

Binary search

Recursive binary search.

$$\text{Eq. } \begin{bmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 3 & 4 & 6 & 7 & 9 & 12 & 16 & 17 \end{bmatrix} \quad n=8 \quad \text{target} = 18.$$

f (arr, low, high) <

IF (low > high) ↴ Base
return -1; Case

`mid = (low + high) / 2;` → odd way

If (a const) == target)

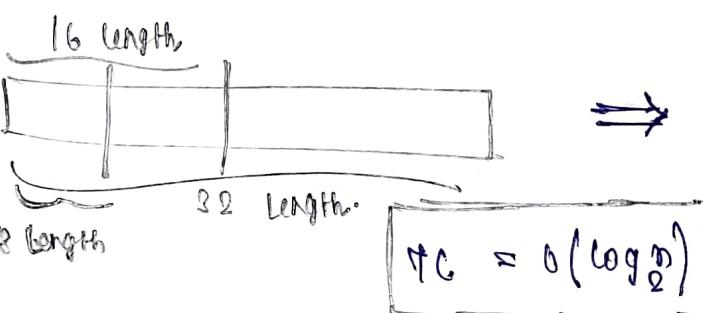
return mid;

else if (target > a[mid])

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return p (arr, mid + 1, high, target);
```

else

```
return & (arr, low, mid - 1, target);
```



$$B \cap n \neq \emptyset$$

$$64 = 2^6$$

MW overflow cases.

A diagram of a binary search tree with 5 nodes. The root node contains the value 'how'. The left child of the root contains 'low' and the right child contains 'high'. The left child of 'low' is '1' and its right child is '2'. The right child of 'high' is '3' and its left child is '4'. The right child of '3' is '5'. Arrows point from the labels 'height' and 'MAX-KEY' to the rightmost node '5'.

$$mid = (low + high) / 2;$$

YNY-MAN

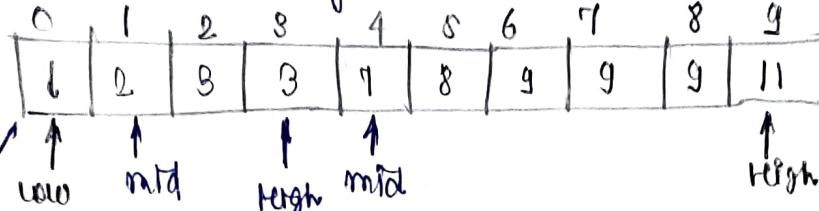
$$\text{mid} = \text{low} + \frac{(\text{high} - \text{low})}{2}$$

Q2 Lower Bound \rightarrow smallest index such that $arr[low] \geq n$

$$arr[7] = [8 \ 5 \ 2 \ B \ 4 \ 1 \ 5 \ 1 \ 9] \quad n=5$$

$$\begin{array}{c|c|c|c|c} x=8 & x=9 & x=19 & x=16 & x=20 \\ lb=2 & lb=3 & lb=4 & lb=4 & lb=\infty \end{array}$$

If no. is repeating so we return lower index.

Eg. 

$$x=1$$

$$ans = 16 / 4 \ \text{Ans} \quad \stackrel{\text{Eq-1}}{mid} = 0 + 9 / 2 = 4$$

If no element possible $a[0] = 7 > 4$ yes possible.

$$mid = 0 + 3 / 2 = 1$$

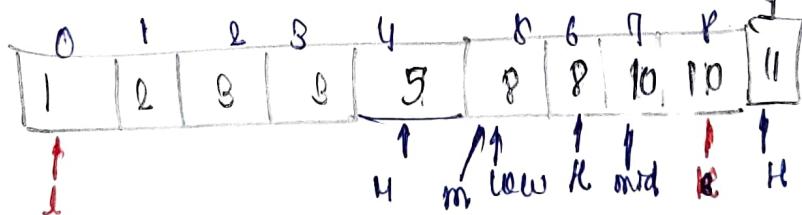
$a[1] = 2 > 1$ yes possible.

$$mid = 0 + 0 / 2 = 0$$

$a[0] = 1 = 1$ yes possible.

Now, $low = 0 \ high = -ve$

$x=9$ So, loop end so $ans \Rightarrow 0$ Ans.

Eg-2 

$$ans = 10 \ \text{Ans}$$

$$mid = 0 + 9 / 2 = 4 \quad a[4] = 5 < 9 \quad \text{No.}$$

$$mid = 8 + 9 / 2 = 9 \quad a[9] = 10 > 9 \quad \text{Yes.}$$

$$mid = 5 + 6 / 2 = 8 \quad a[5] = 8 < 9 \quad \text{No.}$$

$$mid = 6 + 6 / 2 = 6 \quad a[6] = 8 < 9 \quad \text{No.}$$

Now, search space ended so $ans = 9$ Ans

→ Upper bound → smallest index s.t $\text{arr}[md] > n$

$\text{arr}[] = \begin{matrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 \end{matrix}$

$x = 6$	$x = 11$	$n = 12$	$n = 18$	$x = 0$
$UB = 3$	$UB = 9$	$UB = 10$	$LB = 10$	$UB = 0$

→ Search insert position.

$\text{arr}[] = [1 \ 2 \ 4 \ 7]$

same as lower bound.

→ Floor and ceil in sorted array

largest no
 n array $\leq x$

smallest no
in array $>= x$

• $\text{arr}[] = [10 \ 20 \ 30 \ 40 \ 50]$ | • $\text{arr}[] = [10, 20, 28, 30, 40]$

$x = 25$

$\text{Floor} = 20$

$\text{ceil} = 30$.

$x = 28$

$\text{Floor} = 20$

$\text{ceil} = 30$.

$\text{Floor}(\text{arr}, n) \leftarrow$

$\text{ans} = -1$

$\text{low} = 0$

$\text{high} = n-1$

while ($\text{low} \leq \text{high}$) {

$\text{mid} = (\text{low} + \text{high}) / 2$;

if ($\text{arr}[\text{mid}] \leq \text{mid}$) {

$\text{ans} = \text{arr}[\text{mid}]$;

$\text{low} = \text{mid} + 1$;

else {

$\text{high} = \text{mid} - 1$;

} return ans ;

Q3 Find first and last occurrences of x

0	1	2	3	4	5	6	7
2	4	6	8	7	8	11	13

$$x = 8$$

$$\text{first} = 3$$

$$\text{last} = 5$$

$$x = 11$$

$$\text{first} = 6$$

$$\text{last} = 6$$

We lower bound $\rightarrow a[\text{first}] >= n$

and upper bound $\rightarrow a[\text{last}] > n$

$$\text{so } \lfloor b(8) \rfloor = 8 \text{ make } x = 3$$

$$\text{UB}(8) = 8 \text{ make } x = 6 - 1$$

$$= 5$$

$$TC = O(\log_2 n)$$

$$SC = O(1)$$

Hand wall

2	8	8	1	8	8	11	13
---	---	---	---	---	---	----	----

Here we use simple
BS for first and last
occurrences and return it.

but it will not work for all.

$$\text{so, } \lfloor b(10) \rfloor = \underline{\text{index}} = 6 \quad X$$

$$\text{so, } \lfloor b(14) \rfloor = \underline{\text{index}} = 8 \quad X$$

$$\text{so, if } (\lfloor b(x) \rfloor = n \text{ || arr}[\lfloor b \rfloor] \neq x)$$

so, no. not present.

Q4 search in rotated sorted array ↴

arr = [7 8 9 1 2 3 4 5 6] target $\rightarrow 8$

use OR binary search (because OR search + sorted).

1	2	3	4	5
↑	↑	↑		
1	m	high		

$$\text{so, } l < g$$

$$\underline{\text{so, }} 0 + 8/2 = 4 \quad a[4] = 2 \leq \text{target}$$

$$\text{so, } a[\text{low}] = a[0] = 7 \quad \underline{\text{so, }} l > 1 \quad \text{left}$$

$$a[\text{mid}] = a[4] = 2$$

not sorted.

\rightarrow yaha pe done

left ko check kro

right ek ko check

karne se nahi

hoga.

\rightarrow so, we eliminate

right because it

is sorted.

$$\text{so, } \text{low} = 0$$

$$\text{high} = \text{mid} - 1 = 3$$

again do these things.

15 search in rotated sorted array II

$\text{arr} = [9, 8, 1, 2, 3, 3, 3, 4, 8, 6]$ Identify sorted half

target = 3

Eg. $\begin{bmatrix} 6 & 7 & 1 & 2 & 3 & 4 & 4 & 5 \end{bmatrix}$ Ag $\begin{bmatrix} 4 & 5 & 8 & 2 & 3 & 3 & 4 & 4 \end{bmatrix}$

$\begin{array}{c} l \\ m \\ h \end{array}$ Yes sorted

$\begin{bmatrix} 4 & 5 & 8 & 2 & 3 & 3 & 4 & 4 \end{bmatrix}$

$\begin{array}{c} l \\ m \\ h \end{array}$ sorted

try to bring down this condition.

$$\text{arr}[l] \leq \text{arr}[m] \leq \text{arr}[h]$$

problem.

from search space

$$T.C = O(\log_2 n)$$

$\rightarrow \begin{bmatrix} 3 & 3 & 1 & 3 & 3 & 3 & 3 \end{bmatrix}$

$\begin{array}{c} l \\ \downarrow \\ \text{low} \end{array} \quad \begin{array}{c} \uparrow \\ \text{mid} \end{array} \quad \begin{array}{c} \uparrow \\ \text{high} \end{array}$

so,

$\begin{bmatrix} 8 & 1 & 2 & 3 & 3 & 3 & 3 \end{bmatrix}$

$\begin{array}{c} \uparrow \\ \text{low} \end{array} \quad \begin{array}{c} \uparrow \\ \text{mid} \end{array} \quad \begin{array}{c} \uparrow \\ \text{high} \end{array}$

again we shrink.

Here, what pt search space half
will say.

Minimum in rotated sorted array (unique).

$\begin{bmatrix} 0 & 1 & 2 & 8 & 4 & 5 & 6 \end{bmatrix}$

$\begin{bmatrix} 4 & 8 & 6 & 7 & 0 & 1 & 2 \end{bmatrix}$

$\begin{array}{c} \downarrow \\ \text{low} \end{array} \quad \begin{array}{c} \uparrow \\ \text{mid} \end{array} \quad \begin{array}{c} \uparrow \\ \text{high} \end{array}$

return 0.

$\begin{array}{c} l \\ \text{sorted} \end{array} \quad \begin{array}{c} m \\ \text{not sorted} \end{array}$

best answer
is here.

→ eliminates left/right.
→ identify sorted half.

Eg

$\begin{bmatrix} 4 & 8 & 1 & 2 & 3 \end{bmatrix}$

$\begin{array}{c} l \\ m \\ h \end{array}$

Here right half
contains ans.

Eg.

9	8	1	2	3	4	5	6
↓	↓	↓	↓		↑		
1	m	h	on		h		

ans = ~~INT MAX~~

①

done,

So, now

sorted

not sorted

① ↗ low
high
mid

Q7 How many times has an array been visited here slight variation in find min code and manipulated with index.

keep track of index only (Kthm qw).

Q8 single element in sorted array

Eg. [1 1 2 2 3 3 4 8 8 6 6]. → elimination will be done.



$$T_C = O(\log_2 n)$$

element

- (8, 0) → element in right half
- (0, 8) → element in left half.

Q9 Find peak element.

$$\text{arr}[] = [1 2 8 4 8 6 9 8 5 1]$$

$$\text{ans} = 8$$

$$\text{arr}[i-1] < \text{arr}[i] > \text{arr}[i+1]$$

$$\text{Eg. } [1 2 1 3 8 6 4]$$

$$\text{ans} = 2$$

$$\text{Eg. } [8 4 3 2 1] \rightarrow \infty$$

$$\text{ans} = 5$$

$$\text{Eg. } \rightarrow \infty [1 2 8 4 8] \rightarrow \infty$$

$$\text{ans} = 5$$

assume away has peaked.

$$a_m = \begin{bmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 1 \end{bmatrix}$$

$$\text{Ans} = \underline{8}. \quad \underline{\text{Relation 4}} \quad - \quad l = 0 \quad | \quad m = 4 \\ h = 9 \quad | \quad \text{ans}[9] = 5 \quad | \quad \text{check for } 6 \\ \text{anti hal peak.}$$

so left me chafe job

$$\text{14-2} \quad - \quad l=5 \quad | \quad m=7 \\ h=9 \quad | \quad \text{ans } [7] = 8 \quad | \quad \text{yes peak. done.}$$

e.g. $\begin{bmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 1 & 10 & 13 & 9 & 6 & 8 & 4 & 2 & 1 & 0 \end{bmatrix}$

$\frac{11-1}{2} - l=0$	$m=4$	not peak	so, peak on the
$h = g$	$am + b \approx 6$	$7 > 6 > 5$ decreasing	left.

$$\text{II-11} = l=0 \quad m=j \\ h=3 \quad am[1]=10$$

not peak.
1 < 10 < 18
Increasing

So, peak on right

$$\begin{array}{c|c|c} \text{---} & l=2 & m=2 \\ & h=3 & \text{arr}[2] = 18 \end{array} \quad \begin{array}{l} \text{yes peak} \\ \text{done.} \end{array}$$

What is first and last will be peak

$$\text{eg. } \begin{bmatrix} 1 & 2 & 3 & 4 & 5 \end{bmatrix} \quad \text{eg } \begin{bmatrix} 0 & 4 & 3 & 2 & 1 \end{bmatrix}$$

Hence Peak = 5. Hence Peak = 5.

$$\arctan \theta > \arctan (\theta + 1)$$

and $\text{arr}[\text{mid}] > \text{arr}[\text{mid} - 1]$

1

$$m = 0$$

- you help others directly return don't check for others

- Simple manually write code for 1st and last elements.

How to deal
with that.

→ Binary search on answers.

10 Find sqrt(N) | low = 1 | mid = 14
n = 28 | high = 28 | $14 \times 14 > 28$ X

Eq n = 28

ans = 5

| l = 1 | mid = 7
h = 18 | $7 \times 7 > 28$ X

ans = 7.5

| l = 1 | mid = 9
h = 6 | $9 \times 9 < 28$ ✓

so, ans = 9

| l = 4 | mid = 8
h = 6 | $8 \times 8 = 28 < 28$ ✓

| l = 6 | m = 6
h = 6 | $6 \times 6 > 28$ X

11 Find Nth root OR on integers

N = 9 | N = 3 | so, l = 1 | m = 14
m = 69 | M = 27 | h = 27 | $14 \times 14 \times 14 > 27$ X

so, l = 1 | m = 30 | so, l = 1 | m = 7
h = 69 | $8 \times 3 \times 3 \times 3 > 69$ X | h = 13 | $7 \times 7 \times 7 > 27$ X

so, l = 1 | m = 87
m = 84 | X

so, l = 1 | m = 8
h = 16 | X

so, l = 1 | m = 4 | so, l = 1 | m = 1
h = 7 | X | h = 4 | X

so, l = 1 | m = 2
h = 3 | $2 \times 2 \times 2 \times 2 = 16 < 69$ ✓

so, l = 3 | m = 8
h = 8 | $8 \times 3 \times 3 \times 3 > 69$ X

13

Koko eating bananas. → return the minimum integer 05

$$\text{Eg. } \text{M}(1, 9) = [3, 6, 9, 11]$$

$$h = 8$$

h such that koko can eat all bananas within h hours.
Here, k = banana / hr

If take 2 banana / hr.

So,

3	6	7	11
↓	↓	↓	↓
2	3	4	6

 $\rightarrow 2$

 $= 18 \text{ hr}$

So,

3	6	7	11
↓	↓	↓	↓
3	2	3	4

 $\rightarrow 3$

 $= 12 \text{ hr.}$

Eg. 2 banana / hr.

3	6	7	11
↓	↓	↓	↓
1	2	2	3

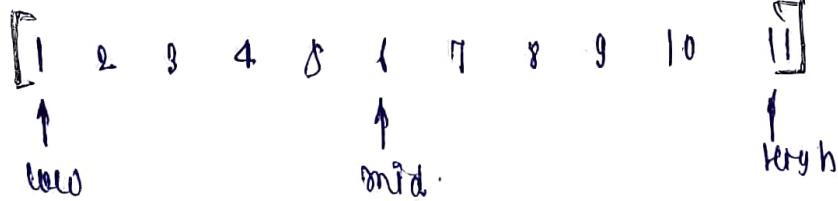
 $\rightarrow 4 \text{ banana / hr}$

 $= 8 \text{ hrs}$

So, $\text{ans} = 4$ hrs

Help us we binary search on answers.

So, answer will be 1 to 11 → max of array.



$$\text{ans} = \cancel{8} / \cancel{p}$$

So, $l = 1$, do, $1 + 1 + 2 + 2$
 $h = 11$ $= 6 < 8$ Yes Possible
 $m = 6$

So,

So, $l = 1$, do $1 + 2 + 3 + 4$
 $h = 8$ No
 $m = 8$ $= 10 > 8$

$$TC = O(N) \times \log_2 \max \{ \cdot \}$$

So, $l = 4$, do $1 + 2 + 2 + 3$, Yes dome.
 $h = 8$
 $m = 4$ $= 8 = 8$

14) minimum no. of days to make m bouquets.

$$\text{bloom day} = \{7, 9, 9, 11, 13, 11, 12, 7\} \quad m=2 \quad k=3$$

18th day $\underbrace{7 \ 9 \ 9}_{\text{no. of bouquets}}$ $\underbrace{11 \ 12 \ 7}_{\text{adjacent flowers required.}}$

no. of bouquets

18th day

17th day: X not possible.

$$\text{so, ans} = 15 \text{ (12.)}$$

12th day $\{v \ v \ v \ v \ X \ v \ v \ v\}$ Yes.

Yes done.

11th day $\{v \ v \ v \ v \ X \ X \ X \ v\}$ No X

eg. $\{1, 10, 3, 10, 2\} \quad m=3 \quad k=2$

10th day $\{v \ v \ v \ v \ v\}$ No 3 bouquets never possible
so, ans = -1.

Now $\boxed{7 \ | \ 9 \ | \ 7 \ | \ 7 \ | \ 13 \ | \ 11 \ | \ 12 \ | \ 7}$

So, as on
answer $\boxed{7 \ | \ 8 \ | \ 9 \ | \ 10 \ | \ 11 \ | \ 12 \ | \ 13}$

$l=7$ | No not possible.
 $h=13$
 $m=10$

$l=11$ | Yes
 $h=13$
 $m=12$

$l=11$ | No
 $h=11$
 $m=11$

$$\text{ans} = 12$$

$$TC = O(N \times \log_{\frac{1}{2}}(\max - \min + 1))$$

15 Find the smallest divisor given a threshold

$$\text{arr}[] = [1 \ 2 \ 3 \ 9] \quad \text{threshold} = 6.$$

Sol. $\frac{1}{4} \frac{2}{4} \frac{3}{4} \frac{9}{4}$

$$1 + 1 + 2 + 3 = 7 > 6$$

ans = 5

Sol. $\frac{1}{8} \frac{2}{8} \frac{3}{8} \frac{9}{8}$

$$1 + 1 + 1 + 2 = 5 < 6 \quad (\text{Yes})$$

16 capacity to ship packages within R days

$$\text{arr}[] = [1 \ 2 \ 3 \ 4 \ 8 \ 6 \ 7 \ 8 \ 9 \ 10] \quad \text{days} = 8.$$

ship will run once per day.

<u>capacity</u> <u>10</u>	day 1 → 1 2 3 4 = 10. → 1
	day 2 → 5 = 5 → 2
	day 3 → 6 = 6 → 4
	day 4 → 7
	day 5 → 8
	day 6 → 9
	day 7 → 10
total 7 days	

ans = 15

Now
capacity 10.
so, day - 1 = 1 2 3 4 5 }
day - 2 = 6 7 }
day - 3 = 8 }
day - 4 = 9 }
day - 5 = 10 }
} Yes
} 5 days
} completed

So,

min	max
Capacity	Capacity

= # max_element
of array

max capacity
of ship to be hold

be sum or all
array.

$$\text{so, Range} = [10, 58]$$

$so, l = 10$	yes	$arr[10] = 18 \%$	$so, l = 18$	yes
$h = 19$		18	$h = 19$	
$m = 17$			$m = 17$	
$so, l = 10$	X	$so, l = 18$	yes	
$h = 19$		$h = 16$		
$m = 14$		$m = 18$		

[17] Find k^{th} missing number

$arr[] = [8, 1, 2, 3, 4, 7, 11] \quad n=8.$

$so, 1 \underline{(2 \ 3 \ 4)} \underline{5} \underline{6} \underline{7} \underline{1} \boxed{9} \ 10 \ 11$

return 9

$$arr[7] = [4 \ 7 \ 9]$$

$$k = 8.$$

$so, \text{high is pointing}$
to \rightarrow index

$so, \text{not good way.}$

Always not good.

$arr[\text{high}] + \text{more.}$

$arr[\text{high}] + \underbrace{(k - \text{missing})}_{\text{now}}$

$arr[\text{high}] + k - (arr[\text{high}] - \text{high} + 1)$

$(arr[\text{high}] + k - arr[\text{high}] + \text{high} - 1)$

$\boxed{k + \text{high} - 1} \quad \underline{\text{Ans}}$

$so, l = 10$	$m = 18$	$h = 16$	<u>yes</u>
$m = 18$	$h = 16$	$l = 10$	
$h = 15$		$m = 18$	
$o \ 1 \ 2 \ 3$	$4 \ 7 \ 11$		
$2 \ 3$	4	7	
\downarrow	\downarrow	\downarrow	
$1 \ 1 \ 1$	$3 \ 6$		
1	2		
	h		

missing no.

$$so, l = 0 \quad | \quad m[2] = 1 \quad X$$

$$h = 4$$

$$m = 2$$

$$so, l = 3 \quad | \quad m[3] = 3 \quad X$$

$$h = 4$$

$$m = 3$$

$$so, l = 4 \quad | \quad m[4] = 6 \cdot 1 \quad X$$

$$h = 4$$

$$m = 4$$

$$l = 5 \quad | \quad \text{find no.}$$

$$h = 5$$

$$m = 5$$

$$arr[high] = 9$$

and missing no. = 3.

$$\text{diff} = 2.$$

$$so, \boxed{1 \rightarrow 2 \rightarrow 3}.$$

one

high + how many more
you need.

18 Aggregative Cores (min diff b/w cores) is maximum

$\text{arr}[] = [0 \ 3 \ 4 \ 7 \ 10 \ 9]$

cores = 4.

→ 0 3 4 7 9 10.

0 3 4 7 9 10 = 1

- see iterative part cores and how much cores we part is selected by binary search.

$$0_1 + 0_2 - 5 \ 0_3 + 0_4 = 1$$

$$0_1 + 0_2 + 0_3 - 3 \ 0_4 = 9$$

→ same as 1st type so ans = 8.

Magnetic force b/w two balls.

19 Allocate Books.

$$\text{arr}[] = [20 \ 46 \ 28 \ 49 \ 24]$$

Students = 4.

$$20 | 46 | 28 | 49 + 24. \rightarrow \text{max} = 73$$

$$20 | 46 | 28 + 49 | 24 \rightarrow \text{max} = 77$$

$$20 | 46 + 28 | 49 | 24 \rightarrow \text{max} = 74$$

We should start from 49 → max. element. (100)
 $\text{high} \approx \text{sum of array.}$

$$[49 \ \dots \ \dots \ \dots \ \dots \ \dots \ \dots \ 172]$$

$$\begin{array}{l|l} \text{Q. } l = 49 & 1 \rightarrow 20 \ 46 \ 28 \\ h = 172 & 2 \rightarrow 49 \ 24 \\ m = 110 & \end{array} \times \text{only 2 students.}$$

$$\begin{array}{l|l} \text{Q. } l = 49 & 1 \rightarrow 20 \ 46 \\ h = 109 & 2 \rightarrow 28 \ 49 \\ m = 78 & 8 \rightarrow 24 \end{array} \times \text{only 8 students}$$

$S_0, l = 49$	$l \rightarrow$	$\textcircled{10}$ students	$S_0, l = 70$	<u>4 students</u>
$h = 79$	$e \rightarrow$		$h = 79$	
$m = 68$	$3 \rightarrow$		$m = 73$	
	$4 \rightarrow$			$S_0, \text{ one } = \frac{73}{7} = 10 \dots 3$

$S_0, l = 62$	$l \rightarrow$	$\textcircled{8}$ students	$S_0, l = 70$	<u>4 students</u>
$h = 79$			$h = 72$	
$m = 69$			$m = 71$	

2n Pointers Partition
 and, split array largest sum,
 $\min = [10 20 30 40]$

$S_0, l = 70$

$h = 70$

$m = 70$

5 students

not possible

Now,

$h = 71$

$h = 70$

$\min = [10 20 30 40]$

$l = 2$

$S_0, [10] + [20 30 40] S_0, \textcircled{90} P_2$

min (all max)

ans.

$[10 20] + [30 40] S_0, \textcircled{50} P_2$

$[10 20 30] + [40] S_0, \textcircled{60} P_1$

min (60) point

Now split array largest sum.

Split array in k subarrays such each one has one element and max subarray sum is minimum.

same

as Pointers Partition.

Q1 Minimize maximum distance b/w gas stations. 08.

Ex. $\text{arr1} = [1, 2, 3, 4, 5]$ Ans width 10⁶

$k = 4$ or actual ans will be accepted.

$\rightarrow \text{ans} = [1, 2, 3, 4, 5, 6, 7, 8, 9]$.

Ex. $\text{arr1} = [1, 7]$ Optimal 6 decimal places.

$k = 2$.

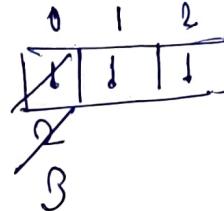
Now, $[1, 2, 4, 7] = 9$

$[1, 2, 5, 7] \rightarrow ②$ \leftarrow better.

Ex. $[1, 13, 17, 23]$ $k = 5$

$\downarrow \downarrow \downarrow \downarrow$

$3 \quad 6 \quad 1$



$\text{ans} = [1, 4, 7, 10, 13, 18, 17, 20, 23]$

Q2 Median of two sorted arrays

$\text{arr1} = [1, 3, 4, 7, 10, 12] \rightarrow m1s = [1, 2, 3, 3, 4, 6, 7, 10, 12, 15]$

$\text{arr2} = [2, 8, 6, 18]$

$$n=10 \quad 4+6/2 = 5$$

Ex. $\text{arr1} = [2, 3, 4]$

$\text{arr2} = [1, 3]$

$\boxed{[2, 3, 3, 4]}$

↑
ans

Now, how we use binary search?

$\text{arr1} = [1, 2, 3, 3, 4] \quad | \quad 6, 7, 10, 12, 18] \quad \left\{ \begin{array}{l} \text{on hypothetical} \\ \text{scenarios.} \end{array} \right.$

pick 1 ↓ | 3, 4, 7, 10, 12 } not valid
element 2, 8, 6, 18 |

pick 2 ↓ | 4, 7, 10, 12 } not valid.
element 2, 8, 6 | 18

Pick 8 element $\{1, 8, 4, 11, 10, 12, 2, 9, 12, 8, 1, 8\}$

→ 8 element form arr 1 and 2 form another

Pick 4 element $\{1, 3, 4, 7, 2, 6, 8, 10, 12, 18\}$

$4 < 6$
 $8 < 7$

Yes

$\{0, 1, 2, 3, 4, 5, 6\}$

• we check only cross guys.

test for validation.

Now

greater in $(11, 12)$ + ~~greatest~~ small

2

$$\text{so, } \frac{4+6}{2} = 10/2 = 5 \text{ Ans}$$

Now $\{0, 1, 2, 3, 4, 5, 6\}$ do BS on that array.

we do BS on shelter array too.

23 kth element of two sorted arrays.

$$arr1[] = [2, 3, 6, 7, 9]$$

$$arr2[] = [1, 4, 8, 10]$$

$$k = 4$$

	4	n
2	3	6 7 9
1	4	8 10
\circlearrowleft	\circlearrowright	$(n-k)$

24

Row with maximum 1's.

0	0	0	1	1	1
1	0	0	0	0	0
2	0	1	1	1	1
3	0	0	0	0	0
4	0	1	1	1	1

- = 3 • Brute way we done in $O(n \times m)$
- = 0
- = 4 now use binary search.
- = 0 we found only first occurrence of one's
- = 4 or LB OR 1
and UB OR 0

return $\underline{\text{②}}$ min \underline{m}
row

Now, mat[0] = [0 0 1 1 1]

25

Search in 2D matrix

mat[1][1] =	$\begin{bmatrix} 3 & 4 & 7 & 9 \\ 12 & 13 & 16 & 18 \\ 20 & 21 & 28 & 29 \end{bmatrix}$	target = 28
		<u>yes</u>

Simply flatten this in a single array.

So, [3 4 7 9 12 13 16 18 20 21 28 29].

How? Formula $1D \rightarrow 2D$.

$$\text{mid} = 10 \xrightarrow{\frac{10}{4}} \text{mat}[2, 2]$$

So, $(\text{mid} / \text{mat}[0].\text{size}(), \text{mid} \% \text{mat}[0].\text{size}())$.

→ vertical coordinate hai yehi formula hai.

26

Search in 2D matrix

1	4	7	11	18
2	5	8	12	19
3	6	9	16	22
10	13	14	17	24
18	21	23	26	30

Here every row and column are sorted.

1	4	7	11	18
2	5	8	12	19
3	6	9	16	22
10	18	14	17	24
19	21	28	26	30

These all
selected,

target = 14.

$$m_{\text{eff}} = \cancel{0} \cancel{1} \cancel{2} \cancel{3}$$

$$|CD| = \cancel{4}$$

Time $\rightarrow O(n+m)$

$$\text{row}(8,2) = 14$$

卷二

Find peak element

→ 2 adjacent cells are
not same.

0	4	2	5	1	4
1	2	9	3	2	3
2	1	7	6	0	1
3	3	6	2	3	7

- There are multiple peak elements but you have to return only single one {row, column}.

$$\begin{array}{l} \text{low} = 0 \quad \text{high} = n. \\ \text{mid} = 2 \end{array} \quad \left| \begin{array}{l} \text{max} = 6 \\ \text{but not peak.} \end{array} \right.$$

now delete right half

Find low = 0 high = 1 | max = 4 so return this.
mid = 0 | but it's peak

$$T_C = O\left(\log_2 m \times n\right).$$

$$S_C \leq O(1).$$

28 Median or row wise sorted matrix.

$$\left[\begin{array}{ccccc} 1 & 8 & 7 & 9 & 11 \\ 2 & 3 & 4 & 8 & 10 \\ 9 & 10 & 12 & 14 & 16 \end{array} \right] \xrightarrow{\quad} \left[\begin{array}{ccccccccc} 1 & 2 & 3 & 4 & 5 & 8 & 7 & 9 & 9 \\ 2 & 3 & 4 & 6 & 6 & 7 & 7 \end{array} \right]$$

 A off-tough
Resistor.

BINARY SEARCH QUESTIONS

01.) IMPLEMENT LOWER BOUND

CODE

```
int lowerBound(vector<int> arr, int n, int x) {  
    int low=0;  
    int high=n-1;  
    int ans=n;  
  
    while(low<=high){  
        int mid=(low+high)/2;  
        if(arr[mid]>=x){  
            ans=mid;  
            high=mid-1;  
        }  
        else{  
            low=mid+1;  
        }  
    }  
    return ans;  
}
```

02.) IMPLEMENT UPPER BOUND

CODE

```
int upperBound(vector<int> &arr, int x, int n){  
    int low=0;  
    int high=n-1;  
    int ans=n;  
  
    while(low<=high){  
        int mid=(low+high)/2;  
        if(arr[mid]>x){  
            ans=mid;  
            high=mid-1;  
        }  
        else  
            low=mid+1;  
    }  
    return ans;  
}
```

BINARY SEARCH QUESTIONS

03.) SEARCH INSERT POSITION

CODE

Same as lower bound question that is done above

```
int searchInsert(vector<int>& arr, int x)
{
    int n=arr.size();
    int low=0;
    int high=n-1;
    int ans=n;
    while(low<=high) {
        int mid=(low+high)/2;
        if(arr[mid]>=x) {
            ans=mid;
            high=mid-1;
        }
        else{
            low=mid+1;
        }
    }
    return ans;
}
```

BINARY SEARCH QUESTIONS

04.) FLOOR AND CEIL

CODE

```
int floor(vector<int> &a, int n, int x) {
    int ans=-1;
    int low=0;
    int high=n-1;
    while(low<=high) {
        int mid=(low+high)/2;
        if(a[mid]<=x) {
            ans=a[mid];
            low=mid+1;
        }
        else{
            high=mid-1;
        }
    }
    return ans;
}

int ceil(vector<int> a, int n, int x) {
    int low=0;
    int high=n-1;
    int ans=-1;

    while(low<=high) {
        int mid=(low+high)/2;
        if(a[mid]>=x) {
            ans=a[mid];
            high=mid-1;
        }
        else{
            low=mid+1;
        }
    }
    return ans;
}

pair<int, int> getFloorAndCeil(vector<int> &a, int n, int x) {
    int ans1=floor(a, n, x);
    int ans2=ceil(a, n, x);
    pair<int,int> p=make_pair(ans1, ans2);
    return p;
}
```

BINARY SEARCH QUESTIONS

05.) FIND FIRST AND LAST OCCURENCES IN AN ARRAY

CODE

```
class Solution {
public:
    int lowerBound(vector<int> arr, int n, int x) {
        int low=0;
        int high=n-1;
        int ans=-1;

        while(low<=high){
            int mid=(low+high)/2;
            if(arr[mid]==x){
                ans=mid;
                high=mid-1;
            }
            else if(arr[mid]<x){
                low=mid+1;
            }
            else{
                high=mid-1;
            }
        }
        return ans;
    }

    int upperBound(vector<int> &arr, int n, int x){
        int low=0;
        int high=n-1;
        int ans=n;

        while(low<=high){
            int mid=(low+high)/2;
            if(arr[mid]>x){
                ans=mid;
                high=mid-1;
            }
            else
                low=mid+1;
        }
        return ans;
    }

    Vector<int> searchRange(vector<int>& nums, int target) {
        int n=nums.size();
        vector<int> ans;
        int lb=lowerBound(nums, n, target);
        ans.push_back(lb);
```

BINARY SEARCH QUESTIONS

```
        if(lb==-1)
            ans.push_back(-1);
        else
            ans.push_back(upperBound(nums, n, target)-1);
        return ans;
    }
};
```

06.) SEARCH IN ROTATED SORTED ARRAY

CODE

```
class Solution {
public:
    int search(vector<int>& nums, int target) {
        int low=0;
        int high=nums.size()-1;

        while(low<=high){
            int mid=(low+high)/2;
            if(nums[mid]==target)
                return mid;

            if(nums[low]<=nums[mid]){
                if(nums[low]<=target && target<=nums[mid]){
                    high=mid-1;
                }
                else{
                    low=mid+1;
                }
            }
            else{
                if(nums[mid]<=target && target<=nums[high]){
                    low=mid+1;
                }
                else{
                    high=mid-1;
                }
            }
        }
        return -1;
    }
};
```

BINARY SEARCH QUESTIONS

07.) SEARCH IN ROTATED SORTED ARRAY II

CODE

```
class Solution {
public:
    bool search(vector<int>& nums, int target) {
        int low=0;
        int high=nums.size()-1;

        while(low<=high){
            int mid=(low+high)/2;
            if(nums[mid]==target)
                return true;

            // EXTRA CASE ONLY
            if(nums[low]==nums[mid] && nums[mid]==nums[high]){
                low++;
                high--;
                continue;
            }
            if(nums[low]<=nums[mid]){
                if(nums[low]<=target && target<=nums[mid]){
                    high=mid-1;
                }
                else{
                    low=mid+1;
                }
            }
            else{
                if(nums[mid]<=target && target<=nums[high]){
                    low=mid+1;
                }
                else{
                    high=mid-1;
                }
            }
        }
        return false;
    }
};
```

BINARY SEARCH QUESTIONS

8.) FIND MINIMUM IN ROTATED SORTED ARRAY

CODE

```
class Solution {  
public:  
    int findMin(vector<int>& nums) {  
        int low=0;  
        int high=nums.size()-1;  
        int ans=INT_MAX;  
  
        while(low<=high){  
            int mid=(low+high)/2;  
            if(nums[low]<=nums[mid]){  
                ans=min(ans, nums[low]);  
                low=mid+1;  
            }  
            else{  
                high=mid-1;  
                ans=min(ans, nums[mid]);  
            }  
        }  
        return ans;  
    }  
};
```

BINARY SEARCH QUESTIONS

9.) FIND MINIMUM IN ROTATED SORTED ARRAY II

CODE

```
class Solution {
public:
    int findMin(vector<int>& nums) {
        int low=0;
        int high=nums.size()-1;
        int ans=INT_MAX;

        while(low<=high){
            int mid=(low+high)/2;
            // EXTRA CASE ONLY
            if(nums[low]==nums[mid] && nums[mid]==nums[high]){
                ans=min(ans, nums[low]);
                low++;
                high--;
                continue;
            }
            if(nums[low]<=nums[mid]){
                ans=min(ans, nums[low]);
                low=mid+1;
            }
            else{
                high=mid-1;
                ans=min(ans, nums[mid]);
            }
        }
        return ans;
    }
};
```

BINARY SEARCH QUESTIONS

10.) HOW MANY TIMES ARRAY IS ROTATED

CODE

Slight variation in find minimum in sorted array code

```
#include <bits/stdc++.h>
int findMin(vector<int>& nums) {
    int low=0;
    int high=nums.size()-1;
    int ans=INT_MAX;
    int index=-1;
    while(low<=high) {
        int mid=(low+high)/2;
        if(nums[low]<=nums[high]) {
            if(nums[low]<ans) {
                index=low;
                ans=nums[low];
            }
            break;
        }
        if(nums[low]<=nums[mid]) {
            if(nums[low]<ans) {
                index=low;
                ans=nums[low];
            }
            low=mid+1;
        }
        else{
            if(nums[mid]<ans) {
                index=mid;
                ans=nums[mid];
            }
            high=mid-1;
        }
    }
    return index;
}

int findKRotation(vector<int> &arr) {
    int ans=findMin(arr);
    return ans;
}
```

BINARY SEARCH QUESTIONS

11.) SINGLE ELEMENT IN SORTED ARRAY

CODE

```
class Solution {
public:
    int singleNonDuplicate(vector<int>& nums) {
        if(nums.size()==1)
            return nums[0];

        if(nums[0]!=nums[1])
            return nums[0];

        if(nums[nums.size()-1]!=nums[nums.size()-2])
            return nums[nums.size()-1];

        int low=1;
        int high=nums.size()-1;

        while(low<=high){
            int mid=(low+high)/2;
            if(nums[mid]!=nums[mid+1] && nums[mid]!=nums[mid-1])
                return nums[mid];
            if((mid%2==1) && nums[mid-1]==nums[mid]
               || (mid%2==0) && nums[mid]==nums[mid+1]){
                low=mid+1;
            }
            else{
                high=mid-1;
            }
        }
        return -1;
    }
};
```

BINARY SEARCH QUESTIONS

12.) FIND PEAK ELEMENT

CODE

```
class Solution {  
public:  
    int findPeakElement(vector<int>& nums) {  
        int n=nums.size();  
        if(n==1)  
            return 0;  
        if(nums[0]>nums[1])  
            return 0;  
        if(nums[n-1]>nums[n-2])  
            return n-1;  
  
        int low=1;  
        int high=n-2;  
  
        while(low<=high){  
            int mid=(low+high)/2;  
            if(nums[mid]>nums[mid-1] && nums[mid]>nums[mid+1])  
                return mid;  
            }  
            else if(nums[mid]>nums[mid-1])  
                low=mid+1;  
            }  
            else{  
                high=mid-1;  
            }  
        }  
        return -1;  
    }  
};
```

BINARY SEARCH QUESTIONS

13.) PEAK INDEX IN A MOUNTAIN ARRAY

CODE

```
class Solution {
public:
    int peakIndexInMountainArray(vector<int>& nums) {
        int n=nums.size();
        if(n==1)
            return 0;
        if(nums[0]>nums[1])
            return 0;
        if(nums[n-1]>nums[n-2])
            return n-1;

        int low=1;
        int high=n-2;

        while(low<=high){
            int mid=(low+high)/2;
            if(nums[mid]>nums[mid-1] && nums[mid]>nums[mid+1]){
                return mid;
            }
            else if(nums[mid]>nums[mid-1]){
                low=mid+1;
            }
            else{
                high=mid-1;
            }
        }
        return -1;
    }
};
```

BINARY SEARCH QUESTIONS

14.) FIND SQRT

CODE

```
class Solution {  
public:  
    int mySqrt(int x) {  
        if(x==0)  
            return 0;  
  
        int low=1;  
        int high=x;  
        int ans=INT_MIN;  
  
        while(low<=high){  
            int mid=low+((high-low)/2);  
            if(mid>x/mid){  
                high=mid-1;  
            }  
            else if(mid<=x/mid){  
                ans=max(ans, mid);  
                low=mid+1;  
            }  
        }  
        return ans;  
    }  
};
```

BINARY SEARCH QUESTIONS

15.) FIND THE Nth ROOT OF AN INTEGER

CODE

```
#include <bits/stdc++.h>
using namespace std;

int func(int mid, int n, int m) {
    long long ans = 1;
    for (int i = 1; i <= n; i++) {
        ans = ans * mid;
        if (ans > m) return 2;
    }
    if (ans == m) return 1;
    return 0;
}

int NthRoot(int n, int m) {
    int low = 1, high = m;
    while (low <= high) {
        int mid = (low + high) / 2;
        int midN = func(mid, n, m);
        if (midN == 1) {
            return mid;
        } else if (midN == 0) {
            low = mid + 1;
        } else {
            high = mid - 1;
        }
    }
    return -1;
}
```

BINARY SEARCH QUESTIONS

16.) KOKO EATING BANANAS

CODE

```
class Solution {  
public:  
    long long func(vector<int> &p, int mid){  
        long long sum=0;  
        for(int i=0;i<p.size();i++){  
            sum=sum+ceil((double)p[i]/(double)mid);  
        }  
        return sum;  
    }  
  
    int minEatingSpeed(vector<int>& p, int h) {  
        long long int low=1;  
        long long high=*max_element(p.begin(), p.end());  
        int ans=-1;  
  
        while(low<=high){  
            long long int mid=low+((high-low)/2);  
            long long int tot_hr=func(p, mid);  
            if(tot_hr<=h){  
                ans=mid;  
                high=mid-1;  
            }  
            else{  
                low=mid+1;  
            }  
        }  
        return ans;  
    }  
};
```

BINARY SEARCH QUESTIONS

17.) MINIMUM NUMBER OF DAYS TO MAKE M BOUQUETS

CODE

```
class Solution {  
  
public:  
    bool func(vector<int> &b, int day, int m, int k){  
        int count=0;  
        int num=0;  
        for(int i=0;i<b.size();i++){  
            if(b[i]<=day){  
                count++;  
            }  
            else{  
                num=num+(count/k);  
                count=0;  
            }  
        }  
        num=num+(count/k);  
        return num>=m;  
    }  
  
    int minDays(vector<int>& b, int m, int k) {  
        if(b.size()/k<m)  
            return -1;  
  
        int low=*min_element(b.begin(), b.end());  
        int high=*max_element(b.begin(), b.end());  
        int ans=INT_MAX;  
  
        while(low<=high){  
            int mid=(low+high)/2;  
            if(func(b, mid, m, k)==true){  
                ans=mid;  
                high=mid-1;  
            }  
            else{  
                low=mid+1;  
            }  
        }  
        return ans;  
    }  
};
```

BINARY SEARCH QUESTIONS

18.) MAXIMUM CANDIES ALLOCATED TO K CHILDREN

CODE

```
class Solution {  
public:  
    bool func(vector<int> &c, int mid, long long k){  
        long long int count=0;  
        for(int i=0;i<c.size();i++){  
            count=count+(c[i]/mid);  
        }  
        return count>=k;  
    }  
  
    int maximumCandies(vector<int>& c, long long k) {  
        int low=1;  
        int high=*max_element(c.begin(), c.end());  
  
        while(low<=high){  
            int mid=low+((high-low)/2);  
            if(func(c, mid, k)){  
                low=mid+1;  
            }  
            else{  
                high=mid-1;  
            }  
        }  
        return high;  
    }  
};
```

BINARY SEARCH QUESTIONS

19.) FIND THE SMALLEST DIVISOR GIVEN A THRESHOLD VALUES

CODE

```
class Solution {  
public:  
    int func(vector<int> &nums, int mid){  
        int count=0;  
        for(int i=0;i<nums.size();i++){  
            count=count+ceil((double)nums[i]/(double)mid);  
        }  
        return count;  
    }  
  
    int smallestDivisor(vector<int>& nums, int th) {  
        int low=1;  
        int high=*max_element(nums.begin(), nums.end());  
  
        while(low<=high){  
            int mid=(low+high)/2;  
            if(func(nums, mid)<=th){  
                high=mid-1;  
            }  
            else{  
                low=mid+1;  
            }  
        }  
        return low;  
    }  
};
```

BINARY SEARCH QUESTIONS

20.) CAPACITY TO SHIP PACKAGES WITHIN D DAYS

CODE

```
class Solution {  
public:  
    int func(vector<int> &w, int cap){  
        int day=1;  
        int load=0;  
        for(int i=0;i<w.size();i++){  
            if(w[i]+load>cap){  
                day=day+1;  
                load=w[i];  
            }  
            else{  
                load=load+w[i];  
            }  
        }  
        return day;  
    }  
  
    int shipWithinDays(vector<int>& w, int days) {  
        int low=*max_element(w.begin(), w.end());  
        int high=accumulate(w.begin(), w.end(), 0);  
        while(low<=high){  
            int mid=(low+high)/2;  
            int day=func(w, mid);  
            if(day<=days)  
                high=mid-1;  
            else  
                low=mid+1;  
        }  
        return low;  
    }  
};
```

BINARY SEARCH QUESTIONS

21.) Kth MISSING POSITIVE NUMBERS

CODE

```
class Solution {  
public:  
    int findKthPositive(vector<int>& arr, int k) {  
        int low=0;  
        int high=arr.size()-1;  
  
        while(low<=high){  
            int mid=(low+high)/2;  
            int miss=arr[mid]-(mid+1);  
            if(miss<k)  
                low=mid+1;  
            else  
                high=mid-1;  
        }  
        return k+high+1;  
    }  
};
```

BINARY SEARCH QUESTIONS

22.) AGGRESSIVE COWS

CODE

```
bool ispossible(vector<int> &stalls, int k, int mid) {  
    int cowcount = 1;  
    int lastpos = stalls[0];  
  
    for (int i = 0; i < stalls.size(); i++) {  
        if (stalls[i] - lastpos >= mid) {  
            cowcount++;  
            if (cowcount == k) {  
                return true;  
            }  
            lastpos = stalls[i];  
        }  
    }  
    return false;  
}  
  
int aggressiveCows(vector<int> &stalls, int k) {  
    sort(stalls.begin(), stalls.end());  
    int s = 0;  
    int e = *max_element(stalls.begin(), stalls.end());  
    int ans = -1;  
  
    while (s <= e) {  
        int mid = s + (e - s) / 2;  
        if (ispossible(stalls, k, mid)) {  
            ans = mid;  
            s = mid + 1;  
        } else {  
            e = mid - 1;  
        }  
    }  
    return ans;  
}
```

BINARY SEARCH QUESTIONS

23.) MAGNETIC FORCE BETWEEN TWO BALLS

CODE

Same as aggressive cow questions

```
class Solution {
public:
    bool ispossible(vector<int> &stalls, int k, int mid) {
        int cowcount = 1;
        int lastpos = stalls[0];

        for (int i = 0; i < stalls.size(); i++) {
            if (stalls[i] - lastpos >= mid) {
                cowcount++;
                if (cowcount == k) {
                    return true;
                }
                lastpos = stalls[i];
            }
        }
        return false;
    }

    int maxDistance(vector<int>& stalls, int k) {
        sort(stalls.begin(), stalls.end());
        int s = 0;
        int e = *max_element(stalls.begin(), stalls.end());
        int ans = -1;

        while (s <= e) {
            int mid = s + (e - s) / 2;
            if (ispossible(stalls, k, mid)) {
                ans = mid;
                s = mid + 1;
            } else {
                e = mid - 1;
            }
        }
        return ans;
    }
};
```

BINARY SEARCH QUESTIONS

24.) ALLOCATE BOOKS

CODE

```
#include <bits/stdc++.h>
using namespace std;

int func(vector<int> &arr, int pages) {
    int stud=1;
    long long pageStud=0;

    for(int i=0;i<arr.size();i++) {
        if(pageStud+arr[i]<=pages) {
            pageStud=pageStud+arr[i];
        }
        else{
            stud++;
            pageStud=arr[i];
        }
    }
    return stud;
}

int findPages(vector<int>& arr, int n, int m) {
    if(m>n)
        return -1;

    int low=*max_element(arr.begin(), arr.end());
    int high=accumulate(arr.begin(), arr.end(), 0);

    while(low<=high) {
        int mid=(low+high)/2;
        int students=func(arr, mid);
        if(students>m)
            low=mid+1;
        else
            high=mid-1;
    }
    return low;
}
```

BINARY SEARCH QUESTIONS

25.) SPLIT ARRAY LARGEST SUM

CODE

Same as ALLOCATE BOOKS

```
class Solution {
public:
    int func(vector<int> &arr, int pages){
        int stud=1;
        long long pageStud=0;

        for(int i=0;i<arr.size();i++){
            if(pageStud+arr[i]<=pages){
                pageStud=pageStud+arr[i];
            }
            else{
                stud++;
                pageStud=arr[i];
            }
        }
        return stud;
    }

    int splitArray(vector<int>& arr, int m) {
        int n=arr.size();
        if(m>n)
            return -1;

        int low=*max_element(arr.begin(), arr.end());
        int high=accumulate(arr.begin(), arr.end(), 0);

        while(low<=high){
            int mid=(low+high)/2;
            int students=func(arr, mid);
            if(students>m)
                low=mid+1;
            else
                high=mid-1;
        }
        return low;
    }
};
```

BINARY SEARCH QUESTIONS

26.) PAINTER'S PARTITION PROBLEM

CODE

Same as ALLOCATED BOOKS

```
#include <bits/stdc++.h>
using namespace std;

int func(vector<int> &arr, int pages) {
    int stud=1;
    long long pageStud=0;

    for(int i=0;i<arr.size();i++) {
        if(pageStud+arr[i]<=pages) {
            pageStud=pageStud+arr[i];
        }
        else{
            stud++;
            pageStud=arr[i];
        }
    }
    return stud;
}

int findLargestMinDistance(vector<int> &arr, int m)
{
    int n=arr.size();
    if(m>n)
        return -1;

    int low=*max_element(arr.begin(), arr.end());
    int high=accumulate(arr.begin(), arr.end(), 0);

    while(low<=high) {
        int mid=(low+high)/2;
        int students=func(arr, mid);
        if(students>m)
            low=mid+1;
        else
            high=mid-1;
    }
    return low;
}
```

BINARY SEARCH QUESTIONS

27.) FIND Kth SMALLEST PAIR DISTANCE

CODE

```
class Solution {
public:
    bool isPossible(vector<int>& nums, int k, int mid) {
        int count = 0;
        int left = 0;

        for (int right = 0; right < nums.size(); right++) {
            while (nums[right] - nums[left] > mid) {
                left++;
            }
            count += right - left;
        }
        return count >= k;
    }

    int smallestDistancePair(vector<int>& nums, int k) {
        sort(nums.begin(), nums.end());
        int low = 0, high = nums.back() - nums.front();

        while (low < high) {
            int mid = low + (high - low) / 2;
            if (isPossible(nums, k, mid)) {
                high = mid;
            } else {
                low = mid + 1;
            }
        }
        return low;
    }
};
```

BINARY SEARCH QUESTIONS

28.) MEDIAN OF TWO SORTED ARRAYS

CODE 01

```
class Solution {  
public:  
    double findMedianSortedArrays(vector<int>& a, vector<int>& b) {  
        int n1=a.size();  
        int n2=b.size();  
  
        int i=0;  
        int j=0;  
        int n=(n1+n2);  
  
        int ind2=n/2;  
        int ind1=ind2-1;  
  
        int count=0;  
        int ind1el=-1;  
        int ind2el=-1;  
  
        while(i<n1 && j<n2){  
            if(a[i]<b[j]){  
                if(count==ind1)  
                    ind1el=a[i];  
                if(count==ind2)  
                    ind2el=a[i];  
                count++;  
                i++;  
            }  
            else{  
                if(count==ind1)  
                    ind1el=b[j];  
                if(count==ind2)  
                    ind2el=b[j];  
                count++;  
                j++;  
            }  
        }  
        while(i<n1){  
            if(count==ind1)  
                ind1el=a[i];  
            if(count==ind2)  
                ind2el=a[i];  
            count++;  
            i++;  
        }  
    }  
}
```

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```
while(j<n2){  
    if(count==ind1)  
        ind1el=b[j];  
    if(count==ind2)  
        ind2el=b[j];  
    count++;  
    j++;  
}  
  
if(n%2==1)  
    return ind2el;  
return (double)((double)(ind1el+ind2el)/2.0);  
}  
};
```

CODE 02

```
class Solution {  
public:  
    double findMedianSortedArrays(vector<int>& a, vector<int>& b) {  
        int n1=a.size();  
        int n2=b.size();  
  
        if(n1>n2)  
            return findMedianSortedArrays(b, a);  
  
        int low=0;  
        int high=n1;  
  
        int left=(n1+n2+1)/2;  
        int n=n1+n2;  
  
        while(low<=high){  
            int mid1=(low+high)>>1;  
            int mid2=left-mid1;  
  
            int l1=INT_MIN;  
            int l2=INT_MIN;  
            int r1=INT_MAX;  
            int r2=INT_MAX;  
  
            if(mid1<n1)  
                r1=a[mid1];  
            if(mid2<n2)  
                r2=b[mid2];  
            if(mid1-1>=0)  
                l1=a[mid1-1];  
        }  
    }  
};
```

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```
if(mid2-1>=0)
l2=b[mid2-1];

if(l1<=r2 && l2<=r1){
    if(n%2==1)
        return max(l1, l2);
    return ((double)(max(l1, l2)+min(r1, r2)))/2.0;
}
else if(l1>r2)
high=mid1-1;
else
low=mid1+1;
}
return 0;
};

};
```

BINARY SEARCH QUESTIONS

29.) Kth ELEMENT OF TWO SORTED ARRAYS

CODE

```
#include <bits/stdc++.h>
int kthElement(vector<int> &a, vector<int> &b, int n1, int n2, int k) {
    if(n1>n2)
        return kthElement(b, a, n2, n1, k);

    int low=max(k-n2, 0);
    int high=min(k, n1);

    int left=k;
    int n=n1+n2;

    while(low<=high) {
        int mid1=(low+high)>>1;
        int mid2=left-mid1;

        int l1=INT_MIN;
        int l2=INT_MIN;
        int r1=INT_MAX;
        int r2=INT_MAX;

        if(mid1<n1)
            r1=a[mid1];
        if(mid2<n2)
            r2=b[mid2];
        if(mid1-1>=0)
            l1=a[mid1-1];
        if(mid2-1>=0)
            l2=b[mid2-1];

        if(l1<=r2 && l2<=r1) {
            return max(l1, l2);
        }
        else if(l1>r2)
            high=mid1-1;
        else
            low=mid1+1;
    }
    return 0;
}
```

BINARY SEARCH QUESTIONS

30.) ROW WITH MAXIMUM 1's

CODE

```
class Solution {
public:
    int lowerBound(vector<int> &arr, int n, int x) {
        int low=0;
        int high=n-1;
        int ans=n;
        sort(arr.begin(), arr.end());

        while(low<=high){
            int mid=(low+high)/2;
            if(arr[mid]>=x){
                ans=mid;
                high=mid-1;
            }
            else{
                low=mid+1;
            }
        }
        return ans;
    }

    vector<int> rowAndMaximumOnes(vector<vector<int>>& mat) {
        int count_max=0;
        int index=0;

        for(int i=0;i<mat.size();i++){
            int m=mat[i].size();
            int count=m-lowerBound(mat[i], m, 1);
            if(count>count_max){
                count_max=count;
                index=i;
            }
        }
        return {index, count_max};
    }
};
```

BINARY SEARCH QUESTIONS

31.) SEARCH IN 2D MATRIX I

CODE

```
class Solution {  
public:  
    bool searchMatrix(vector<vector<int>>& mat, int target) {  
        int n=mat.size();  
        int m=mat[0].size();  
  
        int low=0;  
        int high=(n*m)-1;  
  
        while(low<=high){  
            int mid=(low+high)/2;  
            int row=mid/m;  
            int col=mid%m;  
  
            if(mat[row][col]==target)  
                return true;  
            else if(mat[row][col]<target)  
                low=mid+1;  
            else  
                high=mid-1;  
        }  
        return false;  
    }  
};
```

BINARY SEARCH QUESTIONS

32.) SEARCH IN 2D MATRIX II

CODE

```
class Solution {
public:
    bool searchMatrix(vector<vector<int>>& mat, int target) {
        int n=mat.size();
        int m=mat[0].size();

        int row=0;
        int col=m-1;

        while(row<n && col>=0){
            if(mat[row][col]==target)
                return true;
            else if(mat[row][col]<target)
                row++;
            else
                col--;
        }
        return false;
    }
};
```

BINARY SEARCH QUESTIONS

33.) FIND PEAK ELEMENT II

CODE

```
class Solution {
public:
    vector<int> findPeakGrid(vector<vector<int>>& mat) {
        int m = mat.size();
        int n = mat[0].size();
        int low = 0;
        int high = n - 1;

        while (low <= high) {
            int mid = low + (high - low) / 2;
            int maxRow = 0;

            for (int i = 0; i < m; i++) {
                if (mat[i][mid] > mat[maxRow][mid]) {
                    maxRow = i;
                }
            }

            bool isLeftValid = false;
            bool isRightValid = false;

            if (mid - 1 >= 0 && mat[maxRow][mid] < mat[maxRow][mid - 1]) {
                isLeftValid = true;
            }

            if (mid + 1 < n && mat[maxRow][mid] < mat[maxRow][mid + 1]) {
                isRightValid = true;
            }

            if (!isLeftValid && !isRightValid) {
                return {maxRow, mid};
            }

            if (isLeftValid) {
                high = mid - 1;
            }
            else {
                low = mid + 1;
            }
        }
        return {-1, -1};
    };
}
```

BINARY SEARCH QUESTIONS

34.) MEDIAN OF A ROW WISE SORTED MATRIX

CODE

```
class Solution {
public:
    int countLessEqual(vector<vector<int>>& matrix, int mid) {
        int count = 0;
        int rows = matrix.size();
        int cols = matrix[0].size();

        for (int i = 0; i < rows; ++i) {
            count += upper_bound(matrix[i].begin(), matrix[i].end(), mid)
                    - matrix[i].begin();
        }
        return count;
    }

    int findMedian(vector<vector<int>>& matrix) {
        int rows = matrix.size();
        int cols = matrix[0].size();

        int low = INT_MAX, high = INT_MIN;
        for (int i = 0; i < rows; ++i) {
            low = min(low, matrix[i][0]);
            high = max(high, matrix[i][cols - 1]);
        }

        int desired = (rows * cols + 1) / 2;

        while (low < high) {
            int mid = low + (high - low) / 2;
            if (countLessEqual(matrix, mid) < desired) {
                low = mid + 1;
            } else {
                high = mid;
            }
        }

        return low;
    }
};
```

THANK YOU !