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Pledge: "I pledge my honor that I have abided by the Stevens Honor System" – Himanshu Rana(hrana2)

Use the Master Theorem to find the complexity of each recurrence relation listed below.

1. $T(n) = T\left(\frac{n}{2}\right) + n^2$
 $a = 1, b = 2, d = 2$
 $1 < 2^2$
 Complexity: $\Theta(n^2)$
2. $T(n) = 4T\left(\frac{n}{2}\right) + n^2$
 $a = 4, b = 2, d = 2$
 $4 = 2^2$
 Complexity: $\Theta(n^2 \log_2 n)$
3. $T(n) = 3T\left(\frac{n}{3}\right) + \sqrt{n}$
 $a = 3, b = 3, d = \frac{1}{2}$
 $3 > 3^{(1/2)}$
 Complexity: $\Theta(n^{\log_3 3}) = \Theta(n)$

For each function below, write the recurrence relation for its running time and then use the Master Theorem to find its complexity.

```
4. int f(int arr[], int n) {
    if (n == 0) {
        return 0;
    }
    int sum = 0;
    for (int j = 0; j < n; ++j) {
        sum += arr[j];
    }
    return f(arr, n / 2) + sum + f(arr, n / 2);
}
```

Recurrence: $T(n) = 2T(n/2) + n$ $a = 2, b = 2, d = 1$ $2 = 2^1$ Complexity: $\Theta(n \log_2 n)$

```
5. void g(int n, int arrA[], int arrB[]) {  
    if (n == 0) {  
        return;  
    }  
    for (int i = 0; i < n; ++i) {  
        for (int j = 0; j < n; ++j) {  
            arrB[j] += arrA[i];  
        }  
    }  
    g(n / 2, arrA, arrB);  
}
```

Recurrence: $T(n) = T(n/2) + (n^2)$

$a = 1, b = 2, d = 2$

$1 < 2^2$

Complexity: $\Theta(n^2)$