

ASCON

Authenticated Encryption
CAESAR and NIST LWC Winner

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Outline

- 1 T1 - Construction
- 2 T2 - Cryptanalysis
- 3 T3 - Verilog
- 4 T4 - Automated Analysis
- 5 Conclusion

Abstract

Abstract. Authenticated encryption satisfies the basic need for authenticity and confidentiality in our information infrastructure. In this paper, we provide the specification of ASCON-128 and ASCON-128a. Both authenticated encryption algorithms provide efficient authenticated encryption on resource-constrained devices and on high-end CPUs. Furthermore, they have been selected as the “primary choice” for lightweight authenticated encryption in the final portfolio of the CAESAR competition. In addition, we specify the hash function ASCON-HASH, and the extendable output function ASCON-XOF. Moreover, we complement the specification by providing a detailed overview of existing cryptanalysis and implementation results.

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Key Points

- **Authentication** - AEAD, Single Pass, Online
- **Device Versatility** - High End CPU, Small Devices
- **Primitive Options** - Encryption, Hash, XOF
- **Cryptanalysis** - Linear, Differential, Zero-Sum, Integral

Cipher Specifications

Recommended parameters for ASCON authenticated encryption

Cipher	Bit size of					Rounds	
	key	nonce	tag	rate	capacity	p^a	p^b
ASCON-128	128	128	128	64	256	12	6
ASCON-128a	128	128	128	128	192	12	8

Sponge Construction

- **State (s)** - The sponge operates on a state of 320 bits
- **Rate (r)** - Rate at which message is consumed
- **Capacity (c)** - Encryption, Hash, XOF
- **Tag (t)**- Associated with Authentication.
- **a, b** - Number of Rounds in **End** and **Core** Permutations

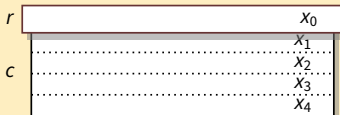
State (S)

$$S = x_0 || x_1 || x_2 || x_3 || x_4$$



$$S = S_r || S_c$$

ASCON-128



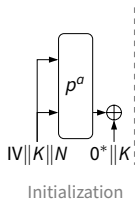
- **Rate (r)** - 64 Bits Block
- **Capacity (c)** - 256 Bits

ASCON-128a



- **Rate (r)** - 128 Bits Block
- **Capacity (c)** - 192 Bits

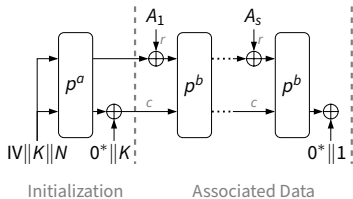
Authenticated Encryption



Duplex-Sponge Construction

- **Initialization** - Initialize the 320-Bit State S with Initialization Vector IV , key K , and Nonce N . Run it through **End Permutation** p^a and inject padded K .

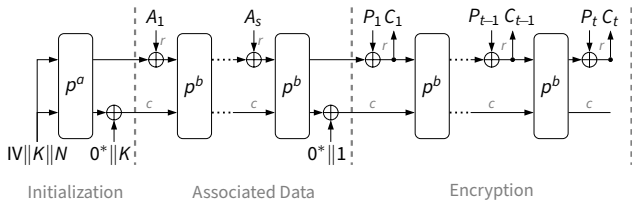
Authenticated Encryption



Duplex-Sponge Construction

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- **Associated Data Processing** - Inject S with Associated Data blocks A_i and digest through **Core Permutation** p^b .

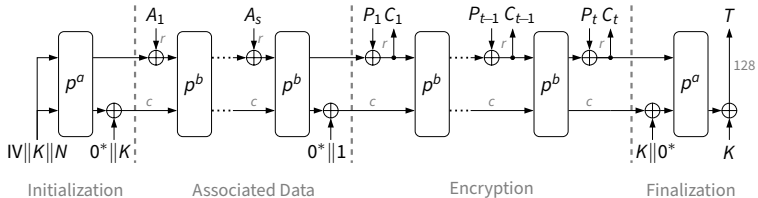
Authenticated Encryption



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- **Encryption** - Inject S with plaintext blocks P_i and extract ciphertext blocks C_i after each run of p_b .

Authenticated Encryption



Duplex-Sponge Construction

- **Initialization** - Initialize the 320-Bit State S with Initialization Vector IV , key K , and Nonce N . Run it through **End Permutation** p^a and inject padded K .
- **Associated Data Processing** - Inject S with Associated Data blocks A_i and digest through **Core Permutation** p^b .
- **Encryption** - Inject S with plaintext blocks P_i and extract ciphertext blocks C_i after each run of p_b .
- **Finalization** - Inject padded K again and extract the Tag T for authentication after running through p^a .

Initialization and Associated Data Processing

Initialization

IV	k	r	a	b	0^{160-k}	x_0
K						x_1
						x_2
N						x_3
						x_4

$$IV_{k,r,a,b} \leftarrow k || r || a || b || 0^{160-k}$$

$$S \leftarrow IV || K || V$$

$$S \leftarrow p^a(S) \oplus (0^{320-k} || K)$$

- $|IV| - 64, |K| - 128, |N| - 128$
- ASCON-128 - 80400c0600000000
- ASCON-128a - 80800c0800000000

Initialization and Associated Data Processing

Initialization

IV	k	r	a	b	0^{160-k}	x_0
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- $|IV| - 64, |K| - 128, |N| - 128$
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Associated Data Processing

- Padding the Associated Data A and splitting s r -Bit Blocks

$$A_1, \dots, A_s \leftarrow A || 1 || 0^{r-1-(|A| \bmod r)}$$

- Digesting Associated Data Blocks

$$S \leftarrow p^b((S_r \oplus A_i) || S_c), 1 \leq i \leq s$$

- Adding 1-bit Domain Separation Constant

$$S \leftarrow S \oplus (0^{319} || 1)$$

Encryption and Finalisation

Encryption

- Padding the Plaintext P and splitting t r -Bit Blocks

$$P_1, \dots, P_t \leftarrow P || 1 || 0^{r-1-(|P| \bmod r)}$$

- Injecting Plaintext and Extracting Ciphertext

$$C_i \leftarrow S_r \oplus P_i \quad 1 \leq i \leq t$$

$$S \leftarrow \begin{cases} p^b(C_i || S_c) & \text{if } 1 \leq i < t \\ C_i || S_c & \text{if } i = t \end{cases}$$

- Unpadding Last Ciphertext C_t

$$C'_t \leftarrow \lfloor C_t \rfloor_{|P| \bmod r}$$

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- Unpadding Last Ciphertext C_t

$$C'_t \leftarrow \lfloor C_t \rfloor_{|P| \bmod r}$$

Finalization

- Adding Padded Key to State

$$S \leftarrow S \oplus (0^r || K || 0^{c-k})$$

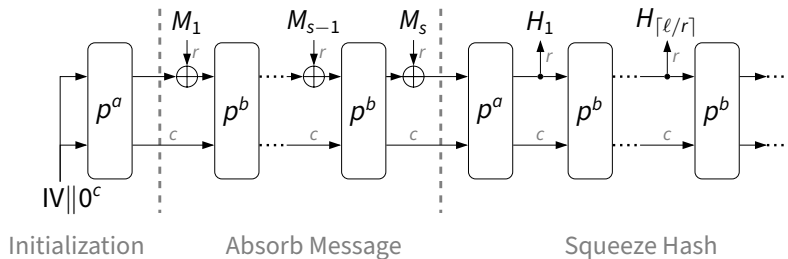
- End Permutation

$$S \leftarrow p^a(S)$$

- Extracting Tag T from the Least Significant 128 Bits after adding Key

$$T \leftarrow [S]^{128} \oplus [K]^{128}$$

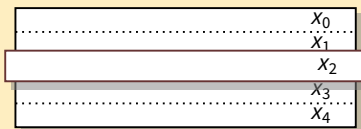
Hash



Round Constant (p_C)

Permutation ($p = p_L \circ p_S \circ p_C$)

1 Byte Constant

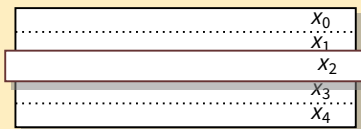


- $x_2 := x_2 \oplus c_r$
- 12 Round Constants
- $p^a \rightarrow c_r$ and $p^b \rightarrow c_{a-b+r}$

Round Constant (p_C)

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1 Byte Constant



- $x_2 := x_2 \oplus c_r$
- 12 Round Constants
- $p^a \rightarrow c_r$ and $p^b \rightarrow c_{a-b+r}$

p^{12}	p^8	p^6	Constant	p^{12}	p^8	p^6	Constant
0			000000000000000000f0	6	2	0	00000000000000000096
1			000000000000000000e1	7	3	1	00000000000000000087
2			000000000000000000d2	8	4	2	00000000000000000078
3			000000000000000000c3	9	5	3	00000000000000000069
4	0		000000000000000000b4	10	6	4	0000000000000000005a
5	1		000000000000000000a5	11	7	5	0000000000000000004b

Sbox (p_S) and Linear Layer (p_L)

Permutation ($p = p_L \circ p_S \circ p_C$)

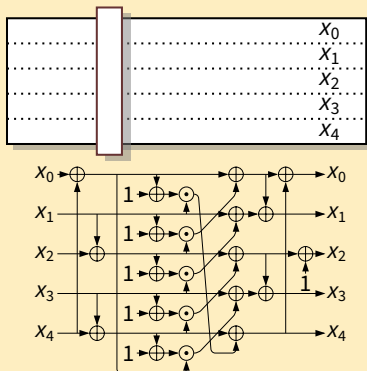
x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f	10	11	12	13	14	15	16	17	18	19	1a	1b	1c	1d	1e	1f
$S(x)$	4	b	1f	14	1a	15	9	2	1b	5	8	12	1d	3	6	1c	1e	13	7	e	0	d	11	18	10	c	1	19	16	a	f	17

Sbox (p_S) and Linear Layer (p_L)

Permutation ($p = p_L \circ p_S \circ p_C$)

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f	10	11	12	13	14	15	16	17	18	19	1a	1b	1c	1d	1e	1f
$S(x)$	4	b	1f	14	1a	15	9	2	1b	5	8	12	1d	3	6	1c	1e	13	7	e	0	d	11	18	10	c	1	19	16	a	f	17

Sbox

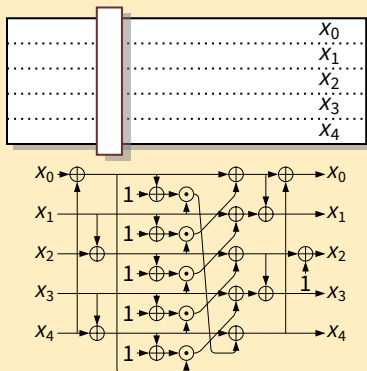


Sbox (p_S) and Linear Layer (p_L)

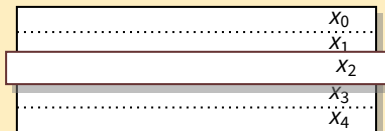
Permutation ($p = p_L \circ p_S \circ p_C$)

x	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f	10	11	12	13	14	15	16	17	18	19	1a	1b	1c	1d	1e	1f
$S(x)$	4	b	1f	14	1a	15	9	2	1b	5	8	12	1d	3	6	1c	1e	13	7	e	0	d	11	18	10	c	1	19	16	a	f	17

Sbox



Linear Layer



$$X_0 := X_0 \oplus (X_0 \ggg 19) \oplus (X_0 \ggg 28)$$

$$X_1 := X_1 \oplus (X_1 \ggg 61) \oplus (X_1 \ggg 39)$$

$$X_2 := X_2 \oplus (X_2 \ggg 1) \oplus (X_2 \ggg 6)$$

$$X_3 := X_3 \oplus (X_3 \ggg 10) \oplus (X_3 \ggg 17)$$

$$X_4 := X_4 \oplus (X_4 \ggg 7) \oplus (X_4 \ggg 41)$$

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DDT

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f	10	11	12	13	14	15	16	17	18	19	1a	1b	1c	1d	1e	1f		
0	32	
1	4	.	4	.	4	.	4	4	.	4	.	4	.	4	.	
2	4	.	4	.	4	.	4	.	4	.	4	.	4	.	4	.	
3	.	4	.	.	.	4	.	.	4	.	.	.	4	.	.	4	.	4	.	.	4	.	.	4	4	
4	8	8	8	8	.	
5	4	.	4	.	4	.	4	.	4	.	4	.	.	4	.	4	
6	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	
7	.	.	4	4	.	.	4	4	.	.	4	4	.	.	4	4	
8	4	4	4	4	4	4	4	4	
9	.	2	.	2	2	.	2	.	2	.	2	.	.	2	2	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	
10	.	2	2	.	2	.	2	.	2	2	.	2	.	2	.	2	.	2	2	.	2	.	.	2	.	2	2	.	2	.	2	.	2	
11	.	.	2	2	.	.	2	2	.	.	2	2	.	.	2	2	.	.	2	2	.	.	2	2	.	.	2	2	.	.	2	2	.	
12	.	8	8	8	8	
13	.	2	.	2	.	2	.	2	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	
14	.	4	4	.	4	.	.	4	4	4	.	4	.	.	4	
15	4	4	.	.	4	4	4	4	.	.	4	4	.	.	.	
16	8	.	8	8	.	8	
17	8	.	8	.	8	.	8	.	8	
18	.	2	.	2	.	2	.	2	.	2	.	2	.	2	2	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	
19	.	.	8	.	8	8	.	8
20	.	.	.	4	4	4	4	4	4	4	4	
21	4	.	4	.	4	.	4	4	.	4	4	.	4	
22	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	
23	.	.	4	.	4	4	.	4	4	.	4	4	.	4	.	.	.	
24	.	.	.	2	2	2	2	2	2	2	2	2	2	2	2	2	.	.	.	2	2	2	2	2	
25	.	.	.	4	.	.	4	.	4	4	.	.	4	4	4	.	.	4	.	.	
26	.	2	2	.	2	2	.	2	.	.	2	2	.	.	2	.	2	2	.	2	2	.	2	2	.	2	.	.	2	2	.	2	.	
27	.	.	2	2	2	2	.	.	.	2	2	2	2	2	2	2	2	2	2	2	2	2	.	.	
28	.	4	.	4	.	.	.	4	.	4	4	.	4	4	.	4	
29	.	.	.	4	.	4	.	.	4	4	.	4	4	4	.	4	.	.	
30	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	
31	.	.	4	4	4	4	4	4	4	4	

LAT

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f	10	11	12	13	14	15	16	17	18	19	1a	1b	1c	1d	1e	1f		
0	16	
1	8	.	.	.	4	4	.	-4	4	.	.	.	4	4	.	.	4	-4	4	.	-4	.	-4	.	-4	.	.		
2	-8	8	.	.	4	4	.	4	4	.	4	4	.	4	4	.	.	-4	-4		
3	.	8	4	4	.	4	4	.	-4	-8	4	.	4	.	4	.	-4	.		
4	.	.	.	4	.	-4	.	.	.	4	.	.	.	4	-4	-4	.	4	.	-4	.	-4	.	.	.	-8	.	-4	.	-4	.	4	-4	
5	.	.	.	4	.	4	.	.	.	-4	-4	.	.	.	-4	4	.	-4	-4	4	.	-4	4	.	-4	.	-8	.	-4	
6	.	.	.	4	.	-4	-4	.	.	4	-4	-4	.	-4	-4	.	8	.	-4	-4	.	-4	4	.		
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12	.	.	-8	4	-8	-4	4	.	-4	.	-4	.	.	.	-4	.	4	4	.	-4	
13	.	.	.	-4	-8	4	.	.	.	4	-4	-4	.	.	-4	.	.	.	4	-4	.	-4	-4	4	4	.	.	
14	.	.	.	-4	8	-4	-4	.	.	-4	-4	-4	.	.	.	4	4	.	-4	-4	.	.	4	.	-4	
15	.	.	.	8	-4	-8	-4	.	.	-4	-4	.	.	4	.	4	.	.	.	-4	-4	.	
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19	-8	-8	4	-4	4	4	-4	.	.	.	-4	4	4	-4	
20	.	.	.	4	.	4	.	.	4	4	-4	-4	-4	-4	.	-4	.	4	.	.	4	-4	4	-4	-4	.	4	4	.	.	.	4	.	
21	.	.	.	4	.	-4	-4	4	.	-4	4	.	8	.	4	.	4	4	4	.	-4	-4	.	
22	.	.	.	-4	.	-4	.	.	.	4	.	-4	4	4	4	.	8	.	-4	.	4	.	.	4	.	.	4	4	.	.	.	-4	.	
23	.	.	.	4	.	-4	.	8	.	-4	.	-4	4	.	.	-4	4	-4	4	4	.	4	4	.	
24	-8	.	4	4	.	-4	.	.	4	4	-4	-4	-4	-4	.	.	-4	4	.	.	.	-4	.	
25	4	-4	-4	4	.	-8	4	-4	-4	-4	4	4	.	.	-4	-4	.	
26	.	8	-4	.	-4	-4	.	4	.	.	.	-4	4	-4	-4	.	.	-4	.	.	4	.	-4	.	-4	.	-4	
27	8	.	-4	4	-4	-4	-4	4	-4	4	.	.	-4	-4	.	.	-4	-4	.	-4	
28	.	.	8	4	.	-4	.	.	4	.	4	.	4	-4	4	.	.	-4	.	.	-4	-4	4	4	4	.	.	
29	.	.	.	-4	.	4	.	8	.	4	.	4	8	.	-4	.	-4	4	.	-4	
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31	.	.	8	4	.	4	4	-4	.	-4	4	.	.	-4	.	.	.	4	4	-4	.	.	4	.	-4	.	.	.	

BCT

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f	10	11	12	13	14	15	16	17	18	19	1a	1b	1c	1d	1e	1f	
0	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	32	
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4	32	.	.	.	4	8	12	8	.	16	.	16	4	8	12	8	.	16	.	16	4	8	12	8	16	.	16	.	4	8	12	8	
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7	32	.	8	4	4	.	4	4	.	8	8	12	4	.	4	4	.	.	4	.	4	.	.	.	8	.	12	.	4	.	.	.	
8	32	.	.	4	.	8	4	4	4	4	4	4	8	.	8	.	12	4	12	.	.	.	4	.	.	8	4	
9	32	2	.	2	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	2	.	2	.
10	32	2	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	2	.	2	.	2	.	2	.	2	.	2	.	2	.
11	32	.	2	2	.	.	2	2	.	.	2	2	.	.	2	2	.	.	2	2	.	2	.	2	.	2	.	.	2	2	.	2	.
12	32	12	16	8	16	8	.	4	12	.	8	.	8	.	4	.	12	16	8	16	8	16	4	16	.	12	.	8	.	8	.	4	.
13	32	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.
14	32	12	4	.	4	.	4	8	8	8	4	4	4	4	.	8	.	12	.	4	.	4	.	4	.
15	32	4	8	4	8	.	.	.	8	4	12	.	12	4	.	.	.	4	.	4	4	8	.	4	4	4	.	.	.
16	32	4	16	4	16	8	16	8	.	12	16	12	16	8	16	8	4	.	4	.	8	.	8	.	12	.	12	.	8	.	8	.	.
17	32	8	.	8	.	8	16	8	16	4	.	4	.	4	16	4	16	12	.	12	.	12	16	12	.	8	.	8	.	8	16	8	8
18	32	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.	2	.
19	32	16	12	8	12	8	.	.	16	16	12	8	12	8	.	.	16	.	8	4	8	4	.	.	16	16	8	4	8	4	.	.	
20	32	.	.	.	4	4	12	4	.	8	.	8	4	4	12	4	8	.	8	.	8	8	.
21	32	4	8	4	.	4	.	4	.	.	8	.	.	12	.	12	.	8	8	8	4	8	4	4
22	32	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	
23	32	.	8	4	4	.	4	4	.	8	8	12	4	.	4	4	.	.	4	.	4	.	.	.	8	.	12	.	4	.	.	.	
24	32	.	.	.	2	2	2	2	2	2	2	2	2	2	2	2	.	.	.	2	2	2	2	.
25	32	.	.	4	.	8	4	4	4	4	4	4	8	.	8	.	12	4	12	.	.	.	4	.	.	8	4	.
26	32	2	2	.	.	2	2	.	2	.	.	2	2	.	.	2	.	2	.	2	.	2	2	.	2	.	2	.	2	.	2	.	2
27	32	.	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	.	.
28	32	4	8	4	8	.	.	.	8	4	12	.	12	4	.	.	4	.	4	4	8	.	4	4	4	.	.	.
29	32	8	.	4	.	4	.	.	12	4	.	12	8	.	8	.	8	4	8	.	8	.	4	.	4	.	.	
30	32	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
31	32	8	12	4	12	4	.	.	8	.	8	.	8	.	.	.	8	.	4	4	4	4	.	.	.	8

Sbox Analysis

- **Differential Branch Number - 3**
- **Differential Uniformity - 8**

DDT Frequency Analysis

$f(\Delta_i, \Delta_o)$	S
0	707
2	176
4	120
6	0
8	20

Sbox Analysis

- **Differential Branch Number** - 3
- **Differential Uniformity** - 8

DDT Frequency Analysis

$f(\Delta_i, \Delta_o)$	S
0	707
2	176
4	120
6	0
8	20

- **Linear Branch Number** - 3
- **Maximum Absolute Linear Bias** - 8

LAT Frequency Analysis

ϵ	S
-8	18
-4	174
0	647
4	162
8	22

BCT BDT

Sbox Analysis

● Boomerang Uniformity - 16

BCT Frequency Analysis

$f(\Delta_i, \nabla_o)$	S
0	445
2	176
4	150
8	110
12	50
16	30

BDT Frequency Analysis

$f(\Delta_i, \Delta_o, \nabla_o)$	S
0	31624
2	352
4	720
8	40
32	32

Differential and Linear Cryptanalysis

- **Linear Layer** (Σ_i) - Differential and Linear Branch Number is 4.
- The minimum number of active S-boxes after 3 rounds - (p^3)
 - **Differential Characteristics** - 15
 - **Linear Characteristics** - 13

Differential and Linear Cryptanalysis

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Rounds (R)	1	2	3	4	5
Minimum # Active Sboxes (Differential)	1	4	15	≤ 44	≤ 78
Minimum # Active Sboxes (Linear)	1	4	13	≤ 43	≤ 67

Differential and Linear Cryptanalysis

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Collision Producing Differentials - Forgery Attack

- Differentials with differences in the rate part (S_r) of State at the input (Δ_I^r) and output (Δ_O^r) of p^b .

$$\Delta_I^r \xrightarrow{p^b} \Delta_O^r$$

- Might be useful in Forgery attack on the AEAD scheme.

Truncated Differential over Sbox

Undisturbed Bits

For a specific input difference Δ_i of an S-box, if some bits of the output difference Δ_o^* remain invariant, then we call such bits **undisturbed**.

$$Pr \left[\Delta_i \xrightarrow{S} \Delta_o^* \right] = 1$$

Truncated Differential over Sbox

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Δ_i	Δ_o^*
00001	*1***
00010	1***1
00011	***0*
00100	**110
00101	1****
01111	*1*0*

Δ_i	Δ_o^*
10000	*10**
10001	10**1
10011	0***0
10100	0*1**
10101	****1
10110	1****

Δ_i	Δ_o^*
00110	****1
00111	0**1*
01000	**11*
01011	***1*
01100	**00*
01110	*0***

Δ_i	Δ_o^*
10111	****0
11000	**1**
11100	**0**
11110	*1***
11111	*0***

Truncated Differential over Sbox

Undisturbed Bits

For a specific input difference Δ_i of an S-box, if some bits of the output difference Δ_o^* remain invariant, then we call such bits **undisturbed**.

$$Pr \left[\Delta_i \xrightarrow{S} \Delta_o^* \right] = 1$$

Note - Inverse S-Box only has 2 undisturbed bits.

Δ_i	Δ_o^*
00001	*1***
00010	1***1
00011	***0*
00100	**110
00101	1****
01111	*1*0*

Δ_i	Δ_o^*
10000	*10**
10001	10**1
10011	0***0
10100	0*1**
10101	****1
10110	1****

Δ_i	Δ_o^*
00110	****1
00111	0**1*
01000	**11*
01011	***1*
01100	**00*
01110	*0***

Δ_i	Δ_o^*
10111	****0
11000	**1**
11100	**0**
11110	*1***
11111	*0***

3.5 Round Truncated Differential Distinguisher

[illegible]

A set of navigation icons typically found in Beamer presentations, including symbols for back, forward, search, and other slide controls.

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Impossible Differential

5 Round Differential

x_0	0000000000000000
x_1	0000000000000000
x_2	0000000000000000
x_3	0000000000000000
x_4	8000000000000000
\nrightarrow	
x_0	01000000000100002
x_1	0000000000000000
x_2	0000000000000000
x_3	0000000000000000
x_4	0000000000000000

Impossible Differential

5 Round Impossible Truncated Differential

3.5 Truncated Differential Distinguisher

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x_2	0000000000000000
x_3	0000000000000000
x_4	0000000000000000

Impossible Differential

5 Round Differential

x_0	0000000000000000
x_1	0000000000000000
x_2	0000000000000000
x_3	0000000000000000
x_4	8000000000000000
\nrightarrow	
x_0	0100000000100002
x_1	0000000000000000
x_2	0000000000000000
x_3	0000000000000000
x_4	0000000000000000

5 Round Impossible Truncated Differential	
3.5 Truncated Differential Distinguisher	
S_4	*****0*****
	*****0*****
	*****0*****
	*****0*****
	*****0*****
1.5 Round Impossible Backward Differential	

5 Round Differential

x_0	0000000000000000
x_1	0000000000000000
x_2	0000000000000000
x_3	0000000000000000
x_4	8000000000000000
\nearrow	
x_0	0100000000100002
x_1	0000000000000000
x_2	0000000000000000
x_3	0000000000000000
x_4	0000000000000000

[illegible]

Zero Sum Distinguisher

ASCON Sbox - χ Keccak Mapping

- χ has branch number 2, Fix Point at 0, Outputs depend on 3 Inputs
- ASCON S-box is designed by adding lightweight affine transformations to the input and output of .

Zero Sum Distinguisher

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Algebraic Degree

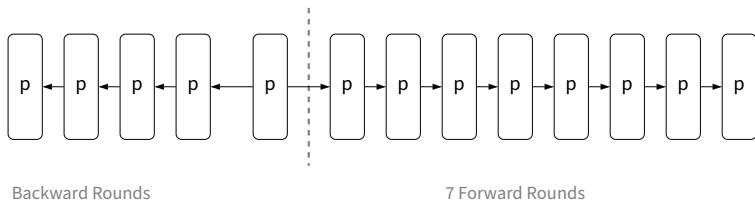
- ASCON S-box (S) \rightarrow ASCON Permutation (p)

$$d(S) = 2 \rightarrow d(p) \leq 2 \rightarrow d(p^r) \leq 2^r$$

- ASCON S-Box Inverse (S^{-1}) \rightarrow ASCON Permutation Inverse (p^{-1})

$$d(S^{-1}) = 3 \rightarrow d(p^{-1}) \leq 3 \rightarrow d((p^{-1})^r) = 3^r$$

12 Round (p^a) - Zero Sum Distinguisher (2^d)

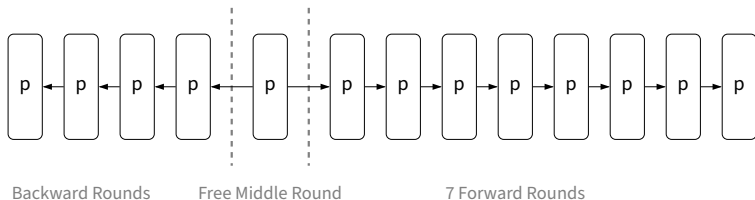


Basic Distinguisher (2^{244})

- 5 Backward Rounds - (3^5)
- 7 Forward Rounds - (2^7)

$$d = \max\{3^5, 2^7\} + 1 = 243 + 1 = 244$$

12 Round (p^a) - Zero Sum Distinguisher (2^d)



Basic Distinguisher (2^{244})

- 5 Backward Rounds - (3^5)
 - 7 Forward Rounds - (2^7)
- $$d = \max\{3^5, 2^7\} + 1 = 243 + 1 = 244$$

Free Middle Round (2^{130})

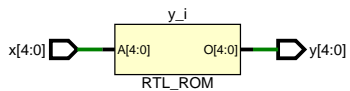
- d is a multiple of 5-Bit Sbox
- All d variables positioned in S-box
- S-Box Output also contains d variable and $320 - d$ constant bits
- 4 Backward, 7 Forward Rounds

$$d = \max\{3^4, 2^7\} + 1 = 128 + 1 = 129 \rightarrow 130$$

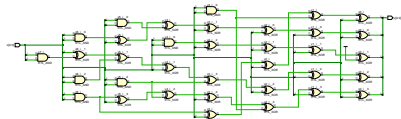
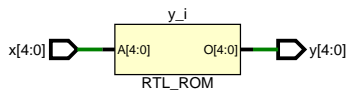
Outline

- 1 T1 - Construction
- 2 T2 - Cryptanalysis
- 3 T3 - Verilog**
- 4 T4 - Automated Analysis
- 5 Conclusion

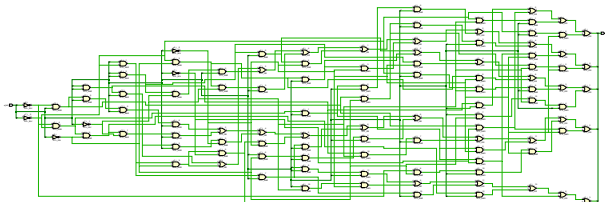
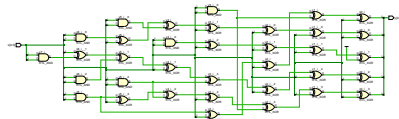
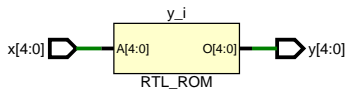
Sbox



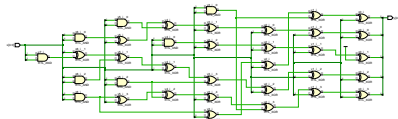
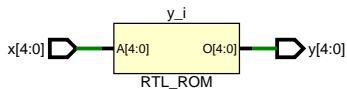
Sbox



Sbox

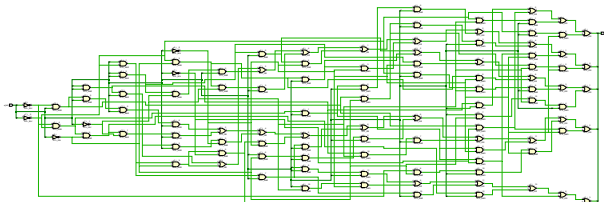


Sbox

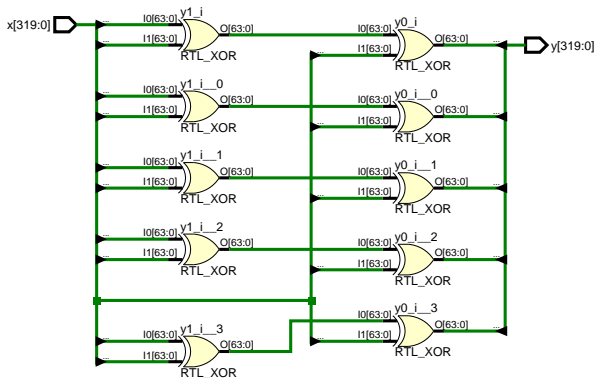


LuT	ANF	K-MAP
43.920	42.480	39.240

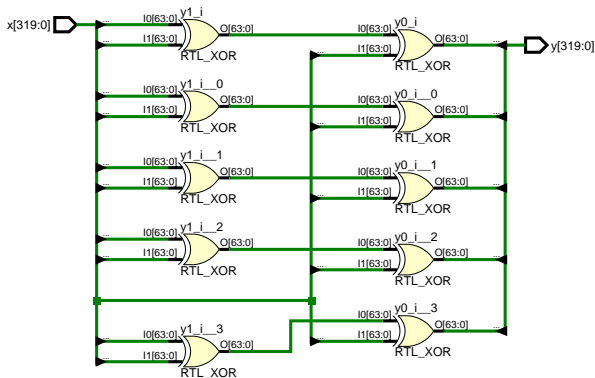
Component	Slice LuT (Logic)	OBUF	IBUF	LUT5	LUT4
Number Used	3	5	5	3	2



Linear Diffusion - 1958.400 GE

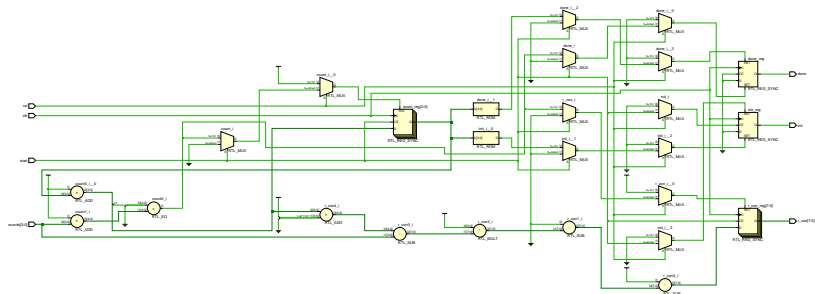


Linear Diffusion - 1958.400 GE

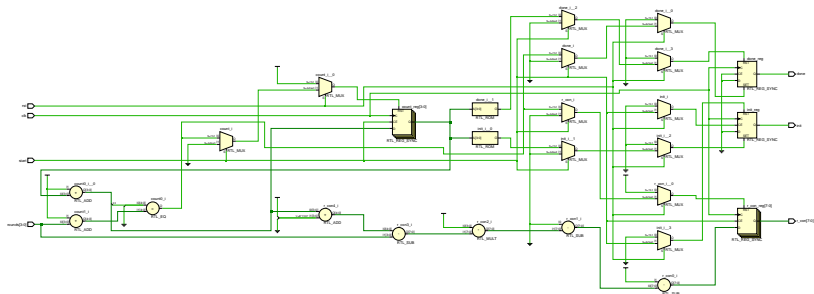


Components	Number Used
Slice LuT (Logic)	169
obuf	320
lut3	320
ibuf	320

Round Constant - 321.840 GE



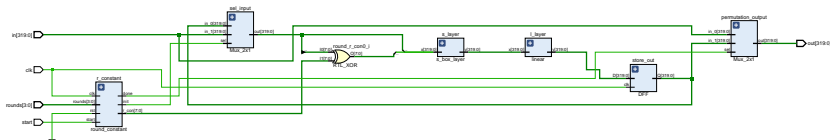
Round Constant - 321.840 GE



Components	Slice LuT (Logic)	Slice Register (Flip Flop)	FDRE	FDRE	OBUF	LUT1	LUT6	IBUF	LUT4	LUT2	LUT3	FDSE	LUT5	BUFG (Clock)
Number Used	29	14	12	10	9	7	7	7	6	5	3	2	1	1



Permutation



● LuT Sbox - 9148.320 GE

Components	Slice LuT (Logic)	Slice Registers (Flip Flop)	LUT3	FDRE	IBUF	OBUF	LUT5	LUT4	LUT6	LUT	1LUT2	BUFG (Clock)
Number Used	740	334	862	334	326	320	196	110	99	9	5	1

● KMAP Sbox - 8848.800 GE

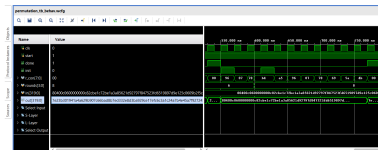
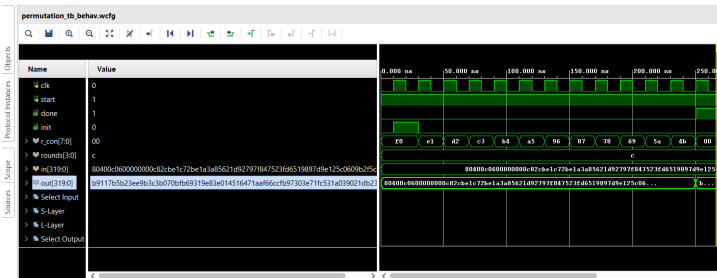
Components	Slice LuT (Logic)	Slice Registers (Flip Flop)	LUT3	FDRE	IBUF	OBUF	LUT5	LUT6	LUT4	LUT1	LUT2	BUFG (Clock)
Number Used	777	334	955	334	326	320	196	134	13	9	5	1

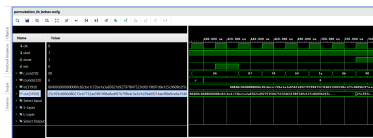
● ANF Sbox - 9056.160 GE

Components	Slice LuT (Logic)	Slice Registers (Flip Flop)	LUT3	FDRE	IBUF	OBUF	LUT6	LUT5	LUT4	LUT1	LUT2	BUFG (Clock)
Number Used	883	334	955	334	326	320	246	140	13	9	5	1

A set of small navigation icons typically found in Beamer presentations, including symbols for back, forward, search, and other slide controls.

Test Bench - p^{12}, p^8, p^6





Outline

- 1 T1 - Construction
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MILP

Number of Active Sboxes - Basic

Variables

MILP

Number of Active Sboxes - Basic

Variables

- $x_{r,w,b} \in \{0, 1\}$ - S-box input Bit b ($0 \leq b \leq 63$) of word X_w ($0 \leq w \leq 4$) of round r active or not.

MILP

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- $d_{r,b} \in \{0, 1\}$ - b^{th} ($0 \leq b \leq 63$) S-box of round r is active

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- $d_{r,b} \in \{0, 1\}$ - b^{th} ($0 \leq b \leq 63$) S-box of round r is active
- $u_{r,w,b} \in \{0, 1\}$ - Linear layer model in word x_w ($0 \leq w \leq 4$) of round r .

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- $y_{r,w,b} \in \{0, 1\}$ - S-box output Bit b ($0 \leq b \leq 63$) of word X_w ($0 \leq w \leq 4$) of round r active or not.
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Objective Function

- Minimize the Active S-boxes

$$\min \sum_{r=1}^R \sum_{b=0}^{63} d_{r,b}$$

MILP

Number of Active Sboxes - Basic

Constraints

- **Non-Triviality** - At least 1 Active Input Bit at Start $\sum_{w=0}^4 \sum_{b=0}^{63} x_{0,w,b} \geq 1$

MILP

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- **Active S-box** - Minimum 1 and Maximum 5 Active Input and Output Bit(s).

$$d_{r,b} \leq \sum_{w=0}^4 x_{r,w,b} \leq 5d_{r,b}$$

$$d_{r,b} \leq \sum_{w=0}^4 y_{r,w,b} \leq 5d_{r,b}$$

MILP

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- **S-box Branch Number** - 3. $\sum_{w=0}^4 (x_{r,w,b} + y_{r,w,b}) \geq 3d_{r,b}$

MILP

Number of Active Sboxes - Basic

Constraints

- **Non-Triviality** - At least 1 Active Input Bit at Start $\sum_{w=0}^4 \sum_{b=0}^{63} x_{0,w,b} \geq 1$
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$$d_{r,b} \leq \sum_{w=0}^4 x_{r,w,b} \leq 5d_{r,b}$$

$$d_{r,b} \leq \sum_{w=0}^4 y_{r,w,b} \leq 5d_{r,b}$$

- **S-box Branch Number** - 3. $\sum_{w=0}^4 (x_{r,w,b} + y_{r,w,b}) \geq 3d_{r,b}$
- **Linear Layer - XOR Operation Model**

$$y_{r,0,b} + y_{r,0,b-19} + y_{r,0,b-28} + x_{r+1,0,b} = 2u_{r,0,b}$$

$$y_{r,2,b} + y_{r,2,b-1} + y_{r,2,b-6} + x_{r+1,2,b} = 2u_{r,2,b}$$

$$y_{r,1,b} + y_{r,1,b-61} + y_{r,1,b-39} + x_{r+1,1,b} = 2u_{r,1,b}$$

$$y_{r,3,b} + y_{r,3,b-10} + y_{r,3,b-17} + x_{r+1,3,b} = 2u_{r,3,b}$$

$$y_{r,4,b} + y_{r,4,b-7} + y_{r,4,b-41} + x_{r+1,4,b} = 2u_{r,4,b}$$

Logical Conditional Modelling

MILP

$$(x_{r,0,b}, x_{r,1,b}, x_{r,2,b}, x_{r,3,b}, x_{r,4,b}) = (\delta_0, \delta_1, \delta_2, \delta_3, \delta_4) \Rightarrow y_{r,w,b} = \delta$$

$$\sum_{w'=0}^4 (-1)^{\delta_i} x_{r,w',b} + (-1)^{\delta+1} y_{r,w,b} - \delta + \sum_{w'=0}^4 \delta_i \geq 0$$

Logical Conditional Modelling

MILP

$$(x_{r,0,b}, x_{r,1,b}, x_{r,2,b}, x_{r,3,b}, x_{r,4,b}) = (\delta_0, \delta_1, \delta_2, \delta_3, \delta_4) \Rightarrow y_{r,w,b} = \delta$$

$$\sum_{w'=0}^4 (-1)^{\delta_i} x_{r,w',b} + (-1)^{\delta+1} y_{r,w,b} - \delta + \sum_{w'=0}^4 \delta_i \geq 0$$

Recall - Undisturbed Bits

Δ_i	Δ_o^*
00001	*1***
00011	***0*
00101	1****
11111	*0***
01011	***1*

Δ_i	Δ_o^*
10101	****1
10110	1****
11110	*1***
01110	*0***

Δ_i	Δ_o^*
10111	****0
11000	**1**
11100	**0**
00110	****1

Gurobi Implementation

Rounds	Active S-Boxes	Variables (Real)	Variables (Binary)	Inequalities	Time Taken
1	1	960	384	2561	0.03 s
2	4	1920	448	5121	0.51 s
3	12	2880	512	7681	2 min 52.47 s

Table: Basic MILP

Rounds	Active S-Boxes	Variables (Real)	Variables (Binary)	Inequalities	Time Taken
1	1	965	384	3393	0.02 s
2	4	1925	448	6785	1.32 s
3	12	2885	512	10177	3 min 28.54 s

Table: Logical Conditional Modeling

Convex Hull

MILP

```
ascon_hull = list(Polyhedron(vertices=generate_vertices(Ascon)).inequality_generator())
```

```
ascon_hull
```

```
[An inequality (-1, 0, 0, 0, 0, 0, 0, 0, 0) x + 1 >= 0,
An inequality (0, -1, 0, 0, 0, 0, 0, 0, 0) x + 1 >= 0,
An inequality (0, 0, -1, 0, 0, 0, 0, 0, 0) x + 1 >= 0,
An inequality (0, 0, 0, -1, 0, 0, 0, 0, 0) x + 1 >= 0,
An inequality (-1, 0, 0, -1, -1, 0, 0, -1, -1, 0) x + 4 >= 0,
An inequality (0, -1, 0, -1, -1, 0, 1, 1, 1, 0) x + 2 >= 0,
An inequality (0, 0, 0, 0, -1, 0, 0, 0, 0, 0) x + 1 >= 0,
An inequality (1, -1, -1, -1, -1, 0, 1, 0, 0, 0) x + 3 >= 0,
An inequality (0, 0, 0, 0, 0, 0, 1, 0, 0, 0) x + 0 >= 0,
An inequality (-1, 0, 0, -1, -1, 0, 0, 1, 1, 0) x + 2 >= 0,
An inequality (0, 0, -1, -1, -1, 0, 1, 1, 1, 0) x + 2 >= 0,
An inequality (-3, -2, -3, -4, -6, -1, -3, -3, -3, -1) x + 23 >= 0,
An inequality (-1, -1, -1, -1, -1, 0, -1, 0, 0, 0) x + 5 >= 0,
An inequality (-1, -1, -1, -2, -2, 0, -1, -1, -1, 0) x + 8 >= 0,
An inequality (-1, -1, -1, -1, 1, 0, 1, 0, 0, 0) x + 3 >= 0,
An inequality (0, -1, -1, -1, -1, 0, 0, -1, -1, 0) x + 5 >= 0,
An inequality (-2, -2, -2, -1, 1, 0, 2, -1, -1, 0) x + 7 >= 0,
An inequality (-1, -1, -1, -2, -1, 0, 1, -2, -2, 0) x + 8 >= 0,
An inequality (-1, -1, -1, -2, -1, 0, 1, 2, 2, 0) x + 4 >= 0,
An inequality (-1, -1, -1, -3, -2, 0, 2, 3, 3, 0) x + 5 >= 0,
An inequality (-1, -1, -1, -1, 0, 0, 1, 1, 1, 0) x + 3 >= 0,
An inequality (1, -1, -1, -2, -2, 0, 2, 1, 1, 0) x + 4 >= 0,
An inequality (0, 0, 0, -1, -1, -1, 1, 1, 1, -1) x + 3 >= 0,
An inequality (0, 0, 0, 0, 0, 0, 0, 0, 0, -1) x + 1 >= 0,
An inequality (0, 0, -1, -1, -1, -1, 0, -1, -1, 0) x + 5 >= 0,
```

Figure: Number of Inequalities generated - 2415

Outline

- 1 T1 - Construction
- 2 T2 - Cryptanalysis
- 3 T3 - Verilog
- 4 T4 - Automated Analysis
- 5 Conclusion**

References

- <https://ascon.iaik.tugraz.at>
- **[DEMS15]** *Cryptanalysis of Ascon* - Christoph Dobraunig, Maria Eichlseder, Florian Mendel, and Martin Schläffer. **IACR** - 2015/030 (pp. 28, 31–33, 35).
- **[Tez16]** *Truncated, Impossible, and Improbable Differential Analysis of Ascon* - Cihangir Tezcan. **IACR** - 2016/490 (pp. 28, 33).

Thanks

Brownie Points - LaTeX-Tikz Block Diagrams

- State Design
- Authenticated Encryption
- Zero Sum Distinguisher

Implementation Info

- Github Link: <https://github.com/highgroundmaster/ASCON>