

MATH 601 (DUE 11/13)

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1. FACTORING POLYNOMIALS WITH COEFFICIENTS IN FINITE FIELDS

Exercise. (Problem 14) For $a \in \mathbb{F}_q$, what are the possible values for $a^{(q-1)/2}$? How many different a take each value?

Proof. Let $\langle \alpha \rangle = (\mathbb{F}_q)^*$. Let $k \in \mathbb{Z}$. If k is even, then $(\alpha^k)^{(q-1)/2} = (\alpha^{k/2})^{q-1} = 1$. If $k = 2l+1$ for some l , then $(\alpha^k)^{(q-1)/2} = \alpha^{l(q-1)} \cdot \alpha^{(q-1)/2} = \alpha^{(q-1)/2} = -1$ because -1 has degree 2 and $\alpha^{(q-1)/2}$ is the only element in $\langle \alpha \rangle$ of degree 2. Therefore,

$$a^{(q-1)/2} = \begin{cases} 0 & (a = 0) \\ 1 & (\exists l \in \mathbb{Z}, a = \alpha^{2l}) \\ -1 & (\exists l \in \mathbb{Z}, a = \alpha^{2l+1}). \end{cases}$$

This is well defined because every nonzero element in \mathbb{F}_q is in $\langle \alpha \rangle$ and $2 \mid |\langle \alpha \rangle| = q - 1$, so the parity of the exponent does not depend on the choice of k . Hence, 1 value gives 0, $(q-1)/2$ values give 1, and $(q-1)/2$ values give -1 . \square

Exercise. (Problem 15) Let $f(x)$ be as in problem 13 and let $h \in \mathbb{F}_q[x]$ be a randomly chosen polynomial. What is the probability that $h^{(q^r-1)/2} = \pm 1$ in the ring $\mathbb{F}_q[x]/(f(x))$.

Proof. As shown in Problem 13 last week, there exists an isomorphism $\Phi : \mathbb{F}_q[x]/(f(x)) \rightarrow \mathbb{F}_q[x]/(f_1(x)) \times \cdots \times \mathbb{F}_q[x]/(f_m(x))$ by the Chinese Remainder Theorem. For any $h \in \mathbb{F}_q[x]$, $\Phi(h + (f)) = (h + (f_1), \dots, h + (f_m))$. Moreover, $\Phi(h^{(q-1)/2} + (f)) = (h^{(q-1)/2} + (f_1), \dots, h^{(q-1)/2} + (f_m))$. Therefore, $h^{(q-1)/2} + (f) = 1$ if and only if $h^{(q-1)/2} + (f_1), \dots, h^{(q-1)/2} + (f_m)$ all equal 1.

Let $\alpha_1, \dots, \alpha_m$ be generators of $(\mathbb{F}_q[x]/(f_1(x)))^*, \dots, (\mathbb{F}_q[x]/(f_m(x)))^*$. For each i , $h^{(q-1)/2} + (f_i) = 1$ if and only if $h \in \langle \alpha_i^2 \rangle$ by Problem 14. Therefore, $h^{(q-1)/2} + (f) = 1$ if and only if $(h + (f_1), \dots, h + (f_m)) \in \langle \alpha_1^2 \rangle \times \cdots \times \langle \alpha_m^2 \rangle$. There are exactly $((q^r - 1)/2)^m$ elements that satisfy that. Therefore,

$$\frac{\left(\frac{q^r-1}{2}\right)^m}{(q^r)^m} = \left(\frac{q^r-1}{2q^r}\right)^m = \left(\frac{1}{2} - \frac{1}{2q^r}\right)^m.$$

is the probability that $h^{(q^r-1)/2} = 1$ in $\mathbb{F}_q[x]/(f(x))$.

Using the exact same argument, we can derive that the probability that $h^{(q^r-1)/2} = -1$ is exactly the same value. \square

Exercise. (Problem 16) With $f(x)$ as in problem 13, write $f(x) = g_1(x) \cdots g_m(x)$ for the factorization into irreducible factors. Express $\gcd(f(x), h^{(q^r-1)/2} - 1)$ in terms of the $g_i(x)$'s.

Proof. $\gcd(f(x), h^{(q^r-1)/2} - 1)$ is the product of $g_i(x)$'s that divide $h^{(q^r-1)/2} - 1$. It is divisible by $g_i(x)$ if and only if $h \in \langle \alpha_i^2 \rangle$ from Problem 15. \square

Exercise. (Problem 17) Describe a probabilistic factoring algorithm which has a very high probability of finding the irreducible factors of a polynomial $f(x) \in \mathbb{F}_q[x]$, provided one knows ahead of time that $f(x)$ is a product of m distinct irreducible polynomials of degree r .

Proof. Let i_0 be fixed. Given a random $h(x) \in \mathbb{F}_q[x]$, the probability that $h^{(q-1)/2} - 1 \in (f_{i_0})$ is $1/2 - 1/(2q^r)$, which is slightly smaller than 50%. Therefore, it is likely that given a random $h(x) \in \mathbb{F}_q[x]$, the probability that $h^{(q-1)/2} - 1 \in (f_i)$ for *some* i 's is high. However, the probability that $h^{(q-1)/2} - 1 \in (f_i)$ in *all* i 's is low.

In other words, the probability that $h^{(q-1)/2} - 1$ is a proper divisor of f is high. Therefore, we can expect to factor $f(x)$ by

- Step 1: Generate a random polynomial $h(x) \in \mathbb{F}_q[x]/(f(x))$.
- Step 2: Calculate $h^{(q^r-1)/2} - 1$. This step can be done efficiently by exponentiation by squaring.
- Step 3: Calculate $d(x) = \gcd(f(x), h^{(q^r-1)/2} - 1)$. This step can be done efficiently by the Euclid algorithm.
- Step 4: If $1 \leq \deg(d(x)) < \deg(f(x))$, then factorize $f(x)/d(x)$ and $d(x)$ further by going back to Step 1 unless it is degree r . Otherwise, we were unlucky, so we go back to Step 1.

\square

Exercise. (Problem 18, 19, 20)

- Problem 18: $(x^2 + x - 1)^4$
- Problem 19: $(x^3 - 25x^2 - 35x + 3)(x^4 + 4x^2 + 5x + 3)(x^5 + 4x^2 + 8x + 3)$.
- Problem 20: $(x^4 + 4x^2 + 5x + 3)(x^4 + 15x^3 - 16x^2 - 27x - 26)(x^4 - 3x^3 + 9x^2 - 23x + 1)$.

I used the following Python code to factorize. The idea is to use the methods developed in Problem 11 and Problem 17. Later, I noticed that I should have added code to check if $f(x)$ is square free, but for some reason, the code was still able to factorize the polynomial for Problem 18.

```
from sympy import *
from random import *

x = symbols('x')

# Find a random polynomial of degree <= deg in Z_{mod}.
def randpoly(deg, mod):
    p = poly(0, x, modulus = mod)
```

```

    for d in range(deg):
        p = x * p + randint(0, mod - 1)
    return poly(p, x, modulus = mod)

# Find  $f^{\text{exp}} \bmod f$  in  $\mathbb{Z}_{\text{mod}}$ .
def polypow(f, exp, modf, mod):
    res = poly(1, x, modulus = mod)
    while exp > 0:
        if exp % 2 == 1:
            quotient, res = div(res * f, modf, modulus = mod)
            quotient, f = div(f * f, modf, modulus = mod)
            exp = exp // 2
    return res

# Calculate  $x^{(p^n)} - x \bmod f$ .
def xqd(p, n, modf):
    res = polypow(x, p**n, modf, p)
    res -= poly(x, x, modulus = p)
    return res

def factor(f, p, originaldegree, factors):
    # Problem 11
    for n in range(2, originaldegree):
        g = xqd(p, n, f)
        d = gcd(f, g)
        if 1 <= d.degree() < f.degree():
            # We found a proper factor.
            # Factorize further.
            factor(d, p, originaldegree, factors)
            quotient, remainder = div(f, d, modulus = p)
            factor(quotient, p, originaldegree, factors)
    return

# Problem 17
for r in range(2, f.degree()):
    if f.degree() % r != 0: continue
    for i in range(10):
        h = randpoly(r, p)
        # Raise h to the power of  $(p^r - 1)/2$ .
        h = polypow(h, (p**r - 1) // 2, f, p)
        h = h - poly(1, x, modulus = p)

```

```

        d = gcd(f, h)
        if d.degree() == 0 or d.degree() == f.degree():
            continue
        else:
            # We found a proper factor.
            # Factorize further.
            factor(d, p, originaldegree, factors)
            quotient, remainder = div(f, d)
            factor(quotient, p, originaldegree, factors)
        return
    factors.append(f)

def factorizepoly(f, mod):
    print("Factorize %s" % f)
    factors = []
    factor(f, mod, f.degree(), factors)
    prod = poly(1, x, modulus = mod)
    for fac in factors:
        prod *= fac
        print(latex(fac))
    if prod != f:
        print("*****ERROR!*****")
    print()
    return

f = poly(x**8 + x**7 - x**6 + x**5 + x**4 - x**3 - x**2 - x + 1, x, modulus = 3)
factorizepoly(f, 3)

f = poly((x**12+48*x**11+42*x**10+58*x**9+11*x**8+25*x**7+22*x**6+30*x**5+15*x**4+6*x**3+3*x**2+2*x+1), x, modulus = 73)
factorizepoly(f, 73)

f = poly((x**12+12*x**11 + 25*x**10 + 40*x**9 + 6*x**8 + 15*x**7 + 24*x**6 + 15*x**5 + 6*x**4 + 3*x**3 + 3*x**2 + 2*x + 1), x, modulus = 73)
factorizepoly(f, 73)

```

2. GALOIS THEORY III

Exercise. (Problem 1) Prove Proposition 23 part (ii).

Proof. Clearly, $F \subset gK \subset L$ because $g \in \text{Aut}(L/F)$. gK is a subfield because g preserves addition, multiplication and multiplicative inverse, so gK is closed under addition, multiplication and multiplicative inverse.

Let $\phi \in \text{Aut}(L/gK)$. Then clearly, $g^{-1}\phi g \in \text{Aut}(L)$. $g^{-1}\phi g$ fixes K because $\forall x \in K, (g^{-1}\phi g)(x) = g^{-1}(g(x)) = x$. Therefore, $\phi \in g \text{Aut}(L/K)g^{-1}$.

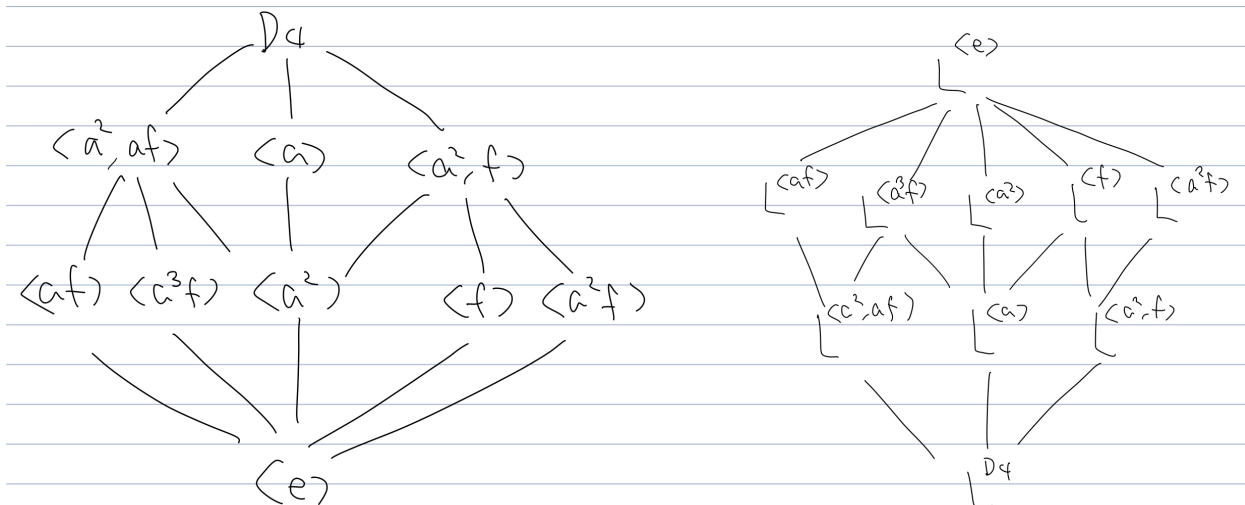


FIGURE 1. Problem 3

Let $g\psi g^{-1} \in g \operatorname{Aut}(L/K)g^{-1}$. Then $g\psi g^{-1} \in \operatorname{Aut}(L)$. For all $g(k) \in g(K)$, $(g\psi g^{-1})(g(k)) = g(\psi(k)) = g(k)$. Therefore, $g\psi g^{-1} \in \operatorname{Aut}(L/gK)$. \square

Exercise. (Problem 2) Show that the Galois correspondence is order reversing.

Proof. Let $H_1 \subset H_2$ be given. Let $x \in K^{H_2}$. Then x is fixed by every element in H_2 . Then x is clearly fixed by every element in H_1 . Thus $x \in K^{H_1}$.

Let $K_1 \subset K_2$. Let $\sigma \in \operatorname{Aut}(L/K_2)$. Then σ clearly fixes K_1 . Thus $\sigma \in \operatorname{Aut}(L/K_1)$. \square

Exercise. (Problem 3) Draw a picture showing all the subgroups of the dihedral group with eight elements, $D_4 := \langle a, f : a^4 = 1 = f^2, f a f = a^{-1} \rangle \simeq \langle (1234), (12)(34) \rangle \subset S_4$ showing which are contained in which. Now draw a diagram of the corresponding intermediate fields in a Galois extension, $F \subset L$, with Galois group isomorphic to D_4 indicating which are contained in which.

Proof. Figure 1. \square

Exercise. (Problem 4) Let $F \subset M$ be a Galois extension with Galois group isomorphic to the dihedral group with eight elements (denoted D_4 in class). Show that there is a tower of intermediate fields, $F \subset K \subset L$ such that $F \subset K$ is Galois and $K \subset L$ is Galois, but $F \subset L$ is not Galois.

Proof. $G_1 = \langle a f \rangle$ is a normal subgroup of $G_2 = \{e, a f, a^2, a^3 f\}$ because the index is 2. Similarly, G_2 is a normal subgroup of D_4 because the index is 2. However, G_1 is not a normal subgroup of D_4 . (For instance, $f \langle a f \rangle f^{-1} = \langle f a \rangle$, but $a f \neq f a$.) By the Fundamental Theorem of Galois Theory, L^{G_1} and L^{G_2} are intermediate fields. By Proposition 23(iii), $L^{G_2} \subset L^{G_1}$ and $L^{D_4} \subset L^{G_2}$ is Galois, but $L^{D_4} \subset L^{G_1}$ is not Galois. \square

Exercise. (Problem 5) Let $F \subset M$ be a Galois extension with Galois group isomorphic to the symmetric group S_4 . Let $H = \langle (123) \rangle \subset S_4$. Make a list of the intermediate fields in the extension, $F \subset M^H$. For each intermediate field L indicate whether or not $F \subset L$ is Galois and whether or not $L \subset M^H$ is Galois.

Proof. There are only 4 subgroups of S_4 that contain S_3 . They are H, S_3, A_4, S_4 .

Clearly, $M^H \subset M^H$ is Galois. H is not a normal subgroup of S_4 because $(14)(12)(14) \notin H$. Therefore, $F \subset M^H \subset$ is not Galois.

$S_3 = \{e, (12), (13), (23), (123), (132)\}$ is a proper subgroup of S_4 that contains H properly. Therefore, $F \subsetneq M^{S_3} \subsetneq M$. Since $[S_3 : H] = 2$, H is a normal subgroup of S_3 . Therefore, $M^{S_3} \subset M^H$ is Galois. S_3 is not a normal subgroup of S_4 because $(14)(12)(14) \notin S_3$. Therefore, $F \subset M^{S_3}$ is not Galois. \square