

# MATH 601 (DUE 11/13)

HIDENORI SHINOHARA

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### 1. FACTORING POLYNOMIALS WITH COEFFICIENTS IN FINITE FIELDS

**Exercise.** (Problem 14) For  $a \in \mathbb{F}_q$ , what are the possible values for  $a^{(q-1)/2}$ ? How many different  $a$  take each value?

*Proof.* Let  $\langle \alpha \rangle = (\mathbb{F}_q)^*$ . Let  $k \in \mathbb{Z}$ . If  $k$  is even, then  $(\alpha^k)^{(q-1)/2} = (\alpha^{k/2})^{q-1} = 1$ . If  $k = 2l+1$  for some  $l$ , then  $(\alpha^k)^{(q-1)/2} = \alpha^{l(q-1)} \cdot \alpha^{(q-1)/2} = \alpha^{(q-1)/2} = -1$  because  $-1$  has degree 2 and  $\alpha^{(q-1)/2}$  is the only element in  $\langle \alpha \rangle$  of degree 2. Therefore,

$$a^{(q-1)/2} = \begin{cases} 0 & (a = 0) \\ 1 & (\exists l \in \mathbb{Z}, a = \alpha^{2l}) \\ -1 & (\exists l \in \mathbb{Z}, a = \alpha^{2l+1}). \end{cases}$$

This is well defined because every nonzero element in  $\mathbb{F}_q$  is in  $\langle \alpha \rangle$  and  $2 \mid |\langle \alpha \rangle| = q - 1$ , so the parity of the exponent does not depend on the choice of  $k$ . Hence, 1 value gives 0,  $(q-1)/2$  values give 1, and  $(q-1)/2$  values give  $-1$ .  $\square$

**Exercise.** (Problem 15) Let  $f(x)$  be as in problem 13 and let  $h \in \mathbb{F}_q[x]$  be a randomly chosen polynomial. What is the probability that  $h^{(q^r-1)/2} = \pm 1$  in the ring  $\mathbb{F}_q[x]/(f(x))$ .

*Proof.* As shown in Problem 13 last week, there exists an isomorphism  $\Phi : \mathbb{F}_q[x]/(f(x)) \rightarrow \mathbb{F}_q[x]/(f_1(x)) \times \cdots \times \mathbb{F}_q[x]/(f_m(x))$  by the Chinese Remainder Theorem. For any  $h \in \mathbb{F}_q[x]$ ,  $\Phi(h + (f)) = (h + (f_1), \dots, h + (f_m))$ . Moreover,  $\Phi(h^{(q-1)/2} + (f)) = (h^{(q-1)/2} + (f_1), \dots, h^{(q-1)/2} + (f_m))$ . Therefore,  $h^{(q-1)/2} + (f) = 1$  if and only if  $h^{(q-1)/2} + (f_1), \dots, h^{(q-1)/2} + (f_m)$  all equal 1.

Let  $\alpha_1, \dots, \alpha_m$  be generators of  $(\mathbb{F}_q[x]/(f_1(x)))^*, \dots, (\mathbb{F}_q[x]/(f_m(x)))^*$ . For each  $i$ ,  $h^{(q-1)/2} + (f_i) = 1$  if and only if  $h \in \langle \alpha_i^2 \rangle$  by Problem 14. Therefore,  $h^{(q-1)/2} + (f) = 1$  if and only if  $(h + (f_1), \dots, h + (f_m)) \in \langle \alpha_1^2 \rangle \times \cdots \times \langle \alpha_m^2 \rangle$ . There are exactly  $((q^r - 1)/2)^m$  elements that satisfy that. Therefore,

$$\frac{\left(\frac{q^r-1}{2}\right)^m}{(q^r)^m} = \left(\frac{q^r-1}{2q^r}\right)^m = \left(\frac{1}{2} - \frac{1}{2q^r}\right)^m.$$

is the probability that  $h^{(q^r-1)/2} = 1$  in  $\mathbb{F}_q[x]/(f(x))$ .

Using the exact same argument, we can derive that the probability that  $h^{(q^r-1)/2} = -1$  is exactly the same value.  $\square$

**Exercise.** (Problem 16) With  $f(x)$  as in problem 13, write  $f(x) = g_1(x) \cdots g_m(x)$  for the factorization into irreducible factors. Express  $\gcd(f(x), h^{(q^r-1)/2} - 1)$  in terms of the  $g_i(x)$ 's.

*Proof.*  $\gcd(f(x), h^{(q^r-1)/2} - 1)$  is the product of  $g_i(x)$ 's that divide  $h^{(q^r-1)/2} - 1$ . It is divisible by  $g_i(x)$  if and only if  $h \in \langle \alpha_i^2 \rangle$  from Problem 15.  $\square$

**Exercise.** (Problem 17) Describe a probabilistic factoring algorithm which has a very high probability of finding the irreducible factors of a polynomial  $f(x) \in \mathbb{F}_q[x]$ , provided one knows ahead of time that  $f(x)$  is a product of  $m$  distinct irreducible polynomials of degree  $r$ .

*Proof.* Let  $i_0$  be fixed. Given a random  $h(x) \in \mathbb{F}_q[x]$ , the probability that  $h^{(q-1)/2} - 1 \in (f_{i_0})$  is  $1/2 - 1/(2q^r)$ , which is slightly smaller than 50%. Therefore, it is likely that given a random  $h(x) \in \mathbb{F}_q[x]$ , the probability that  $h^{(q-1)/2} - 1 \in (f_i)$  for *some*  $i$ 's is high. However, the probability that  $h^{(q-1)/2} - 1 \in (f_i)$  in *all*  $i$ 's is low.

In other words, the probability that  $h^{(q-1)/2} - 1$  is a proper divisor of  $f$  is high. Therefore, we can expect to factor  $f(x)$  by

- Step 1: Generate a random polynomial  $h(x) \in \mathbb{F}_q[x]/(f(x))$ .
- Step 2: Calculate  $h^{(q^r-1)/2} - 1$ . This step can be done efficiently by exponentiation by squaring.
- Step 3: Calculate  $d(x) = \gcd(f(x), h^{(q^r-1)/2} - 1)$ . This step can be done efficiently by the Euclid algorithm.
- Step 4: If  $1 \leq \deg(d(x)) \leq \deg(f(x))$ , then factorize  $f(x)/d(x)$  and  $d(x)$  further by going back to Step 1 unless it is degree  $r$ . Otherwise, we were unlucky, so we go back to Step 1.

$\square$

**Exercise.** (Problem 18, 19, 20)

- Problem 18:  $(x^2 + x - 1)^4$
- Problem 19:
- Problem 20:  $(x^4 + 4x^2 + 5x + 3)(x^4 + 15x^3 - 16x^2 - 27x - 26)(x^4 - 3x^3 + 9x^2 - 23x + 1)$ .

## 2. GALOIS THEORY III

**Exercise.** (Problem 1) Prove Proposition 23 part (ii).

*Proof.* Clearly,  $F \subset gK \subset L$  because  $g \in \text{Aut}(L/F)$ .  $gK$  is a subfield because  $g$  preserves addition, multiplication and multiplicative inverse, so  $gK$  is closed under addition, multiplication and multiplicative inverse.

Let  $\phi \in \text{Aut}(L/gK)$ . Then clearly,  $g^{-1}\phi g \in \text{Aut}(L)$ .  $g^{-1}\phi g$  fixes  $K$  because  $\forall x \in K, (g^{-1}\phi g)(x) = g^{-1}(g(x)) = x$ . Therefore,  $\phi \in g \text{Aut}(L/K) g^{-1}$ .

Let  $g\psi g^{-1} \in g \text{Aut}(L/K) g^{-1}$ . Then  $g\psi g^{-1} \in \text{Aut}(L)$ . For all  $g(k) \in g(K)$ ,  $(g\psi g^{-1})(g(k)) = g(\psi(g^{-1}(g(k)))) = g(\psi(k)) = g(k)$ . Therefore,  $g\psi g^{-1} \in \text{Aut}(L/gK)$ .  $\square$

**Exercise.** (Problem 2) Show that the Galois correspondence is order reversing.

*Proof.* Let  $H_1 \subset H_2$  be given. Let  $x \in K^{H_2}$ . Then  $x$  is fixed by every element in  $H_2$ . Then  $x$  is clearly fixed by every element in  $H_1$ . Thus  $x \in K^{H_1}$ .

Let  $K_1 \subset K_2$ . Let  $\sigma \in \text{Aut}(L/K_2)$ . Then  $\sigma$  clearly fixes  $K_1$ . Thus  $\sigma \in \text{Aut}(L/K_1)$ .  $\square$