

Chapter 6: Information channel

Presented notions

- Physical medium to send information from transmitter to receiver
- Channel has additive noise which changes the propagated signal
- Channel has different types:
 - Discrete channel
 - Memoryless channel
 - Memory channel
 - Binary channel
 - Continuous channel
- Input source, noise source, output are considered random variables
 - $\text{Output} = \text{Input} + \text{Noise}$

6.1. Channel model

- Discrete channel: discrete input signal and discrete output signal
 - Channel model:
 - Random variable represents discrete source
 - $X = \{x_1, x_2, \dots, x_n\}$
 - $P(X) = \{P(x_1), P(x_2), \dots, P(x_n)\}$
 - Alphabet of output
 - $Y = \{y_1, y_2, \dots, y_n\}$
 - $Y_i = x_j$ (alphabet of output is identical to alphabet of input)
 - Generally, number of inputs (r) is different with number of outputs (s).
 - This course only focuses: $r = s = n$

6.1. Channel Model (Cont.)

- Discrete channel:

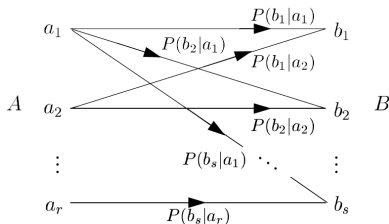
- This course only focus on memoryless channel
- Matrix channel for memoryless channel: consists of all transfer probabilities
 - $P(Y|X) = \{P(y_j|x_i)\}$

$$\begin{bmatrix} P(y_1|x_1) & P(y_2|x_1) & \dots & P(y_n|x_1) \\ P(y_1|x_2) & P(y_2|x_2) & \dots & P(y_n|x_2) \\ \dots & \dots & \dots & \dots \\ P(y_1|x_n) & P(y_2|x_n) & \dots & P(y_n|x_n) \end{bmatrix}$$

- Elements in main diagonal is correct probability
- Elements outside off main diagonal is wrong probability
- If all correct probabilities are equal, all wrong probabilities are equal: uniform channel
- Sum of all elements in one row = 1
- Identity matrix (unit matrix): transmission in channel not wrong
- If all elements are equal, the channel wrong: output independents with input
- Elements of the matrix are symmetric across the main diagonal (symmetric matrix): symmetric channel

6.1. Channel Model (Cont.)

- Transition diagram of the discrete channel:
 - r input points. Each point represents a input signal
 - s output points. Each point represents a output signal
 - Connection line between a input point and a output point represent transfer probability



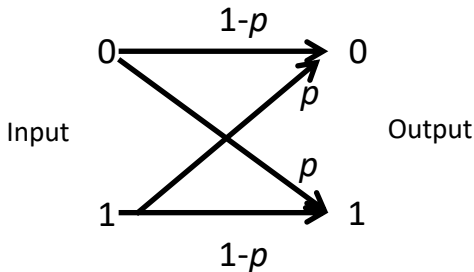
6.1. Channel Model (Cont.)

- Discrete Binary Symmetric Channel (BSC)
 - Binary channel: input signal and output signal are binary
 - For example: 0 and 1
 - Symmetric channel: The channel matrix is symmetric across the main diagonal
 - Channel matrix: contain transmission probability $P(y|x)$
 - $P(y|x)$: conditional probability of receiving output signal y when input signal x is sent
- Binary symmetric channel = binary channel + symmetric channel
 - Example: channel matrix of one BSC

$$P(Y|X) = \begin{array}{cc|c} & & & \\ & 3/4 & 1/4 & x1 \\ & 1/4 & 3/4 & x2 \\ y1 & & y2 & \end{array}$$

6.1. Channel Model (Cont.)

Binary symmetric channel (BSC) model



$$P[Y = 0|X = 1] = P[Y = 1|X = 0] = p$$

$$P[Y = 1|X = 1] = P[Y = 0|X = 0] = 1 - p$$

6.1. Channel Model (Cont.)

- Continuous channel: continuous input signal and continuous output signal
 - Channel model:
 - Random variable represents continuous source
 - $X = \{x\}$ $x_{\min} \leq x \leq x_{\max}$
 - $P_X(x)$ probability density function of source X
 - Output:
 - $Y = \{y\}$ $y_{\min} \leq y \leq y_{\max}$
 - $P(y|x)$ conditional probability density function of output Y when input X has determined values

6.2. Mutual information of discrete channel (Cont.)

- According chapter 2: mutual information between two random variables X, Y:

$$\begin{aligned} I(X; Y) &= \sum \sum P(x,y) \log \frac{P(x,y)}{P(x)P(y)} \\ &= H(X) + H(Y) - H(X,Y) \end{aligned}$$

- Because: $\frac{P(x,y)}{P(y)} = P(x|y)$

$$\begin{aligned} \text{So} \quad : I(X; Y) &= \sum \sum P(x,y) \log \frac{P(x|y)}{P(x)} \\ &= H(X) - H(X|Y) \end{aligned}$$

- Corollary:

- Mutual information is a average information (entropy) that a symbol may transmit through channel
- Mutual information equals average input information ($H(X)$) minus average loss information ($H(X|Y)$).
- Average loss information is generated by noise: $H(X|Y) = H(N)_X$ with noise source N and source X are inputs

6.2. Mutual information of discrete channel (Cont.)

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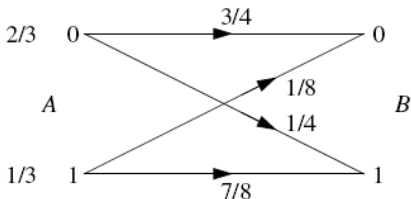
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6.2. Mutual information of discrete channel (Cont.)

- Information transmitted through channel always equals input information minus loss information
- Loss information: noise information generated by noise
- Information transmitted through channel in both ways are equal
- $0 \leq I(X;Y) \leq H(X)$ when no noise, $X = Y$

6.2. Mutual information of discrete channel (Cont.)

- $H(A)$, $H(A|B)$?
- Mutual information?



6.3. Mutual information of continuous channel

- As presented in chapter 2, entropy of continuous source are calculated using integral and probability density function

- So, mutual information is also calculated using integral

$$\begin{aligned} I(X; Y) &= \iint_{x,y} P_{x,y}(x,y) \log \frac{P_{x,y}(x,y)}{P_x(x)P_y(y)} dx dy \\ &= H(X) + H(Y) - H(X,Y) \\ &= \iint_{x,y} P_{x,y}(x,y) \log \frac{P_{x|y}(x|y)}{P_x(x)} dx dy \\ &= H(X) - H(X|Y) \\ &= \iint_{x,y} P_{x,y}(x,y) \log \frac{P_{y|x}(y|x)}{P_y(y)} dx dy \\ &= H(Y) - H(Y|X) \end{aligned}$$

- Properties of mutual information in continuous channel are similar to the properties of mutual information in discrete channel

6.4. Channel capacity

- Maximum average amount of information that channel may transmit in a unit of time without error
- Denoted by C
- Calculated by number of transmitted symbols n_o multiply maximum average amount of information that a symbol can transmit through channel (mutual information)
 - $C = n_o \times I(X;Y)_{\max}$
- Remark: in information theory, n_o is physical parameter so that n_o is considered as an unit
 - $C = I(X;Y)_{\max}$

6.4. Channel capacity (Cont.)

- Discrete channel:

- Number of transmitted symbols (n_o) is bandwidth of channel (Δf)
 - So, $C = \Delta f \times I(X;Y)_{\max}$
- $C = \Delta f \times (H(X) - H(X|Y))_{\max} = \Delta f \times (H(X) - H(N)_X)_{\max}$
- Without noise: $H(N)_X = 0$
 - So: $C = \Delta f \times H(X)_{\max} = \Delta f \times \log |X| = \Delta f \times \log L$ (L is number of different symbols of source X)

6.4. Channel capacity (Cont.)

- Continuous channel:
 - n_o is calculated as the number of samples of the discrete source equivalent to the continuous source
 - According to sampling theorem:
 - $n_o = 2 F_{\max}$ (F_{\max} is maximum frequency of source connected to channel)
 - $F_{\max} = \Delta f$ (bandwidth of channel)
 - $C = 2 \Delta f \times I(X;Y)_{\max}$
 $= 2 \Delta f \times (H(Y) - H(Y|X))_{\max}$

6.4. Channel capacity (Cont.)

- Continuous channel:

- Normally, source has Gaussian distribution and noise has also Gaussian distribution:

- $H(Y|X) = H(N)_Y = \log \sqrt{2\pi e P_N}$

- P_N : average power (variance) of the noise source

- $H(Y) = \log \sqrt{2\pi e P_Y}$

- P_Y : average power (variance) of the output

- $P_Y = P_X + P_N$

- P_X : average power (variance) of input source

- The variance of the sum of the Gaussian random variables is equal to the sum of the variances of each variable

- $C = 2 \Delta f \times (\log \sqrt{2\pi e P_Y} - \log \sqrt{2\pi e P_N})$

$$= 2\Delta f \times \log \sqrt{\frac{P_Y}{P_N}}$$

$$= \Delta f \times \log \frac{P_Y}{P_N} = \Delta f \times \log \frac{P_X + P_N}{P_N} = \Delta f \times \log \left(1 + \frac{P_X}{P_N}\right)$$

$$\rightarrow C = \Delta f \times \log \left(1 + \frac{S}{N}\right)$$

$\frac{S}{N}$: signal/noise rate: calculated by the ratio of the average power of the input signal with the average of the noise measured at the output of the channel