow \text{NIL}$

$Q \leftarrow \text{empty}$; ENQ(Q, s)

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 /* If v is first discovered */

 color[v] \leftarrow grey

$d[v] \leftarrow d[u] + 1$

$p[v] \leftarrow u$

 ENQ(Q, v)

End If

End For

 color[u] \leftarrow black

 /* Done exploring u */

End While

End BFS



BFS(G, s)

/* $G = (V, E)$ and $s \in V$ */

$\text{color}[s] \leftarrow \text{grey}$; $d[s] \leftarrow 0$; $p[s] \leftarrow \text{NIL}$

For each $v \in V - \{s\}$ **do**

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 $p[v] \leftarrow \text{NIL}$

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While Q is not empty **do**

$u \leftarrow \text{DEQ}(Q)$

/* Explore u */

For each $(u, v) \in E$ **do**

/* Explore edge (u, v) */

If $\text{color}[v] = \text{white}$ **then do**

/* If v is first discovered */

$\text{color}[v] \leftarrow \text{grey}$

$d[v] \leftarrow d[u] + 1$

$p[v] \leftarrow u$

$\text{ENQ}(Q, v)$

End If

End For

$\text{color}[u] \leftarrow \text{black}$

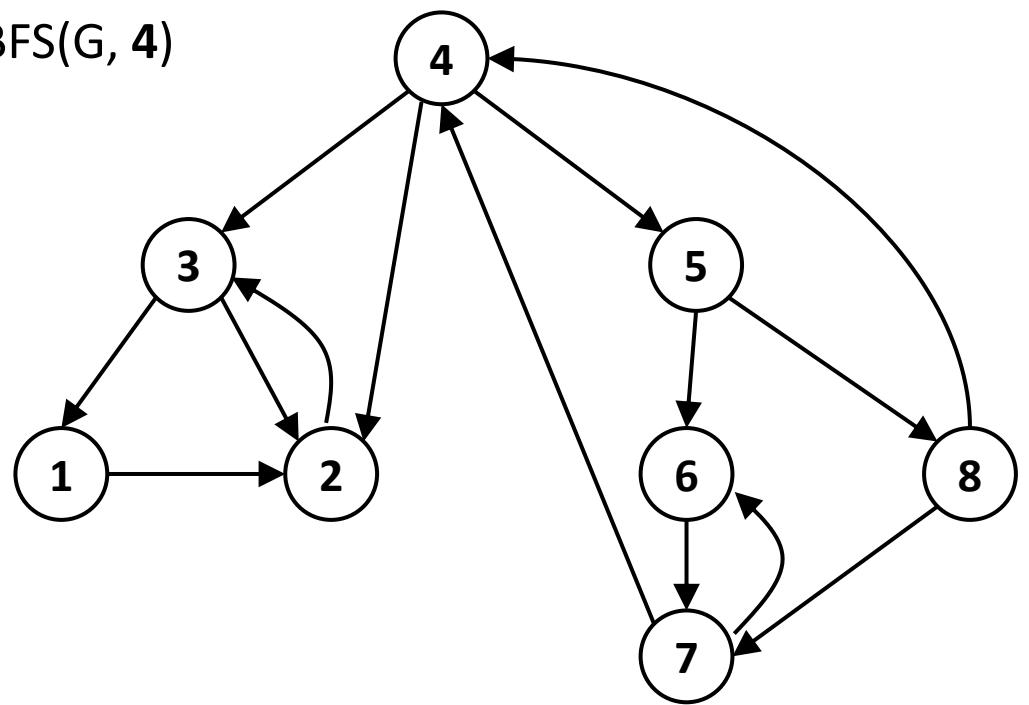
/* Done exploring u */

End While

End BFS



BFS(G, 4)



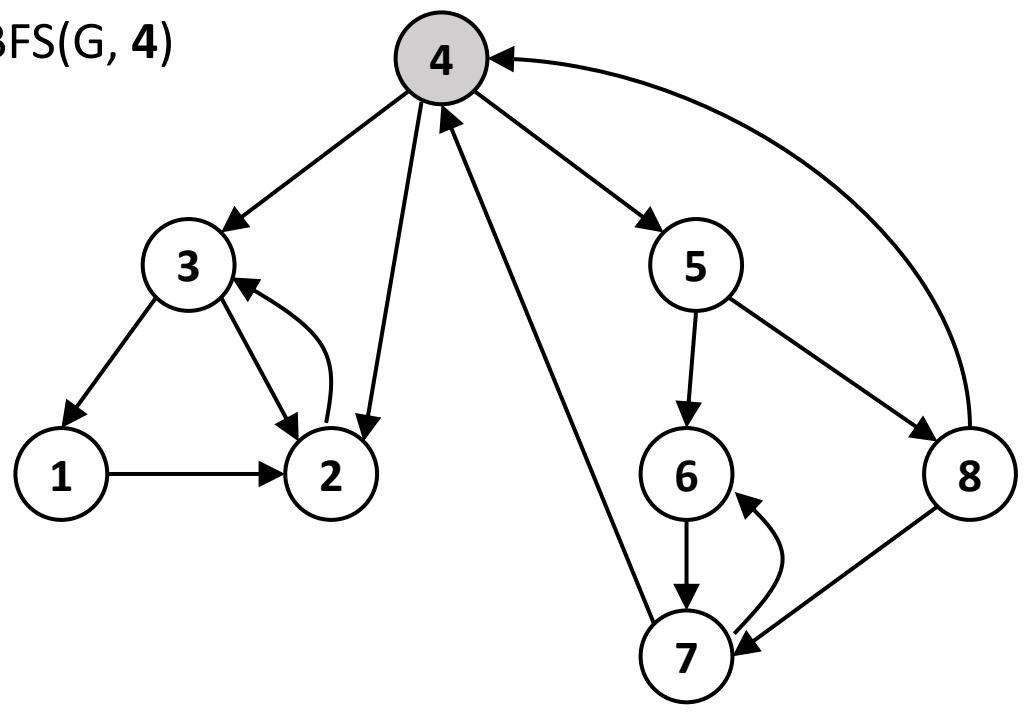
Adj List of G :

| | |
|---|-------------|
| 1 | → 2 |
| 2 | → 3 |
| 3 | → 1 → 2 |
| 4 | → 3 → 2 → 5 |
| 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

Contents of Q :



BFS(G, 4)



4

Adj List of G :

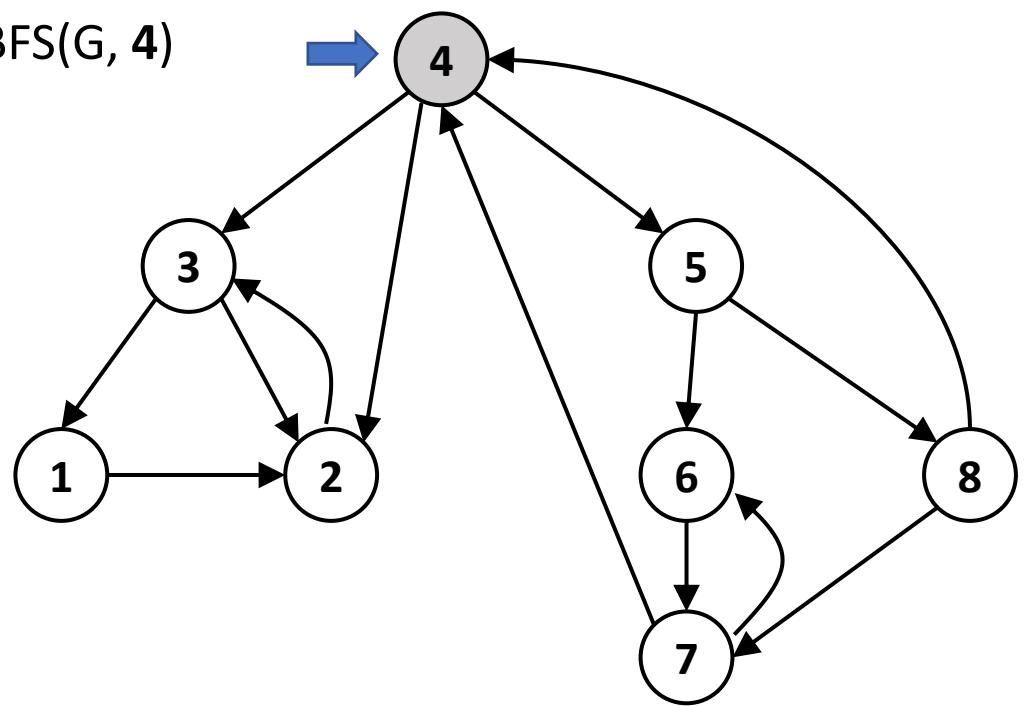
| | | |
|---|---|-----------|
| 1 | → | 2 |
| 2 | → | 3 |
| 3 | → | 1 → 2 |
| 4 | → | 3 → 2 → 5 |
| 5 | → | 6 → 8 |
| 6 | → | 7 |
| 7 | → | 6 → 4 |
| 8 | → | 7 → 4 |

Contents of Q :

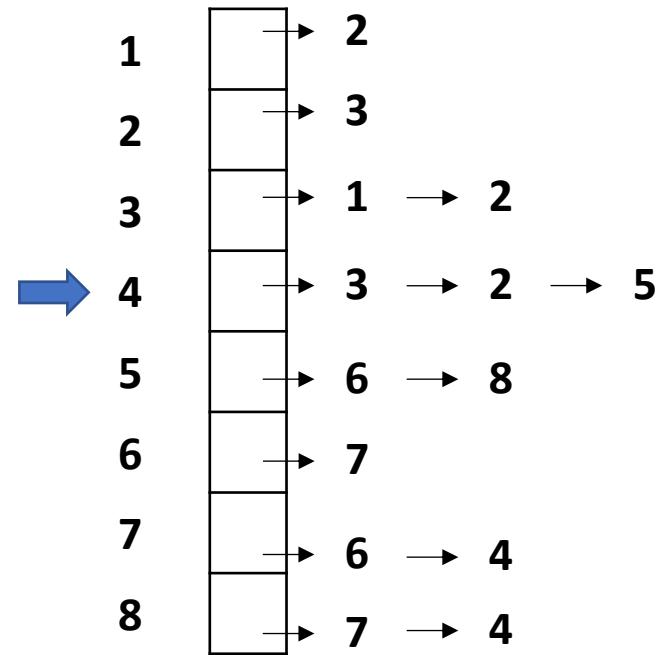
d = 0
4



BFS(G , 4)



Adj List of G :

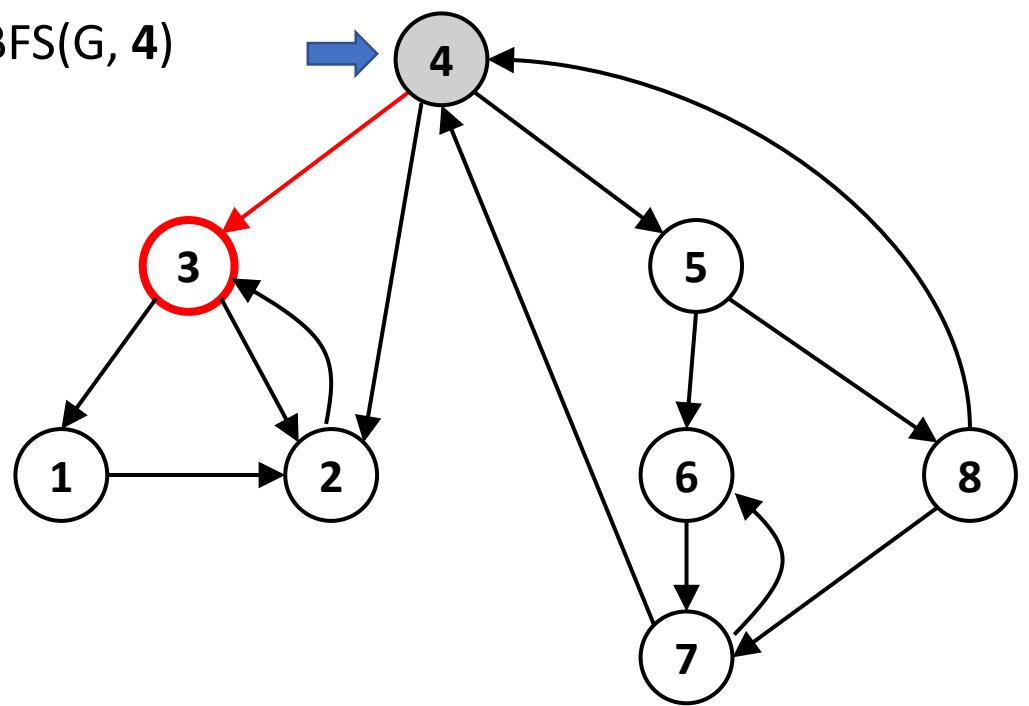


Contents of Q :

$d = 0$
4



BFS(G, 4)



Adj List of G :

→

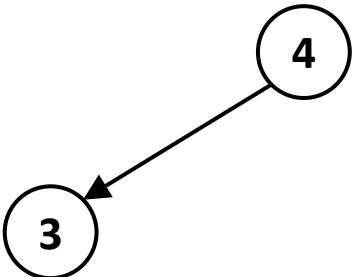
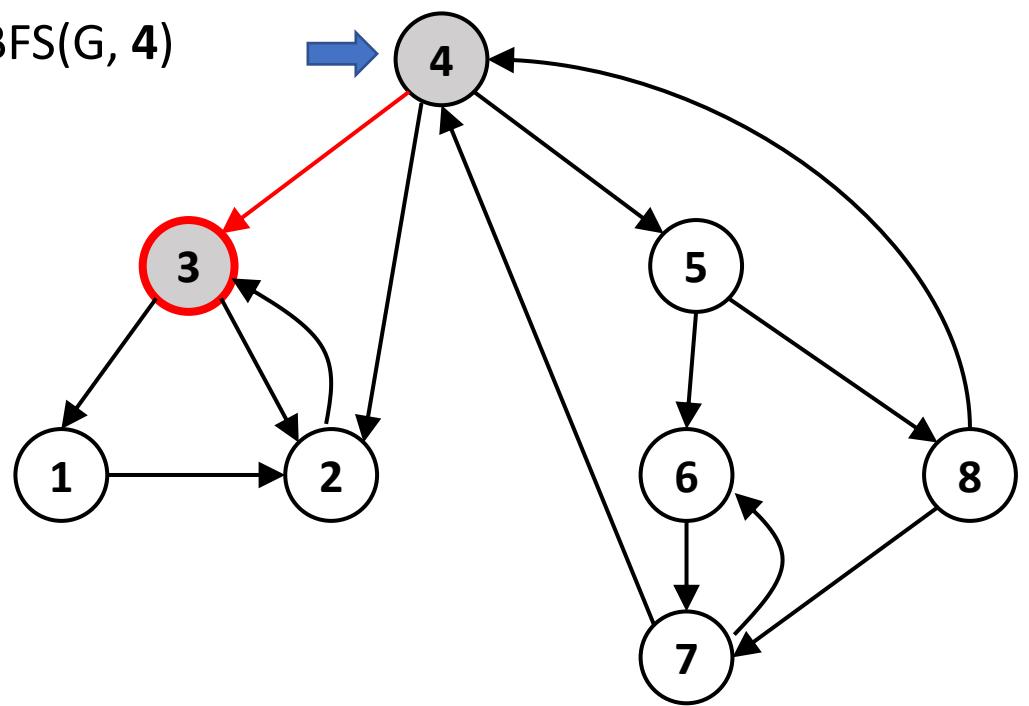
| | | |
|---|---|-----------|
| 1 | → | 2 |
| 2 | → | 3 |
| 3 | → | 1 → 2 |
| 4 | → | 3 → 2 → 5 |
| 5 | → | 6 → 8 |
| 6 | → | 7 |
| 7 | → | 6 → 4 |
| 8 | → | 7 → 4 |

Contents of Q :

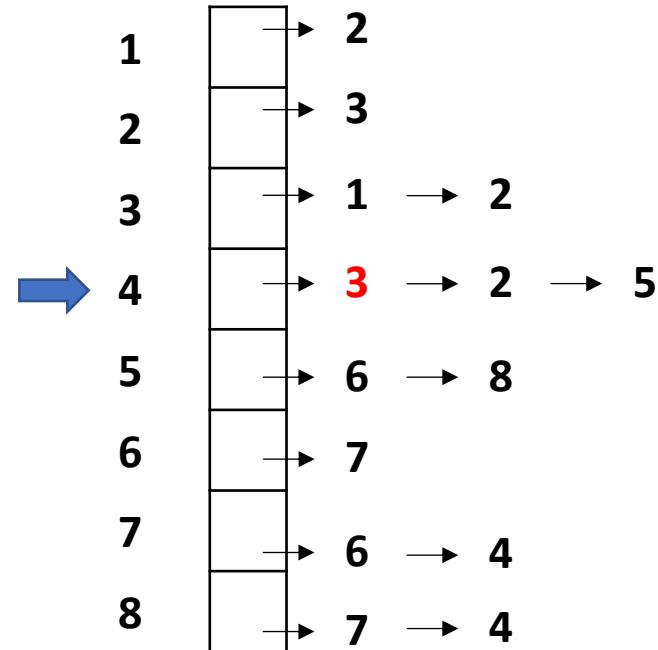
d = 0
4



BFS(G , 4)



Adj List of G :

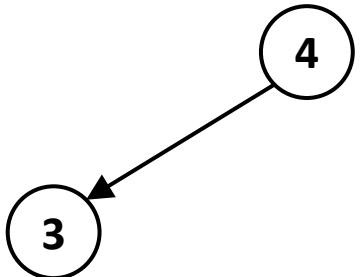
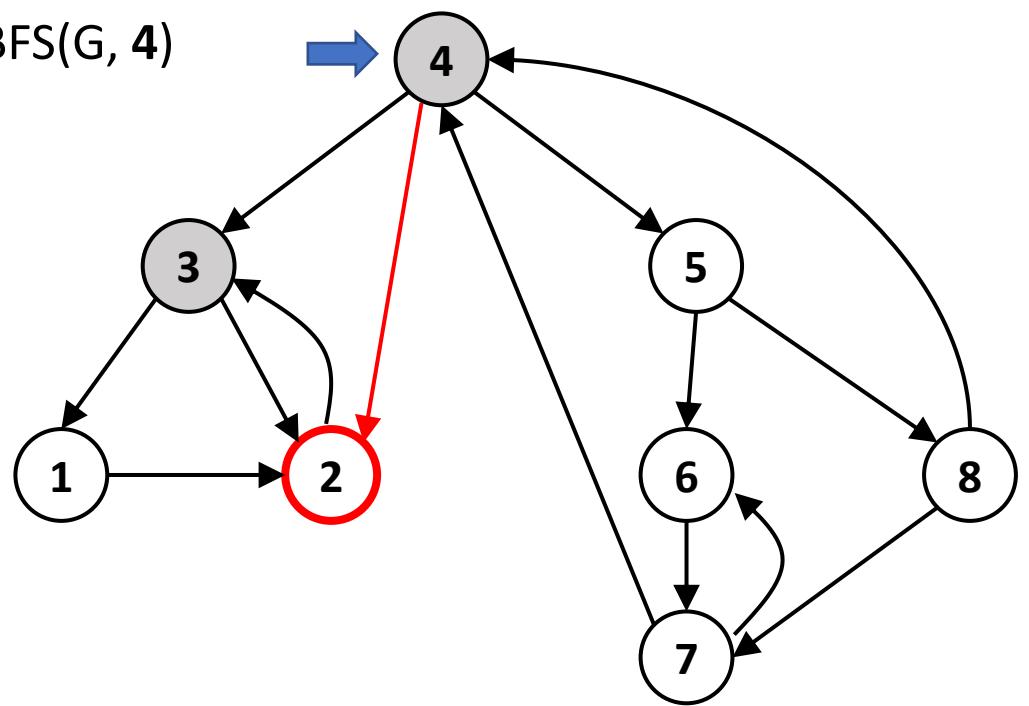


Contents of Q :

| | |
|---------|---|
| $d = 0$ | 4 |
| $d = 1$ | 3 |



BFS(G , 4)



Adj List of G :

→

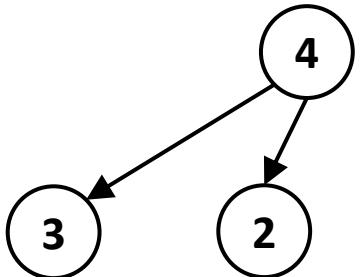
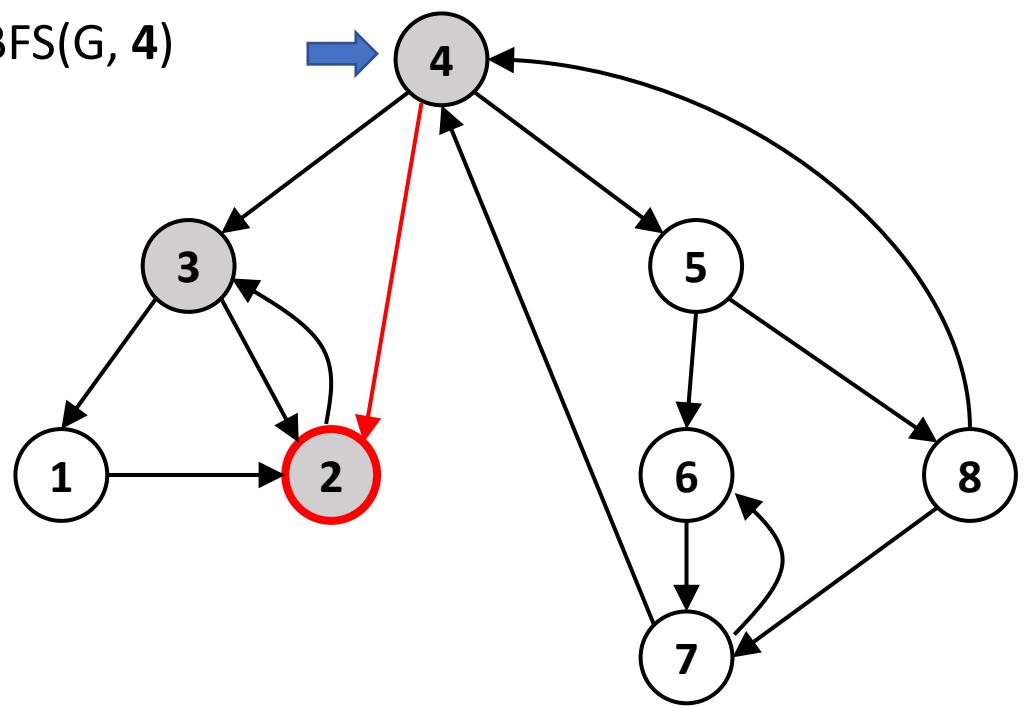
| | | |
|---|---|-----------|
| 1 | → | 2 |
| 2 | → | 3 |
| 3 | → | 1 → 2 |
| 4 | → | 3 → 2 → 5 |
| 5 | → | 6 → 8 |
| 6 | → | 7 |
| 7 | → | 6 → 4 |
| 8 | → | 7 → 4 |

Contents of Q :

| | |
|---------|---|
| $d = 0$ | 4 |
| $d = 1$ | 3 |



BFS(G , 4)



Adj List of G :

The adjacency list representation of the graph G is as follows:

| | |
|---|-------------------|
| 1 | → 2 → 3 |
| 2 | → 1 → 3 |
| 3 | → 1 → 2 → 4 |
| 4 | → 5 → 6 → 7 |
| 5 | → 2 → 8 |
| 6 | → 7 → 4 |
| 7 | → 8 → 4 |
| 8 | → 7 |

Contents of Q :

$d = 0$

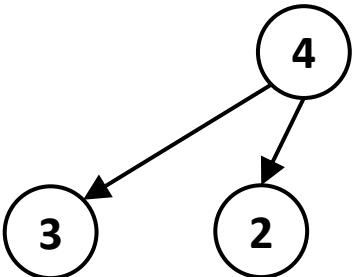
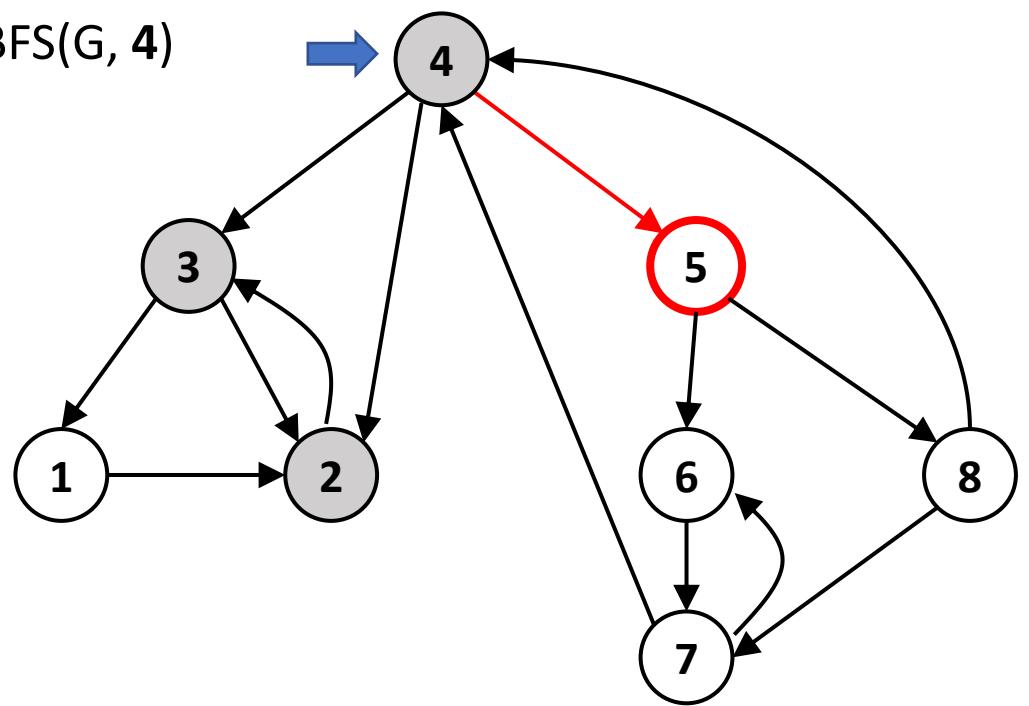
| |
|---|
| 4 |
|---|

$d = 1$

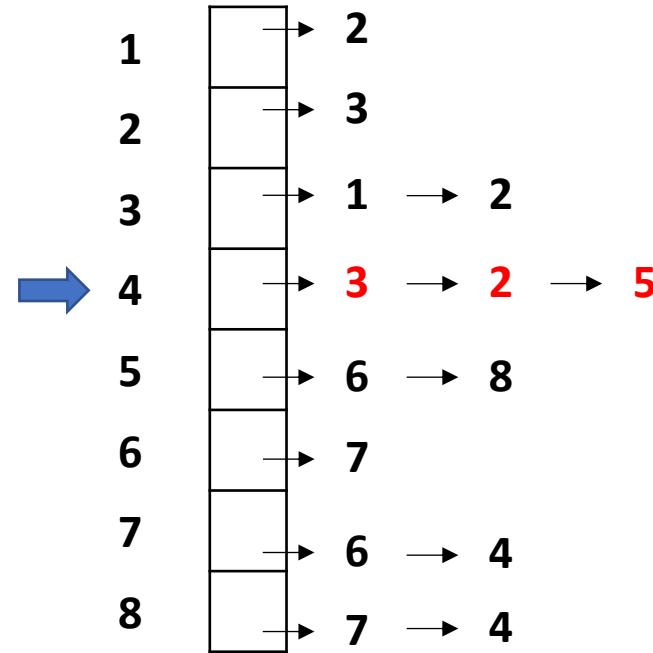
| | |
|---|---|
| 3 | 2 |
|---|---|



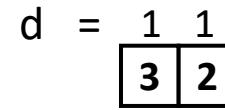
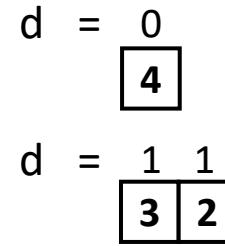
BFS(G , 4)



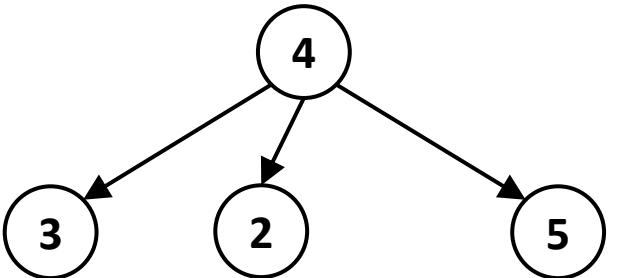
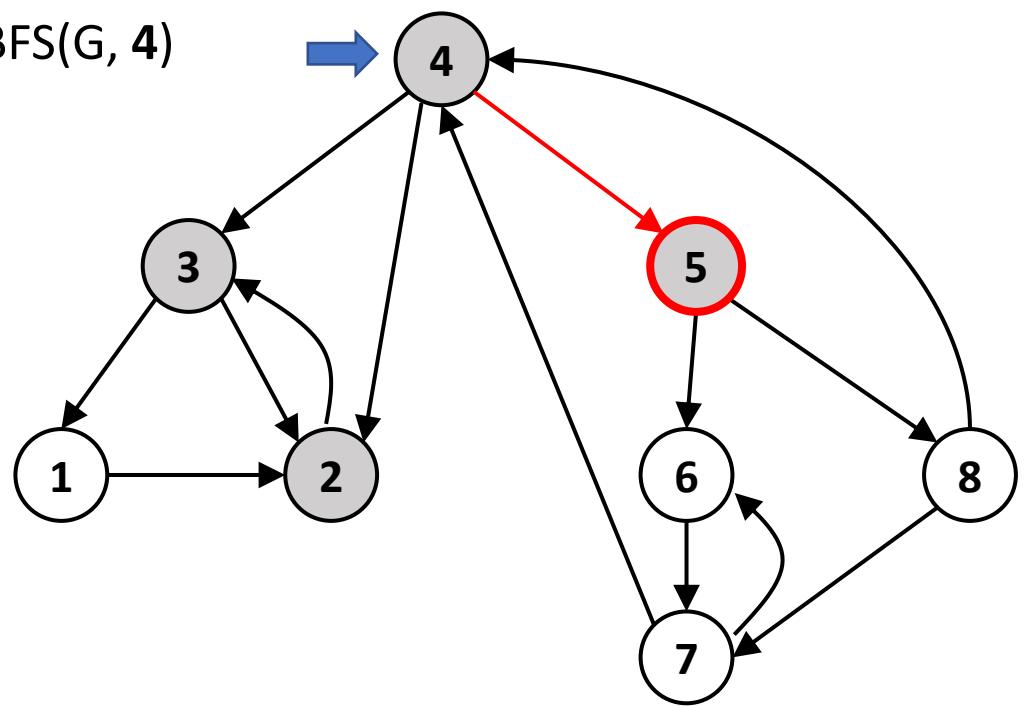
Adj List of G :



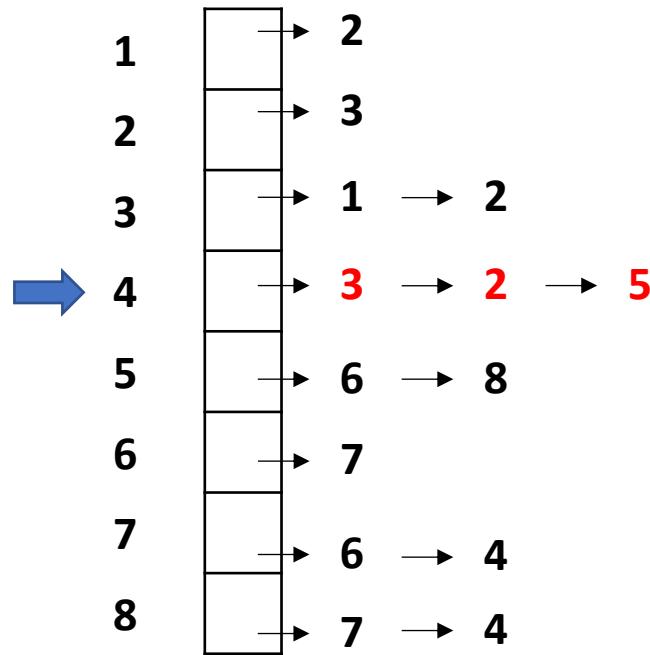
Contents of Q :



BFS(G , 4)



Adj List of G :



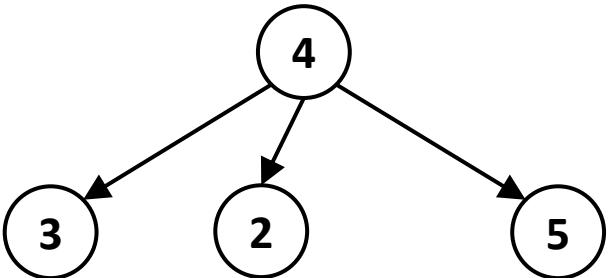
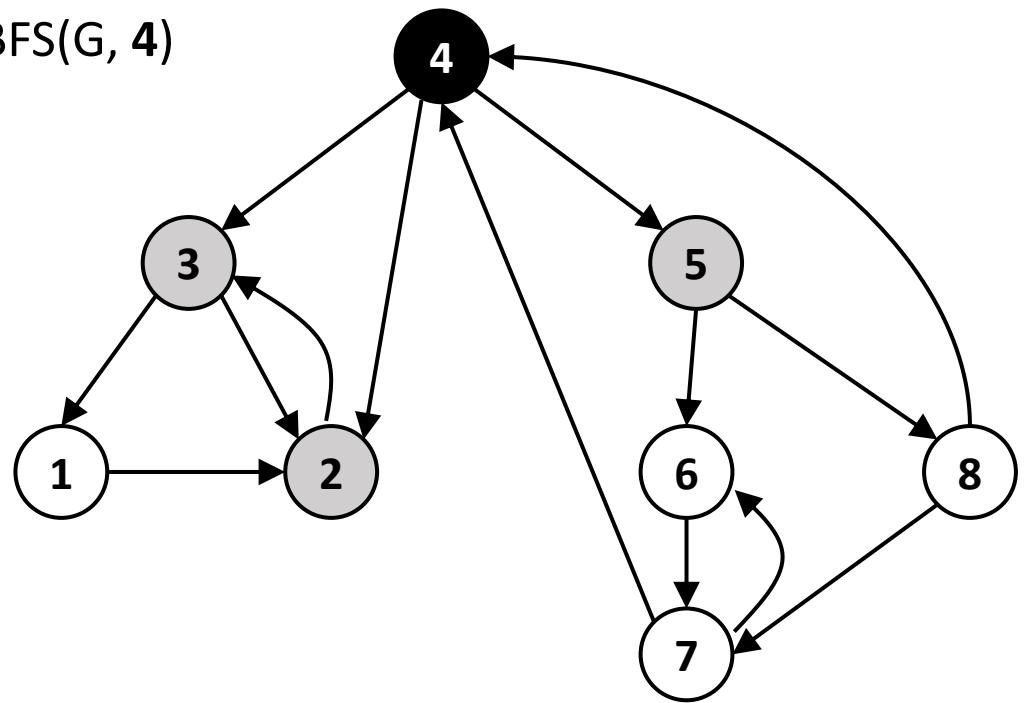
Contents of Q :

$d = 0$
4

$d = 1$ 1 1 1
3 2 5



BFS(G, 4)



Adj List of G :

| | | |
|---|---|-----------|
| 1 | → | 2 |
| 2 | → | 3 |
| 3 | → | 1 → 2 |
| 4 | ✓ | 3 → 2 → 5 |
| 5 | → | 6 → 8 |
| 6 | → | 7 |
| 7 | → | 6 → 4 |
| 8 | → | 7 → 4 |

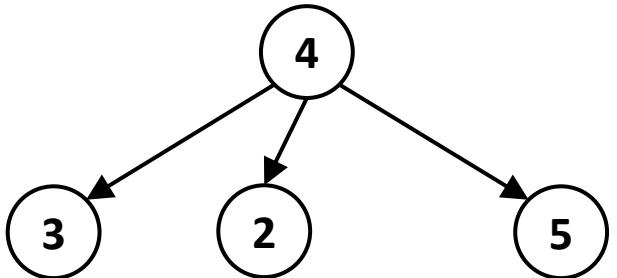
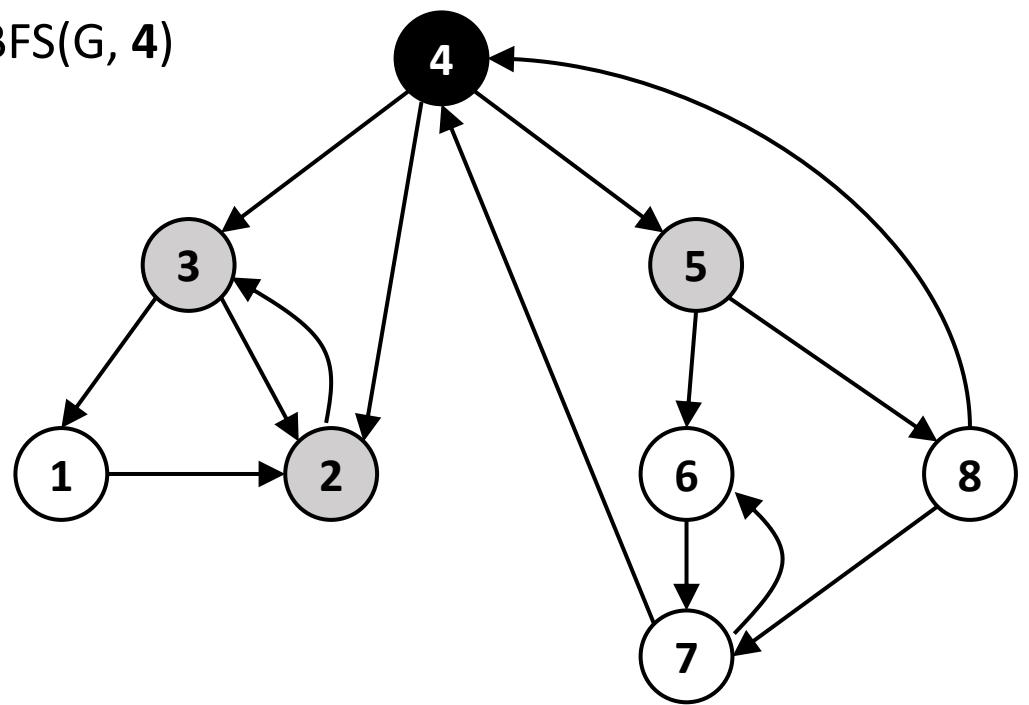
Contents of Q :

d = 0
4

d = 1 1 1
3 2 5



BFS(G, 4)



Adj List of G :

| | | |
|---|---|-----------|
| 1 | → | 2 |
| 2 | → | 3 |
| 3 | → | 1 → 2 |
| 4 | ✓ | 3 → 2 → 5 |
| 5 | → | 6 → 8 |
| 6 | → | 7 |
| 7 | → | 6 → 4 |
| 8 | → | 7 → 4 |

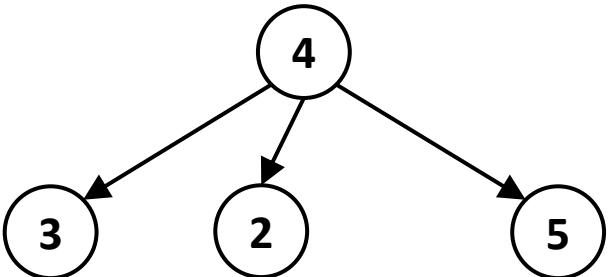
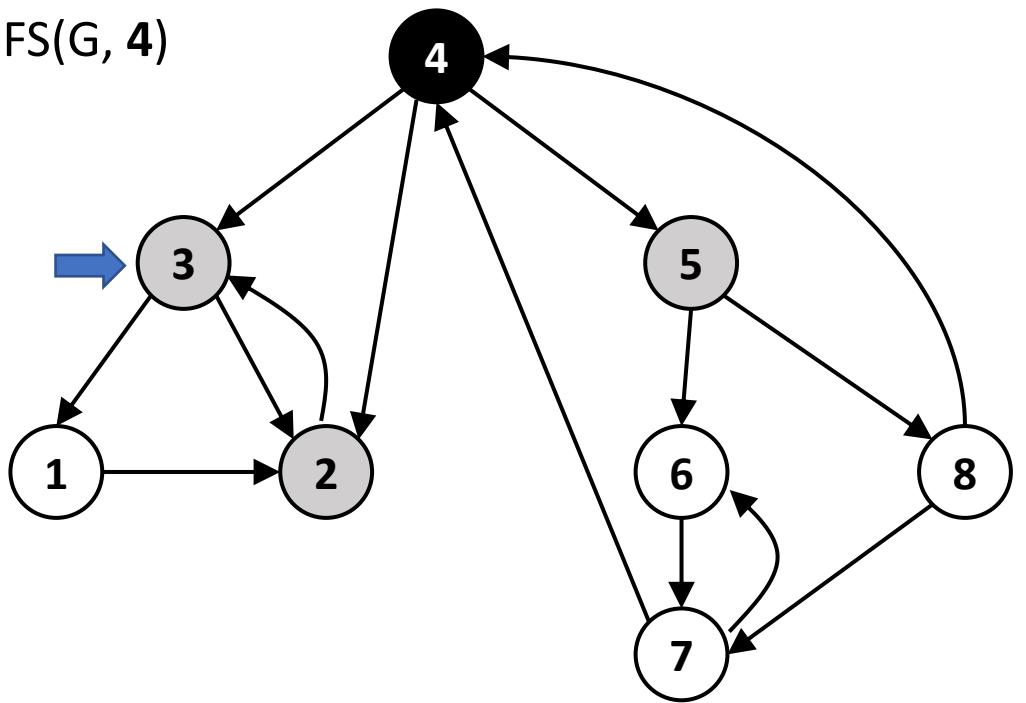
Contents of Q :

d = 0
4

d = 1 1 1
3 2 5



BFS(G , 4)



Adj List of G :

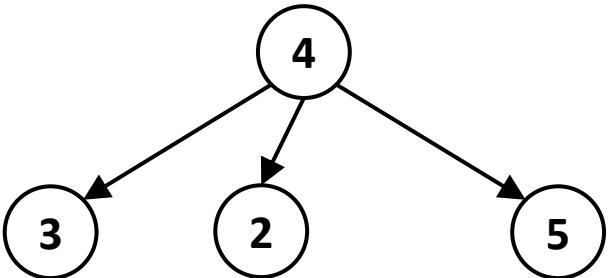
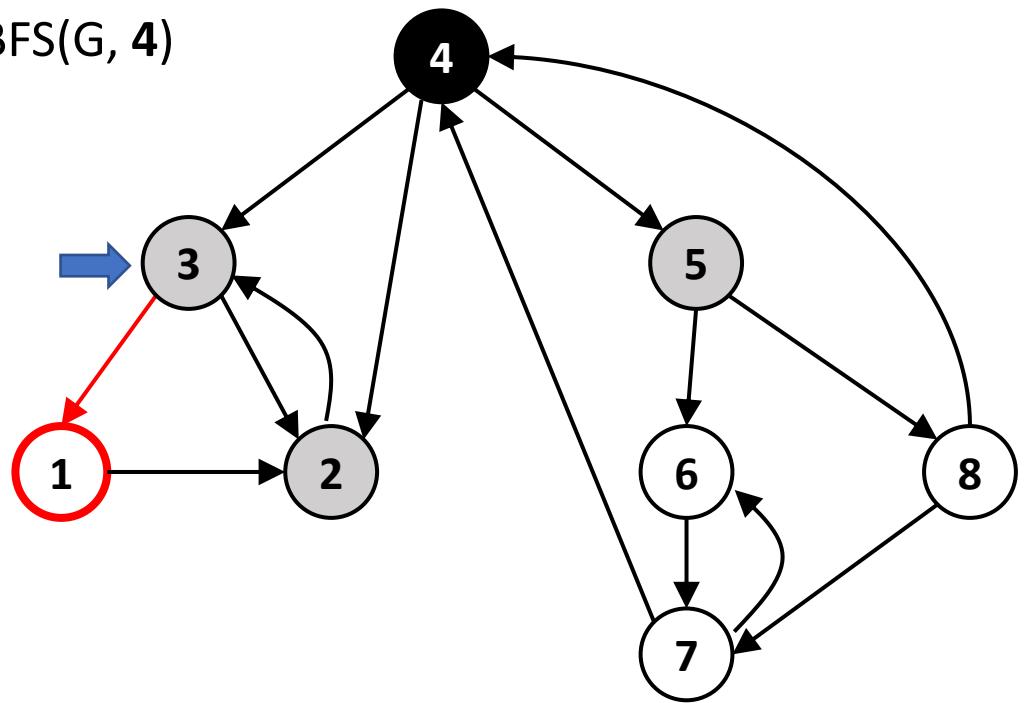
| | | |
|---|---|-----------|
| 1 | → | 2 |
| 2 | → | 3 |
| 3 | → | 1 → 2 |
| 4 | ✓ | 3 → 2 → 5 |
| 5 | → | 6 |
| 6 | → | 8 |
| 7 | → | |
| 8 | → | 6 → 4 |
| | → | 7 → 4 |

Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 |
| | 2 5 |



BFS(G , 4)



Adj List of G :

The adjacency list for graph G is shown as follows:

| | |
|---|-------------|
| 1 | → 2 |
| 2 | → 3 |
| 3 | → 1 → 2 → 5 |
| 4 | ✓ |
| 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

A blue arrow points to the row for node 3, and a checkmark is placed next to the row for node 4.

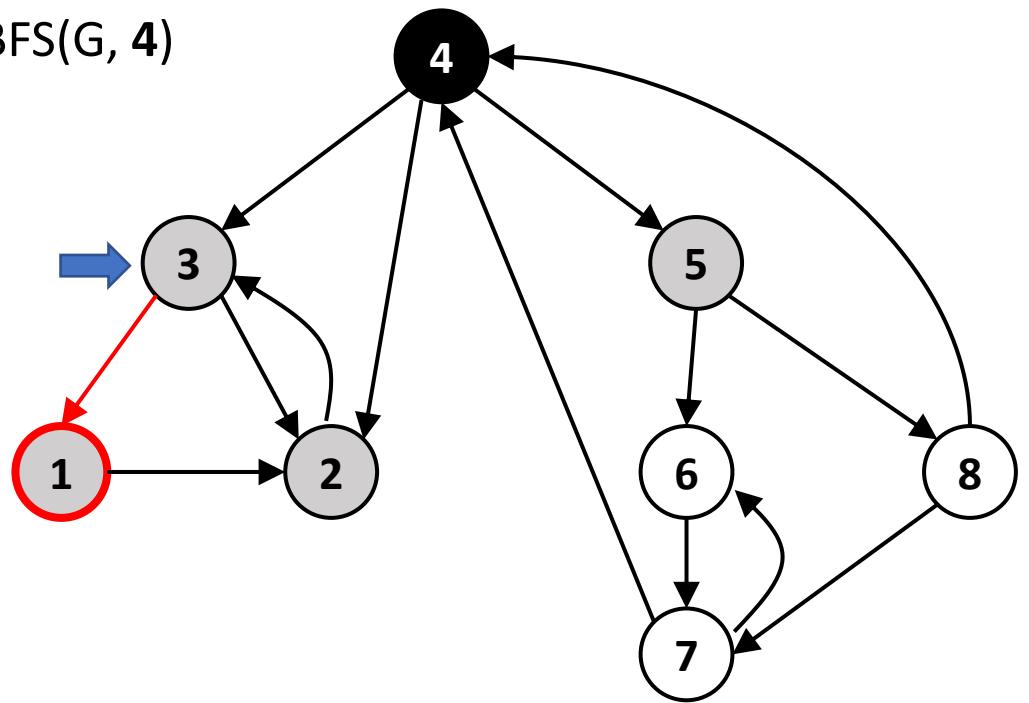
Contents of Q :

The queue Q contains nodes with their depths (d):

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 3 2 5 |
| $d = 1$ | 2 5 |



BFS(G, 4)



Adj List of G :

The adjacency list for graph G is as follows:

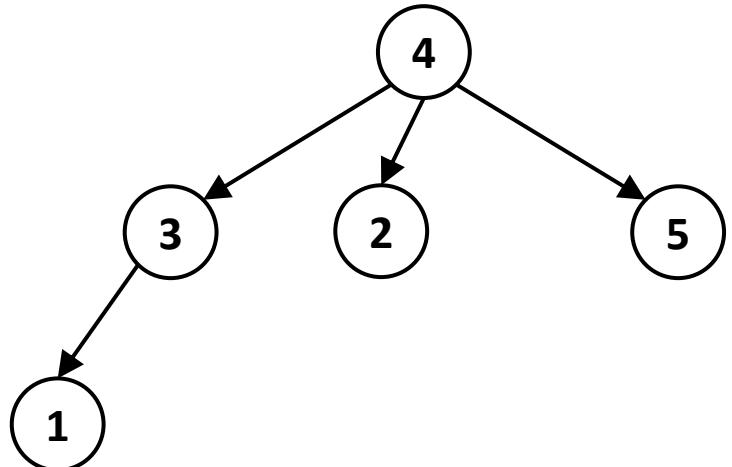
| | |
|---|-------------|
| 1 | → 2 |
| 2 | → 3 |
| 3 | → 1 → 2 → 5 |
| 4 | ✓ |
| 5 | → 3 → 2 → 8 |
| 6 | → 7 → 4 |
| 7 | → 8 → 4 |
| 8 | |

A blue arrow points to node 3, and a checkmark is placed next to node 4.

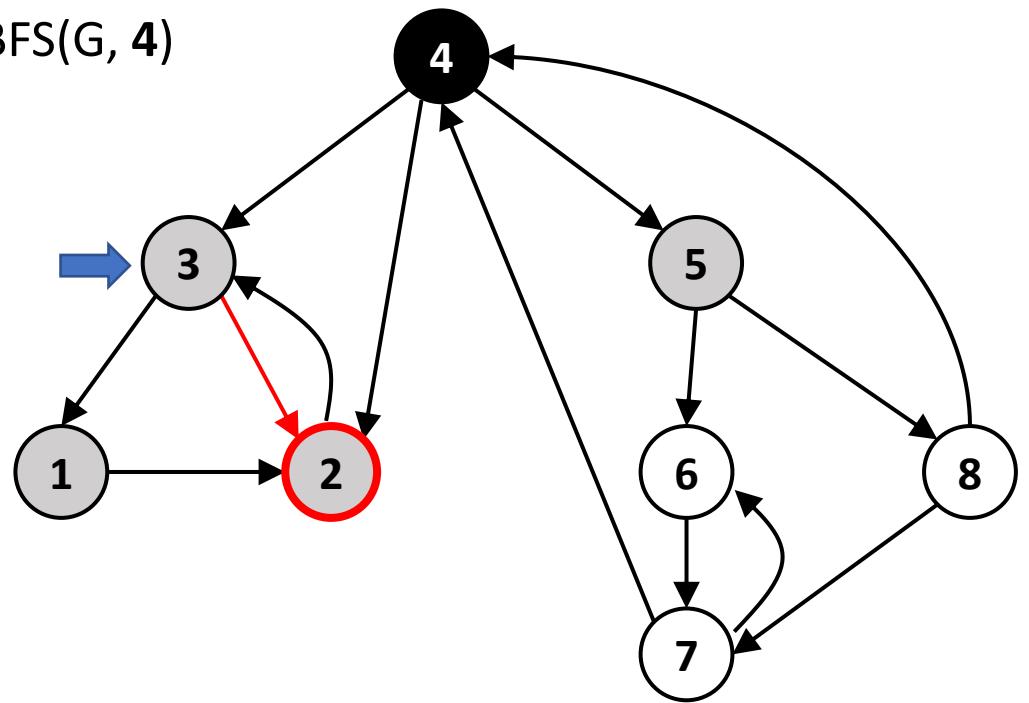
Contents of Q :

The queue Q contains nodes with their depths (d):

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| 1 1 2 | |
| 2 5 1 | |



BFS(G , 4)



Adj List of G :

The adjacency list for graph G is as follows:

| | |
|---|---------------|
| 1 | → 2 |
| 2 | → 3 |
| 3 | → 1 → 2 |
| 4 | ✓ → 3 → 2 → 5 |
| 5 | → 6 → 8 |
| 6 | → 7 → 4 |
| 7 | → 8 → 4 |
| 8 | → 7 |

Contents of Q :

$d = 0$

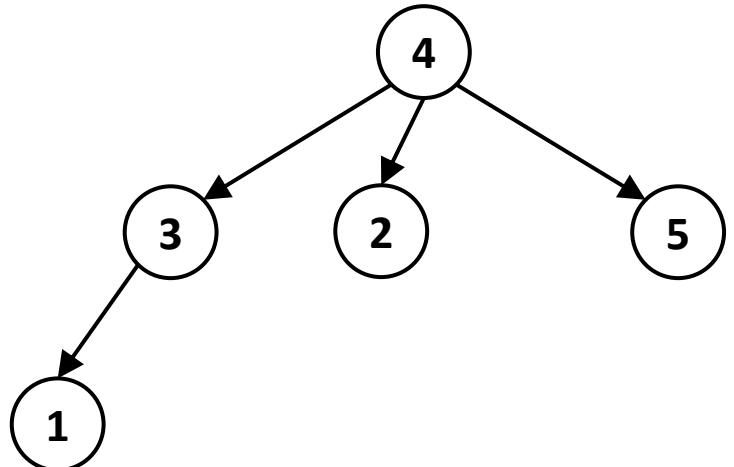
| |
|---|
| 4 |
|---|

$d = 1$

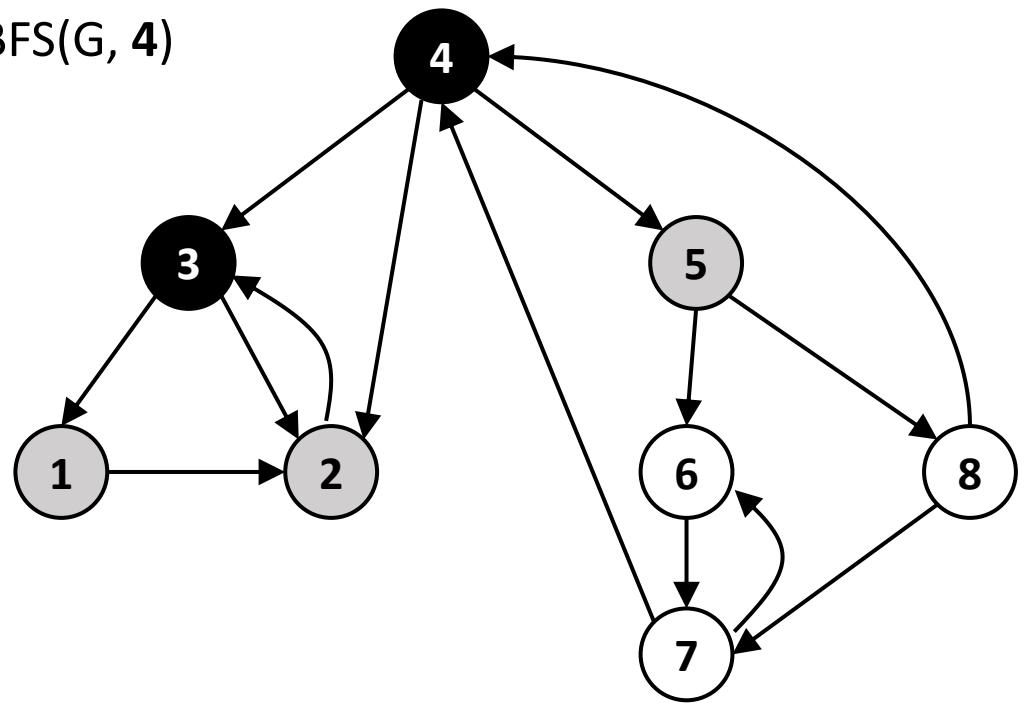
| | | |
|---|---|---|
| 3 | 2 | 5 |
|---|---|---|

$d = 2$

| | | |
|---|---|---|
| 1 | 1 | 2 |
|---|---|---|



BFS(G, 4)

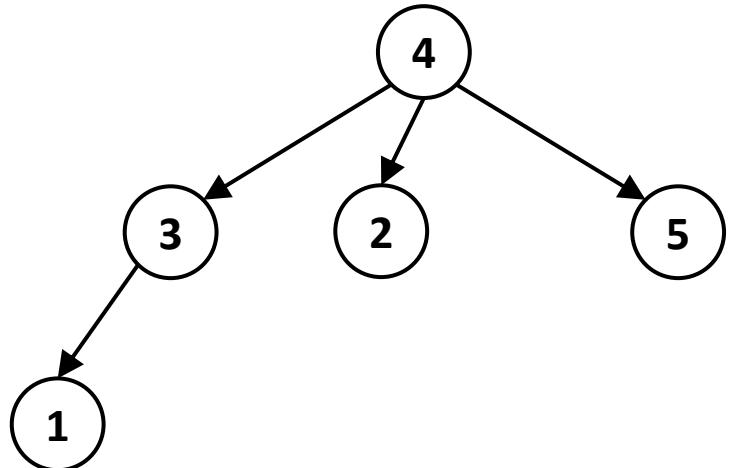


Adj List of G :

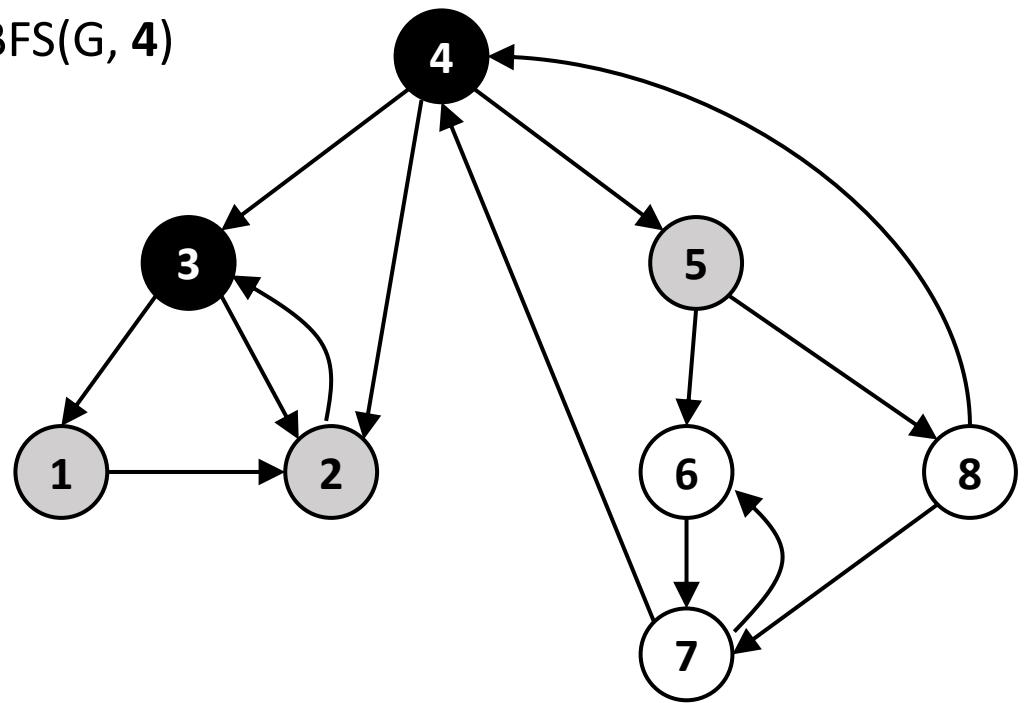
| | | |
|-----|---|-----------|
| 1 | → | 2 |
| 2 | → | 3 |
| ✓ 3 | → | 1 → 2 |
| ✓ 4 | → | 3 → 2 → 5 |
| 5 | → | 6 → 8 |
| 6 | → | 7 |
| 7 | → | 6 → 4 |
| 8 | → | 7 → 4 |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |



BFS(G, 4)

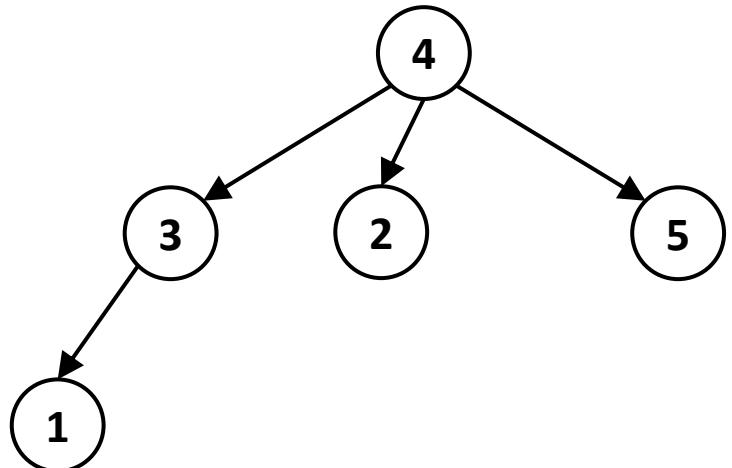


Adj List of G :

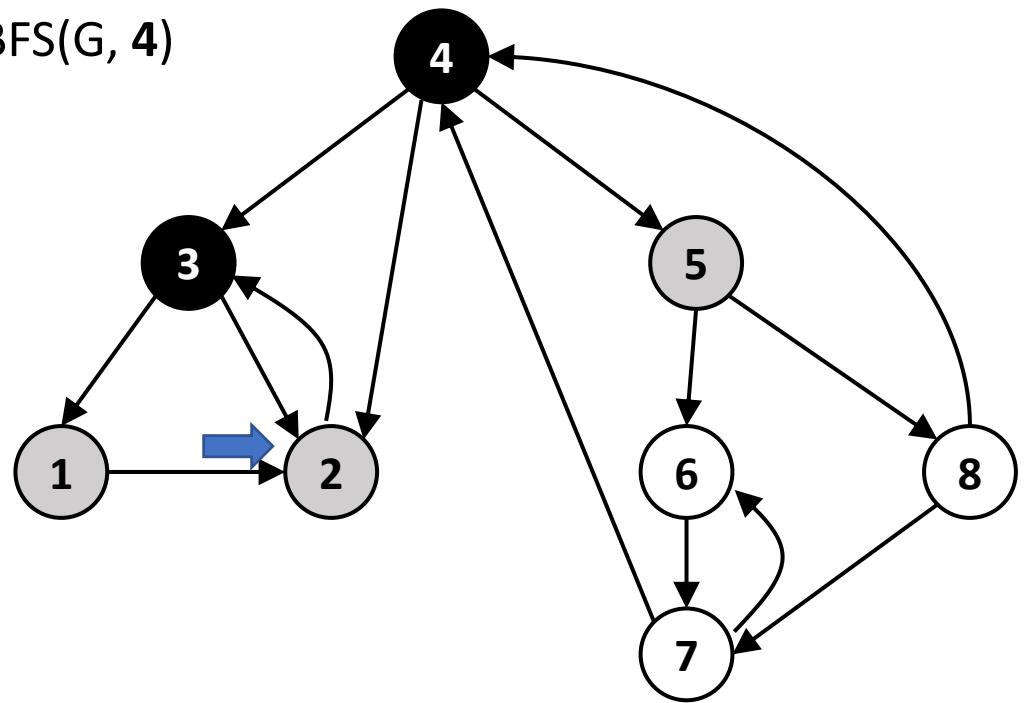
| | |
|-----|-----------|
| 1 | 2 |
| 2 | 3 |
| ✓ 3 | 1 → 2 |
| ✓ 4 | 3 → 2 → 5 |
| 5 | 6 → 8 |
| 6 | 7 |
| 7 | 6 → 4 |
| 8 | 7 → 4 |

Contents of Q :

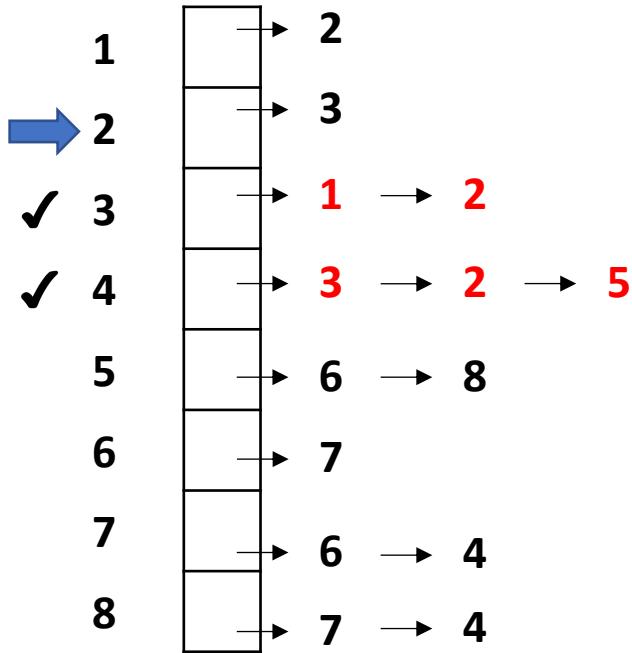
| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |



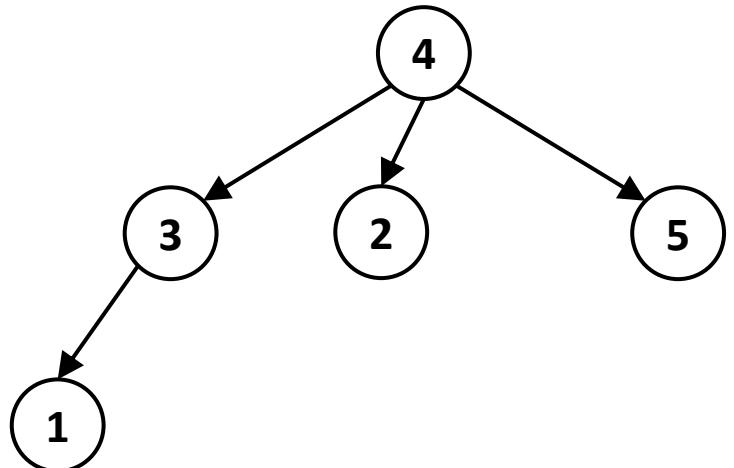
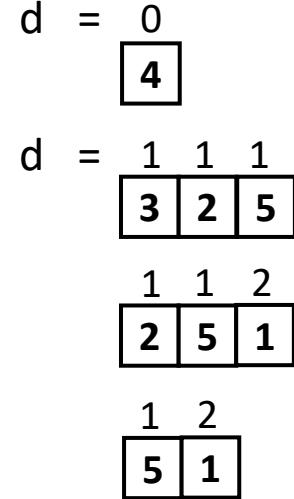
BFS(G, 4)



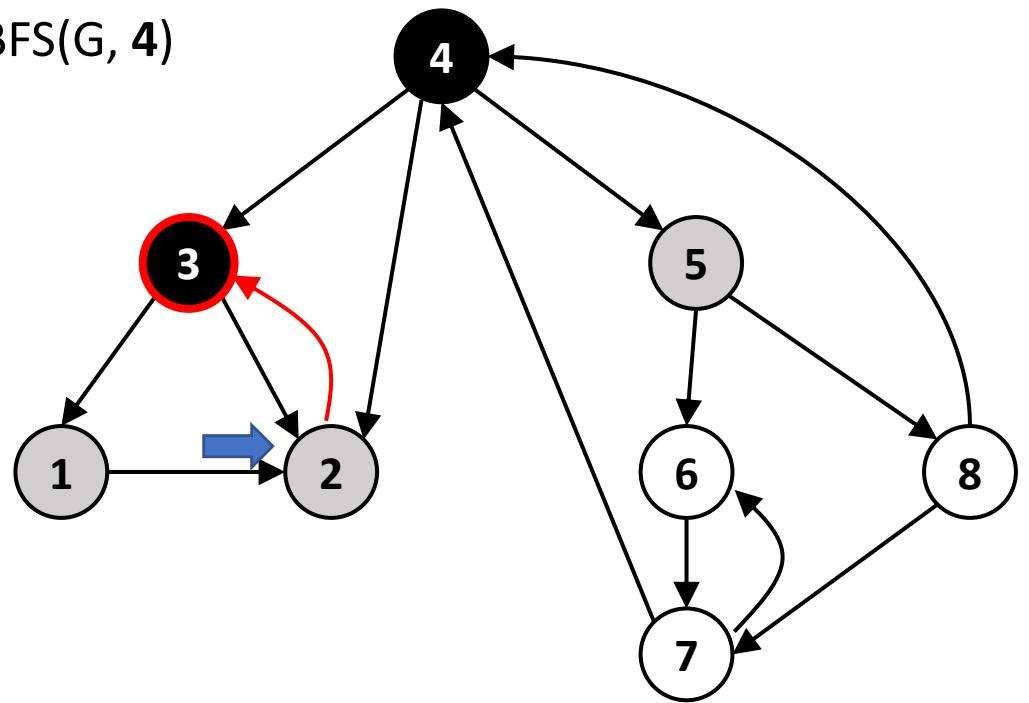
Adj List of G :



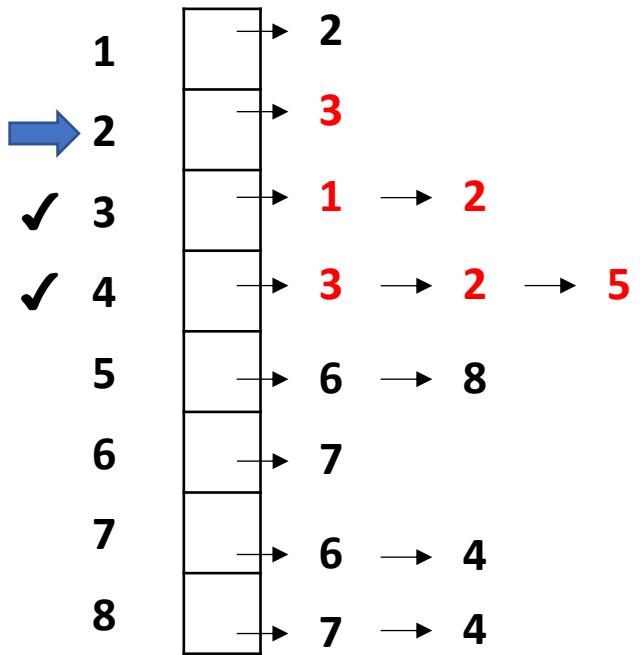
Contents of Q :



BFS(G , 4)

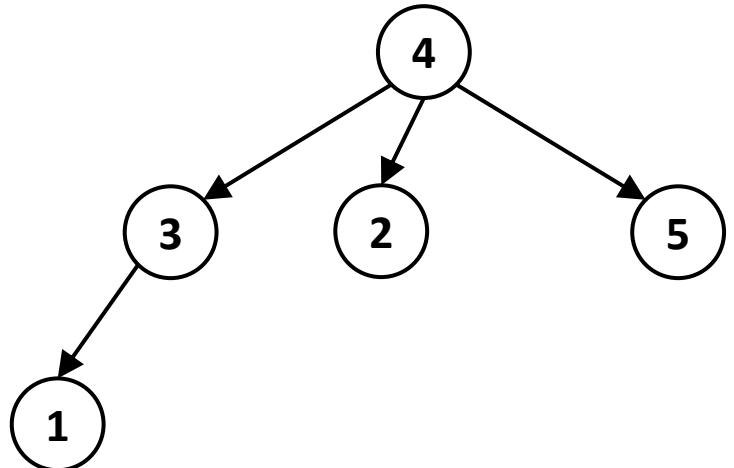


Adj List of G :

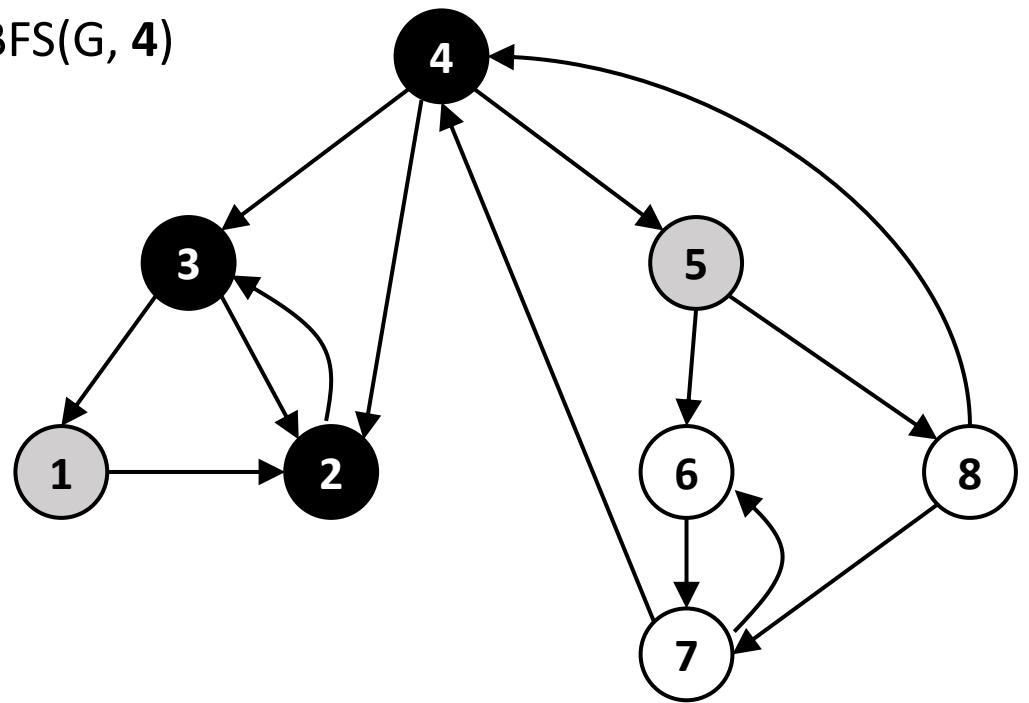


Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |



BFS(G , 4)

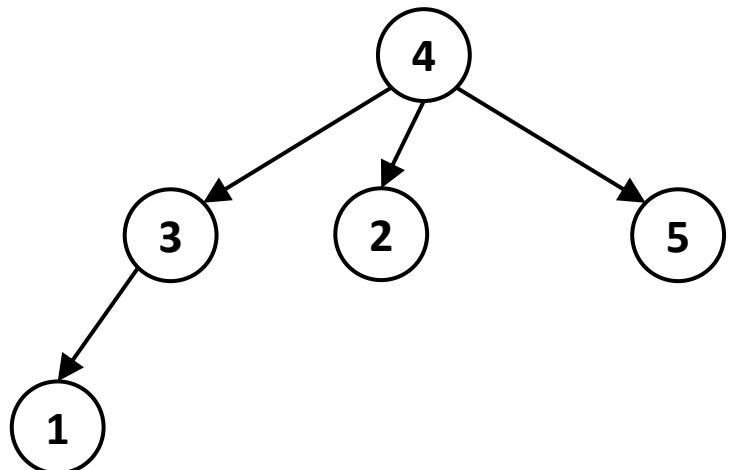


Adj List of G :

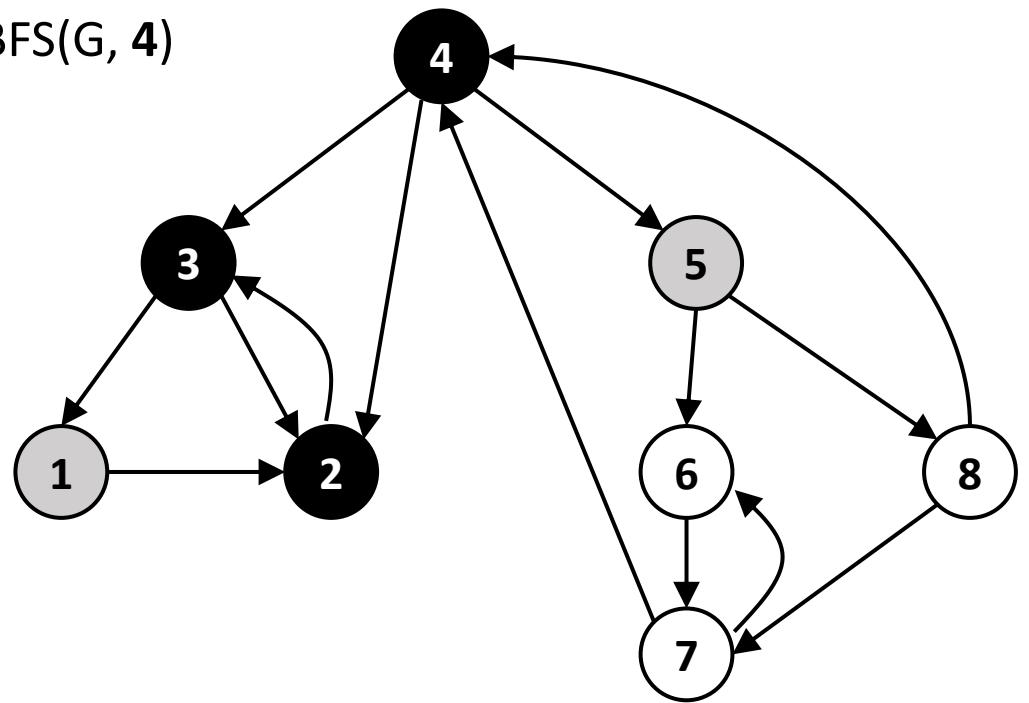
| | | |
|---|---|-----------|
| 1 | ✓ | 2 |
| 2 | ✓ | 3 |
| 3 | ✓ | 1 → 2 |
| 4 | ✓ | 3 → 2 → 5 |
| 5 | | 6 → 8 |
| 6 | | 7 |
| 7 | | 6 → 4 |
| 8 | | 7 → 4 |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |



BFS(G , 4)

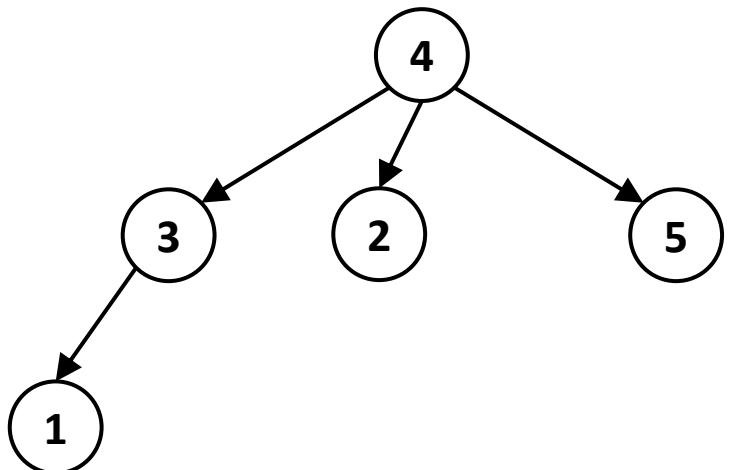


Adj List of G :

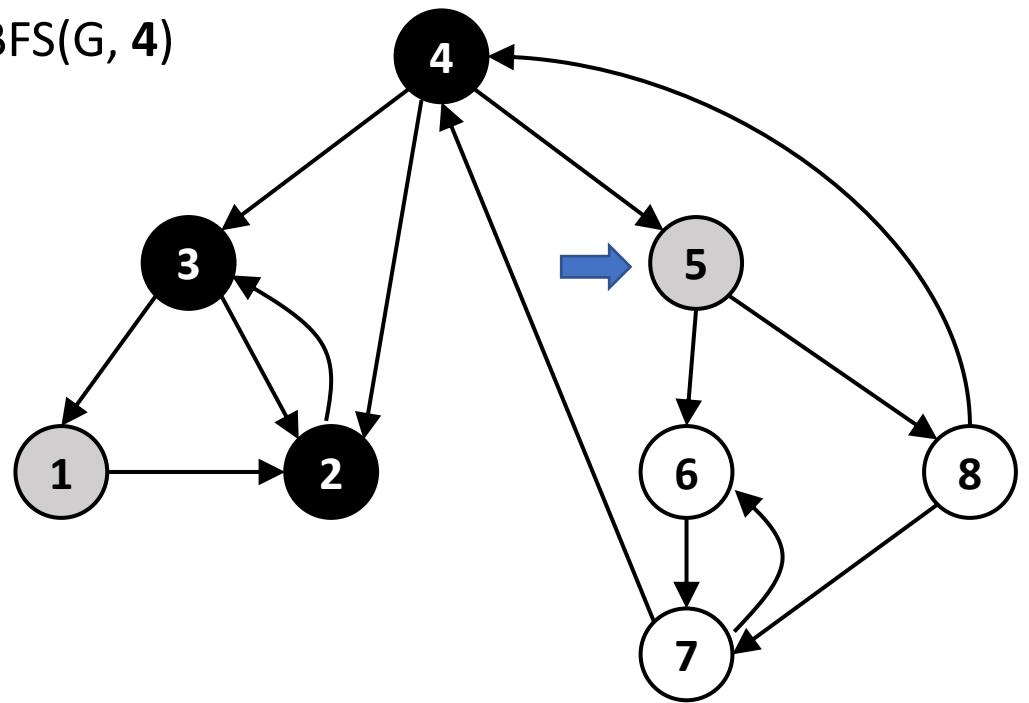
| | | |
|---|---|-----------|
| 1 | ✓ | 2 |
| 2 | ✓ | 3 |
| 3 | ✓ | 1 → 2 |
| 4 | ✓ | 3 → 2 → 5 |
| 5 | | 6 → 8 |
| 6 | | 7 |
| 7 | | 6 → 4 |
| 8 | | 7 → 4 |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |



BFS(G, 4)

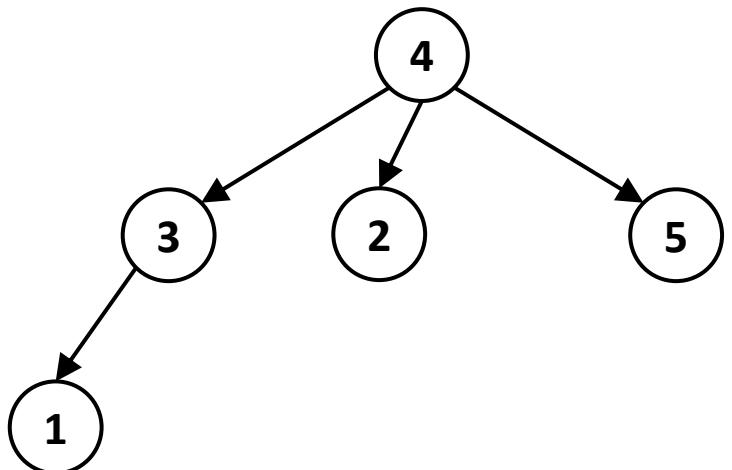


Adj List of G :

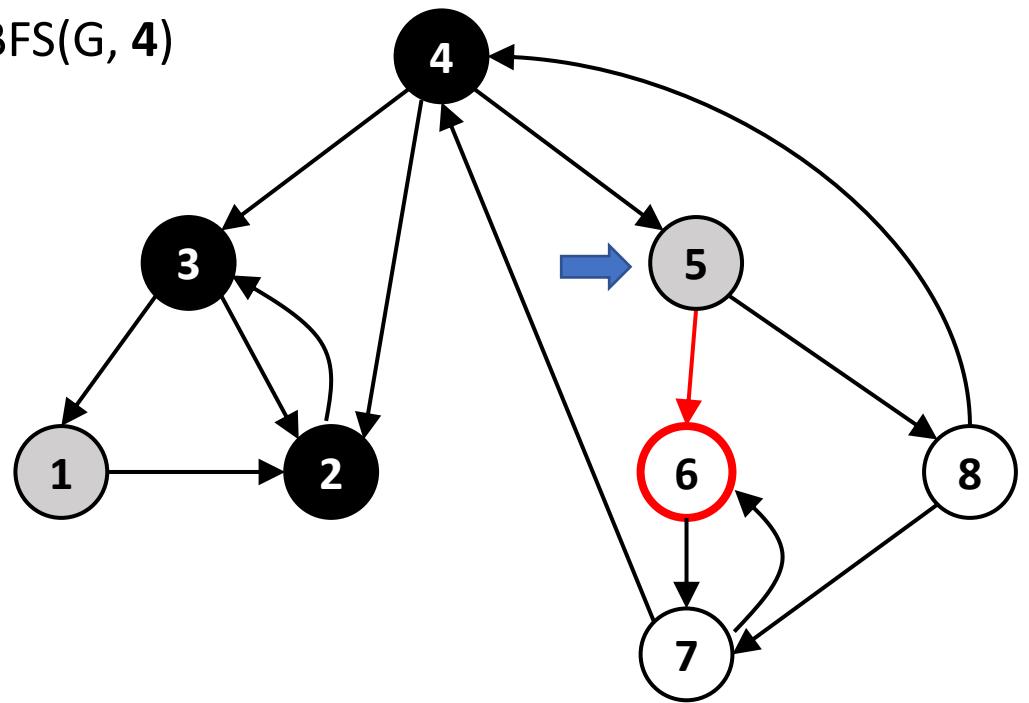
| | |
|-----|-------------|
| 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| → 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 |
| | 1 |



BFS(G, 4)

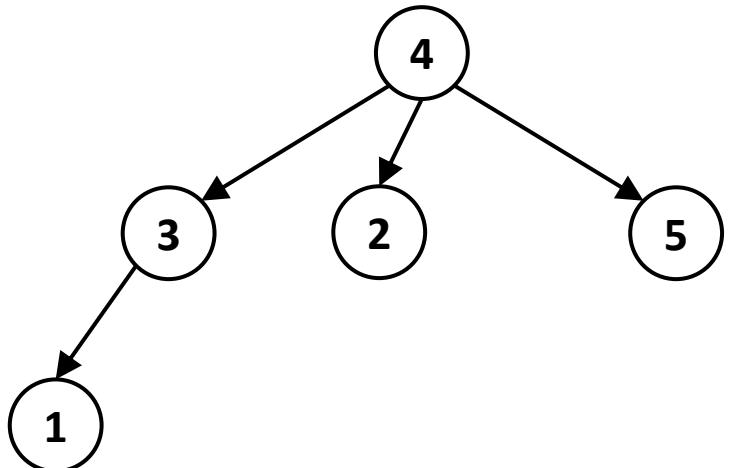


Adj List of G :

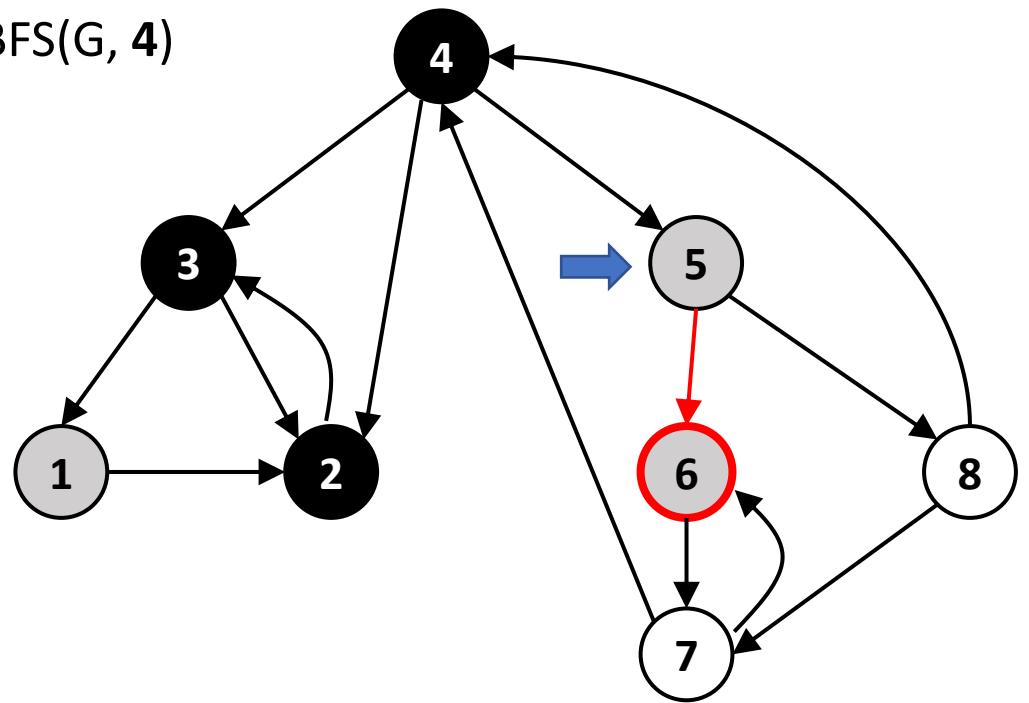
| | |
|-----|-------------|
| 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| → 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 7 | → 7 → 4 |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 |
| | 1 |

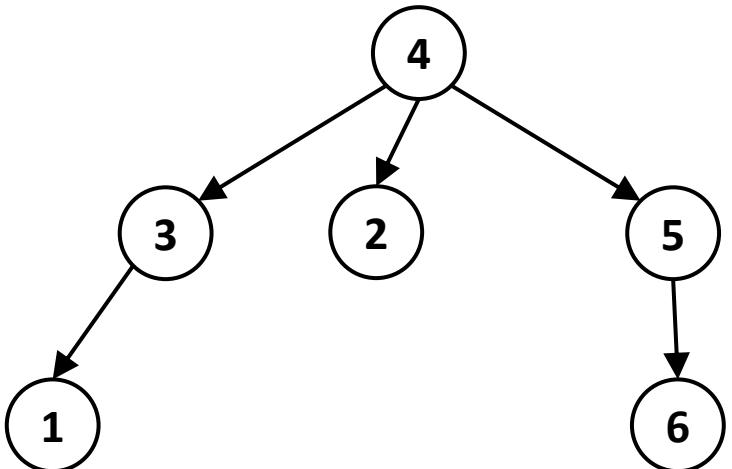


BFS(G , 4)



Adj List of G :

| | |
|-----|-----------|
| 1 | 2 |
| ✓ 2 | 3 |
| ✓ 3 | 1 → 2 |
| ✓ 4 | 3 → 2 → 5 |
| → 5 | 6 → 8 |
| 6 | 7 |
| 7 | 6 → 4 |
| 8 | 7 → 4 |

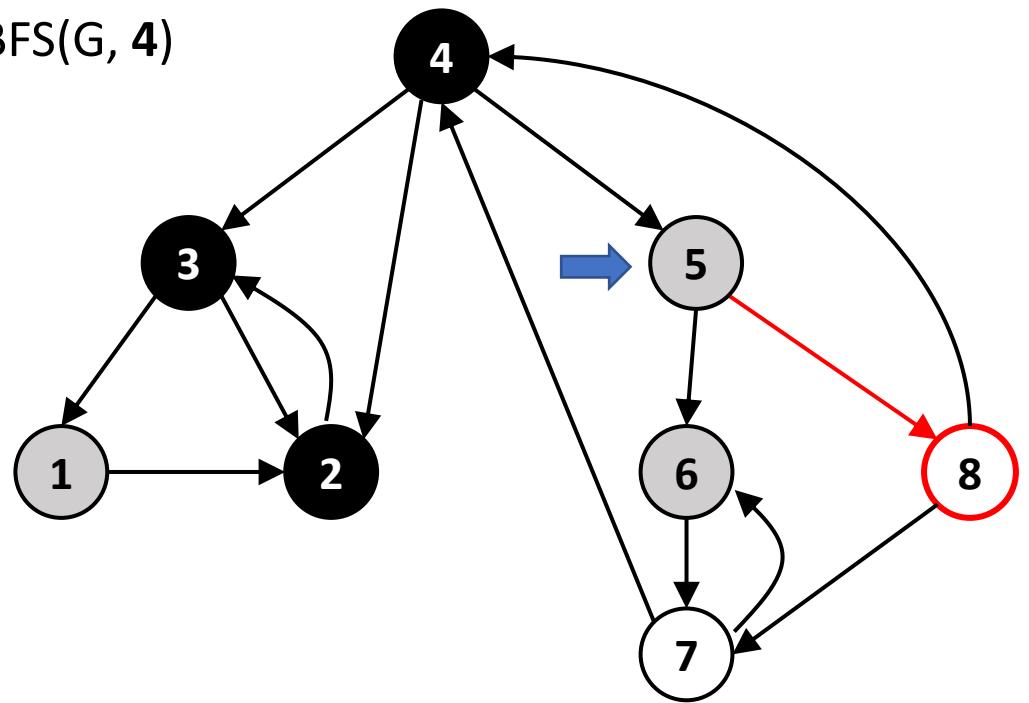


Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 |
| | 1 6 |



BFS(G , 4)

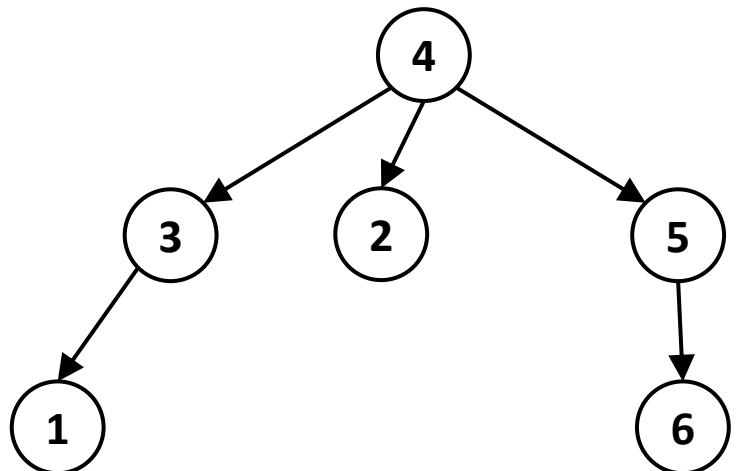


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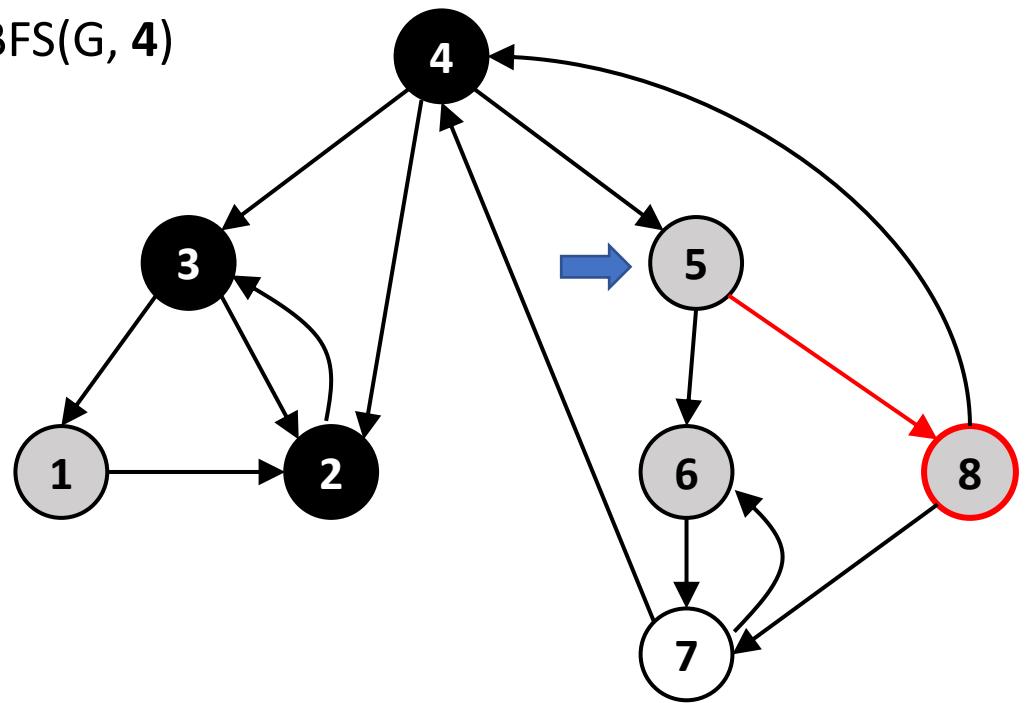
| | | |
|---|---|-----------|
| 1 | ✓ | 2 |
| 2 | ✓ | 3 |
| 3 | ✓ | 1 → 2 |
| 4 | ✓ | 3 → 2 → 5 |
| 5 | ✓ | 6 → 8 |
| 6 | | 7 |
| 7 | | 6 → 4 |
| 8 | | 7 → 4 |

Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 |
| | 1 6 |



BFS(G , 4)

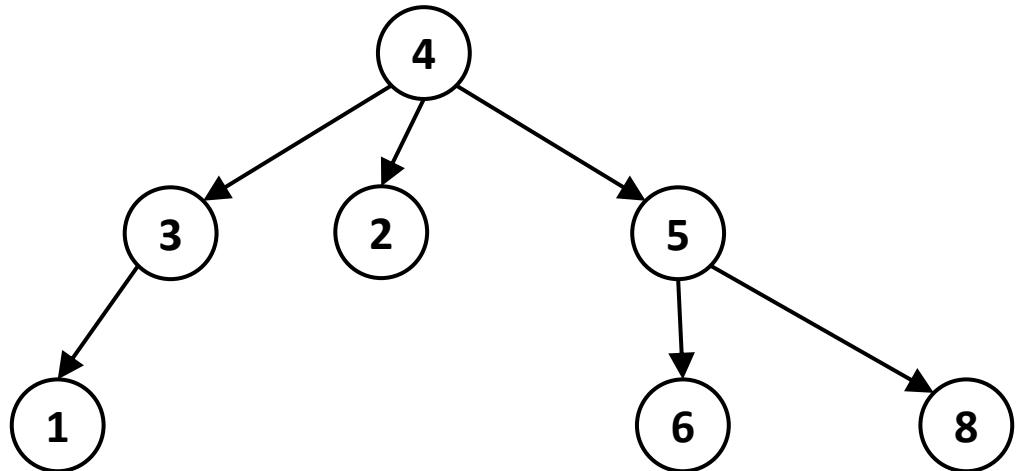


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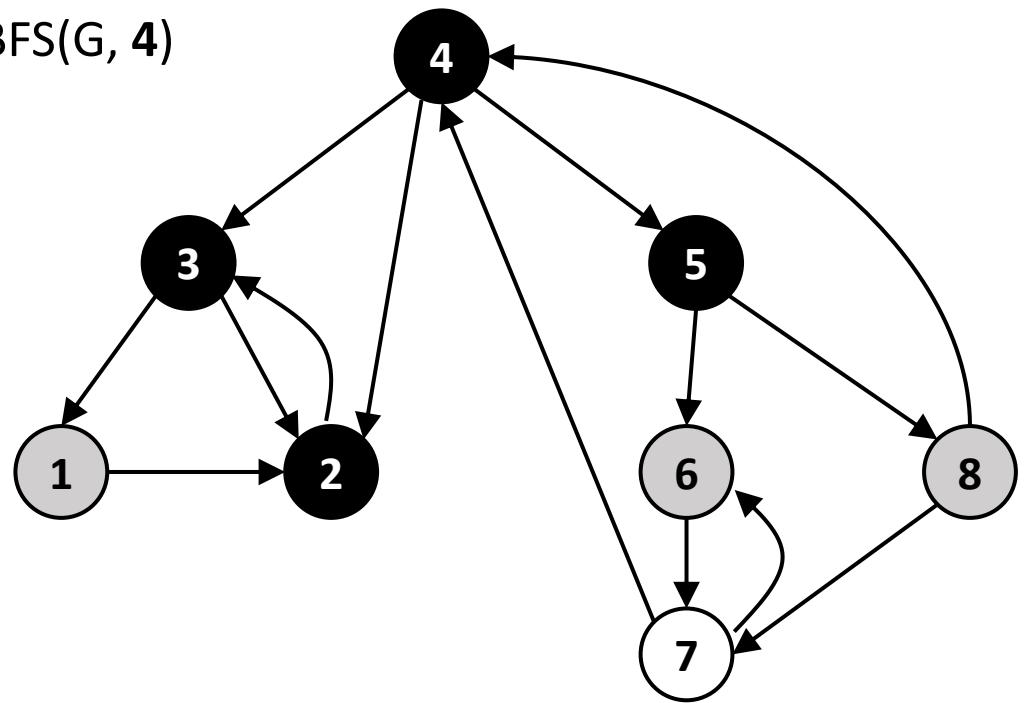
| | | |
|-----|---|-----------|
| 1 | → | 2 |
| ✓ 2 | → | 3 |
| ✓ 3 | → | 1 → 2 |
| ✓ 4 | → | 3 → 2 → 5 |
| → 5 | → | 6 → 8 |
| 6 | → | 7 |
| 7 | → | 6 → 4 |
| 8 | → | 7 → 4 |

Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |



BFS(G , 4)

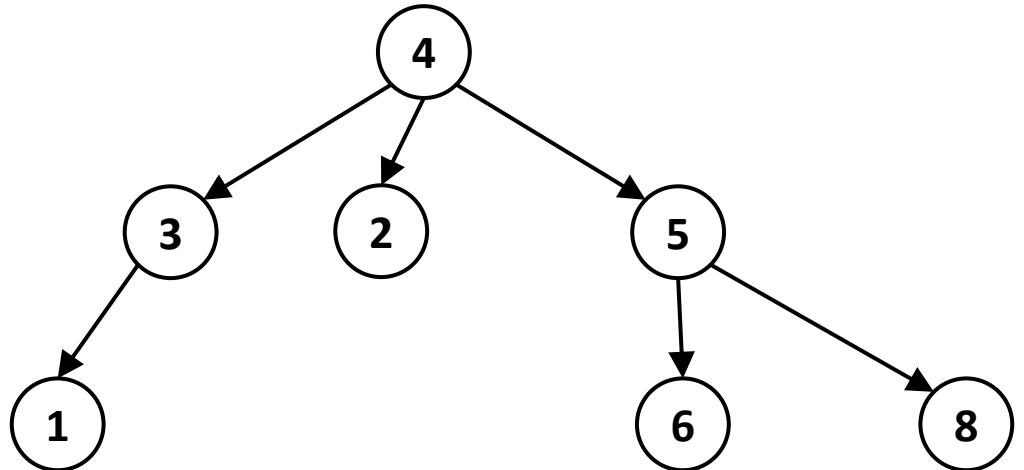


Adj List of G :

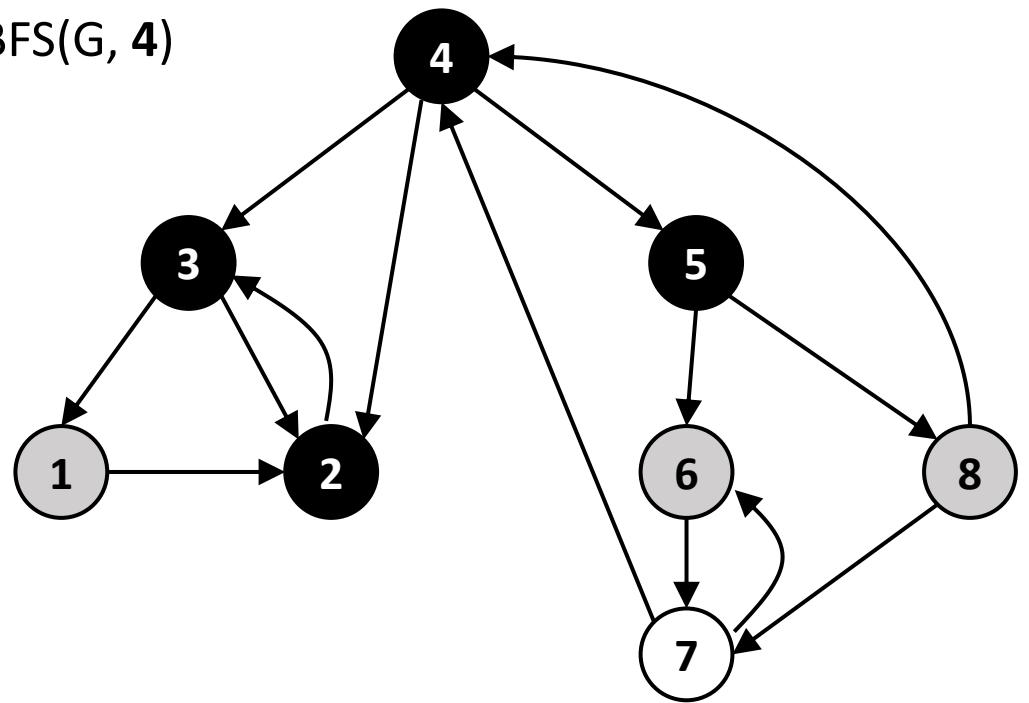
| | |
|-----|-----------|
| 1 | 2 |
| ✓ 2 | 3 |
| ✓ 3 | 1 → 2 |
| ✓ 4 | 3 → 2 → 5 |
| ✓ 5 | 6 → 8 |
| 6 | 7 |
| 7 | 6 → 4 |
| 8 | 7 → 4 |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |



BFS(G , 4)

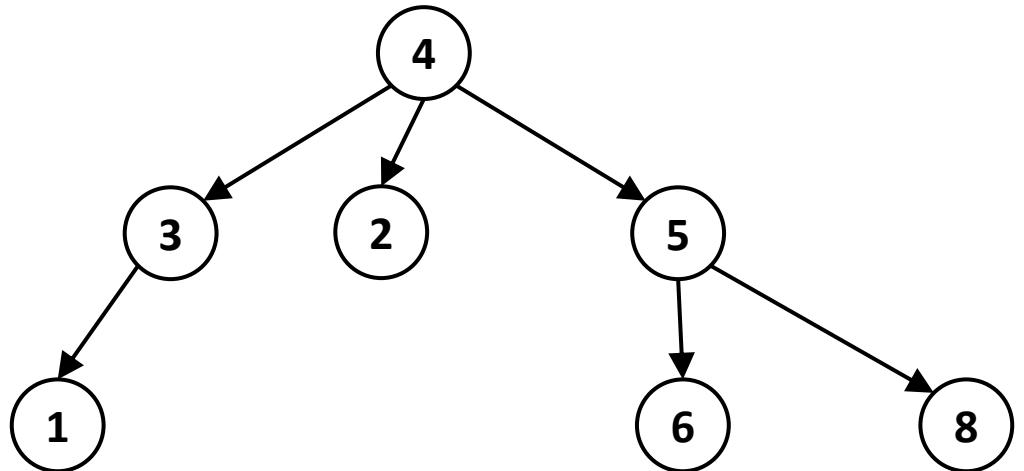


Adj List of G :

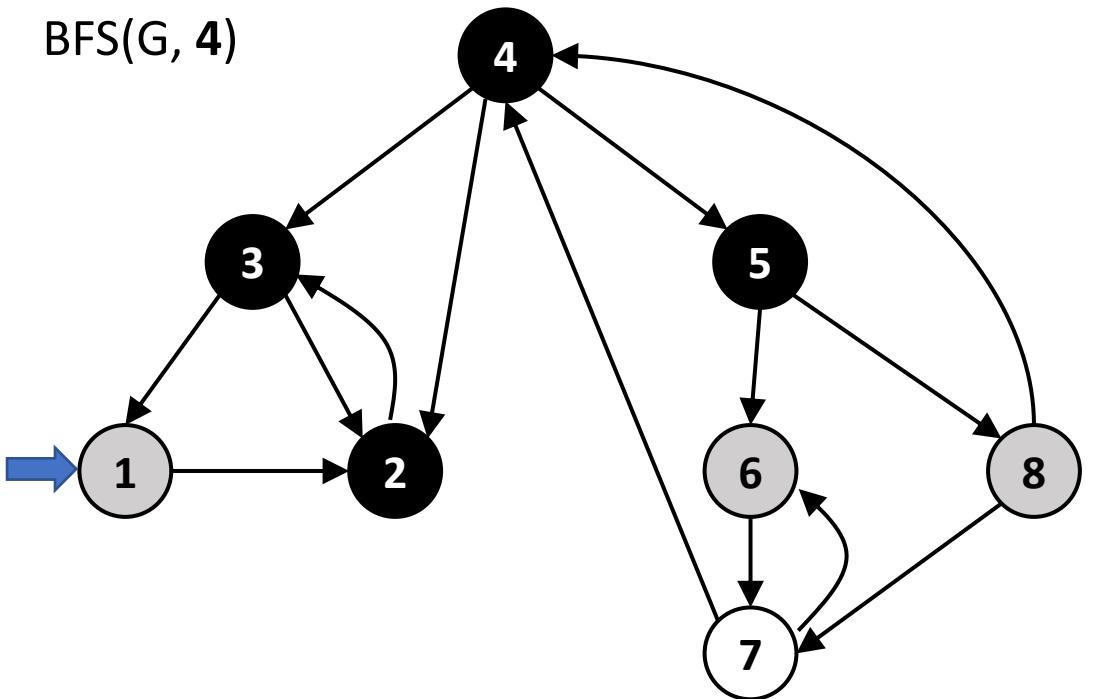
| | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| 1 | ✓ | 2 | ✓ | 3 | ✓ | 4 | ✓ | 5 | ✓ | 6 | ✓ | 7 | ✓ | 8 |
| | | 2 | | 3 | | 4 | | 5 | | 6 | | 7 | | 8 |
| | | 3 | → | 1 | → | 2 | → | 5 | | 6 | → | 8 | | |
| | | | | 1 | → | 2 | → | 5 | | | | | | |
| | | | | | | 3 | → | 2 | → | 5 | | | | |
| | | | | | | | 6 | → | 8 | | | | | |
| | | | | | | | | 7 | | | | | | |
| | | | | | | | | 6 | → | 4 | | | | |
| | | | | | | | | | 7 | → | 4 | | | |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |

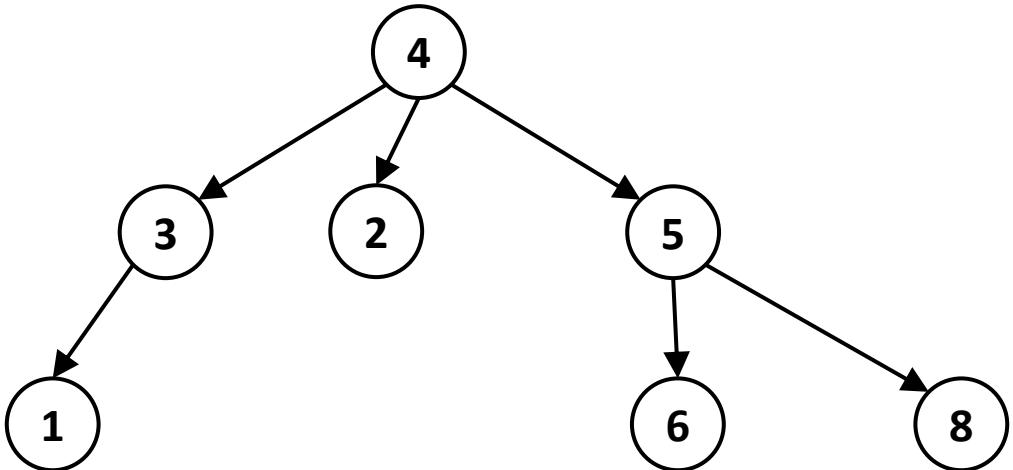


BFS(G , 4)



Adj List of G :

| | |
|-----|-----------|
| → 1 | 2 |
| ✓ 2 | 3 |
| ✓ 3 | 1 → 2 |
| ✓ 4 | 3 → 2 → 5 |
| ✓ 5 | 6 → 8 |
| 6 | 7 |
| 7 | 6 → 4 |
| 8 | 7 → 4 |

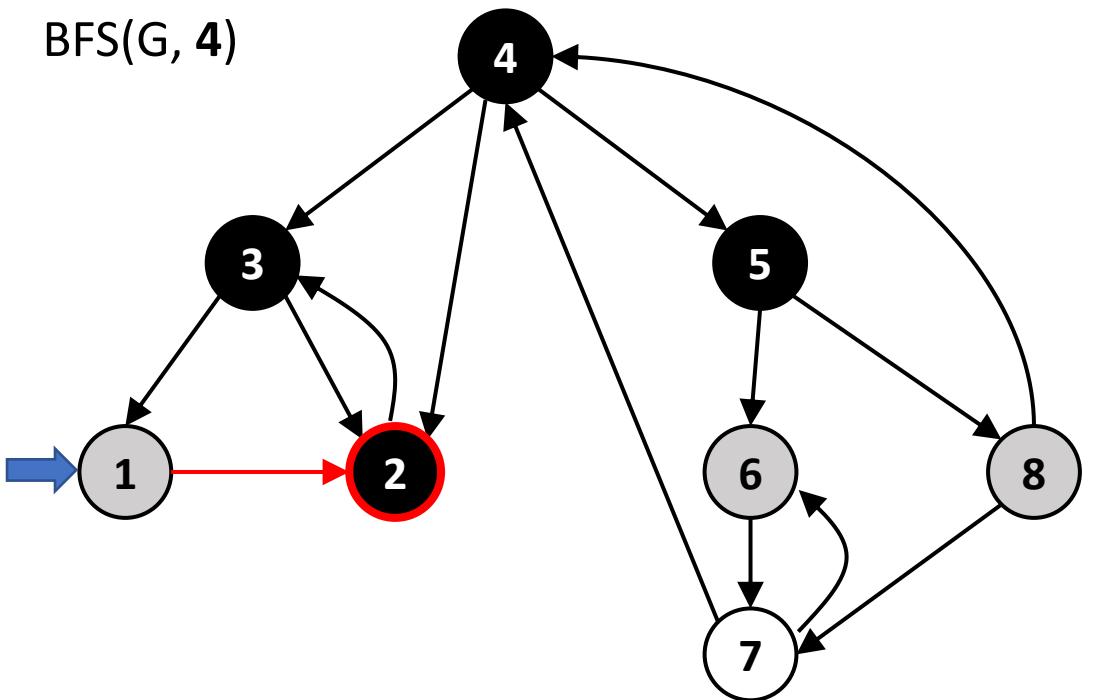


Contents of Q :

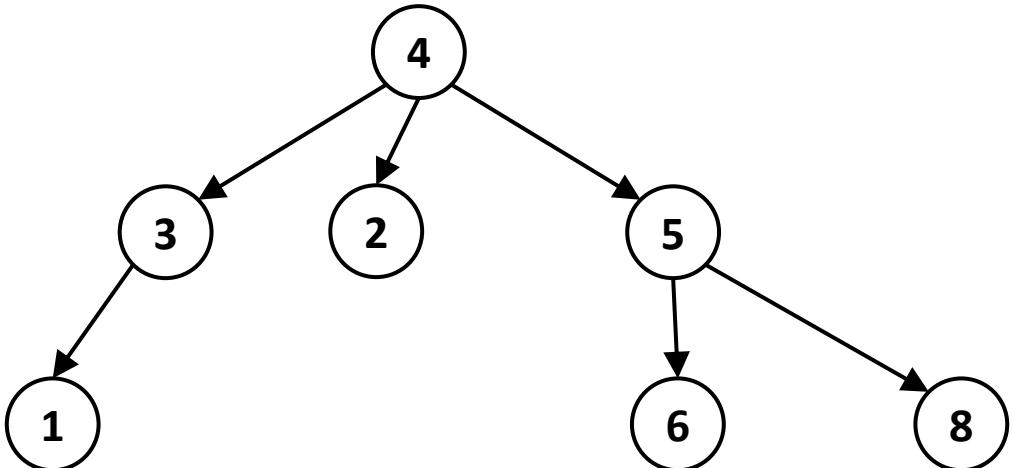
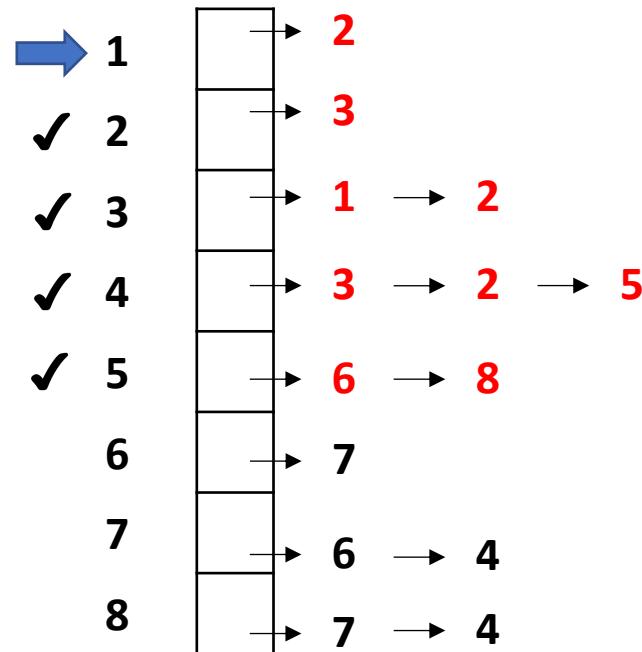
| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |



BFS(G , 4)



Adj List of G :

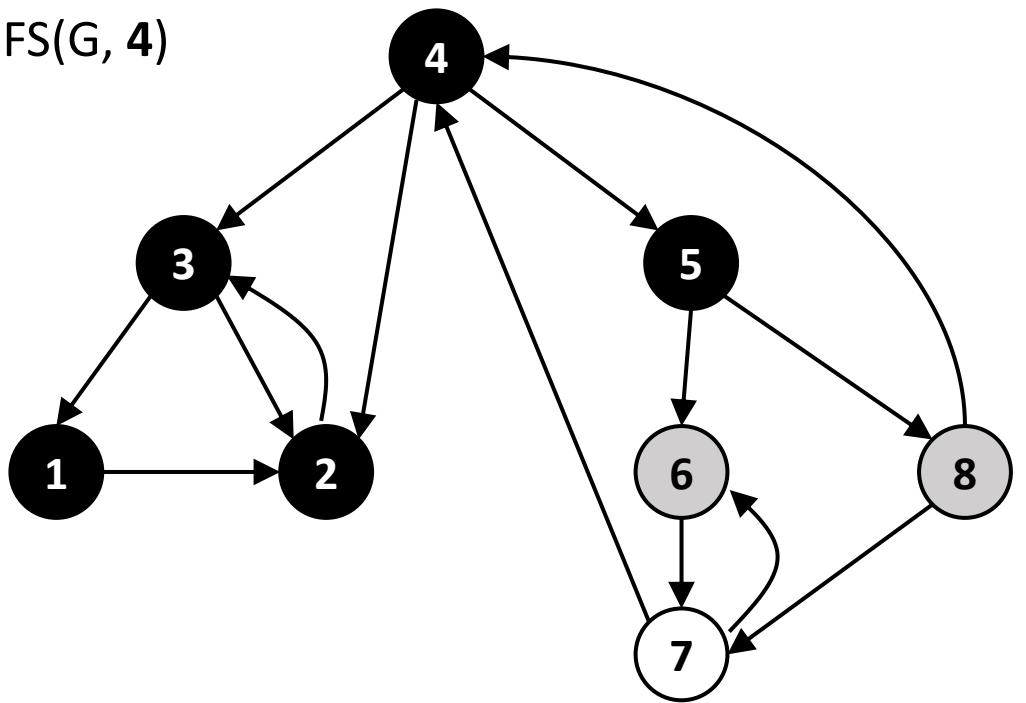


Contents of Q :

| | | | | |
|---------|---|---|---|---|
| $d = 0$ | <table border="1"><tr><td>4</td></tr></table> | 4 | | |
| 4 | | | | |
| $d = 1$ | <table border="1"><tr><td>3</td><td>2</td><td>5</td></tr></table> | 3 | 2 | 5 |
| 3 | 2 | 5 | | |
| $d = 2$ | <table border="1"><tr><td>2</td><td>5</td><td>1</td></tr></table> | 2 | 5 | 1 |
| 2 | 5 | 1 | | |
| $d = 3$ | <table border="1"><tr><td>5</td><td>1</td></tr></table> | 5 | 1 | |
| 5 | 1 | | | |
| $d = 4$ | <table border="1"><tr><td>1</td><td>6</td><td>8</td></tr></table> | 1 | 6 | 8 |
| 1 | 6 | 8 | | |
| $d = 5$ | <table border="1"><tr><td>6</td><td>8</td></tr></table> | 6 | 8 | |
| 6 | 8 | | | |



BFS(G , 4)

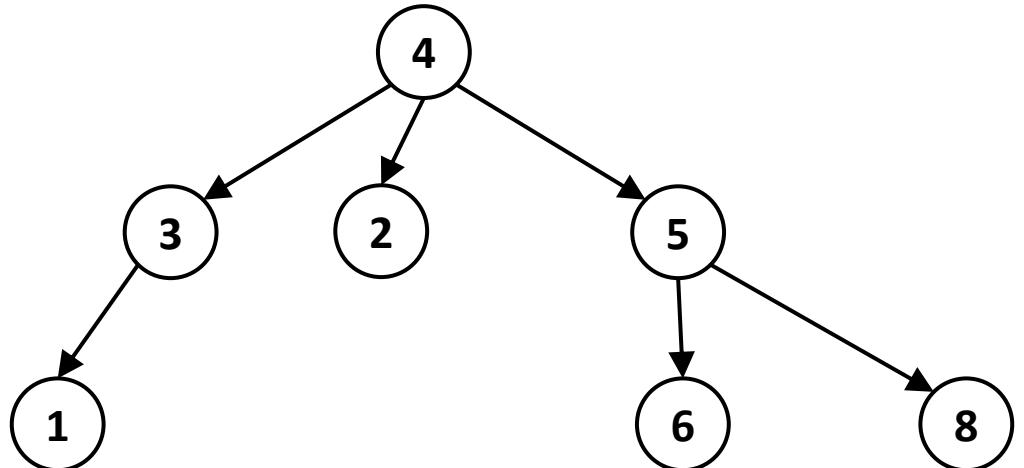


Adj List of G :

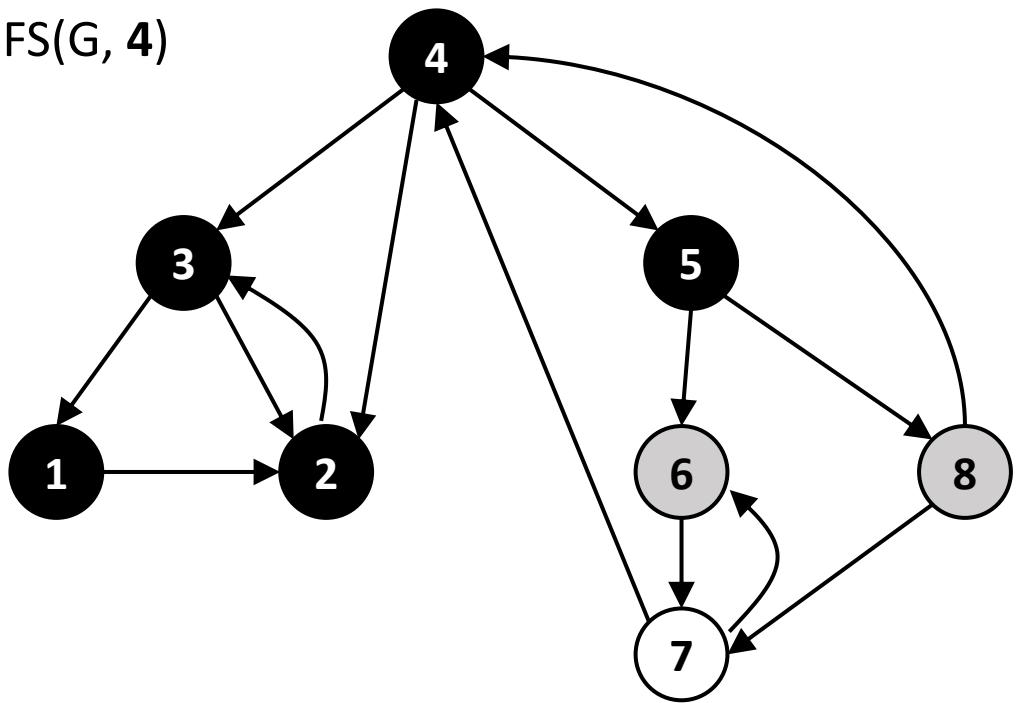
| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |

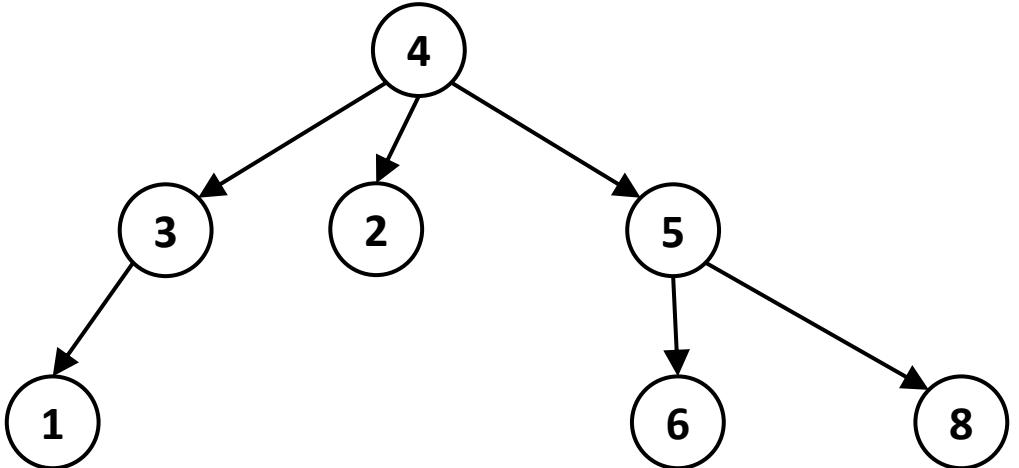


BFS(G , 4)



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

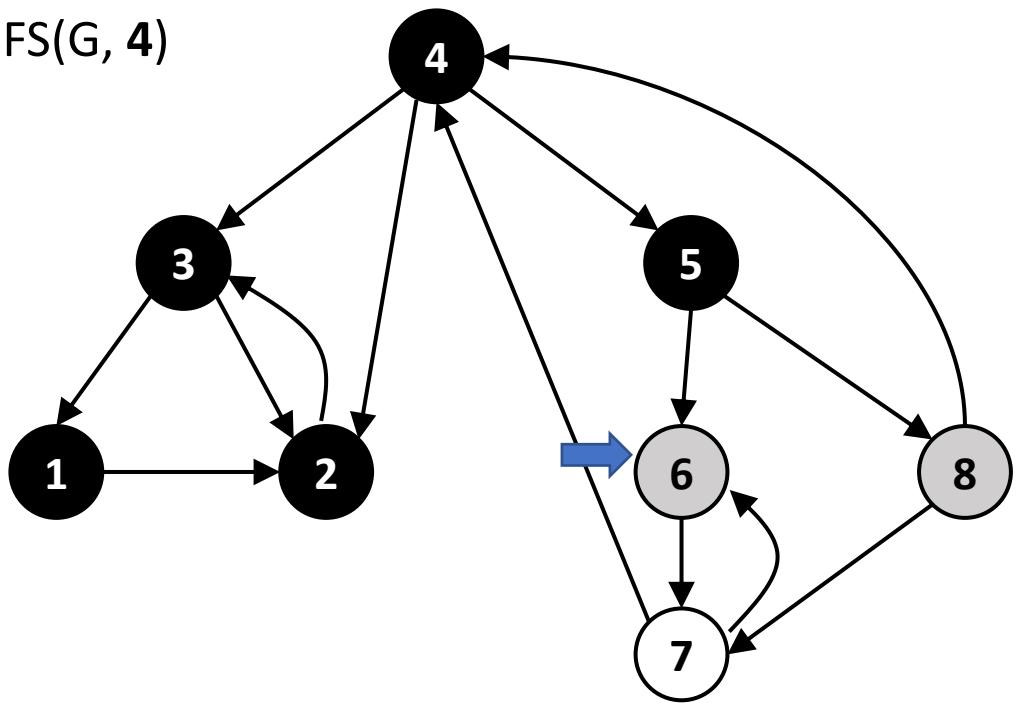


Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |

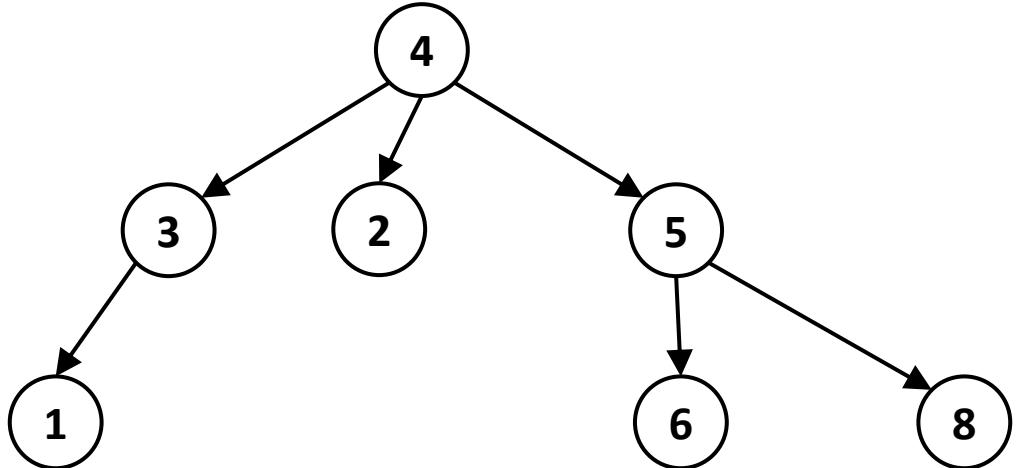


BFS(G , 4)



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

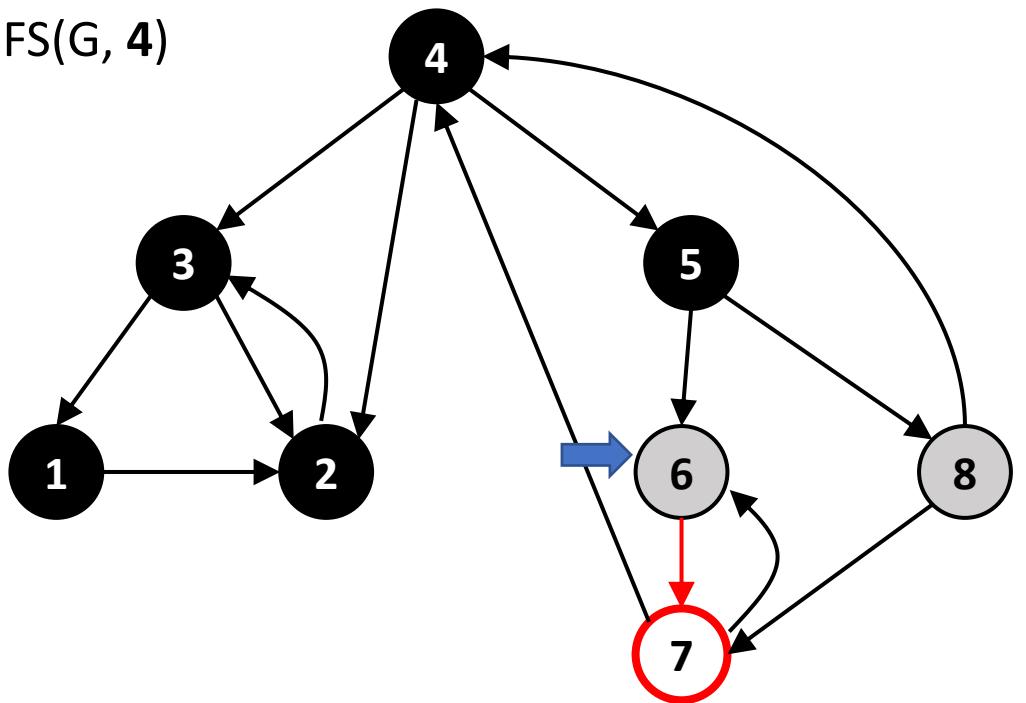


Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 |
| | 8 |

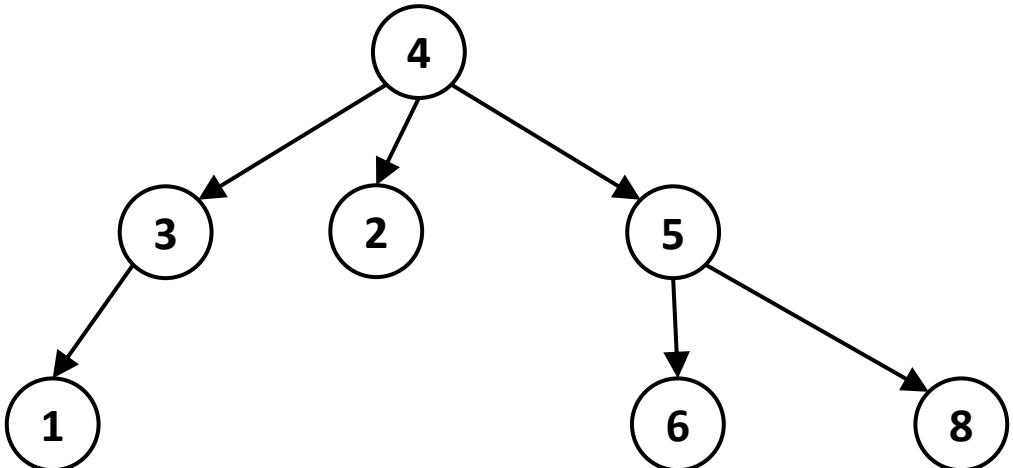


BFS(G , 4)



Adj List of G :

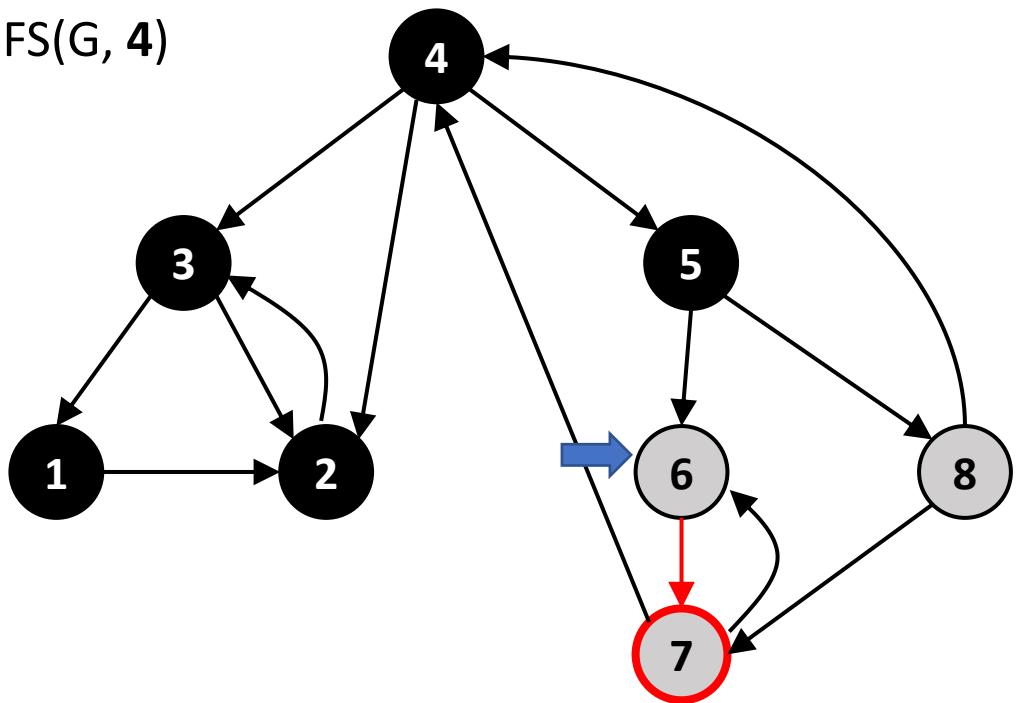
| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |



Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 |
| | 8 |

BFS(G , 4)

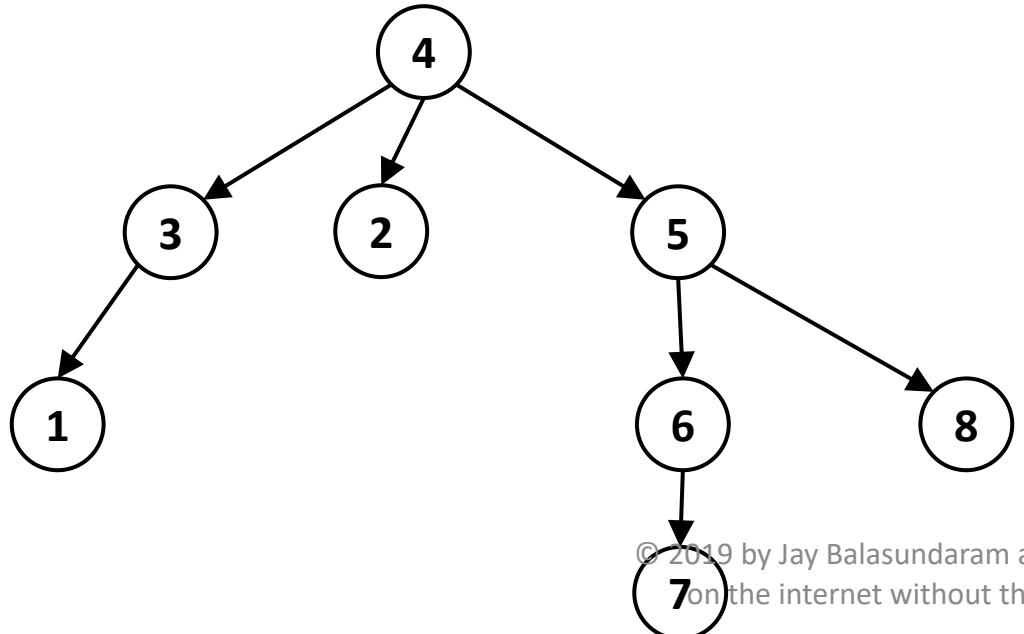


Adj List of G :

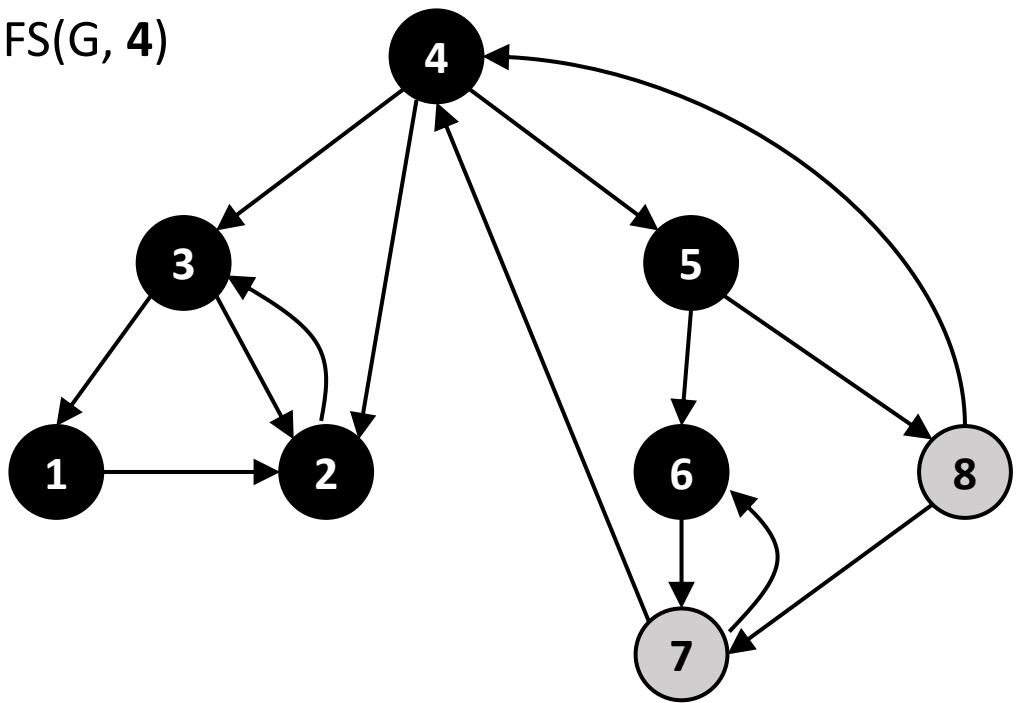
| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
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| | 1 2 |
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| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |



BFS(G, 4)

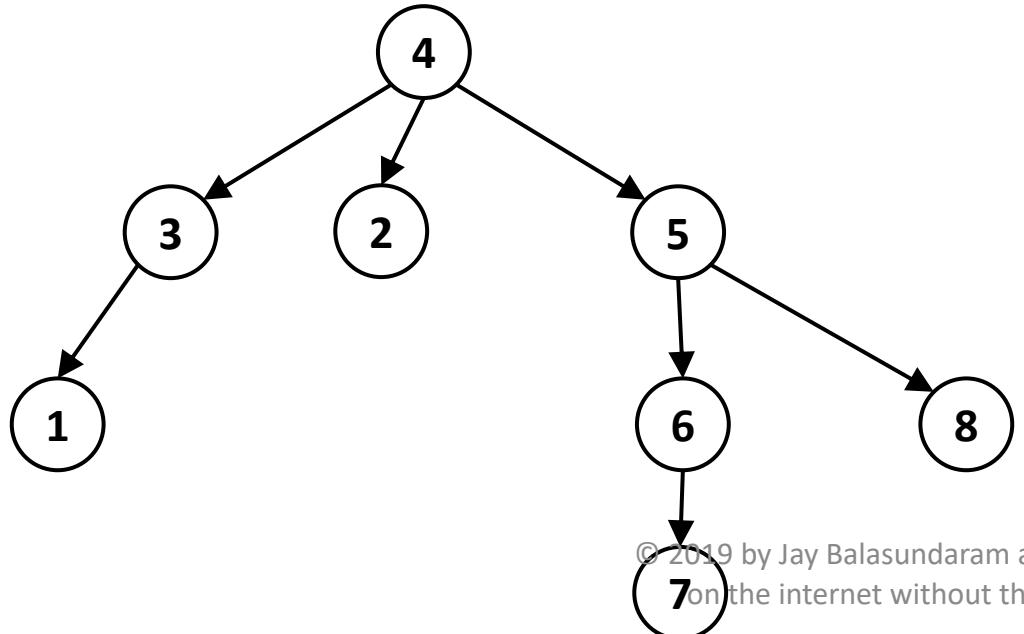


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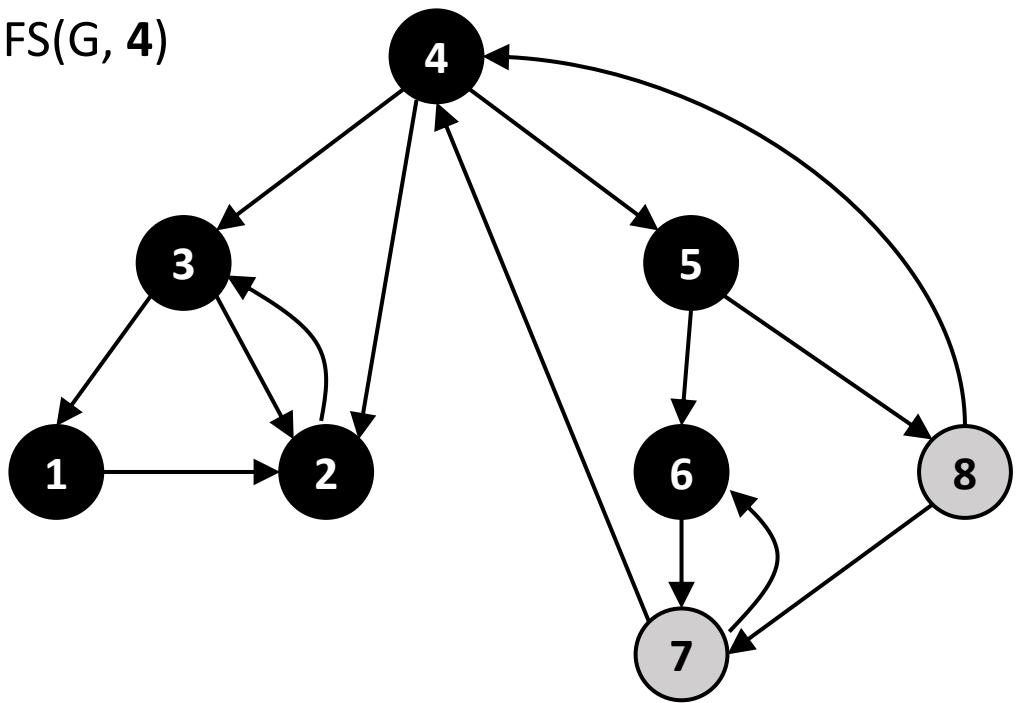
| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| ✓ 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |



BFS(G , 4)

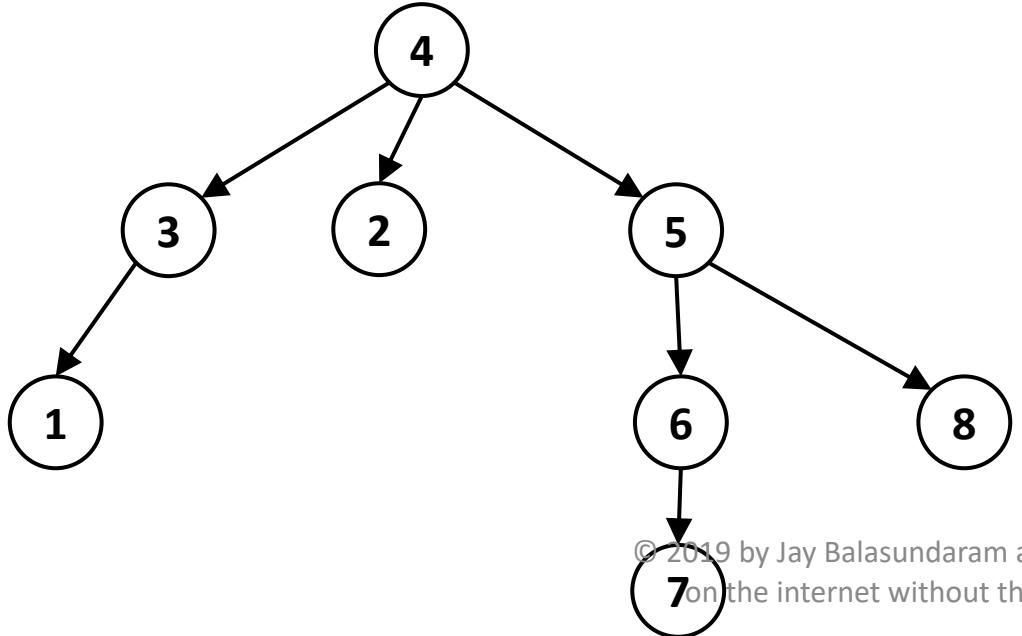


Adj List of G :

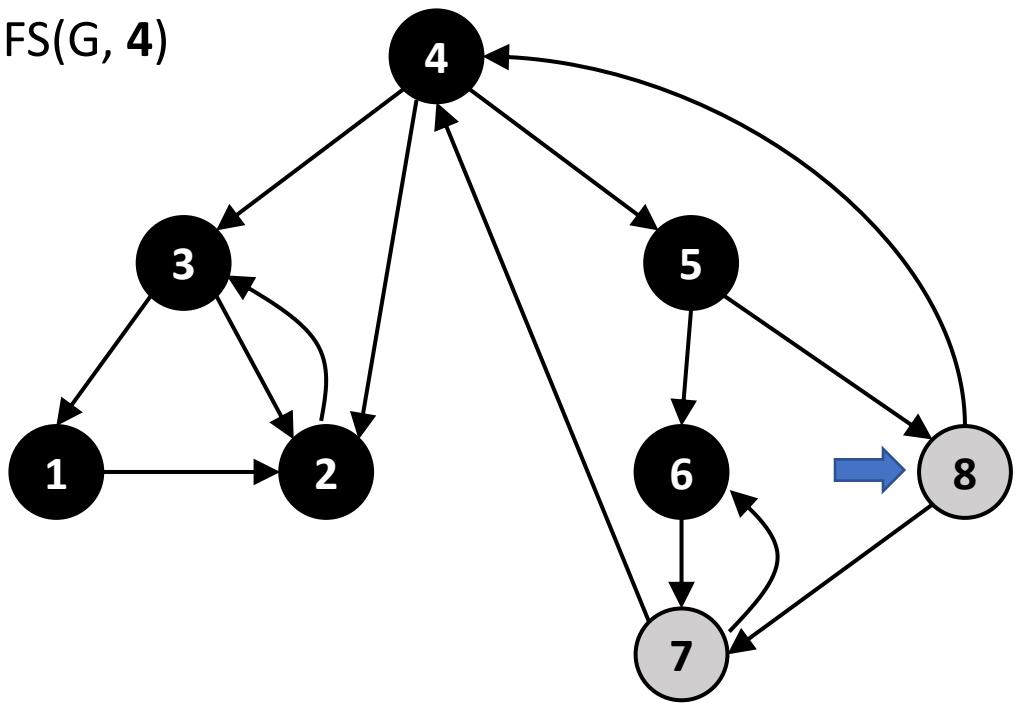
| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| ✓ 6 | → 7 |
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| | |
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| $d = 0$ | 4 |
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| | 1 1 2 |
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| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |

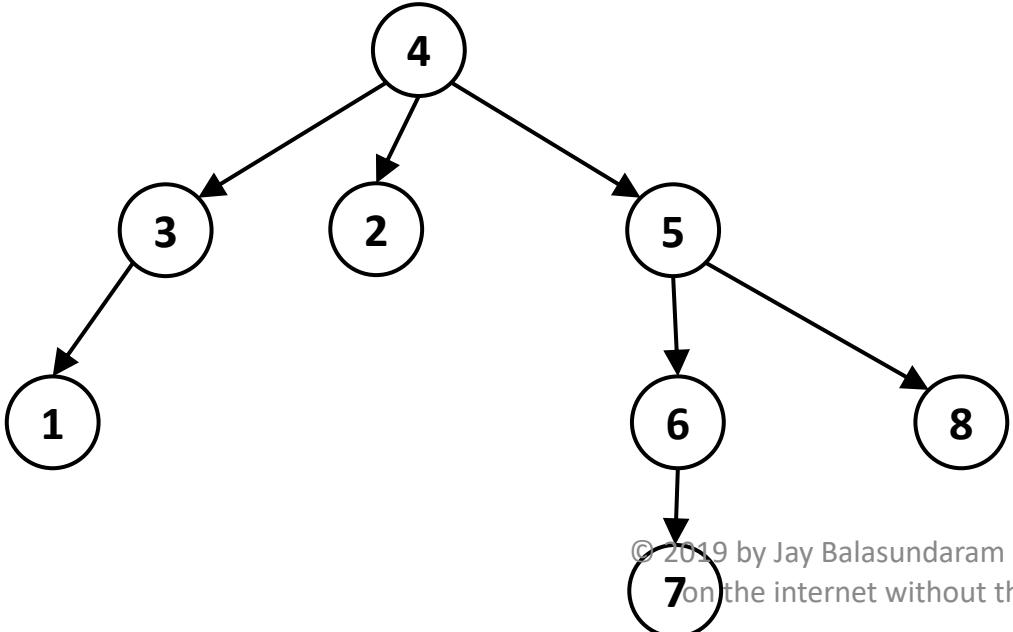


BFS(G, 4)



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| ✓ 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

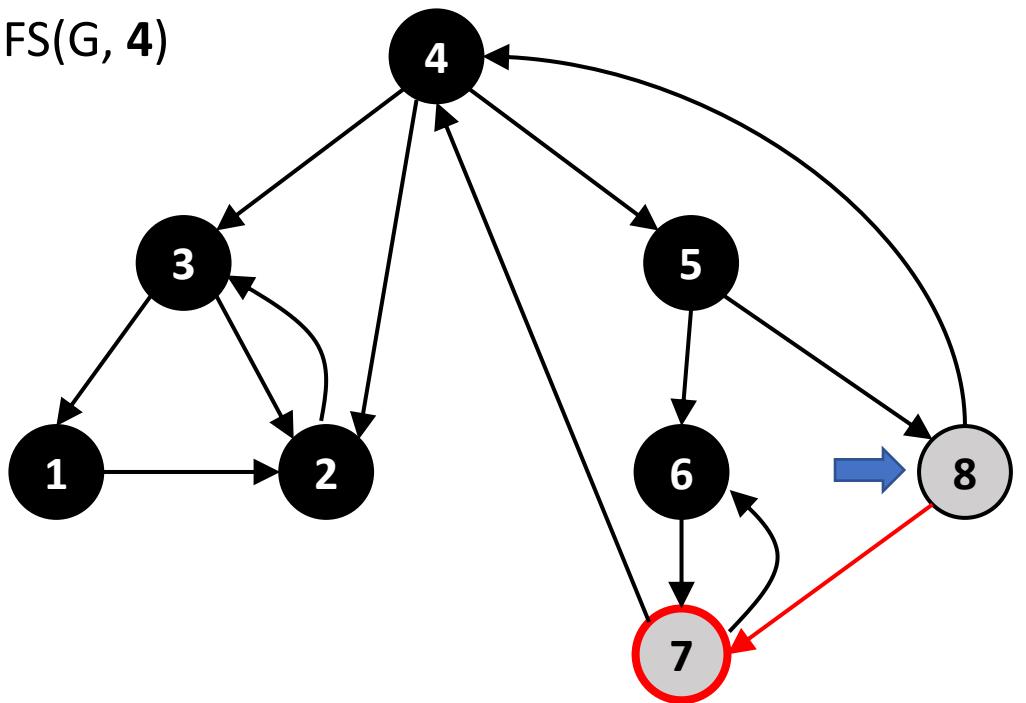


Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |
| | 3 |
| | 7 |

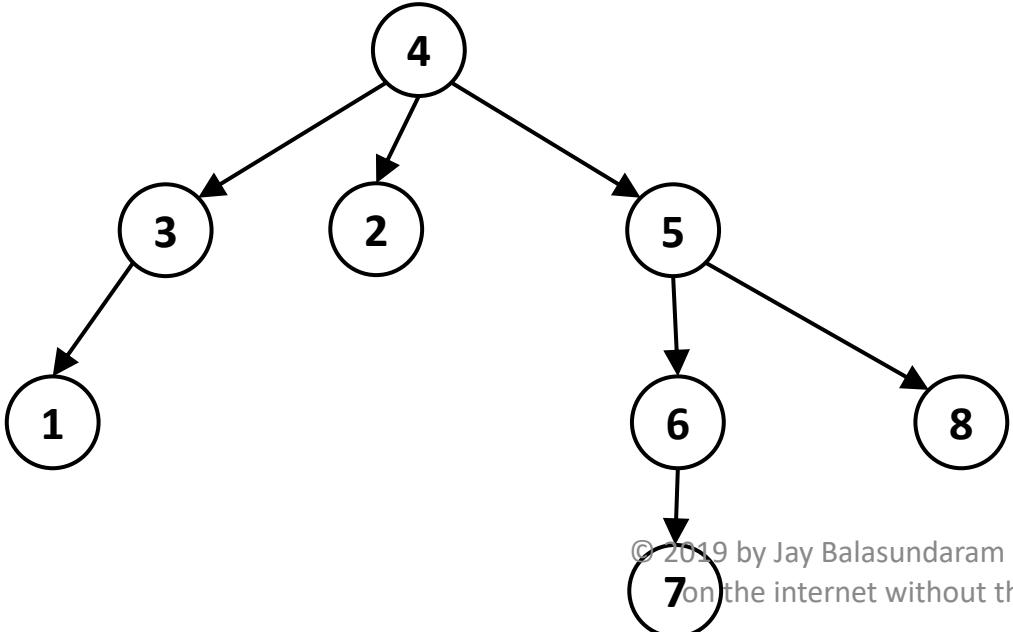


BFS(G, 4)



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| ✓ 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

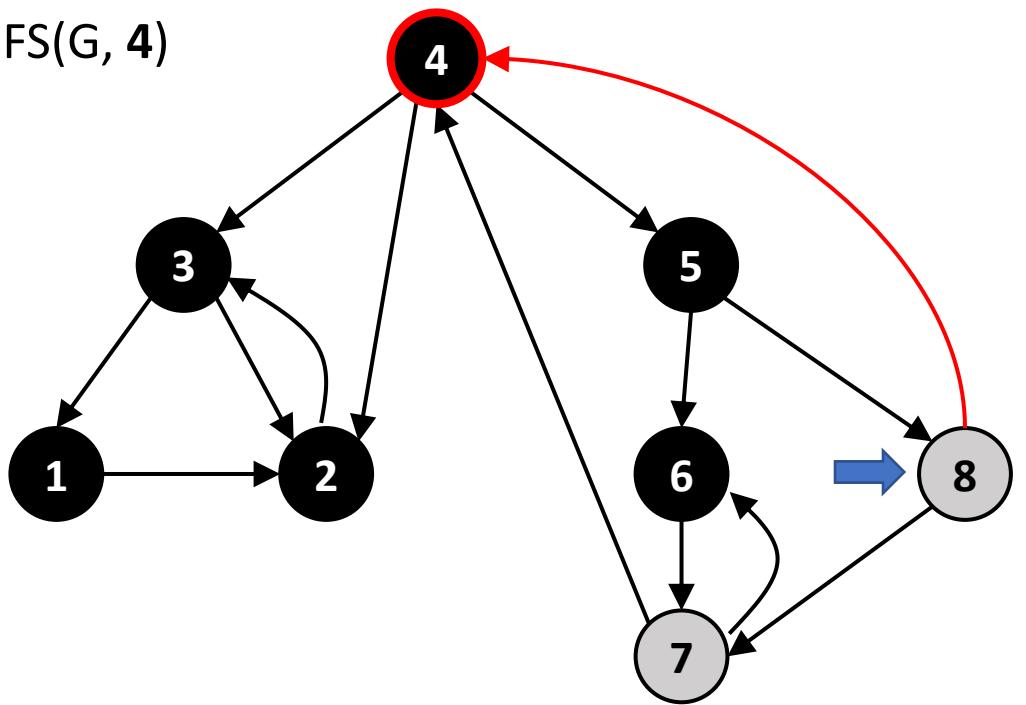


Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |
| | 3 |
| | 7 |

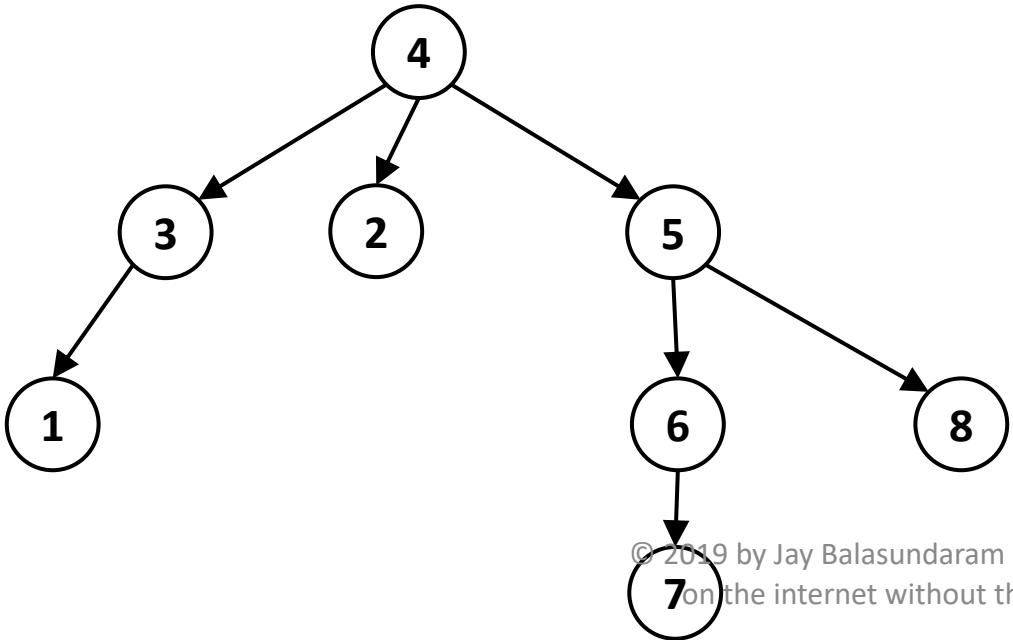


BFS(G , 4)



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| ✓ 6 | → 7 |
| 7 | → 6 → 4 |
| 8 | → 7 → 4 |

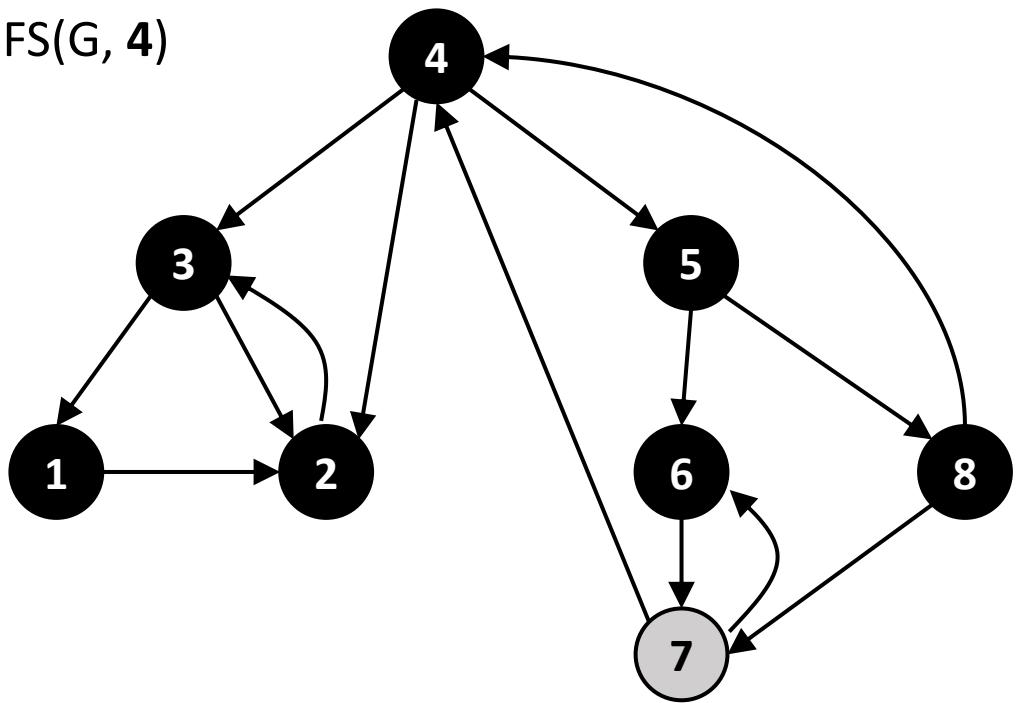


Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |
| | 3 |
| | 7 |



BFS(G , 4)

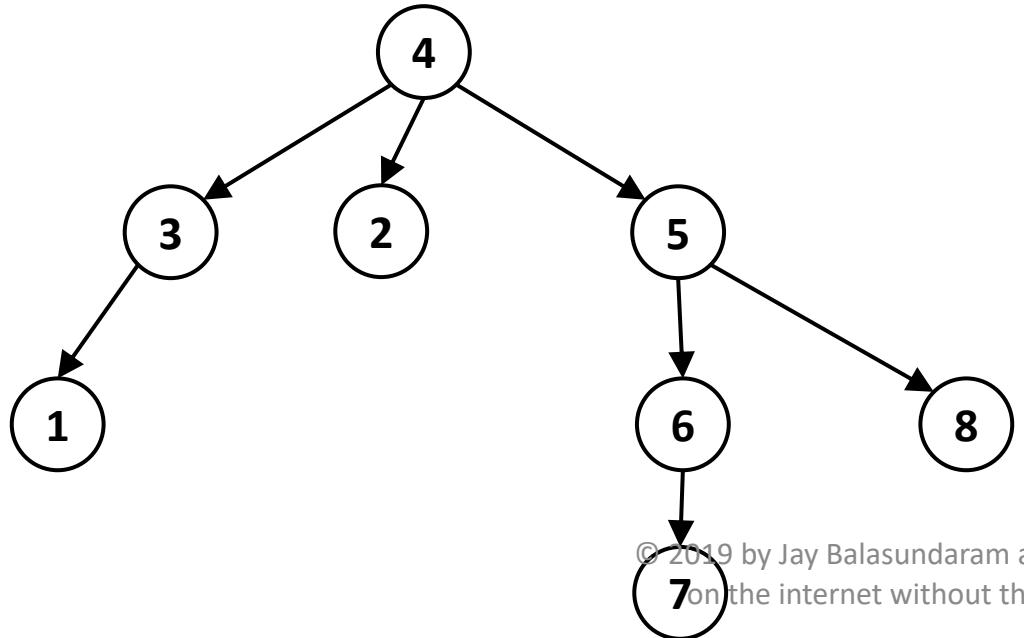


Adj List of G :

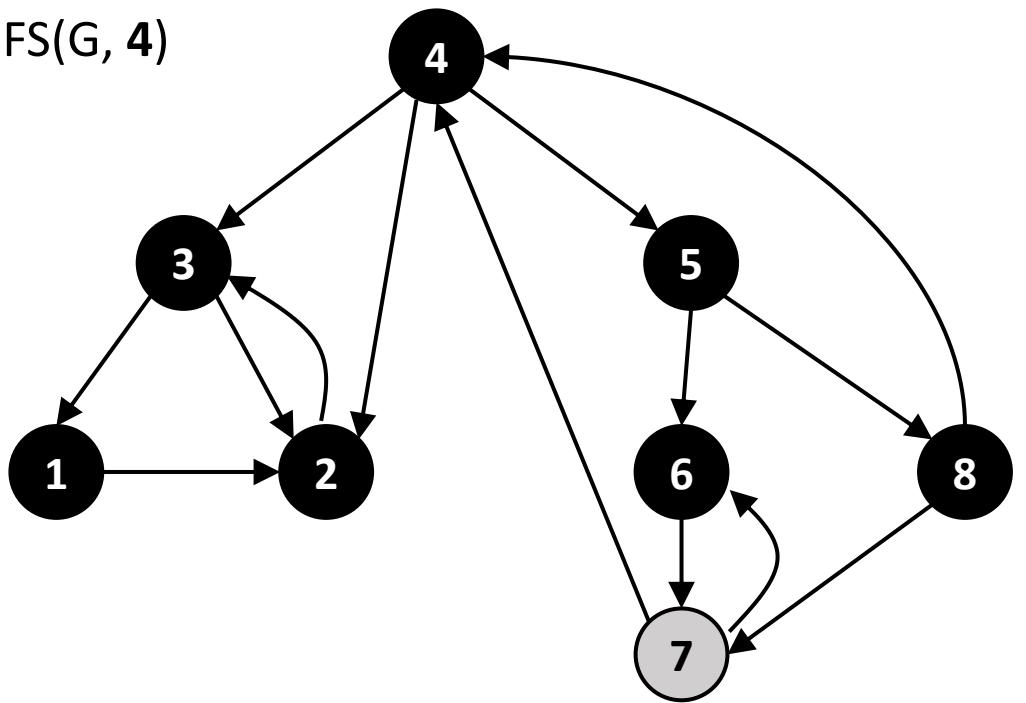
| | |
|-----|-------------|
| ✓ 1 | → 2 |
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| | 7 |



BFS(G , 4)

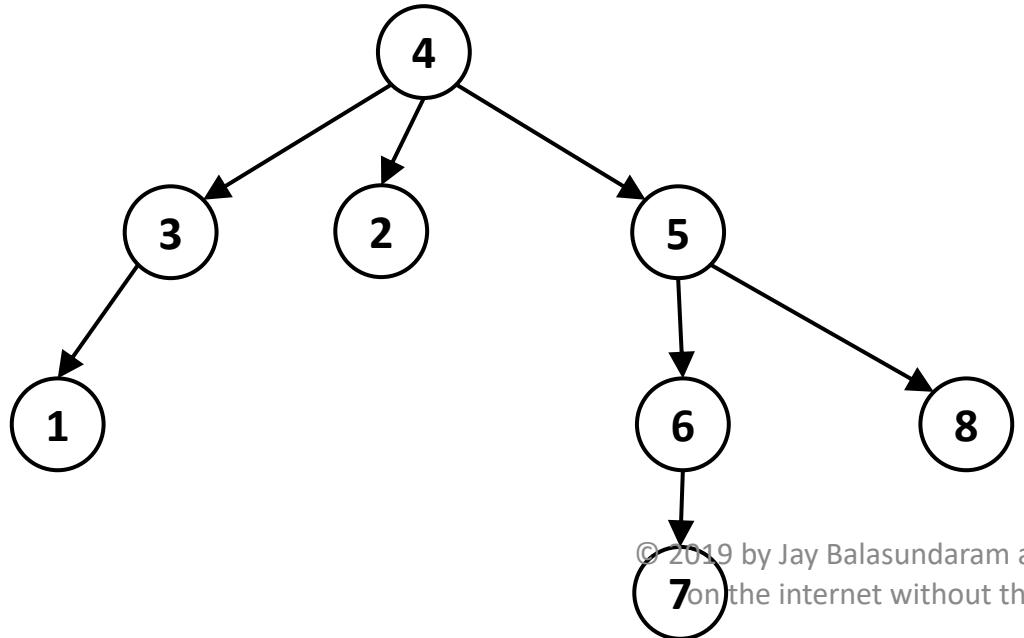


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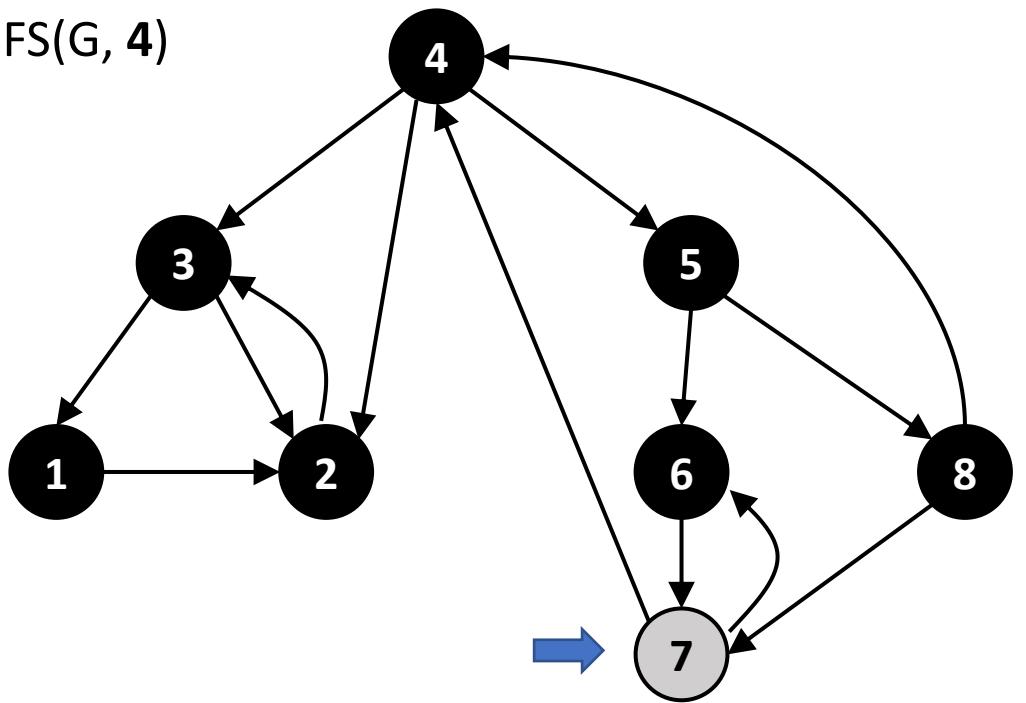
| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
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Contents of Q :

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| | 1 1 2 |
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| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |
| | 3 |
| | 7 |

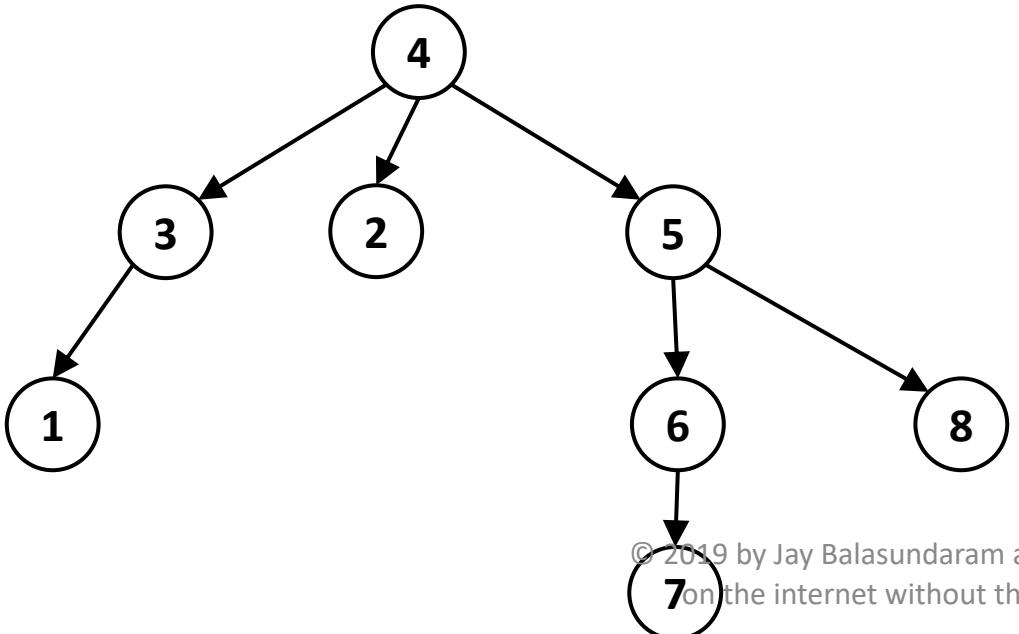


BFS(G , 4)



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
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| ✓ 7 | → 6 → 4 |
| ✓ 8 | → 7 → 4 |

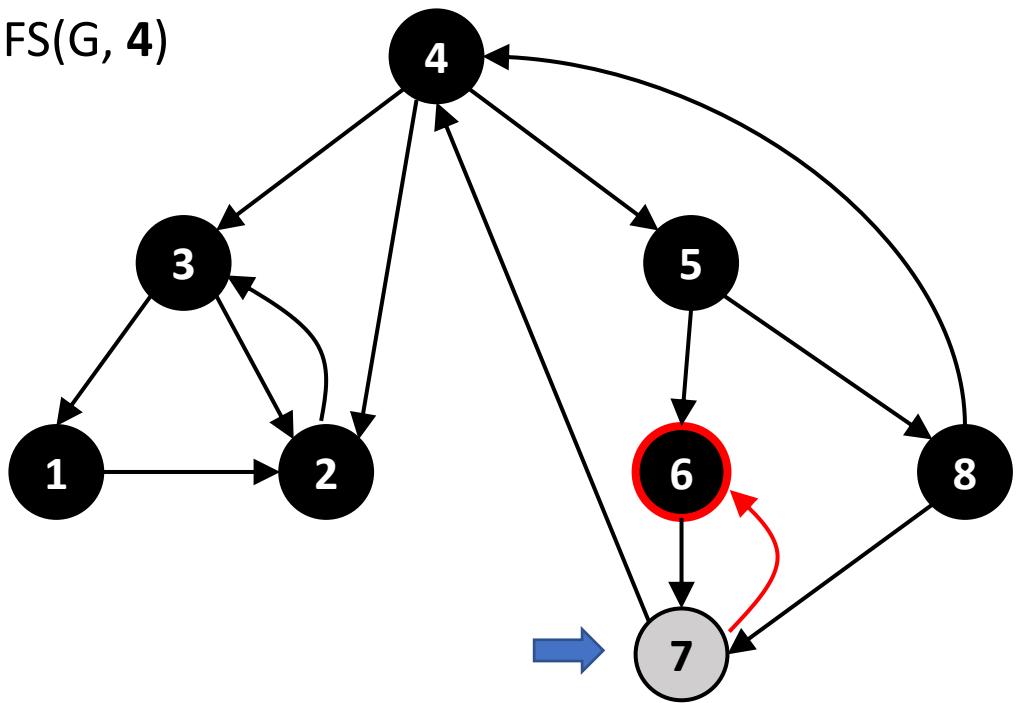


Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
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| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
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| | 2 3 |
| | 8 7 |
| | 3 |
| | 7 |

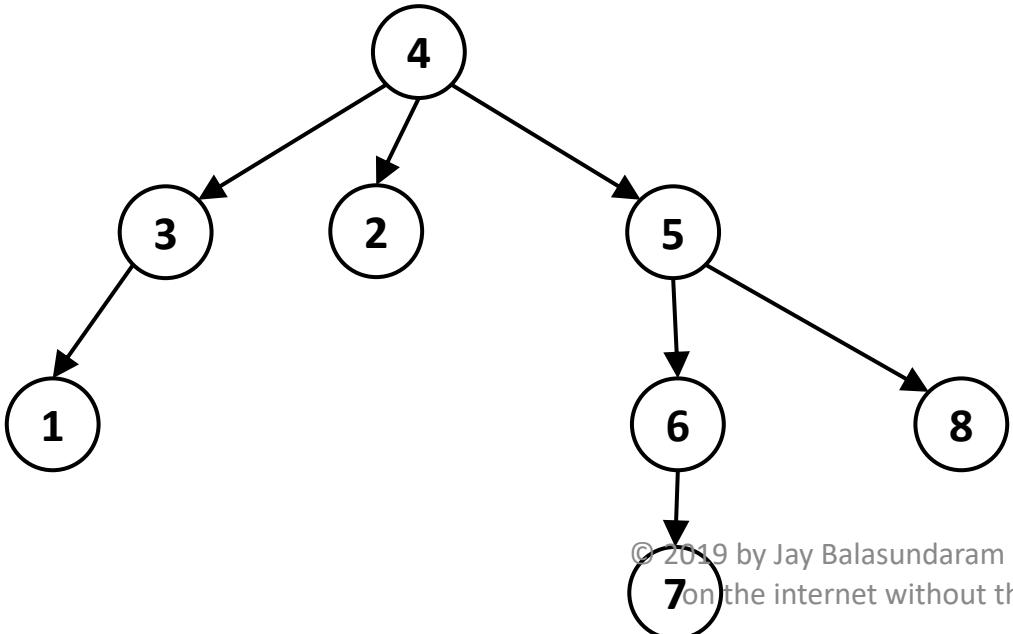


BFS(G , 4)



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
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| ✓ 5 | → 6 → 8 |
| ✓ 6 | → 7 |
| ✓ 7 | → 6 → 4 |
| ✓ 8 | → 7 → 4 |

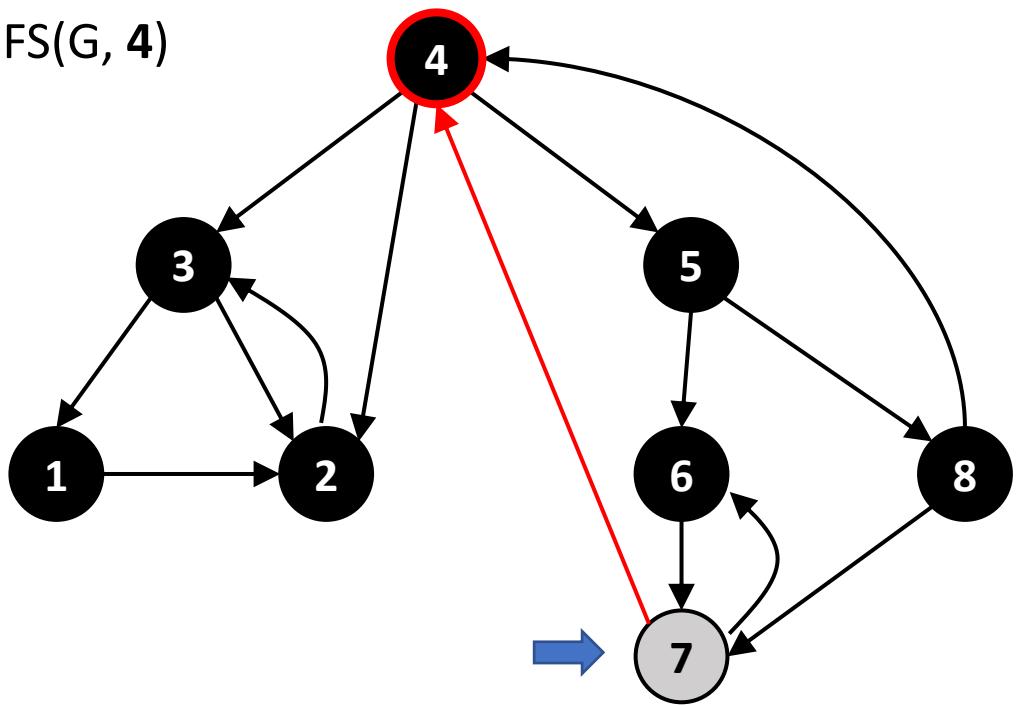


Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
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| | 1 1 2 |
| | 2 5 1 |
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| | 2 2 2 |
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| | 2 3 |
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| | 7 |

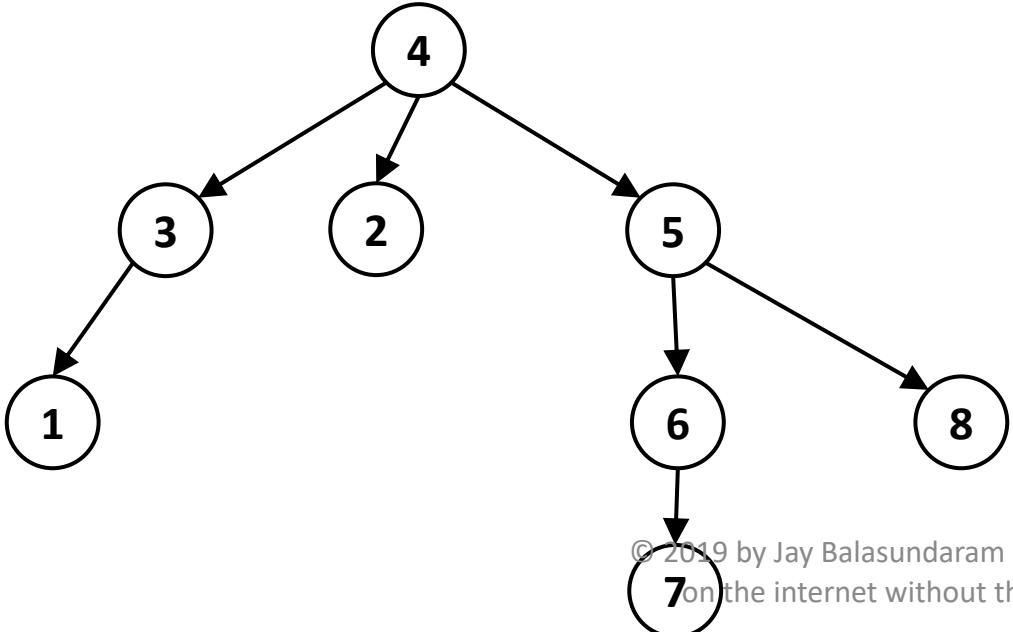


BFS(G , 4)



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| ✓ 6 | → 7 |
| ✓ 7 | → 6 → 4 |
| ✓ 8 | → 7 → 4 |

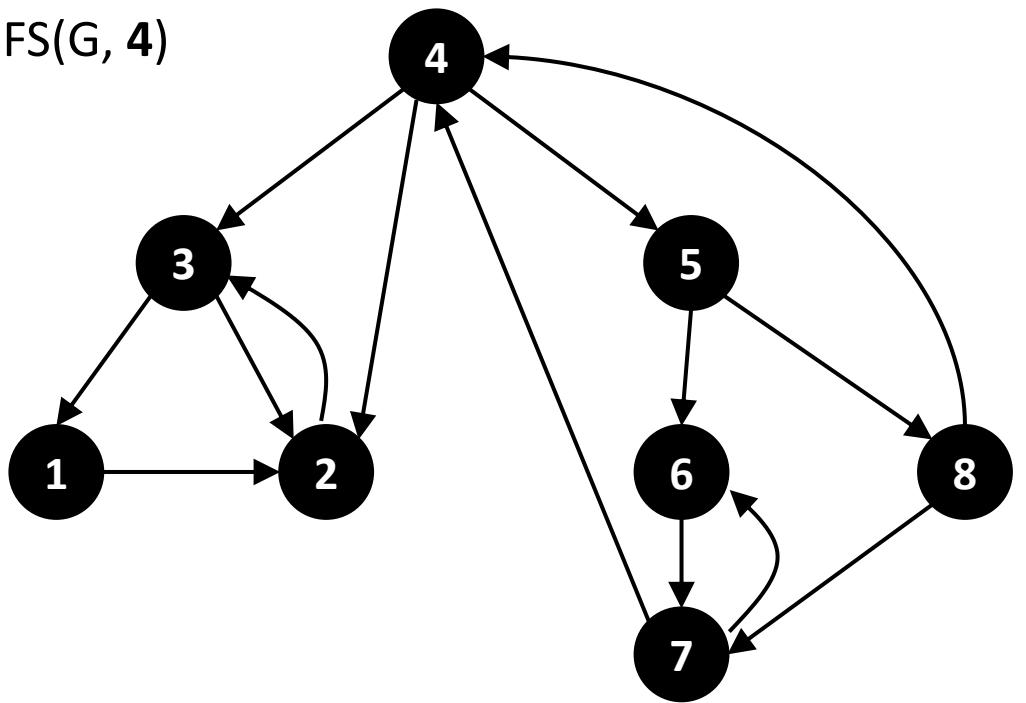


Contents of Q :

| | |
|---------|-------|
| $d = 0$ | 4 |
| $d = 1$ | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |
| | 3 |
| | 7 |



BFS(G, 4)

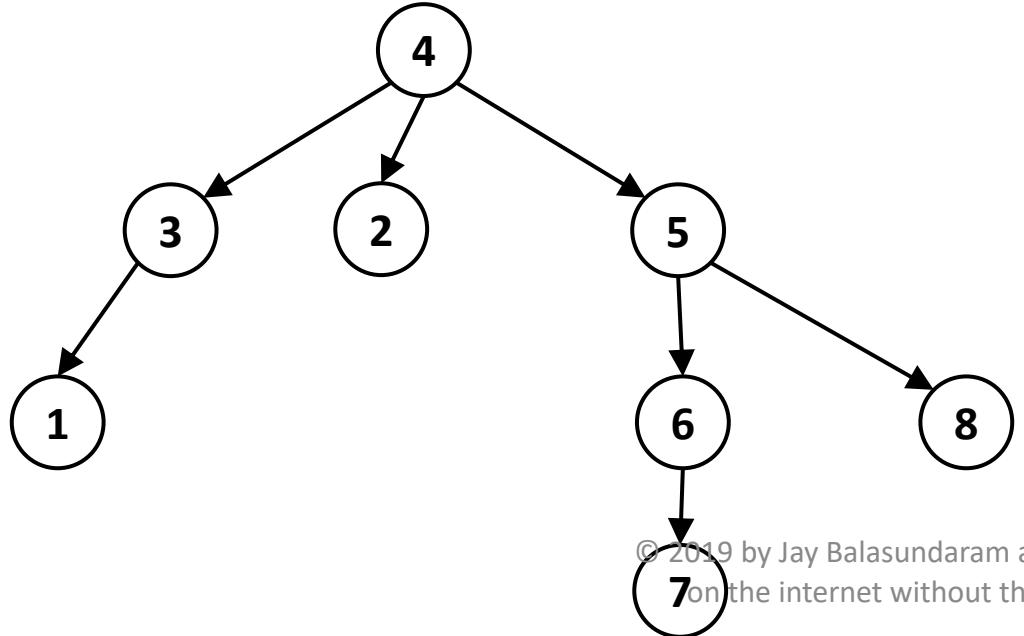


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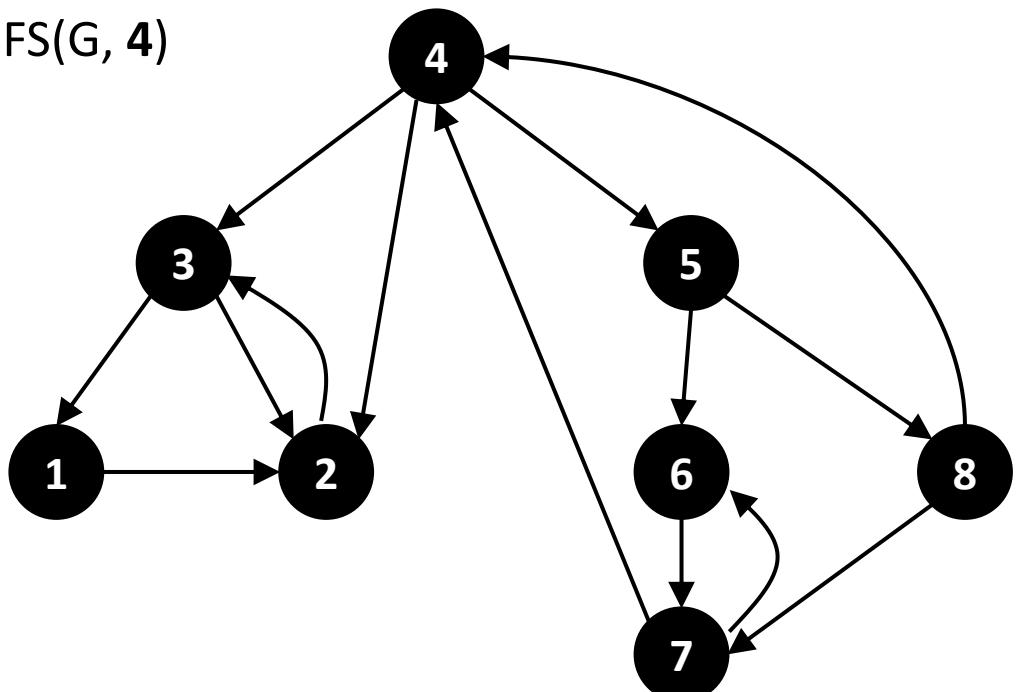
| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
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| ✓ 6 | → 7 |
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Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
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| | 1 1 2 |
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| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
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| | 2 3 |
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| | 7 |



BFS(G, 4)



Adj List of G :

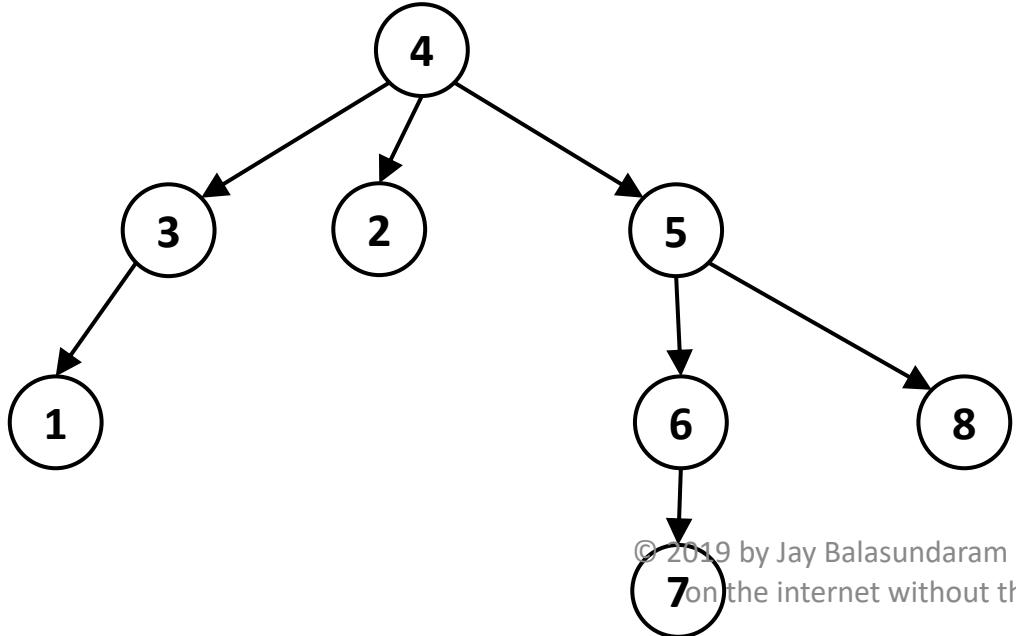
| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
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| ✓ 6 | → 7 |
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Contents of Q :

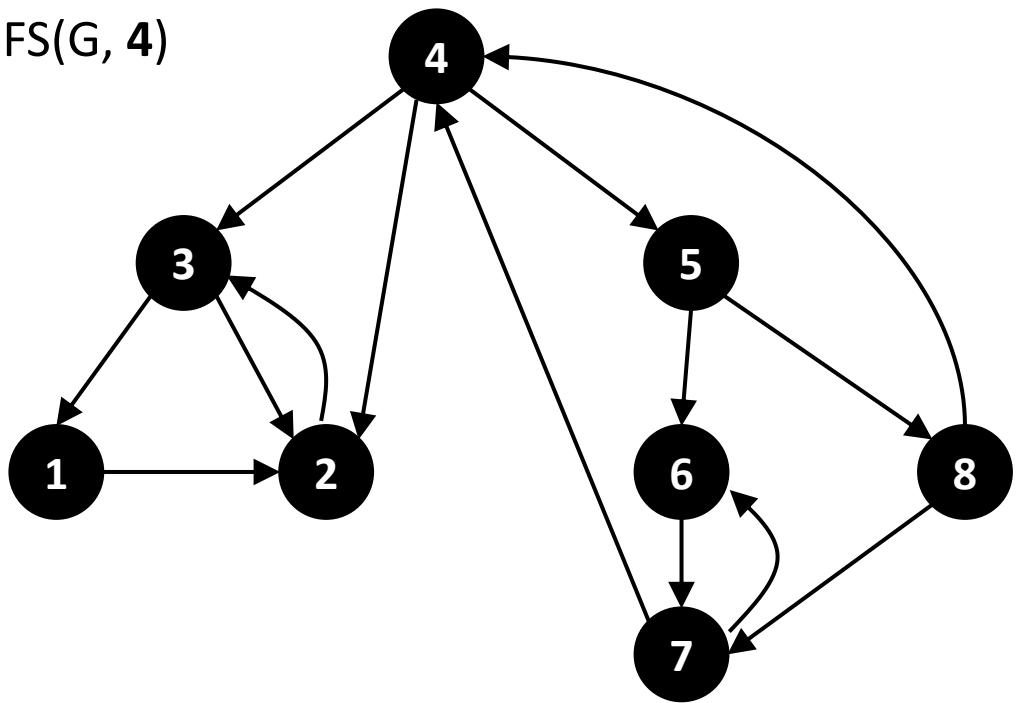
| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |
| | 3 |
| | 7 |

Worst-Case Time Complexity of BFS:

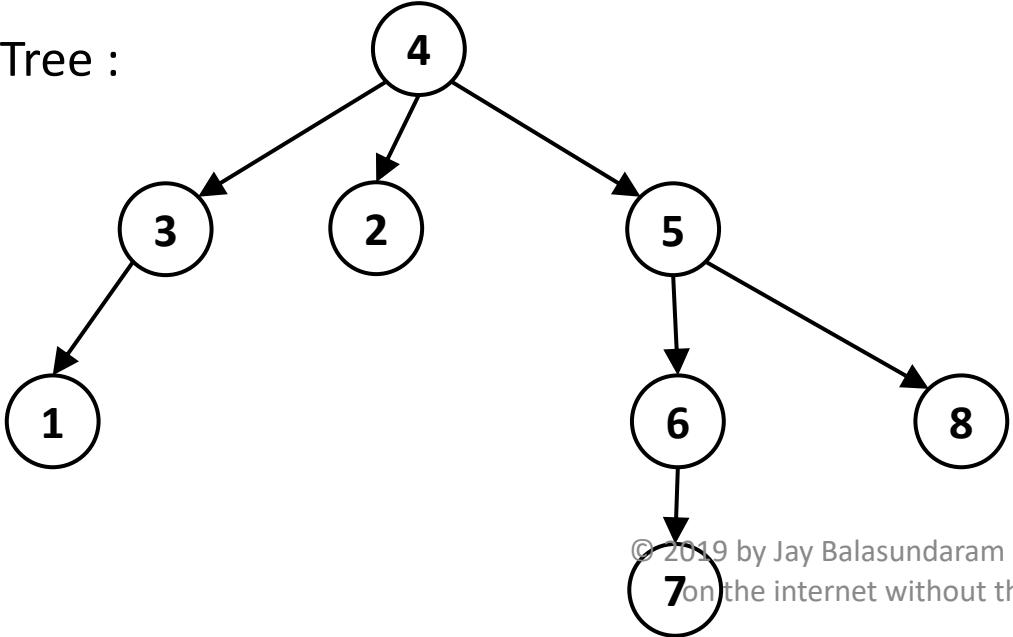
$$O(|V| + |E|)$$



BFS(G, 4)



BFS Tree :



Adj List of G :

| | |
|-----|-------------|
| ✓ 1 | → 2 |
| ✓ 2 | → 3 |
| ✓ 3 | → 1 → 2 |
| ✓ 4 | → 3 → 2 → 5 |
| ✓ 5 | → 6 → 8 |
| ✓ 6 | → 7 |
| ✓ 7 | → 6 → 4 |
| ✓ 8 | → 7 → 4 |

Contents of Q :

| | |
|-------|-------|
| d = 0 | 4 |
| d = 1 | 1 1 1 |
| | 3 2 5 |
| | 1 1 2 |
| | 2 5 1 |
| | 1 2 |
| | 5 1 |
| | 2 2 2 |
| | 1 6 8 |
| | 2 2 |
| | 6 8 |
| | 2 3 |
| | 8 7 |
| | 3 |
| | 7 |

Worst-Case Time Complexity of BFS:

$$O(|V| + |E|)$$

