

## PROBLEMS - SET 8

**Problem 1.** Define a time reversible Markov Chain on the positive integers  $\{1, 2, \dots\}$  with stationary distribution  $\pi(i) \propto i^{-5/4}$  for  $i \in \{1, 2, \dots\}$ .

**Solution 1.** Let  $P_{i,j}$  denote the transition probabilities of a random walk on the positive integers. We assume that

$$\begin{aligned} P_{i,i+1} &= p, \quad P_{i+1,i} = 1 - p \quad \text{for any } i \geq 1 \\ P_{1,1} &= 1 - p \end{aligned}$$

with  $p \in (0, 1)$ .

For  $i \neq j$ , define  $\alpha(i, j) = 0$  if  $j \neq i+1, i-1$ ,

$$\alpha(i, i+1) = \min \left\{ \frac{\pi_{i+1} P_{i+1,i}}{\pi_i P_{i,i+1}}, 1 \right\} = \min \left\{ \frac{1-p}{p} \frac{\pi_{i+1}}{\pi_i}, 1 \right\} = \min \left\{ \frac{(i+1)^{-5/4} (1-p)}{i^{-5/4} p}, 1 \right\}$$

$$\alpha(i, i-1) = \min \left\{ \frac{\pi_{i-1} P_{i-1,i}}{\pi_i P_{i,i-1}}, 1 \right\} = \min \left\{ \frac{p}{1-p} \frac{\pi_{i-1}}{\pi_i}, 1 \right\} = \min \left\{ \frac{(i-1)^{-5/4} p}{i^{-5/4} (1-p)}, 1 \right\}$$

Note that

$$\pi_i P_{i,i+1} \alpha(i, i+1) = \pi_{i+1} P_{i+1,i} \alpha(i+1, i).$$

The transition probabilities for the desired Markov Chain are  $Q_{i,j} = P_{i,j} \alpha(i, j)$  for  $j \neq i$ , by construction indeed  $\pi_i P_{i,j} \alpha(i, j) = \pi_j P_{j,i} \alpha(j, i)$ . When  $i = 1$

$$Q_{1,j} = \begin{cases} p \alpha(1, 2) = p \min\{1, 2^{-5/4} \frac{1-p}{p}\} & \text{if } j = 2 \\ 1 - Q_{1,2} = 1 - p \min\{1, 2^{-5/4} \frac{1-p}{p}\} & \text{if } j = 1 \\ 0 & \text{if } j > 2 \end{cases}$$

When  $i > 1$

$$Q_{i,j} = \begin{cases} (1-p) \alpha(i, i-1) = (1-p) \min\{1, (\frac{i-1}{i})^{-5/4} \frac{p}{1-p}\} = \min\{1-p, p (\frac{i}{i-1})^{5/4}\} & \text{if } j = i-1 \\ p \alpha(i, i+1) = p \min\{1, (\frac{i+1}{i})^{-5/4} \frac{1-p}{p}\} & \text{if } j = i+1 \\ 1 - Q_{i,i-1} - Q_{i,i+1} & \text{if } j = i \\ 0 & \text{if } |i-j| > 1 \end{cases}$$

Note that by choosing  $p = 1/2$ , all the previous expressions can be simplified.

**Problem 2.** Let  $(V, E)$  be a graph. Define a reversible Markov Chain  $X_n$  with state space  $S = \{X, X \subset V\} = 2^V$  the set of subsets of  $V$ , such that its invariant measure is the measure in which the measure of a subset  $X \subset V$  is proportional to  $2^{E|X}$ , where  $E|X$  is the number of edges in  $E$  that have both endpoints in  $X$ .

**Solution 2.** First of all we have to fix a reference transition matrix. To do this, we define a rule to pass from one state  $X$  to another state  $X'$  of the chain. Note that the states of the chain are subsets of  $V$ .

We define the following rule to update a set  $X_n \subset V$ .

**Metropolis Hastings:**

- (i) Randomly select a vertex  $v \in V$ .
- (ii) If  $v \notin X_n$ , set  $X_{n+1} = X_n \cup \{v\}$ .
- (iii) If  $v \in X_n$  set  $X_{n+1} = X_n \setminus \{v\}$  with probability  $p$  and  $X_{n+1} = X_n$  with probability  $1 - p$ .

So, we get that the transition matrix is given by

$$P_{X,X'} = \begin{cases} 0 & \text{if } X, X' \text{ differ on more than 1 vertex} \\ \frac{1}{|V|} & \text{if } X' \setminus X = \{v\} \text{ for exactly one } v \\ \frac{1}{|V|}p & \text{if } X \setminus X' = \{v\} \text{ for exactly one } v \\ \text{generally incomputable} & \text{if } X = X'. \end{cases}$$

As usual, the diagonal elements of the transition matrix are not known, but the algorithm can be simulated without this knowledge. Note that the invariant distribution is given by  $\pi = (\pi_X)_{X \in 2^V}$ , defined as

$$\pi_X = c2^{E|X|}$$

for some normalization constant  $c > 0$ . So, in order for the chain to be reversible, we need that

$$\pi_X P_{X,X'} = \pi_{X'} P_{X',X}.$$

Obviously these equations are trivial if  $X, X'$  differ for more than one vertex. Assume now that  $X' \setminus X = \{v\}$ , then the previous equation reads

$$c2^{E|X|} \frac{1}{|V|} = c2^{E|X'|} \frac{p}{|V|}.$$

Note that

$$2^{E|X'|} = 2^{E|X \cup \{v\}|} = 2^{E|X|} \cdot 2^{|\text{neighbors of } v \text{ in } X|}$$

because  $E|X'| = E|X| + |\text{neighbors of } v \text{ in } X|$ . Therefore, substituting in the balance equation we get

$$p = 2^{-|\text{neighbors of } v \text{ in } X|}.$$

With this choice of  $p$ , we get that  $P_{X,X'}$  is the transition matrix of the required MC. This is just the Metropolis-Hastings algorithm.

**Problem 3.** Let  $X_1, X_2, \dots$  be a sequence of independent exponential random variables with parameter  $\lambda$  and  $N$  be a Geometric random variable of parameter  $\alpha$ . Assume that  $X_1, X_2, \dots$  and  $N$  are also independent. Prove that the random variable

$$T = \sum_{i=1}^N X_i$$

is an exponential distribution of parameter  $\alpha\lambda$ .

**Solution 3.** It is sufficient to show that the moment generating function of  $T$  coincides with the moment generating function of an exponential random variable of parameter  $\alpha\lambda$ .

We compute the moment generation function of  $T$ . We fix  $t \in \mathbb{R}$  and we define the function  $f : \mathbb{N} \rightarrow \mathbb{R}$  as  $f(n) = \mathbb{E}[e^{t(X_1 + \dots + X_n)}]$ . Observe that  $\varphi_T(t) = \mathbb{E}[e^{t(X_1 + \dots + X_N)}] = \mathbb{E}(f(N))$ . Therefore

$$\varphi_T(t) = \sum_{n=1}^{\infty} \mathbb{E}[e^{t(X_1 + \dots + X_n)}] \mathbb{P}[N = n].$$

Since the random variables  $X_i$  are iid, and the moment generation function of  $X_1$  is defined for  $t < \lambda$  as

$$\varphi_{X_1}(t) = \mathbb{E}[e^{tX_1}] = \frac{\lambda}{\lambda - t}$$

we get

$$\mathbb{E}[e^{t(X_1 + \dots + X_n)}] = \prod_{i=1}^n \varphi_{X_1}(t) = \left( \frac{\lambda}{\lambda - t} \right)^n.$$

Then

$$\begin{aligned} \varphi_T(t) &= \sum_{n=1}^{\infty} \left( \frac{\lambda}{\lambda - t} \right)^n \mathbb{P}[N = n] = \frac{\lambda\alpha}{\lambda - t} \sum_{n=1}^{\infty} \left( \frac{\lambda(1 - \alpha)}{\lambda - t} \right)^{n-1} \\ &= \frac{\lambda\alpha}{\lambda - t} \frac{1}{1 - \frac{\lambda(1 - \alpha)}{\lambda - t}} = \frac{\lambda\alpha}{\lambda\alpha - t} \end{aligned}$$

which is the moment generating function of an exponential random variables of parameter  $\alpha\lambda$ .