符号	对应终结符	符号	对应终结符	符号	对应终结符
++ ! ~ + -	ω_0	*/%	ω_1	+ -	ω_2
<< >>	ω_3	<<=>>=	ω_4	== !=	ω_5
= -= += ^=	ω_6	П	" "	&&	'&&'
I	" "	٨	"A"	&	"&"
("(")	")"	;	"."
常数	c	标识符	v	字符串	s
void,int,float,char	type	if	if	else	else
{	"{"	}	"}"	while	while
,	,				

初始文法

我们将实现的的文法设计如下:

```
Program -> External_Declaration | Program External_Declaration
External_Declaration -> Function_Definition [IDF()] {function entry} | Declaration ";"
Function_definition -> type [addT] Direct_declarator Compound_statement
Direct declarator -> v [IDF] '(' ')'
compound_statement -> '{' '}' | '{' Statements'}'
Statements -> Statement | Statements Statement
Statement -> IfStatement | Assignment_expression ";" [clr()] | Declaration ";" [clr()] |
While_Statement | compound_statement | ';' [clr()]
Assignment_expression -> logic_or_expression \omega_6 Assignment_expression (GEQ \omega_6)
logic_or_expression
logic_or_expression -> logic_or_expression '||' logic_and_expression {GEQ || }|
logic_and_expression
```

logic_and_expression -> logic_and_expression "&&" or_expression {GEQ &&}| or_expression or_expression -> or_expression '|' xor_expression {GEQ |} | xor_expression xor_expression -> xor_expression '^' and_expression {GEQ ^}| and_expression and_expression -> and_expression '&' equal_expression {GEQ &} | equal_expression equal_expression -> equal_expression ω_5 relation_expression (GEQ ω_5)|relation_expression relation_expression -> relation_expression ω_4 shift_expression {GEQ ω_4 }| shift_expression shift_expression -> shift_expression ω_3 add_expression |{GEQ ω_3 } add_expression

```
add_expression -> add_expression \omega_2 mul_expression {GEQ \omega_2} | mul_expression
mul_expression -> mul_expression \omega_1 unary_expression {GEQ \omega_1} | unary_expression
unary_expression \rightarrow \omega_0 unary_expression {GEQ \omega_0} | T
T \rightarrow c \{push c\} \mid v \{push v\} \mid s \{push s\} \mid '('Assignment\_expression ')'
IfStatement -> if "(" E ")" {IF(if)} Statement {IE(ie)}
                 \mid if "(" E ")" Statement else {EL(el)} Statement {IE(ie)}
Declaration -> type [add_type] Variable_List
Variable_List -> Variable_List ',' init_declarator | init_declarator
init_declarator \rightarrow v [IDV] | v "=" Assignment_expression [GEQ w6]
While_Statement -> while \{WH()\} '(' Expression ')' {DO(do)} Statement {WE(we)}
```

其中,加粗字体表示为非终结符,花体字体与双引号内容表示为终结符。

LL1文法:

我们使用缩写将上述文法中的非终结符表示出来,其中{}内的内容表示select集, select集后的数字表示

```
产生式的编号,其中使用[]括起来的内容属于动作函数。
P -> ExD P1 \{type\} 58 \mid \epsilon \} 67
P1 -> ExD P1 \{type\} 59 | \epsilon {#} 68
ExD -> tpye [add type] v [add name] ExD1 {type} 60
ExD1 -> Func1 { "(" } 61 | [IDV(v, type, len)] Declaration1 ";" {"=", ";"} 62
Func1 -> "(" ")" [IDF(v, type)] CS { "(" } 63
Declaration1 -> w6 AsE [GEQ w6] 64 | \epsilon {",", ";"} 69
CS -> '{' CS1 '}' { '{'} 65
CS1 -> SS \{\omega_0, c, v, s, '(', '', ';', if, while, type\} 66 | \epsilon \{'\}'\} 70
Ss -> S Ss \{\omega_0, c, v, s, '(', '\{', '\}', if, while, type\}\} 0 | \epsilon \{'\}'\} 41
S -> ASE ";" \{\omega_0, c, v, s, '(')\} 2 | IfS \{if\} 44 | Declaration ";"\{type\} 45 | Wh \{while\} 46 | ";" \{j\}
47 | CS { '{'}}48
ASE -> LOE AsE<sub>1</sub> {\omega_0, c, v, s, '(') 4
\mathrm{AsE}_1 \rightarrow \omega_6 AsE [GEQ \omega_6] {\omega_6} 5 | [deal] \omega LoE {\omega} 6 | \epsilon {";",")", ","} 42
LoE ->LaE LoE_1 {\omega_0, c, v, s, '(' } 7
LoE_1 \rightarrow "||" LaE [GEQ || ] LoE_1 \in "||" 8| "\epsilon" {\omega_6, ';', ')', #}9
LaE -> OrE LaE_1 \{\omega_0, c, v, s, '(')\}10
LaE_1 -> "&&" OrE LaE_1 [GEQ && ] {"&&"}11 | \epsilon {\omega_6,"||",';', ')', #}12
OrE -> XoE OrE<sub>1</sub> {\omega_0, c, v, s, '(')13
OrE_1 -> "|" XoE [GEQ | ] OrE_1 \{"|"\} 14 | \epsilon \{\omega_6, "| |", "\&\&", ';', ")', \#\}15
XoE -> AnE XoE_1 \{\omega_0, c, v, s, '(')\}16
```

 $XoE_1 \rightarrow \text{"^"}$ AnE [GEQ ^] $XoE_1 \in \{\omega_6, \text{"} \mid \text{", "&, ", ", ", ", "}\}$ 18

AnE -> EqE $AnE_1 \{\omega_0, \ c$, v , s, '(' }19

 $AnE_1 -> \text{``&''} \ \textbf{EqE} \ [\mathsf{GEQ} \ \& \] \ AnE_1 \ \ \{\text{``&''}\} \textbf{20} \ \ | \ \epsilon \ \{\omega_6,\text{`'} | \ | \text{''},\text{``&&''},\text{''} | \text{''},\text{''}^{'}, \ ')', \ \#\} \textbf{21}$

EqE -> ReE EqE_1 { ω_0 , c, v, s, '(' } 22

 ${
m EqE}_1 ext{ -> } \omega_5 \ {
m ReE} \ [{
m GEQ} \ \omega_5] \ {
m EqE}_1 \ \{\omega_5\}$ 23 | $\epsilon \ \{\omega_6,..."^", "\&",';', ")', \#\}$ 24

ReE -> ShE ${
m ReE}_1 \ \{\omega_0,\ c$, v , s, '(' }25

 ${\rm ReE}_1 > \omega_4$ **ShE** [GEQ ω_4] ${\rm ReE}_1$ { ω_4 }**26** | ϵ { ω_6 ,..."&", ω_5 ,';', ')', #}**27**

ShE -> AdE ${\rm ShE}_1 \ \{\omega_0,\ c$, v , s, '(' }28

 $ShE_1 \rightarrow \omega_3$ AdE [GEQ ω_3] ShE_1 { ω_3 }29 | ϵ { ω_6 ,... ω_5 , ω_4 ,';', ')', #}30

AdE -> MuE $AdE_1 \{\omega_0, c, v, s, '(')\}$ 31

 $AdE_1 \rightarrow \omega_2$ MuE[GEQ ω_2] AdE_1 { ω_2 }32 | ϵ { $\omega_6,...\omega_4$, ω_3 ,';', ')', #}33

MuE -> UnE MuE_1 { ω_0 , c, v, s, '(' } 34

 $\mathrm{MuE}_1 \rightarrow \omega_1$ UnE [GEQ ω_1] $\mathrm{MuE}_1 \{\omega_1\}$ 35 | $\epsilon \{\omega_6,...\omega_2,\omega_1,';', ')', \#\}$ 36

UnE -> ω_0 UnE [GEQ ω_0] { ω_0 } 37 | T {c , v , s, '(' }38

 $extsf{T} ext{->} c ext{ [push c] } | v ext{ [push v] } | s ext{ [push s] } extsf{39} | '(' extsf{ASE '})' ext{ {"(")}40}$

IfS -> if "(" AsE ")" [IF(if)] S [EL(el)] IFS1 $\{if\}$ [IE(ie)] 49

IFS1 -> else **S** {else} **50** | ' ϵ ' { ω_0 , c, v, s, '(', '\;', if, while, type }**51**

Declaration -> type [add_type] **VL** {type} **52**

 $VL \rightarrow v [IDV()] VL1 \{v\} 53$

VL1 -> **Declaration1 VL1** {"w6"} **54** | "," v [add_name] **VL1** {","} **55** | ϵ {";"} **56**

Wh \rightarrow while [WH()] '(' ASE ')' [DO(do)] **S** [WE(we)] {while} **57**