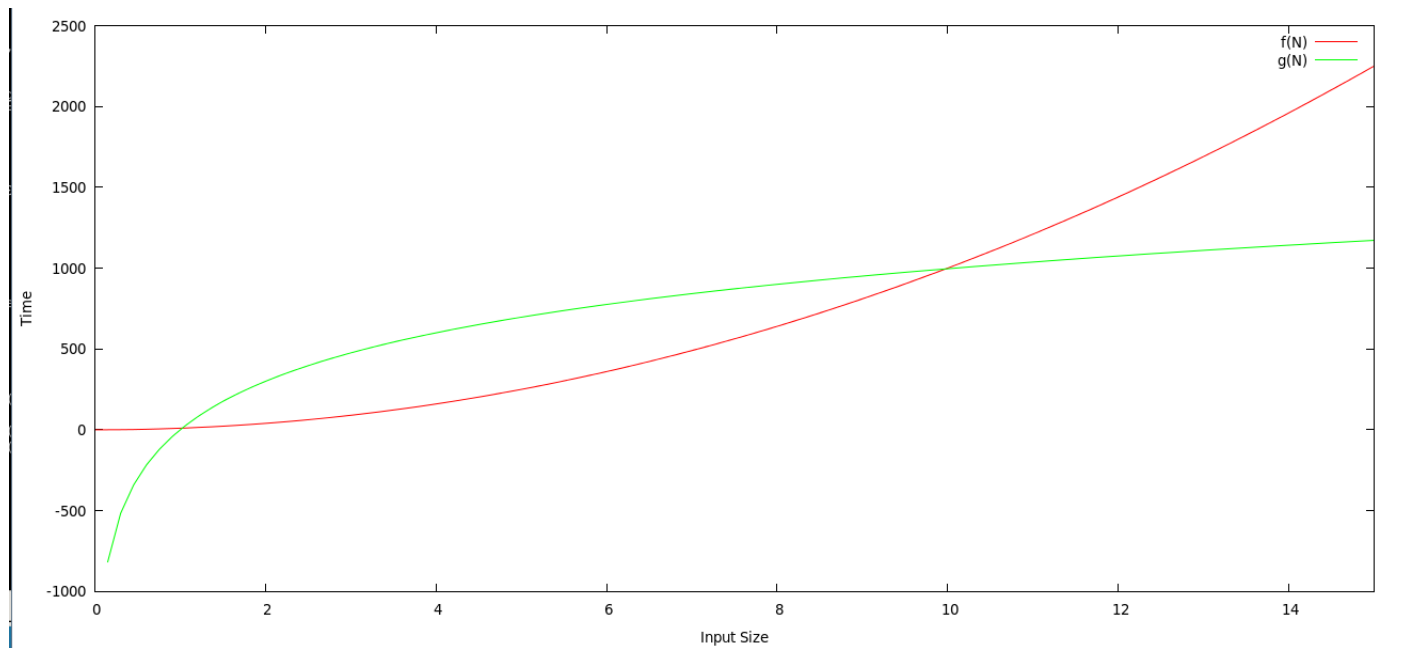


CS 310
Assignment 205

Hieu Le

Problem 1. Algorithm A performs $10n^2$ basic operations, and algorithm B performs $300\lg(n)$ basic operations. Which algorithm is better, and at what value of n does the better algorithm start to show its better performance? Illustrate your answer with a graph generated by a program such as gnuplot or wolfram alpha.

Answer: As depicted in the following graph:



Algorithm B is better than algorithm A. Algorithm B begins to show superior performance at $n = 10$.

Problem 2. Using the definitions, show that

$$6n^2 + 20n \in O(n^3)$$

but

$$6n^2 + 20n \notin \Omega(n^3)$$

Answer:

Let $c = 30$ and $n_0 = 1$. Since $6n^2 + 20n \leq cn^3$ for all $n \geq n_0$, $6n^2 + 20n \in O(n^3)$.

Since $6n^2 + 20n$ is of smaller degree than n^3 , there does not exist any pair of positive numbers c and n_0 such that $6n^2 + 20n \geq cn^3$ for all $n \geq n_0$. Hence, $6n^2 + 20n \notin \Omega(n^3)$.

Problem 3. Given the following algorithm, and assuming that n is an even number, calculate the exact number of times statement `foo` runs, and analyze the algorithm using the the rules (e.g., polynomial).

```

j = 1;
while( j <= n/2 )
{
    i = 1;
    while( i <= j )
    {
        foo;
        i++;
    }
    j++;
}

```

Answer:

The number of times statement `foo` runs is:

$$1 + 2 + 3 + \dots + n/2 = \sum_{k=1}^{n/2} k = (1 + n/2)(n/2)/2 = (2n + n^2)/8.$$

Since $(2n + n^2)/8$ is a polynomial of degree 2, $(2n + n^2)/8 \in \Theta(n^2)$.

Problem 4. Solve the recurrence relation $t_n = 2nt_{n-1}$ where $t_0 = 1$.

Answer: Using the given relation, we have:

$$\begin{aligned}
 t_n &= 2nt_{n-1} \\
 &= 2n(2(n-1)t_{n-2}) \\
 &= 2^2n(n-1)t_{n-2} \\
 &= 2^2n(n-1)(2(n-2)t_{n-3}) \\
 &= 2^3n(n-1)(n-2)t_{n-3} \\
 &\vdots \\
 &= 2^n n(n-1) \dots (2)(1)t_0
 \end{aligned}$$

and therefore $t_n = 2^n n!$ for $n \geq 1$.

Problem 5. Analyze a program whose time complexity is $T(n) = 7T(\frac{n}{4}) + n$.

Answer:

Applying the Master Theorem, we have: $a = 7, b = 4, d = 1$.

Since $a > b^d$, $T(n) \in \Theta(n^{\log_4 7}) \approx \Theta(n^{1.404})$.