

# Part1-Introduction

Completed on	@2020/03/14
Key videos	<a href="https://www.coursera.org/learn/algorithms-divide-conquer/lecture/8vama/about-the-course">https://www.coursera.org/learn/algorithms-divide-conquer/lecture/8vama/about-the-course</a> <a href="https://www.coursera.org/learn/algorithms-divide-conquer/lecture/wKEYL/karatsuba-multiplication">https://www.coursera.org/learn/algorithms-divide-conquer/lecture/wKEYL/karatsuba-multiplication</a> <a href="https://www.coursera.org/learn/algorithms-divide-conquer/lecture/wW9On/merge-sort-analysis">https://www.coursera.org/learn/algorithms-divide-conquer/lecture/wW9On/merge-sort-analysis</a> <a href="https://www.coursera.org/learn/algorithms-divide-conquer/lecture/q5fV4/guiding-principles-for-analysis-of-algorithms">https://www.coursera.org/learn/algorithms-divide-conquer/lecture/q5fV4/guiding-principles-for-analysis-of-algorithms</a>
Note	Introduction to algorithms by two examples, i.e., Karatsuba Multiplication and Merge Sort.

## Note

### ▼ Integer Multiplication

- 最原始的计算：输入：两个 $n$ 位数字；输出：两个数字的乘积；最原始的运算包括加(add)与乘(multiply)
- 一般整数乘法：考虑乘法+加法，每一行地复杂度 $\leq O(2n)$ ；共有 $n$ 行  $\Rightarrow$  总体地算法复杂度为 $O(n^2)$

Handwritten multiplication of 5678 by 1234:

$$\begin{array}{r} 5678 \\ \times 1234 \\ \hline 22712 \\ 17034 \phantom{0} \\ 11356 \phantom{00} \\ 5678 \phantom{000} \\ \hline 7006652 \end{array}$$

Roughly  $n$  operations per row up to a constant

- 能否有更好的大数乘法？

## ▼ Karatsuba Multiplication

$$x \cdot y = 10^n ac + 10^{n/2}(ad + bc) + bd$$

一种快速相乘算法，将两个大数x,y分解成各自一半得到a,b,c,d，然后就可以根据下列公式，得到优化后的成绩，主要是中间部分可以分解为以下方法，节省一次运算，此外可以进行递归，recursive operation进一步分解：

$$ad + bc = (a + b)(c + d) - ac - bd$$

相比之下，只需三次基本操作，而不需要原有的四次操作分别计算ac，ad，bc，bd

## ▼ About the Course

主要关注主题

- Vocabulary for design and analysis of algorithm

"Big-Oh" notation, 应该值得是算法复杂度

"sweet spot"

- Divide and conquer algorithm design paradigm
- Randomization in algorithm design

QuickSort, primality testing, hashing

- Primitives for reasoning about graphs

最短路径，连通性问题

- Use and implementation of data structure

Heaps, balanced **binary search trees**, **hashing tables** (重点讲) and some variants

主要讲解问题

1. Greedy algorithm design paradigm 贪婪算法
2. Dynamic programming algorithm design paradigm 动态规划
3. NP-complete problems and what to do about them NP完备问题
4. Fast heuristics with provable guarantees
5. Fast exact algorithms for special cases

## 6. Exact algorithms that beat brute-force search

### ▼ Merge Sort: Motivation and Example

- Motivation
  - 是对于分而治之方法的范例，比Selection, Insertion, Bubble sorts这些排序算法更为高效 $O(n \log n)$
  - 可以帮助掌握最坏情况分析以及渐进分析的原理 guiding principles for algorithm analysis (worst-case and asymptotic analysis)
- 代码思路：显示把一串数列分为一分为二，然后递归排序之后，再进行合并（Merge）。

### ▼ Merge Sort: Pseudocode

- 伪代码：就是同时从开头遍历两个分数组，若其中一个元素较小，则复制到已声明的空间中并+1，以此类推

#### Pseudocode for Merge:

<pre>C = output [length = n] A = 1<sup>st</sup> sorted array [n/2] B = 2<sup>nd</sup> sorted array [n/2] i = 1 j = 1</pre>	<pre>for k = 1 to n     if A(i) &lt; B(j)         C(k) = A(i)         i++     else [B(j) &lt; A(i)]         C(k) = B(j)         j++ end</pre>
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(ignores end cases)

- 复杂度分析，主要考虑merge这个步骤，合并数组长度为N  
初始复制2次操作，后续循环加判断4N个操作  $\Rightarrow$  总体复杂度  $\leq 4N + 2 \leq 6N$
- 运行时间分析（Running time），精准的上界  $\Rightarrow$  最多以下操作次数：

$$T(n) \leq 6n \log_2 n + 6n = \mathcal{O}(n \log n)$$

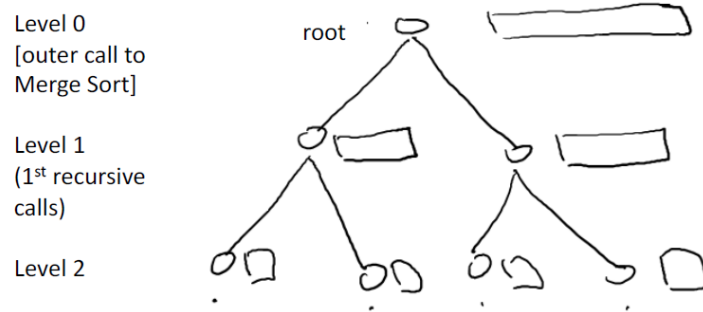
比一般排序的 $O(n^2)$ 复杂度更低

### ▼ Merge Sort: Analysis

- recursion tree 递归跟踪分析方法

一共有  $\log_2 N$  个层次 (this recursion tree has  $\log_2 N$  levels)

Proof of claim (assuming  $n = \text{power of } 2$ ):



对于第 $j$ 层的子问题数量和大小分析：

What is the pattern ? Fill in the blanks in the following statement: at each level  $j=0,1,2,\dots,\log_2(n)$  there are \_\_\_\_ subproblems, each of size \_\_\_\_.

- ☐  $2^j$  and  $2^j$  respectively
- ☐  $\frac{n}{2^j}$  and  $\frac{n}{2^j}$  respectively
- ☒  $2^j$  and  $\frac{n}{2^j}$  respectively

Correct

- 最终复杂度分析

1) 首先对于长度为 $n$ 的数组，每一层都会会花 $6n$ 的时间进行合并  $\Rightarrow$  一共有  $\log_2 N$  个子层+1最终层

## Proof of claim (assuming $n = \text{power of } 2$ ) :

At each level  $j=0,1,2,\dots, \log_2 n$ ,

Total # of operations at level  $j = 0,1,2,\dots, \log_2 n$

$$\leq 2^j * 6\left(\frac{n}{2^j}\right) = 6n$$

# of level- $j$   
subproblems

Size of level- $j$   
subproblem

Work per level -  $j$   
subproblem

Total

$$6n(\log_2 n + 1)$$

Work  
per level

# of  
levels

1

### ▼ Guiding Principles for Analysis of Algorithms

结合归并排序算法进行divide-and-conquer范例算法得学习

- worst-case 最恶劣情况分析：计算整个算法的上界，分析相对较为简单；  
对应的“平均情况”和基准分析需要相关领域的知识

As Opposed to

--"average-case" analysis

--benchmarks

REQUIRES DOMAIN  
KNOWLEDGE

- 分析时会忽略常数项以及低阶项  $\Rightarrow$  真正影响性能的是最高次项
- asymptotic analysis 渐进分析：比较注重分析较为大型连续的数据，而不是简单的单个案例分析！

Asymptotic Analysis : focus on running time for large input sizes  $n$

Eg :  $6n \log_2 + 6n$  "better than"  $\frac{1}{2}n^2$   
MERGE SORT INSERTION SORT

- 什么是快的算法：随着输入大小增长慢，接近常数级复杂度

fast  
algorithm



worst-case running time  
grows slowly with input size

Usually : want as close to linear ( $O(n)$ ) as possible

## Reference

- 大数乘法Karatsuba Multiplication

Karatsuba算法

<https://zh.wikipedia.org/wiki/Karatsuba%E7%AE%97%E6%B3%95>

- Merge Sort

归并排序

<https://zh.wikipedia.org/wiki/%E5%BD%92%E5%B9%B6%E6%8E%92%E5%BA%8F>

插入排序

<https://zh.wikipedia.org/wiki/%E6%8F%92%E5%85%A5%E6%8E%92%E5%BA%8F>

- Sorting Algorithms

Asymptotic Analysis--渐近分析

<https://blog.csdn.net/dingchenxixi/article/details/52449424>