D Kommunikation & Rechnernetze

- D1 Einführung und Motivation
- D2 Technischer Überblick
- D3 Lokale Rechnernetze
- D4 Einige Gemeinsamkeiten von Betriebssystemen und Rechnernetzen: Architekturmodelle und Diensthierarchien

D2 aus: Coulouris/Dollimore/Kindberg: "Distributed Systems - Concepts and Design", 5. Ausg., Pearson, 2012



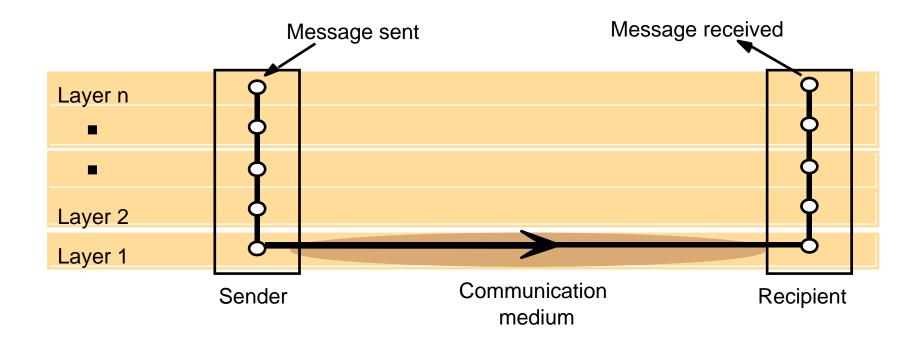
Communication Systems: Network types

	Range	Bandwidth (Mbps)	Latency (ms)
LAN	1-2 kms	10-1000	1-10
WAN	worldwide	0.010-600	100-500
MAN	2-50 kms	1-150	10
Wireless LAN	0.15-1.5 km	2-11	5-20
Wireless WAN	worldwide	0.010-2	100-500
Internet	worldwide	0.05-600	100-500
WPAN (Bluetooth)) 10-30m	0.5-2	5-20
WiFi	0.15-1.5 km	11-108	5-20
WiMAX	5-50 km	1.5-20	5-20
WWAN/GSM3	G worldwide	348-14.4	100-500
LTE	worldwide	100-300	• • •



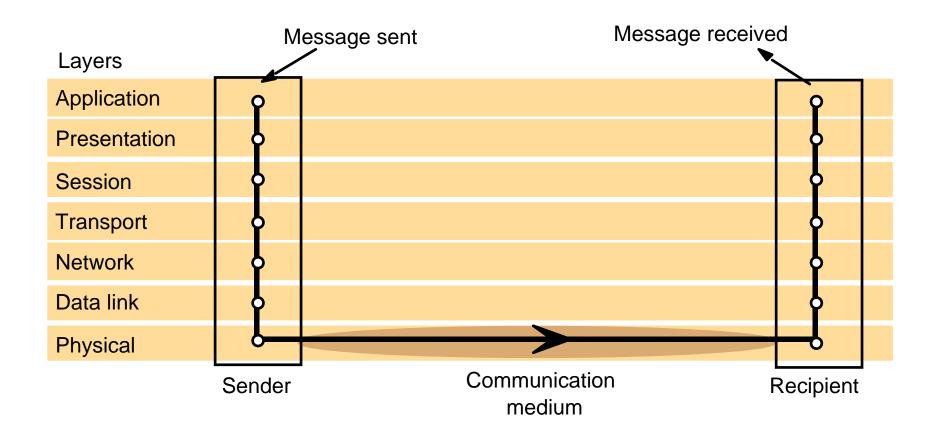
(Stand: ab ca. 2000+)

Principles of Communication Networks (1): Conceptual layering of protocol software





Example: Protocol Layers in the ISO Open Systems Interconnection (ISO/OSI) Model Example



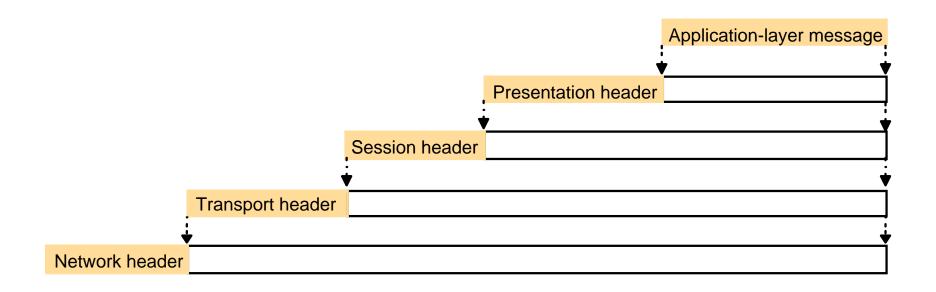


ISO/OSI Protocol Layer Summary

Layer	Description	Examples
Application	Protocols that are designed to meet the communication requirements of specific applications, often defining the interface to a service.	HTTP, FTP, SMTP, CORBA IIOP
Presentation	Protocols at this level transmit data in a network representation that is independent of the representations used in individual computers, which may differ. Encryption is also performed in this layer, if required.	Secure Sockets (SSL),CORBA Data Rep.
Session	At this level reliability and adaptation are performed, such as detection of failures and automatic recovery.	
Transport	This is the lowest level at which messages (rather than packets) are handled. Messages are addressed to communication ports attached to processes, Protocols in this layer may be connection-oriented or connectionless.	TCP, UDP
Network	Transfers data packets between computers in a specific network. In a WAN or an internetwork this involves the generation of a route passing through routers. In a single LAN no routing is required.	IP, ATM virtual circuits
Data link	Responsible for transmission of packets between nodes that are directly connected by a physical link. In a WAN transmission is between pairs of routers or between routers and hosts. In a LAN it is between any pair of hosts.	Ethernet MAC, ATM cell transfer, PPP
Physical	The circuits and hardware that drive the network. It transmits sequences of binary data by analogue signalling, using amplitude or frequency modulation of electrical signals (on cable circuits), light signals (on fibre optic circuits) or other electromagnetic signals (on radio and microwave circuits).	Ethernet base- band signalling, ISDN

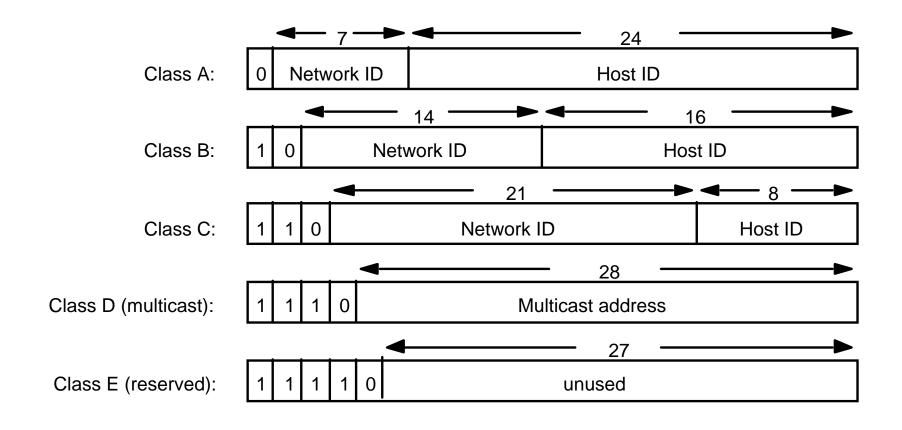


Layering of Protocol Data Units – Encapsulation as it is applied in layered protocols



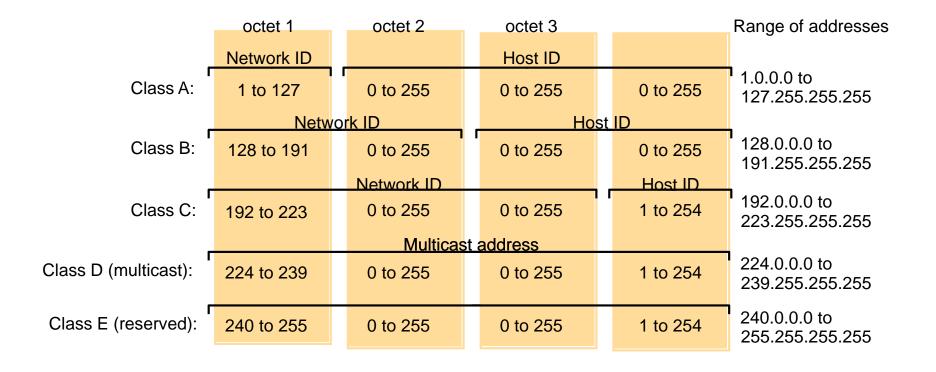


Principles of Communication Networks (2): *Addressing* – Internet address structure, showing field sizes in bits



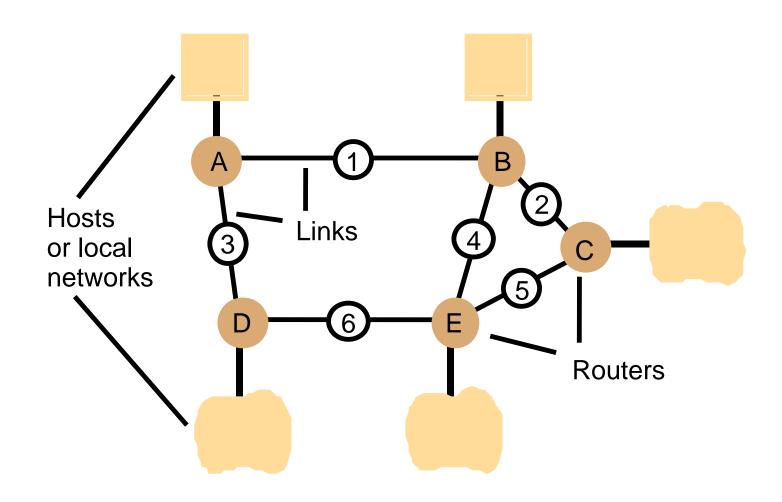


Decimal representation of Internet addresses





Principles of Communication Networks (3): Routing in a wide area network





Routing tables for the network in previous foil

Routings from A			
To	Link	Cost	
A	local	0	
В	1	1	
\mathbf{C}	1	2	
D	3	1	
_ E	1	2	

Routings from B			
То	Link	Cost	
A	1	1	
В	local	0	
C	2	1	
D	1	2	
E	4	1	

Routings from C				
То	Link	Cost		
A	2	2		
В	2	1		
\mathbf{C}	local	0		
D	5	2		
E	5	1		

Routings from D				
То	Link	Cost		
A	3	1		
В	3	2		
C	6	2		
D	local	0		
E	6	1		

Routings from E				
То	Link	Cost		
A	4	2		
В	4	1		
C	5	1		
D	6	1		
Е	local	0		



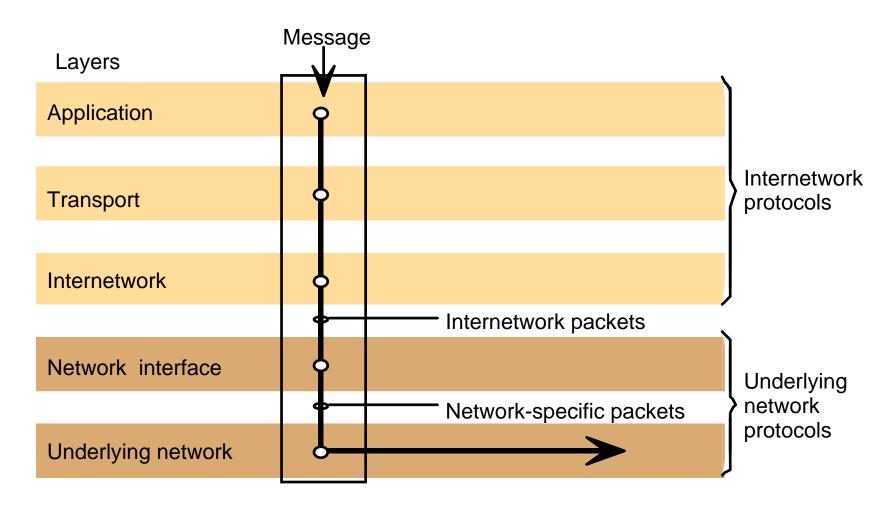
Example routing algorithm (Router Info Protocol, RIP)

Send: Each t seconds or when local table Tl changes, send Tl on each non-faulty outgoing link.

```
Receive: Whenever a routing table Tr is received on link n:
for all rows Rr in Tr {
     if (Rr.link \neq n) {
          Rr.cost = Rr.cost + 1;
          Rr.link = n;
          if (Rr.destination is not in Tl) add Rr to Tl;
           // add new destination to Tl
          else for all rows Rl in Tl {
               if (Rr.destination = Rl.destination) and
                         (Rr.cost < Rl.cost \text{ or } Rl.link = n)) Rl = Rr;
               // Rr.cost < Rl.cost : remote node has better route
               // Rl.link = n: remote node is more authoritative
```

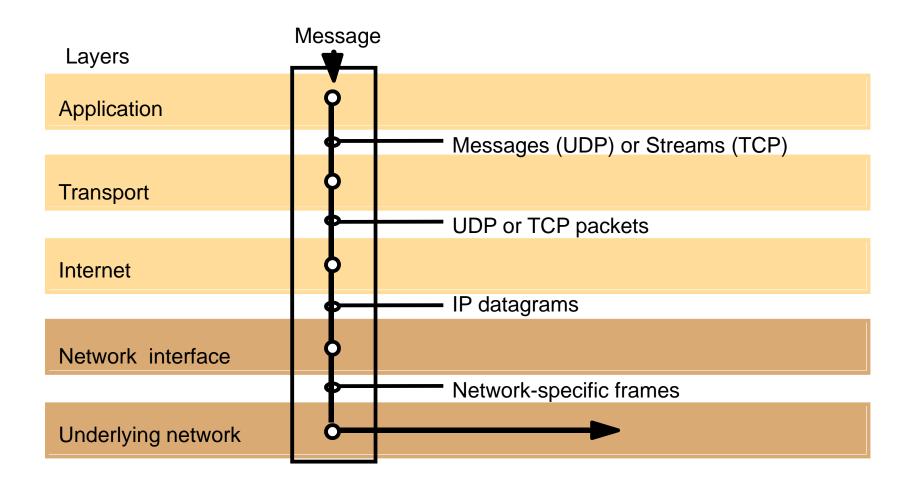


Principles of Communication Networks (4): Simplified Internetwork Layers



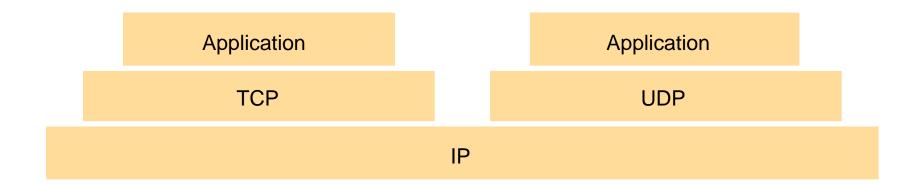


Example: TCP/IP Protocol Layers



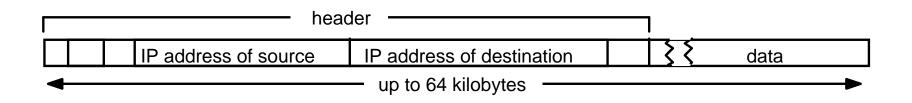


The programmer's conceptual view of a TCP/IP Internet



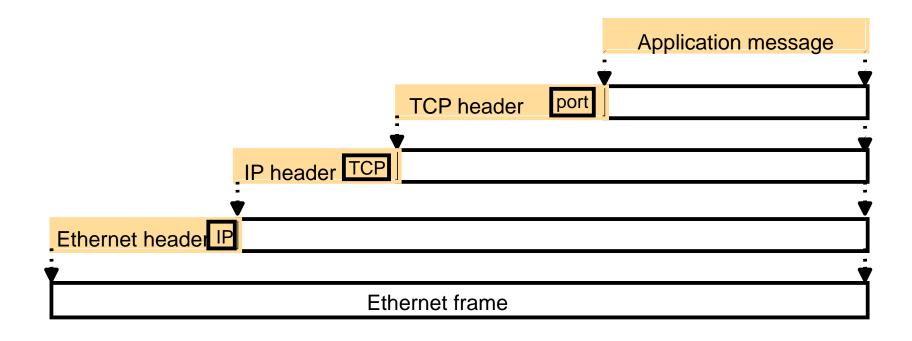


IP packet layout



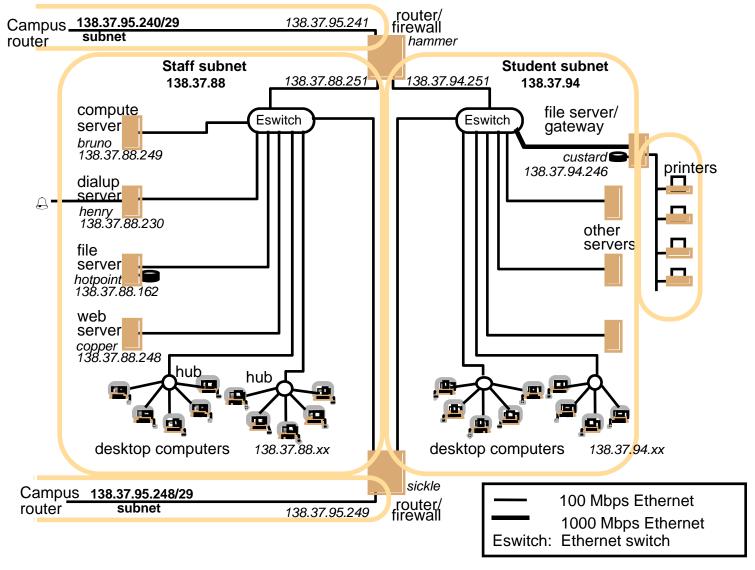


Message encapsulation in data packets transmitted via TCP over an Ethernet





Example: Simplified View of a Computer Science Department Network



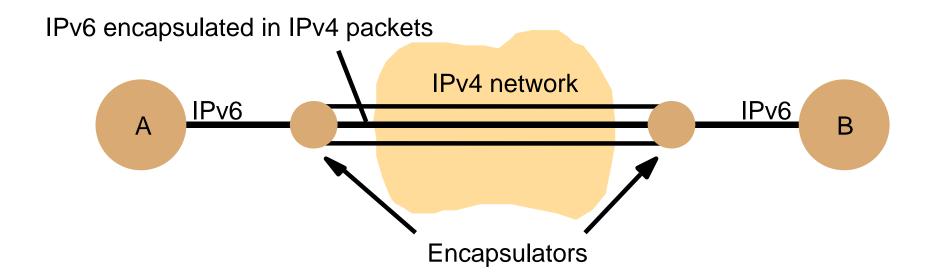


Principles of Communication Networks (5): Extension A: *IPv6 – Header Layout*

Version (4 bits)	Priority (4 bits)	Flow label (24 bits)		
Payload len	gth (16 bits)		Next header (8 bits)	Hop limit (8 bits)
Source address (128 bits)				
Destination address (128 bits)				



Tunnelling for IPv6 migration

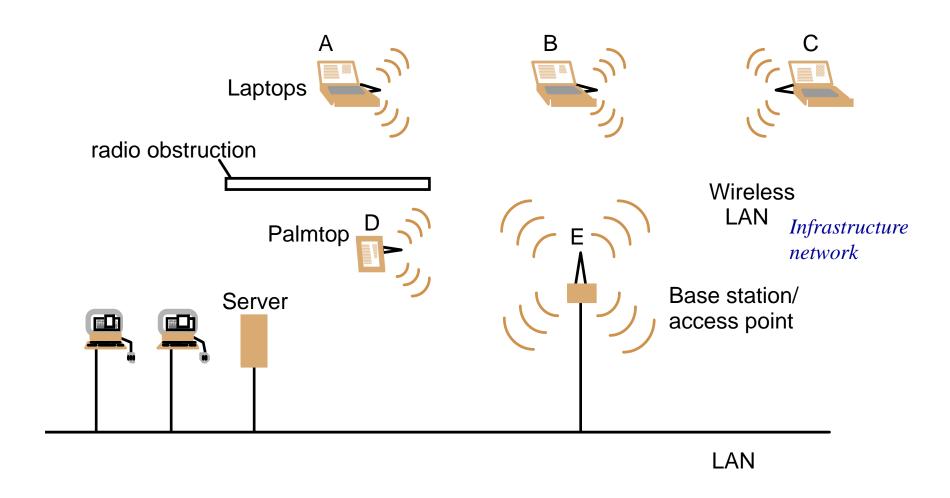


IPv4: seit 1984 – Adressengpass bei 2³² (ca. 4,3 Mrd) Adressen, ab 2011 akut

IPv6: 10 Jahre Einführung – max. 2¹²⁸ Adressen (ca.3,4*10³⁸)

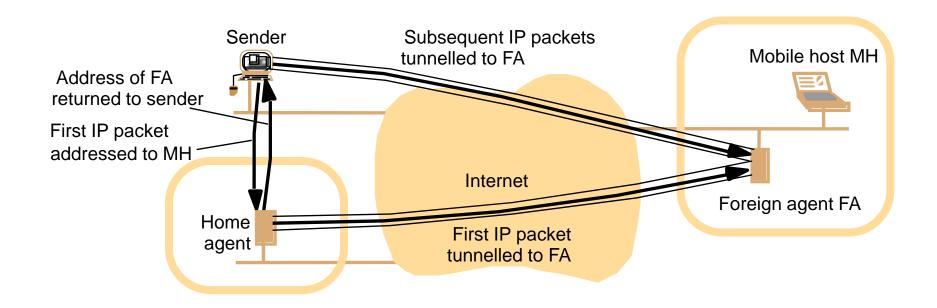


Extension B: Wireless LAN configuration





MobileIP routing mechanism





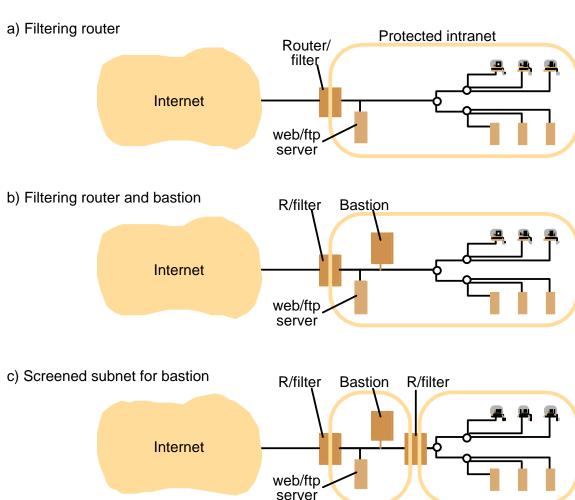
Extension C: Basic Security – Firewall configurations

Firewall goals:

- service control
- behaviour control
- user control

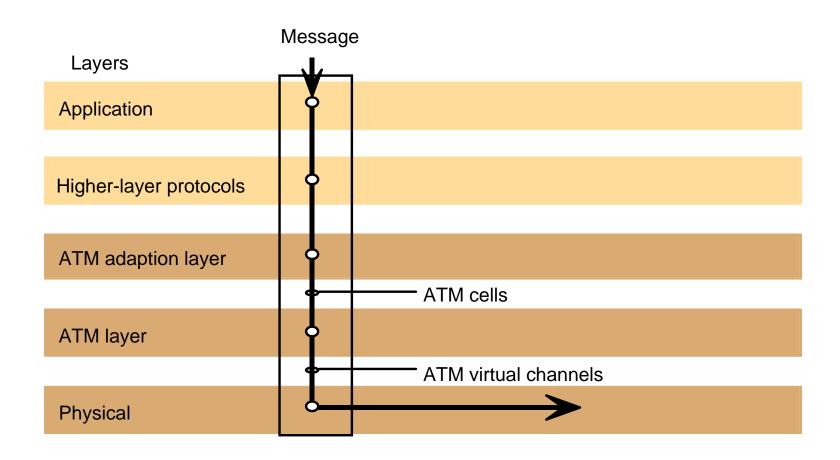
A **bastion host** is a special purpose computer on a network specifically designed and configured to withstand attacks.

The computer generally hosts a single application, for example a proxy server, and all other services are removed or limited to reduce the threat to the computer. It is hardened in this manner primarily due to its location and purpose, which is either on the outside of the firewall or in the DMZ and usually involves access from untrusted networks or computers.



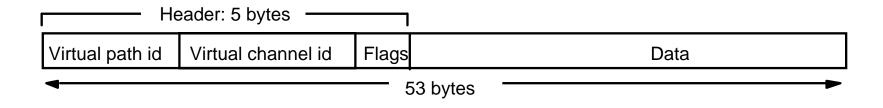


Extension D: Making Communication *faster***: ATM protocol layers**



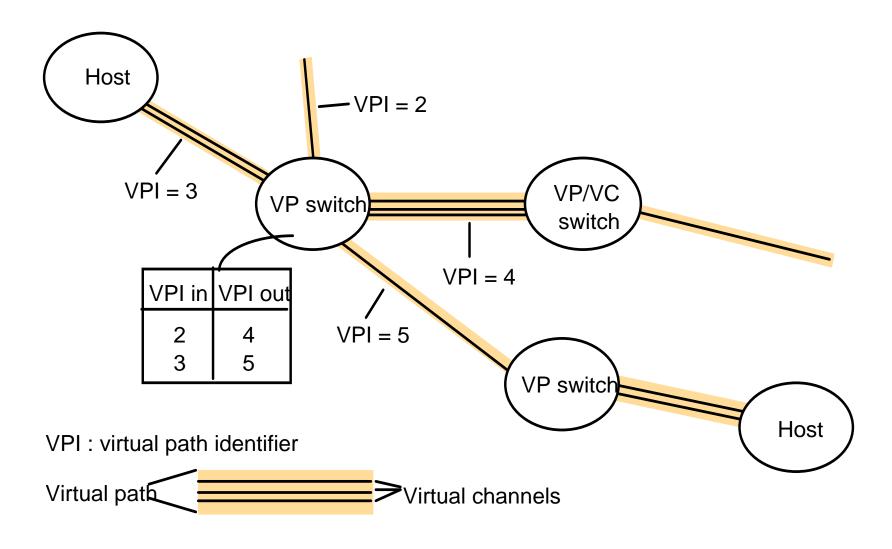


ATM cell layout





Switching virtual paths in an ATM network





Standardisation: IEEE 802.x local network standards

	IEEE No.	Name	Title	Reference
	802.3	Ethernet	CSMA/CD Networks (Ethernet)	[IEEE 1985a]
	802.4		Token Bus Networks	[IEEE 1985b]
	802.5		Token Ring Networks	[IEEE 1985c]
	802.6		Metropolitan Area Networks	[IEEE 1994]
	802.11	WiFi	Wireless Local Area Networks	[IEEE 1999]
	802.15.1	Bluetooth	Wireless Personal Area Networks	[IEEE 2002]
	802.15.4	ZigBee	Wireless Sensor Networks	[IEEE 2003]
_	802.16	WiMAX	Wireless Metropolitan Area Networks	[IEEE 2004a]

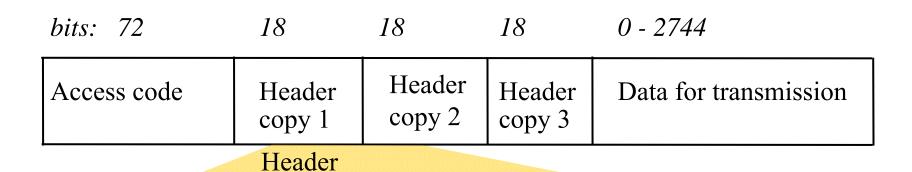


Ethernet ranges and speeds

	10Base5	10BaseT	100BaseT	1000BaseT
Data rate:	10 Mbps	10 Mbps	100 Mbps	1000 Mbps
Max. segment lengths:				
Twisted wire (UTP)	100 m	100 m	100 m	25 m
Coaxial cable (STP)	500 m	500 m	500 m	25 m
Multi-mode fibre	2000 m	2000 m	500 m	500 m
Mono-mode fibre	25000 m	25000 m	20000 m	2000 m



Bluetooth frame structure



bits: 3 1 1 1 4 8

Destination	 Flow	Ack	Seq	Type	Header checksum
Destination	110 **	1 IOK		1990	Treader encekbann

SCO packets (e.g. for voice data) have a 240-bit payload containing 80 bits of data triplicated, filling exactly one timeslot.

