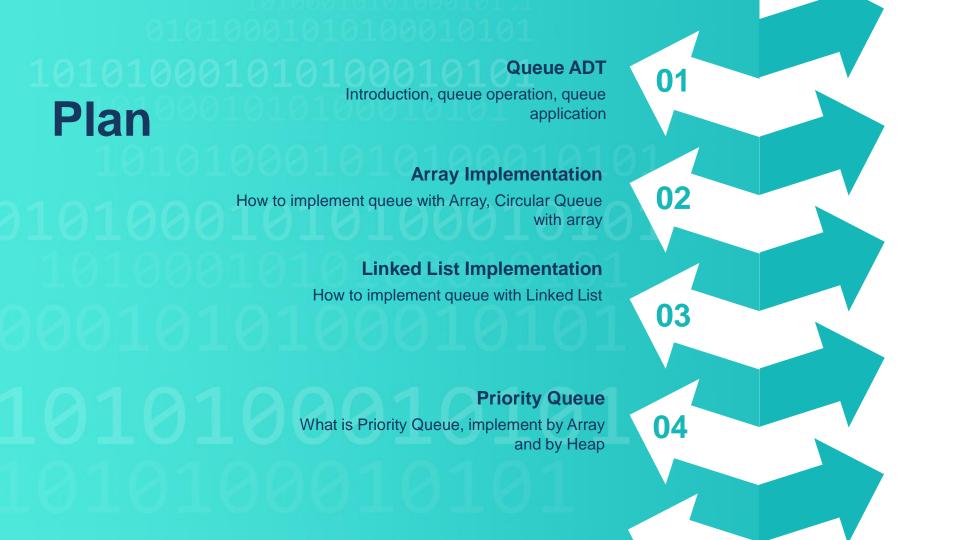
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## data\_tructures(&algorithms, lecture07)

Doan Trung Tung, PhD - University of Greenwich (Vietnam)



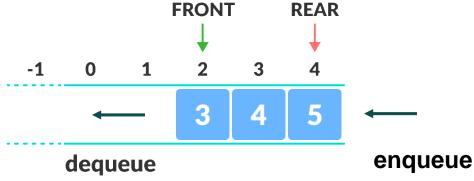
## 

## **Queue ADT**

Introduce Queue data structure, some operations

### What Is A Queue

- A linear data structure that can be accessed in both begins and ends for storing and retrieving data
- A queue is a First In, First Out (FIFO) data structure
  FRONT REAR



## **Operations On Queue**

- clear: clear the queue
- \* is empty: check if a queue is empty
- enqueue: put element on the rear of the queue
- dequeue: take the element on the top out of the queue
- front: get first element without removing it
- rear: get last element without removing it
- \* size: get size of the queue

## **Applications Of Queue**

- Any kind of waiting list
- Operating System:
  - CPU scheduling
  - Disk scheduling
  - Shared resources between processes
- Traverse in tree / graph data structures
- Turned-base games
- Alt-Tab in Windows
- •

## 

## **Array Implementation**

Implement Queue by Array, Circular Queue

## Queue Interface By Array

- Need global variables
- 2 indicators run forward
- Queue size is fixed
- Number of enqueue is limited

```
front rear

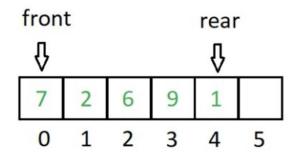
7 2 6 9 1

0 1 2 3 4 5
```

```
#define QUEUE_SIZE 100
extern int front;
extern int rear;
extern int queue[];
void clear_queue(void);
int is_empty(void);
int is_full(void);
void enqueue(const int n);
int dequeue(void);
int size(void);
void overflow(char* msg);
```

## Queue Implementation By Array

- is\_empty: when front = -1
- is\_full: when rear runs to the end of array
- enqueue: need to check full queue before, then increase rear and assign new item to new rear position
- dequeue: need to check empty queue before, then increase front and return the item at old front
- size: distance between front vs rear



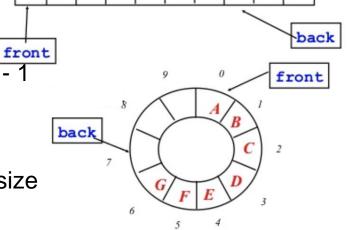


## Queue Implementation By Array

- Number of enqueue is fixed
  - ❖ => When dequeue, if queue has only one item, clear queue to reset front & rear to -1
  - ❖ ⇒ Implement circular queue
- Size is fixed
  - Using dynamic array and grow size if queue is full

## Circular Queue By Array

- Amain idea: front & rear run forward but go back (if possible) to 0 index.
  - is\_empty: front = -1
  - is\_full:
    - ❖ front = 0, rear = max size 1
    - front = rear + 1
  - Using modulo to go back
    - ❖ front = (front + 1) % max size
  - \* real size:
    - ❖ if (front < rear): rear front + 1</p>
    - else: max size (front rear 1)





## Example: Queue of messages

- Simulation of sending message. Commands:
  - +message: enqueue message
    - ❖ When the queue is full, all messages will be sent
  - -: dequeue a message to send
  - a: Sending all messages by dequeue one by one
  - q: Quit simulation, all messages will be sent
  - Using Circular Queue will support unlimited enqueue operation, still having full condition

```
[13]
```

## Example: Queue of messages

- Queue: array of char pointers in heap memory
- Enqueue

```
Check if queue is full
Else

If queue is empty front = 0
Increase rear (go back if needed)
Add message to rear position
```

### Dequeue

```
Check if message is empty
Else
Get message at front
If there is one message, clear queue
Else increase front (go back if needed)
```

## Example: Queue of messages

Running simulation

```
Message sending simulation
>> +hello
Saved hello
>> +world
Saved world
>> +stack and queue
Saved stack and queue
>> -
Sending hello ...
>> +good bye
Saved good bye
>> a
Sending world ...
Sending stack and queue ...
Sending good bye ...
>> a
No message to send!
Program ended with exit code: 0
```

## 

## **Linked List Implementation**

Implement Queue by Linked List

## [16]

### LinkedList vs Queue

Queue is Linked List with limited operations and/or other interfaces

Linked List	Queue
Add to end	Enqueue
Remove from head	Dequeue
Clear list	Clear queue
Is empty	Is empty
Get size	Get size
	To array
	Front / Rear

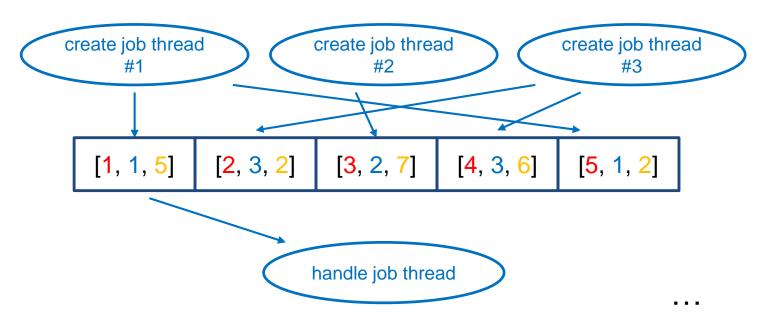
- Simulation a queue of jobs
  - A job has an id, agent id (who creates this job) and duration (in seconds)
  - Several threads will create jobs in parallel and store them in a job queue
  - A thread will get job from that queue to execute it
  - Shared resources: queue of jobs and number of jobs

```
[1, 1, 5] [2, 3, 2] [3, 2, 7] [4, 3, 6] [5, 1, 2] ...
```

## [18]

## Example: Queue of Jobs

Simulation a queue of jobs



[19]

### Example: Queue of Jobs

- Implement queue of jobs
  - Job data

Job node in LL

```
typedef struct
{
    int job_id;
    int agent_id;
    int duration; // in second
} job_data;
```

```
typedef struct str_job job;
struct str_job
{
    job_data data;
    job* next;
};
```

## [20]

- Implement queue of jobs
  - Enqueue operation

```
void enqueue(job** jobs_queue, job_data ajob)
    job* j = create_job(ajob);
    if (is empty(*jobs_queue)) *jobs_queue = j;
    else
        job* rear = *jobs queue;
        while (rear->next != NULL) rear = rear->next;
        rear->next = j;
```

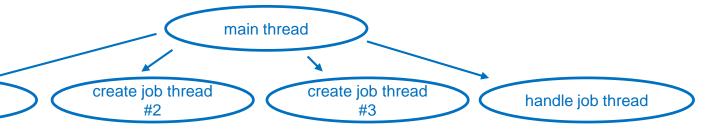
- Implement queue of jobs
  - Dequeue operation

```
job_data dequeue(job** jobs_queue)
    if (is_empty(*jobs_queue))
        printf("Jobs queue is empty!");
        exit(1);
    else
        job* j = *jobs_queue;
        *jobs_queue = j->next;
        job_data data = j->data;
        free(j);
        return data;
```

## Example: Queue of Jobs

- Implement parallel threads
  - pthread\_create: create a new thread to run in parallel

- pthread\_exit: terminate a thread
- sleep: to make current thread idle for a while



create job thread #1

[22]

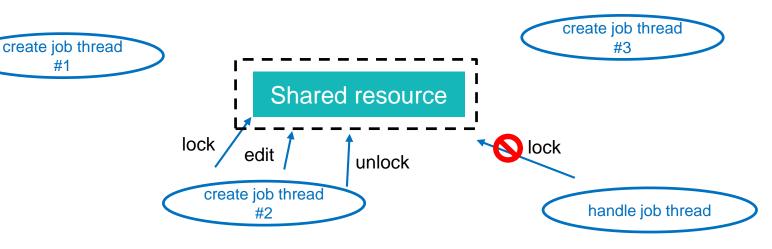
# [23]

## **Example: Queue of Jobs**

### Shared resources

#1

- pthread mutex t locker: declare a locker to lock shared resource
- pthread\_mutex\_init(&locker,NULL): initialize the locker
- pthread\_mutex\_lock(&locker): lock global resources
- pthread mutex unlock(&locker): unlock global resources



- Putting all together
  - main thread

```
// init locker
pthread_mutex_init(&locker,NULL);

// create threads to generate jobs in parallel
for (int agent_id = 1; agent_id <= NAGENTS; agent_id++)
    pthread_create(&tid, NULL, generate_jobs, (void*)agent_id);
// create thread to handle generated jobs
pthread_create(&tid, NULL, handle_jobs, NULL);</pre>
```

## [25]

- Putting all together
  - create job thread

```
int agent id = (int) threadid;
for (int i = 1; i \le NJOBS; i++)
    pthread mutex lock(&locker); // lock the locker
    // increase number of jobs
   // create job data based on agent id, number of jobs
    // enqueue job data
    pthread mutex unlock(&locker); // unlock the locker
    sleep(rand() % WAIT_TIME);
pthread_exit(NULL);
```

[26]

- Putting all together
  - handle job thread

```
Endless loop
Lock the locker

If jobs queue is empty and no more jobs in future
Quit handle thread
Else If queue is not empty
Handle a job
Else
Wait for a job
```

## Example: Queue of Jobs

Putting all together

```
Agent id 1 create job 1
Number of waiting jobs: 1
Handle job 1 of agent 1 in 4s
Agent id 2 create job 2
Agent id 3 create job 3
Agent id 2 create job 4
Agent id 2 create job 5
Agent id 1 create job 6
Agent id 3 create job 7
Number of waiting jobs: 6
Handle job 2 of agent 2 in 0s
Number of waiting jobs: 5
Handle job 3 of agent 3 in 0s
Number of waiting jobs: 4
Handle job 4 of agent 2 in 4s
```

## 

## **Priority Queue**

Introduce Priority Queue and how to implement it



## What is Priority Queue?

- Same as a normal queue but ...
  - Each element has a priority
  - Elements with higher priority will be dequeued before elements with lower priority
  - Elements with same priority will be dequeued in the order of enqueuing

"All animals are equal, but some animals are more equal than others"

George Orwell – Animal Farm

## [30]

## **Example of Priority Queue**

- Enqueue following pairs (key, priority)
  (1, 1), (2, 1), (3, 2), (4, 1), (5, 2), (6, 3), (7, 1)
- Actual queue in building:
  - **\*** (1, 1)
  - **4** (1, 1), (2, 1)
  - **\*** (3, 2), (1, 1), (2, 1)
  - **\*** (3, 2), (1, 1), (2, 1), (4, 1)
  - **\*** (3, 2), (5, 2), (1, 1), (2, 1), (4, 1)
  - **\*** (6, 3), (3, 2), (5, 2), (1, 1), (2, 1), (4, 1)
  - **\*** (6, 3), (3, 2), (5, 2), (1, 1), (2, 1), (4, 1), (7, 1)

# [31]

## Priority Queue by Array

- Enqueue algorithm:
  - Insert new element at the rear
  - If priority of new element is greater than the one ahead then move it up until it's less than or equal

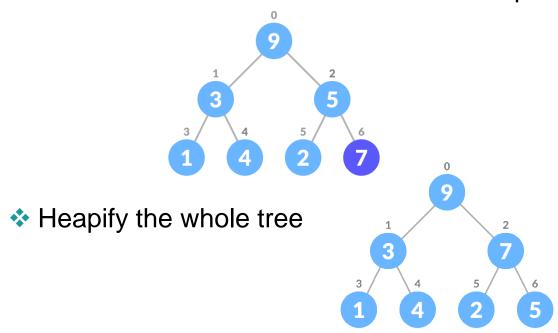
```
4 (3, 2), (1, 1), (2, 1), (4, 1)
```

- **\*** (3, 2), (1, 1), (2, 1), (4, 1), (5, 2)
- **\*** (3, 2), (1, 1), (2, 1), (5, 2), (4, 1)
- **\*** (3, 2), (1, 1), (5, 2), (2, 1), (4, 1)
- **4** (3, 2), (5, 2), (1, 1), (2, 1), (4, 1)
- Complexity:
  - Enqueue: O(n)
  - ❖ Dequeue: O(1)

## [32]

## Priority Queue by Heap

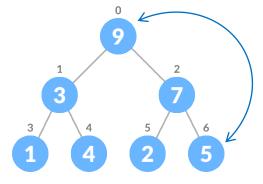
- Enqueue algorithm:
  - Insert new element at the end of the heap



# [33]

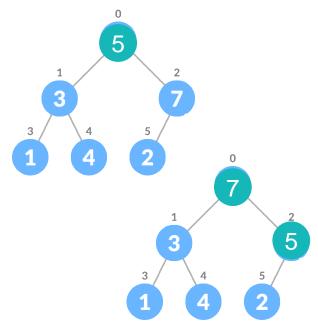
## Priority Queue by Heap

- Dequeue algorithm:
  - Swap the root with the last element then remove it





- Complexity:
  - Enqueue: O(logn)
  - Dequeue: O(logn)



### Priority Queue by Heap

Implement Priority Queue by Heap

```
typedef struct
                     extern int qsize;
    int key;
                      void swap(element *a, element *b);
    int priority;
                     void print_queue(element* queue);
    int order;
                     void heapify(element* queue, int i);
  element;
                      void enqueue(element* queue, element e);
                      element dequeue(element* queue);
                     int greater(element a, element b);
                     int is_empty(void);
```

Fix unstable characteristics of heap

[34]

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[35]

## Priority Queue by Heap

Implement PriorityQueue by Heap

```
void heapify(element* queue, int i)
    int left = 2 * i + 1;
    int right = 2 * i + 2;
    int imax = i;
    if (left < qsize && greater(queue[left], queue[imax]))</pre>
        imax = left;
    if (right < qsize && greater(queue[right], queue[imax]))</pre>
        imax = right;
    if (imax != i)
        swap(&queue[i], &queue[imax]);
        heapify(queue, imax);
```

## Priority Queue by Heap

Implement Priority Queue by Heap

```
int greater(element a, element b)
         // First compare by priority
         // Then compare by order if same priority
             void enqueue(element* queue, element e)
                 // Add e to the end of the queue
                 // Heapify the whole tree bottom-up
                          element dequeue (element* queue)
                             // Get first element, save into e
                             // Swap first element and last element
                             // Heapify the new tree (exclude last element)
                             // Return e
[36]
```