Answers to Caegory Theory by Steve Awodey

Syed Khalid

Nov, 2023

1 Ch1 Categories

1. **Question:** The objects of **Rel** are sets, and an arrow $A \to B$ is a relation from A to B, that is, a subset $R \subseteq A \times B$. The equality relation $\{\langle a, a \rangle \in A \times A | a \in A\}$ is the identity arrow on a set A. Compositoin in **Rel** is to be given by

$$S \circ R = \{ \langle a, c \rangle \in A \times C \mid \exists b (\langle a, b \rangle \in R \& \langle b, c \rangle \in S) \}$$

for $R \subseteq A \times B$ and $S \subseteq B \times C$.

- (a) Show that **Rel** is a category.
- (b) Show also that there is a functor $G : \mathbf{Sets} \to \mathbf{Rel}$ taking objects to themselves and each function $f : A \to B$ to its graph,

$$G(f) = \{ \langle a, f(a) \rangle \in A \times B \mid a \in A \}.$$

(c) Finally, show that there is a functor $C: \mathbf{Rel}^{op} \to \mathbf{Rel}$ taking each relation $R \subseteq A \times B$ to its converse $R^c \subseteq A \times B$, where,

$$\langle a, b \rangle \in R^c \Leftrightarrow \langle b, a \rangle \in R.$$

Answer:

(a) To show that definition of identity when composed with an arbtirary morphism gives back the same morphism.

$$\begin{split} R &= \{\langle b, a \rangle \in B \times A\} \\ I &= \{\langle a, a \rangle \in A \times A | a \in A\} \\ R \circ I &= \{\langle b, a \rangle \in B \times A \mid \exists a (\langle b, a \rangle \in B \times A \ \& \ \langle a, a \rangle \in A \times A)\} \\ &= I \end{split}$$

(b) Identity and composition are already defined. To show that composition is associative.

$$LHS = S \circ (R \circ Q)$$

$$RHS = (S \circ R) \circ Q$$

- 2. **Question:** Consider the following isomorphisms of categories and determine which hold.
 - (a) $\mathbf{Ref} \cong \mathbf{Ref}^{op}$
 - (b) $\mathbf{Sets} \cong \mathbf{Sets}^{op}$
 - (c) For a fixed set X with powerset P(X), as poset categories $P(X) \cong P(X)^{op}$ (the arrows in P(X) are subset inclusions $A \subseteq B$ for all $A, B \subseteq X$).

Answer:

3. Question:

- (a) Show that in **Sets**, the isomorphisms are exactly the bijections.
- (b) Show that in **Monoids**, the isomorphisms are exactly the bijective homomorphisms.
- (c) Show that in **Posets**, the isomorphisms are *not* the same as bijective homomorphisms.

Answer:

4. **Question:** Let X be a topological space and preorder the points by specializatoin: $x \leq y$ iff y is contained in every open set that contains x. Show that this is a preorder, and that it is a poset if X is T_0 (for any two distinct points, there is some open set containing one but not the other). Show that the ordering is trivial is X is T_1 (for any two distinct points, each is contained in an open set not containing the other).

Answer:

5. Question:

For any category \mathbf{C} , define a functor $U: \mathbf{C}/C \to C$ from the slice category over an object C that "forgets about C." Find a functor $F: \mathbf{C}/C \to C^{\to}$ to the arrow category such that $\mathbf{dom} \circ F = U$

Answer:

6. Question:

Construct the "coslice category" C/\mathbf{C} of a category \mathbf{C} under an object C from the slice category \mathbf{C}/C and the "dual category" operation $-^{op}$.

Answer:

7. Question:

Let 2 = a, b be any set with exactly 2 elements a and b. Define a functor $F: \mathbf{Sets}/2 \to \mathbf{Sets} \times \mathbf{Sets}$ with $F(f: X \to 2) = (f^{-1}(a), f^{-1}(b))$. Is this an isomorphism of categories? What about the analogous situation with an one-element set 1 = a instead of 2?

Answer:

8. Question:

Any category \mathbf{C} determines a preorder $P(\mathbf{C})$ by defining a binary relation \leq on the objects by

 $A \leq B$ if and only if there is an arrow $A \to B$

Show that P determines a functor from categories to preorders, by defining its effect on functors between categories and checking the required conditions. Show that P is a (one-sided) inverse to the evident inclusion functor of preorders into categories.

Answer:

9. Question:

Describe the free categories, on the following graphs by determining their objects, arrows, and composition operations.

- (a) $a \xrightarrow{e} b$
- (b) $a \xrightarrow{e} b$
- $(c) \quad a \xrightarrow{\varepsilon} b \\ \downarrow f \\ c$

$$(d) \quad a \xrightarrow{e} b \\ \downarrow f \\ c$$

Answer:

10. Question:

How many free categories on graphs are there which have exactly six arrow? Draw the graphs that generate these categories.

Answer:

11. Question:

Show that the free monoid functor

$$M:\mathbf{Sets} o \mathbf{Mon}$$

(a) Assume the particular choice $M(X) = X^*$ and define its effect

$$M(f): M(A) \to M(B)$$

on a function $f:A\to B$ to be

$$M(f)(a_1...a_k) = f(a_1)...f(a_k), \quad a_1,...a_k \in A$$

(b) Assume only the UMP of the free monoid and use it to determine M on functions, showing the result to be a functor.

Reflect on how these two approaches are related

Answer:

12. Question:

Verify the UMP for free categories on graphs, defined as above with arrows being sequences of edges. Specifically, let $\mathbf{C}(G)$ be the free category on the graph G, so defined, and $i:G\to U(\mathbf{C}(G))$ the graph homomorphism taking vertices and edges to themselves, regarded as objects and arrows

in ${\bf C}(G)$. Show that for any category ${\bf D}$ and graph homomorphism $f:G\to U({\bf D})$, there is a unique functor

$$\overline{h}: \mathbf{C}(G) \to D$$

with

$$U(\overline{(h)}) \circ i = h,$$

where $U: \mathbf{Cat} \to \mathbf{Graph}$ is the underlying graph functor.

Answer:

13. Question:

Use the Cayley representation to show that every small category is isomorphic to a "concrete" one, that is, one in which the objects are sets and the arrows are functions between them.

Answer:

14. Question:

The notion of a category can also be defined with just one sort (arrows) rather than two (arrows and objects); the domains and codomains are taken to be certain *arrows* that act as units under composition, which is partially defined. Read about this definition in section I.1 of Mac Lane's *Categories for the Working Mathematician*, and do the exercise mentioned there, showing that it is equivalent to the usual definition.

Answer:

2 Ch2 Abstract Structures

1. Question:

Show that a function between sets is an epimorphism if and only if it is surjective. Conclude that the isos in **Sets** are exactly the epi-monos.

Answer:

2. Question:

Show that in a poset category, all arrows are both monic and epic.

Answer:

3. Question:

(Inverses are unique.) If an arrow $f:A\to B$ has inverses $g,g':B\to A$ (i.e., $g\circ f=1_A$ and $f\circ g=1_B$, and similarly for g'), then g=g'.

Answer:

4. Question:

With regard to a commutative triangle,

$$\begin{array}{cccc} A & -f \rightarrow B \\ & & \stackrel{|}{\searrow} & \stackrel{|}{\searrow} & \\ & & C \end{array}$$

in any category \mathbf{C} , show

- (a) if f and g are isos (resp. monos, resp. epis), so is h;
- (b) if h is monic, so is f;
- (c) if h is epice, so is g;
- (d) (by example) if h is monic, g need not be.

Answer:

5. Question:

Show that the following are equivalent for an arrow

$$f:A\to B$$

in any category:

- (a) f is an isomorphism.
- (b) f is both a mono and a split epi.
- (c) f is both a split mono and an epi.

(d) f is both a split mono and a split epi.

Answer:

6. Question:

Show that a homomorphism $h:G\to H$ of graphs is monic just if it is injective on both edges and vertices.

Answer:

7. Question:

Show that in any category, any retract of a projective object is also projective.

Answer:

8. Question:

Show that all sets are projective (use the axiom of choice).

Answer:

9. Question:

Show that the epis among posets are the surjections (on elements), and that the one-element poset 1 is projective.

Answer:

10. Question:

Show that sets, regarded as discrete posets, are projective in the category of posets (use the foregoing exercises). Give an example of a poset that is not projective. Show that every projective poset is discrete, that is, a set.

conclude that **Sets** is (isomorphic to) the "full subcategory" of projectives in **Pos**, conisting of all projective posets and all monotone maps between them.

Answer:

11. Question:

Let A be a set. Define an A-monoid to be a monoid M equipped with a function $m:A\to U(M)$ (to the underlying set of M). A morphism $h:(M,m)\to (N,n)$ of A-monoids is to be a monoid homomorphism $h:M\to N$ such that $U(h)\circ m=n$ (a commutative triangle). Together with the evident identities and composites, this defines a category A-Mon of A-monoids.

Show that an initial object in A-Mon is the same thing as a free monoid M(A) on A. (Hint: compare their respective UMPs.)

Answer:

12. Question:

Show that for any Boolean algebra B, Boolean homomorphisms $h: B \to \mathbf{2}$ correspond exactly to ultrafilters in B.

Answer:

13. Question:

In any category with binary products, show directly that

$$A \times (B \times C) \cong (A \times B) \times C$$

Answer:

14. Question:

- (a) For any index set I, define the product $\prod_{i \in I} X_i$ of an I-indexed family of objects $(X_i)_{i \in I}$ in a category, by giving a UMP generalizing that for binary products (the case I = 2).
- (b) Show that in **Sets**, for any set X the set X^I of all functions $f:I\to X$ has this UMP, with respect to the "constant family" where $X_i=X$ for all $i\in I$, and thus

$$X^I \cong \prod_{i \in I} X$$

Answer:

15. Question:

Given a category \mathbb{C} with objects A and B, define the category $\mathbb{C}_{A,B}$ to have objects (X, x_1, x_2) where $x_1 : X \to A, x_2 : X \to B$, and with arrows $f : (X, x_1, x_2) \to (Y, y_1, y_2)$ being arrows $f : X \to Y$ with $y_1 \circ f = x_1$ and $y_2 \circ f = x_2$.

Show that $C_{A,B}$ has a terminal object just in case A and B have a product in ${\bf C}$.

Answer:

16. Question:

In the category of types $\mathbf{C}(\lambda)$ of the λ -calculus, determine the product functor $A, B \mapsto A \times B$ explicitly. Also show that, for any fixed type A, there is a functor $A \to (-) : \mathbf{C}(\lambda) \to \mathbf{C}(\lambda)$, taking any type X to $A \to X$.

Answer:

17. Question:

In any category **C** with products, define the *graph* of an arrow $f:A\to B$ to be the monomorphism

$$\Gamma(f) = \langle 1_A, f \rangle : A \rightarrowtail A \times B$$

(Why is this monic?). Show that for $C = \mathbf{Sets}$ this determines a functor $\Gamma : \mathbf{Sets} \to \mathbf{Rel}$ to the category \mathbf{Rel} of relations, as defined in the exercises

to Chapter 1. (To get an actual relation $R(f)\subseteq A\times B$, take the image of $\Gamma(f):A\rightarrowtail A\times B$.)

Answer:

18. Question:

Show that the forgetful functor $U:\mathbf{Mon}\to\mathbf{Sets}$ from monoids to sets is representable. Infer that U preserves all (small) products.

Answer: