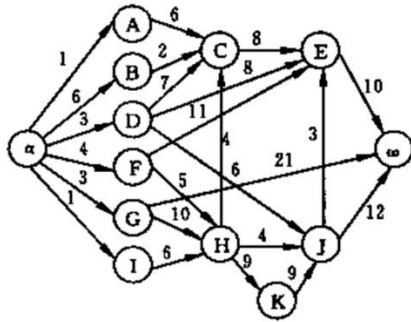


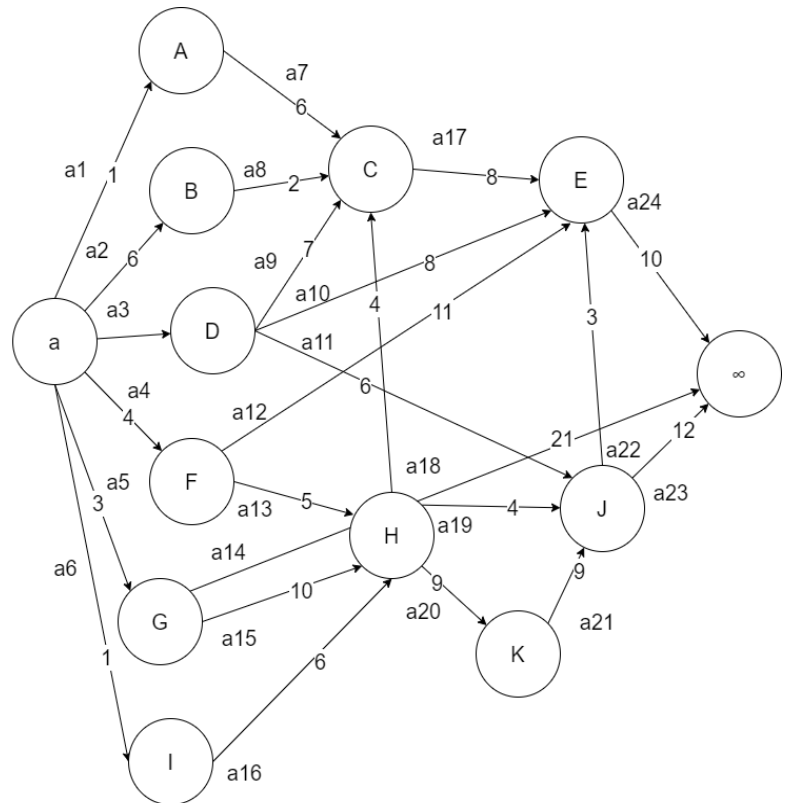
1、针对下图 AOE 网络，计算个活动弧的  $e(a_i)$ ,  $l(a_i)$  的值、各事件（顶点）的  $ve(v_i)$ ,  $l(v_i)$ 。



解：右图为各边的编号以及边和点的信息。

**$Ve(v_i)$  :**

$Ve(A) = 1$   
 $Ve(B) = 6$   
 $Ve(C) = 17(a \rightarrow G \rightarrow H \rightarrow C)$   
 $Ve(D) = 3$   
 $Ve(E) = 25(a \rightarrow G \rightarrow H \rightarrow C \rightarrow E)$   
 $Ve(F) = 4$   
 $Ve(G) = 3$   
 $Ve(H) = 13(a \rightarrow G \rightarrow H)$   
 $Ve(I) = 1$   
 $Ve(J) = 31(a \rightarrow G \rightarrow H \rightarrow K \rightarrow J)$   
 $Ve(K) = 22(a \rightarrow G \rightarrow H \rightarrow K)$   
 $Ve(\infty) = 43$



**$I(v_i)$  :**

$I(A) = 19$   
 $I(B) = 23$   
 $I(C) = 25$   
 $I(D) = 18$   
 $I(E) = 33$   
 $I(F) = 8$   
 $I(G) = 3$

$I(H) = 13$   
 $I(I) = 7$   
 $I(J) = 31$   
 $I(K) = 22$   
 $I(\infty) = 43$

**$e(a_i)$**

$e(a_1) = 0$   
 $e(a_2) = 0$   
 $e(a_3) = 0$   
 $e(a_4) = 0$   
 $e(a_5) = 0$   
 $e(a_6) = 0$

$e(a_7) = 1$   
 $e(a_8) = 6$   
 $e(a_9) = 3$   
 $e(a_{10}) = 3$   
 $e(a_{11}) = 3$   
 $e(a_{12}) = 4$

$e(a_{13}) = 4$   
 $e(a_{14}) = 3$   
 $e(a_{15}) = 3$   
 $e(a_{16}) = 1$   
 $e(a_{17}) = 17$   
 $e(a_{18}) = 13$

$e(a_{19}) = 13$   
 $e(a_{20}) = 13$   
 $e(a_{21}) = 22$   
 $e(a_{22}) = 31$   
 $e(a_{23}) = 31$   
 $e(a_{24}) = 25$

$l(a_i)$

$l(a_1) = 18$

$l(a_2) = 17$

$l(a_3) = 15$

$l(a_4) = 4$

$l(a_5) = 0$

$l(a_6) = 6$

$l(a_7) = 19$

$l(a_8) = 22$

$l(a_9) = 18$

$l(a_{10}) = 25$

$l(a_{11}) = 25$

$l(a_{12}) = 27$

$l(a_{13}) = 8$

$l(a_{14}) = 22$

$l(a_{15}) = 3$

$l(a_{16}) = 7$

$l(a_{17}) = 25$

$l(a_{18}) = 21$

$l(a_{19}) = 27$

$l(a_{20}) = 13$

$l(a_{21}) = 22$

$l(a_{22}) = 30$

$l(a_{23}) = 31$

$l(a_{24}) = 33$

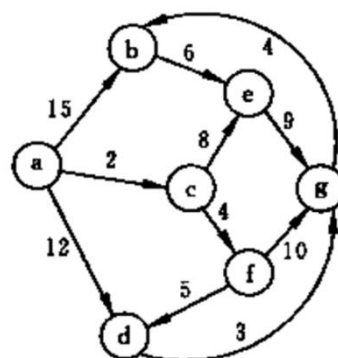
2. 利用 Dijkstra 算法，求图中顶点 a 到其他各顶点的最短路径，写出执行算法过程中各步的状态

解：

步骤一：以 a 为起点初始待定路径表

STEP1

生长点	b	c	d	e	f	g
a	$P(b) = ab$ $D(b) = 15$	$P(c) = ac$ $D(c) = 2$	$P(d) = ad$ $D(d) = 12$	$P(e) = \text{null}$ $D(e) = \infty$	$P(f) = \text{null}$ $D(f) = \infty$	$P(g) = \text{null}$ $D(g) = \infty$



步骤二：从待定路径表中选出一条最短的边，设其顶点为新的生长点，并且对剩下每一个生长点进行比较，若以新生长点为中转的路径短于原始路径则替换。

STEP2

生长点	b	c	d	e	f	g
a	$P(b) = ab$ $D(b) = 15$	$P(c) = ac$ $D(c) = 2$	$P(d) = ad$ $D(d) = 12$	$P(e) = \text{null}$ $D(e) = \infty$	$P(f) = \text{null}$ $D(f) = \infty$	$P(g) = \text{null}$ $D(g) = \infty$
c	$P(b) = ab$ $D(b) = 15$		$P(d) = ad$ $D(d) = 12$	$P(e) = ace$ $D(e) = 10$	$P(f) = acf$ $D(f) = 6$	$P(g) = \text{null}$ $D(g) = \infty$

对步骤二进行  $n-2$  次重复（注：绿色表示当前有新的通路，但路径长度大于原来的路径故步替换，红色为新的通路的路径长度小于原来路径的长度故替换）

STEP3

生长点	b	c	d	e	f	g
a	$P(b) = ab$ $D(b) = 15$	$P(c) = ac$ $D(c) = 2$	$P(d) = ad$ $D(d) = 12$	$P(e) = \text{null}$ $D(e) = \infty$	$P(f) = \text{null}$ $D(f) = \infty$	$P(g) = \text{null}$ $D(g) = \infty$
c	$P(b) = ab$ $D(b) = 15$		$P(d) = ad$ $D(d) = 12$	$P(e) = ace$ $D(e) = 10$	$P(f) = acf$ $D(f) = 6$	$P(g) = \text{null}$ $D(g) = \infty$
f	$P(b) = ab$ $D(b) = 15$		$P(d) = acfd$ $D(d) = 11$	$P(e) = ace$ $D(e) = 10$		$P(g) = acfg$ $D(g) = 16$

## STEP4

生长点	b	c	d	e	f	g
a	P(b) = ab D(b) = 15	P(c) = ac D(c) = 2	P(d) = ad D(d) = 12	P(e) = null D(e) = ∞	P(f) = null D(f) = ∞	P(g) = null D(g) = ∞
c	P(b) = ab D(b) = 15		P(d) = ad D(d) = 12	P(e) = ace D(e) = 10	P(f) = acf D(f) = 6	P(g) = null D(g) = ∞
f	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11	P(e) = ace D(e) = 10		P(g) = acfg D(g) = 16
e	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11			P(g) = acfg D(g) = 16

## STEP5

生长点	b	c	d	e	f	g
a	P(b) = ab D(b) = 15	P(c) = ac D(c) = 2	P(d) = ad D(d) = 12	P(e) = null D(e) = ∞	P(f) = null D(f) = ∞	P(g) = null D(g) = ∞
c	P(b) = ab D(b) = 15		P(d) = ad D(d) = 12	P(e) = ace D(e) = 10	P(f) = acf D(f) = 6	P(g) = null D(g) = ∞
f	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11	P(e) = ace D(e) = 10		P(g) = acfg D(g) = 16
e	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11			P(g) = acfg D(g) = 16
d	P(b) = ab D(b) = 15					P(g) = acfdg D(g) = 14

## STEP6

生长点	b	c	d	e	f	g
a	P(b) = ab D(b) = 15	P(c) = ac D(c) = 2	P(d) = ad D(d) = 12	P(e) = null D(e) = ∞	P(f) = null D(f) = ∞	P(g) = null D(g) = ∞
c	P(b) = ab D(b) = 15		P(d) = ad D(d) = 12	P(e) = ace D(e) = 10	P(f) = acf D(f) = 6	P(g) = null D(g) = ∞
f	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11	P(e) = ace D(e) = 10		P(g) = acfg D(g) = 16
e	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11			P(g) = acfg D(g) = 16
d	P(b) = ab D(b) = 15					P(g) = acfdg D(g) = 14
g	P(b) = ab D(b) = 15					

## STEP7

生长点	b	c	d	e	f	g
a	P(b) = ab D(b) = 15	P(c) = ac D(c) = 2	P(d) = ad D(d) = 12	P(e) = null D(e) = ∞	P(f) = null D(f) = ∞	P(g) = null D(g) = ∞
c	P(b) = ab D(b) = 15		P(d) = ad D(d) = 12	P(e) = ace D(e) = 10	P(f) = acf D(f) = 6	P(g) = null D(g) = ∞
f	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11	P(e) = ace D(e) = 10		P(g) = acfg D(g) = 16
e	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11			P(g) = acfg D(g) = 16
d	P(b) = ab D(b) = 15					P(g) = acfdg D(g) = 14
g	P(b) = ab D(b) = 15					
b				P(e) = ace D(e) = 10		

## STEP8

生长点	b	c	d	e	f	g
a	P(b) = ab D(b) = 15	P(c) = ac D(c) = 2	P(d) = ad D(d) = 12	P(e) = null D(e) = ∞	P(f) = null D(f) = ∞	P(g) = null D(g) = ∞
c	P(b) = ab D(b) = 15		P(d) = ad D(d) = 12	P(e) = ace D(e) = 10	P(f) = acf D(f) = 6	P(g) = null D(g) = ∞
f	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11	P(e) = ace D(e) = 10		P(g) = acfg D(g) = 16
e	P(b) = ab D(b) = 15		P(d) = acfd D(d) = 11			P(g) = acfg D(g) = 16
d	P(b) = ab D(b) = 15					P(g) = acfdg D(g) = 14
g	P(b) = ab D(b) = 15					
b				P(e) = ace D(e) = 10		
a	P(b) = ab D(b) = 15	P(c) = ac D(c) = 2	P(d) = acfd D(d) = 11	P(e) = ace D(e) = 10	P(f) = acf D(f) = 6	P(g) = acfdg D(g) = 14

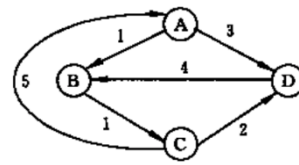
3. 利用 Floyd 算法，求图中各对顶点间的最短路径

解：

STEP1: 初始邻接矩阵 D, 和 Path

	A	B	C	D
A	0	1	$\infty$	3
B	$\infty$	0	1	$\infty$
C	5	$\infty$	0	2
D	$\infty$	4	$\infty$	0

	A	B	C	D
A	-1	-1	-1	-1
B	-1	-1	-1	-1
C	-1	-1	-1	-1
D	-1	-1	-1	-1



STEP2: 以 A 为中转点，进行第一次迭代

	A	B	C	D
A	0	1	$\infty$	3
B	$\infty$	0	1	$\infty$
C	5	6	0	2
D	$\infty$	4	$\infty$	0

	A	B	C	D
A	-1	-1	-1	-1
B	-1	-1	-1	-1
C	-1	A	-1	-1
D	-1	-1	-1	-1

STEP3: 以 B 为中转点，进行第二次迭代

	A	B	C	D
A	0	1	2	3
B	$\infty$	0	1	$\infty$
C	5	6	0	2
D	$\infty$	4	5	0

	A	B	C	D
A	-1	-1	-1	-1
B	-1	-1	-1	-1
C	-1	A	-1	-1
D	-1	-1	B	-1

STEP4: 以 C 为中转点，进行第三次迭代

	A	B	C	D
A	0	1	2	3
B	6	0	1	3
C	5	6	0	2
D	10	4	5	0

	A	B	C	D
A	-1	-1	-1	-1
B	C	-1	-1	C
C	-1	A	-1	-1
D	C	-1	3	-1

STEP5: 以 D 为中转点，进行第四次迭代

	A	B	C	D
A	0	1	2	3
B	6	0	1	3
C	5	6	0	2
D	10	4	5	0

	A	B	C	D
A	-1	-1	-1	-1
B	C	-1	-1	C
C	-1	A	-1	-1
D	C	-1	B	-1

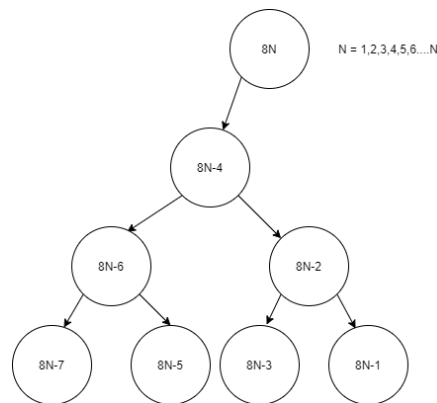
STEP6: 得出最终的 D, 和 Path 矩阵

	A	B	C	D
A	0	1	2	3
B	6	0	1	3
C	5	6	0	2
D	10	4	5	0

	A	B	C	D
A	-1	-1	-1	-1
B	C	-1	-1	C
C	-1	A	-1	-1
D	C	-1	B	-1

4. 已知一个有序表的表长为  $8N$ , 并且表中没有关键字相同的记录。假设按如下所述方法查找一个关键字等于给定值  $K$  的记录: 先在第  $8, 16, 24, \dots, 8K, \dots, 8N$  个记录中进行顺序查找, 或者查找成功, 或者由此确定出一个继续进行折半查找的范围。画出描述上述查找过程的判定树, 并求等概率查找时查找成功的平均查找长度。

解: 该查找过程的判定树如下



$$ASL = \frac{N+1}{2} + \frac{(1+2 \times 2 + 3 \times 4)}{8}$$

-----END-----