CS5005 Parallel Programming Assignment - 2

(Performance profile of synchronization constructs)

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Problem Statement

We discussed various synchronization primitives such as busy-waiting, mutexes, conditional variables, barriers, and read-write locks. You are required to design experiments (or test cases), execute them and measure the performance of each of these synchronization constructs.

Experiment - 1 (array_sum)

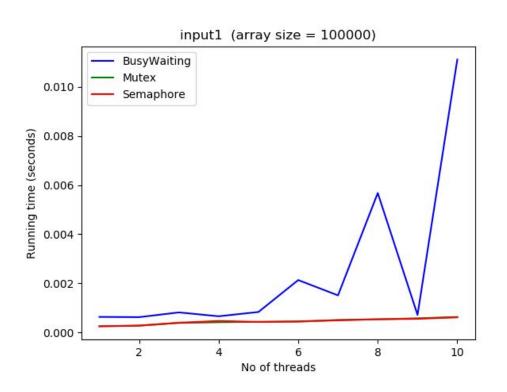
The first experiment is of calculating sum of numbers in an array. Each thread computes a partial sum of the array. These sums are combined into a global sum using critical section. The synchronisation constructs compared using this experiment are

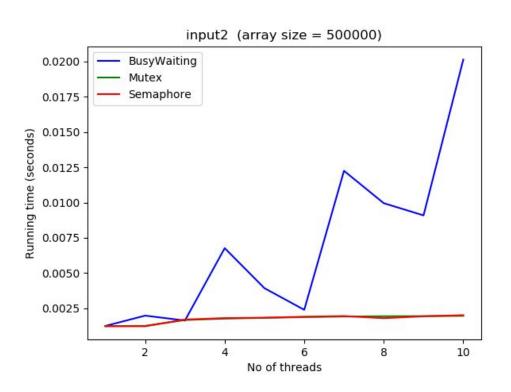
- Busy Waiting
- Mutex
- Semaphore

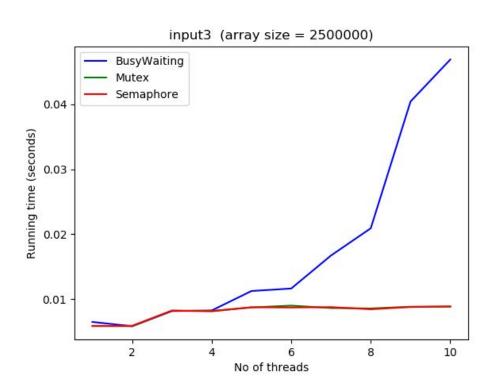
The synchronisation constructs are compared for different array sizes and also for different number of threads.

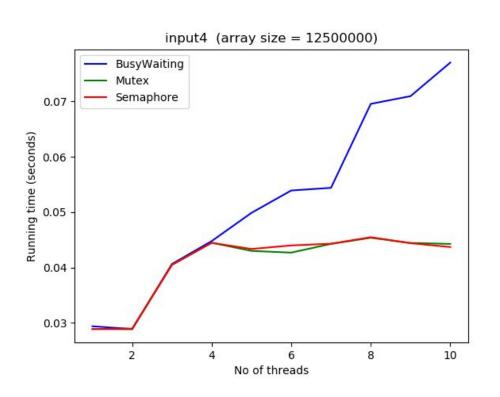
The functions running under the threads for the three constructs are kept similar apart except for the codes specific to them. This ensures uniformity and no extra overhead for any of the three constructs.

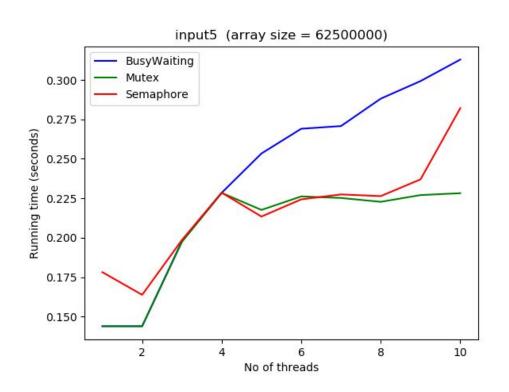
Also since the running time for a program depends on multiple factors an average of running time over various runs is taken. Even then some uneven results are likely to be observed.











Analysis (array_sum)

As the plots on the previous pages shows, the running time for each of the three constructs - busy waiting, mutex, semaphore first decreases when no of threads increases from 1 to 2, because the experiments were performed on a dual core PC. So, at a time maximum of two threads can be run simultaneously. After this the increasing number of threads increases running time as threads with exception of atmost two, must wait to be scheduled at any point of time.

The effect of increasing number of threads is much more pronounced for busy wait than the other two. This is also clear as a thread which is scheduled but is anyway busy waiting wastes processors time, which is not the case with semaphores and mutexes.

The performance of semaphores and mutexes are similar as seen from the plots.

So, we can conclude that both semaphores and mutexes gives almost similar performance and should be used instead of busy waiting to gain performance.

Experiment - 2 (heat_eqlb)

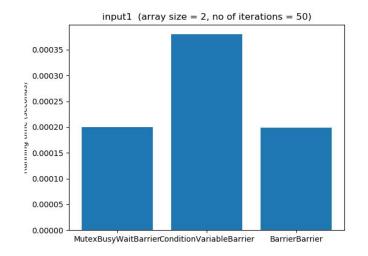
The second experiment is an algorithmic analogy of heat flow through a rod with varying temperature at various points. Ultimately, the rod reaches thermal equilibrium by heat flow from one part of the rod to the other. This experiment simulates the same process using threads. The synchronization constructs compared are

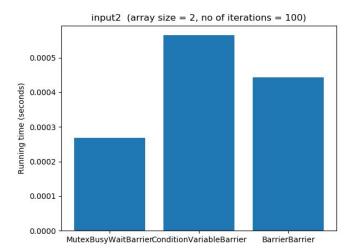
- Synchronisation using mutex and busy waiting
- Condition Variables
- Barriers

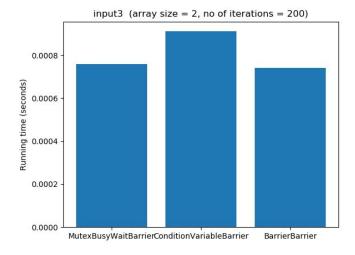
To maintain uniformity each of these is repeated a fixed prespecified number of iterations for an easy comparison.

These are compared for different array sizes and also for different number of iterations. Again the functions running under the threads for the three constructs are kept similar apart except for the codes specific to them. This ensures uniformity and no extra overhead for any of the three constructs. Also an average of running time over various runs is taken.

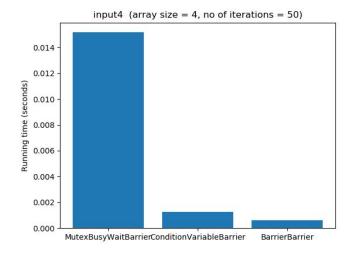
Observations (heat_eqlb)

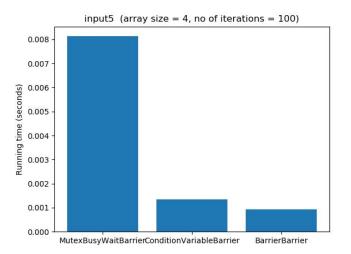


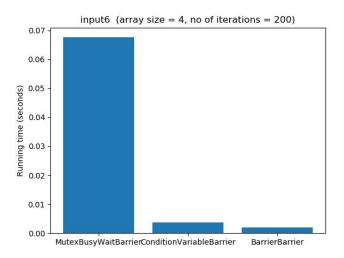




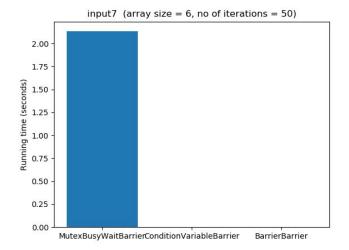
Observations (heat_eqlb)

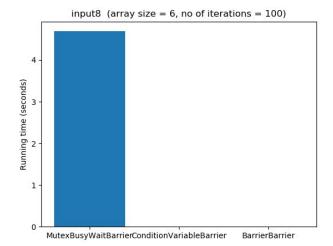


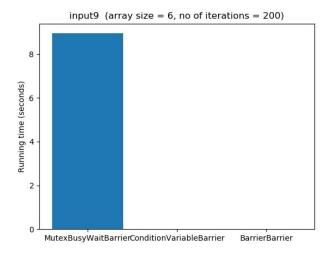




Observations (heat_eqlb)







Analysis (heat_eqlb)

Initially when the number of threads is smaller (2 which is less equal to the number of cores in the PC on which observations was made), the performance of mutex with busy waiting is comparable to that of condition variable and barrier. It is even better than the condition variable. This shows that the overhead involved in implementing the condition variable and barriers is comparable (even we can say it is more) as compared to the wastage of CPU during busy waiting. Also, however this effect is less pronounced as the number of iterations increases and the effect of busy wait becomes apparent.

However, as the number of threads increases (becomes more than available cores in the system) the mutex with busy wait is much badly outperformed by condition variable and barrier. The reason for this is definitely that the threads that are busy waiting also get scheduled by the scheduler, and wastes CPU time.

Also, barriers seem to be performing better compared to condition variables. The reason for this is that condition variables are more general constructs as compared to barriers. So, internal library implementation of condition variables would be much more complex as compared to that of barriers.

So, the conclusion is that if the number of threads involved in a rendezvous is smaller go for mutex with busy wait otherwise use barriers.

Raw Data for first experiment

For input:	1.txt	No accessor							For ohse	rvation - 1
	f the array is of the array i								1010030	<u> </u>
The time s		outing the array	sum (in order of	thread no from :	1 to 10):					
555,5	0.00063	0.00062	0.00081	0.00065	0.00083	0.00213	0.00151	0.00568	0.00071	0.01111
Mutex :	0.00025	0.00027	0.00038	0.00041	0.00043	0.00044	0.00050	0.00053	0.00057	0.00063
Semaphore	: 0.00025	0.00027	0.00039	0.00046	0.00042	0.00044	0.00049	0.00053	0.00055	0.00061
For input	2.txt								For obse	rvation - 2
The sum of	2.txt f the array i of the array								For obse	rvation - 2
The sum of The size of The time :	f the array i of the array spent for com		sum (in order o	f thread no from	1 to 10):				For obse	rvation - 2
The sum of	f the array i of the array spent for com	is: 500000	sum (in order o 0.00163	f thread no from 0.00677	1 to 10): 0.00393	0.00239	0.01225	0.00995	For obse	rvation - 2
The sum of The size of The time :	f the array i of the array spent for com ng :	is: 500000 puting the array				0.00239 0.00188	0.01225 0.00192	0.00995 0.00194	. November 1	

Raw Data for first experiment

	3.txt f the array is: of the array is								<u>For obse</u>	rvation - 3
The time s	spent for compu		sum (in order of	thread no from	1 to 10):					
BusyWaitin	ng: 0.00650	0.00584	0.00817	0.00827	0.01126	0.01164	0.01672	0.02091	0.04041	0.04686
Mutex :	0.00586	0.00586	0.00824	0.00820	0.00873	0.00901	0.00864	0.00858	0.00884	0.00886
Semaphore	: 0.00590	0.00591	0.00824	0.00814	0.00876	0.00872	0.00878	0.00845	0.00880	0.00889
For input	4.txt								For obse	rvation - 4
	f the array is: of the array is								1 01 0056	<u> Ivaliuii - 4</u>
The time s		iting the array	sum (in order o	f thread no from	1 to 10):					
busymattu	0.02937	0.02887	0.04063	0.04477	0.04987	0.05391	0.05440	0.06957	0.07097	0.07704
Mutex :	0.02887	0.02883	0.04052	0.04447	0.04299	0.04268	0.04426	0.04537	0.04443	0.04426
Semaphore	: 0.02886	0.02893	0.04045	0.04446	0.04334	0.04398	0.04429	0.04547	0.04439	0.04370

Raw Data for first experiment

For input5	.txt								Car abaa	rustion E	
The sum of the array is: -32168680 The size of the array is: 62500000									For observation - 5		
The size o	if the array	ls: 62500000									
The time s BusyWaitin		puting the array	sum (in order o	f thread no from	1 to 10):						
busynacea	0.14388	0.14391	0.19773	0.22846	0.25336	0.26903	0.27074	0.28811	0.29924	0.31288	
Mutex :											
	0.14401	0.14392	0.19738	0.22838	0.21760	0.22615	0.22517	0.22273	0.22700	0.22820	
Semaphore											
	0.17814	0.16379	0.19853	0.22853	0.21340	0.22429	0.22741	0.22636	0.23697	0.28209	

For input7.txt

The size of the array is: 6

The number of iterations is: 50

Raw Data for second experiment

For input1.txt	For input2.txt	For input3.txt
The size of the array is: 2	The size of the array is: 2	The size of the array is: 2
The number of iterations is: 50	The number of iterations is: 100	The number of iterations is: 200
The time spent for reaching equilibrium is : MutexBusyWaitBarrier : 0.00020 ConditionVariableBarrier : 0.00038 BarrierBarrier : 0.00020	The time spent for reaching equilibrium is : MutexBusyWaitBarrier : 0.00027 ConditionVariableBarrier : 0.00056 BarrierBarrier : 0.00044	The time spent for reaching equilibrium is : MutexBusyWaitBarrier : 0.00076 ConditionVariableBarrier : 0.00091 BarrierBarrier : 0.00074
For input4.txt	For input5.txt	For input6.txt
The size of the array is: 4	The size of the array is: 4	The size of the array is: 4
The number of iterations is: 50	The number of iterations is: 100	The number of iterations is: 200
The time spent for reaching equilibrium is :	The time spent for reaching equilibrium is	The Aire word for southing and library in

For input9.txt

The size of the array is: 6

The number of iterations is: 200

The time spent for reaching equilibrium is : The time spent for reaching equilibrium is : The time spent for reaching equilibrium is : MutexBusyWaitBarrier : 2.13351 MutexBusyWaitBarrier: 4.68980 MutexBusyWaitBarrier: 8.95425 ConditionVariableBarrier: 0.00317 ConditionVariableBarrier: 0.00627 ConditionVariableBarrier: 0.00167 BarrierBarrier: 0.00310 BarrierBarrier : 0.00086 BarrierBarrier: 0.00163

The size of the array is: 6

The number of iterations is: 100

For input8.txt

Notes

- The experiments were performed on dual core, Intel i5 processor on linux operating system.
- <u>Click here</u> to view the source code used in the analysis.
 - Run "make gentest" to generate test cases file (test case files are not given because of large size).
 - Run "make" to compile the source files.
 - Run "make" to run the file interactively with standard input.
 - Run "make test" to run the code on all input files (input1 to input5 for array_sum and input1 to input9 for heat_eqlb).
 - Run "make <filename>" to run a specific test file (input1 to input5 for array_sum and input1 to input9 for heat_eqlb).
 - Run "make clean" to remove files produced as a result of running the program.
 - Run "make clean" to remove and files produced as a result of running the program and also the input test files.

Thank You

