

# Double Turán Problem

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# Overview

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What is the Turán problem?

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## Question

Given a graph  $F$ , how many edges can an  $n$ -vertex graph have while containing no copy of  $F$  as a subgraph?

What is the hypergraph Turán problem?

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Given an  $r$ -uniform graph  $F$ , how many edges can an  $n$ -vertex  $r$ -uniform hypergraph have while containing no copy of  $F$  as a subgraph?

Let  $F$  be a graph. We call the following quantity the *Turán number* or *extremal number* of  $F$ :

## Definition

$$\text{ex}(n, F) := \max\{e(G) : |V(G)| = n \text{ and } F \not\subseteq G\}$$

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The *Turán graph*  $T_r(n)$  is the complete  $r$ -partite graph on  $n$  vertices with parts of size  $\lfloor n/r \rfloor$  and  $\lceil n/r \rceil$ .

# Introduction

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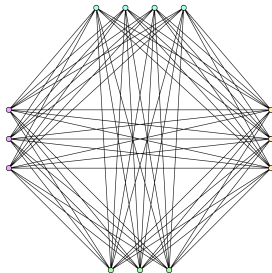


Figure:  $T_4(13)$



## Turán's theorem

The maximum number of edges in an  $n$ -vertex graph containing no clique of order  $r + 1$  is  $e(T_r(n))$ , with equality only for  $T_r(n)$ .

## Erdős-Stone Theorem, Simonovits' Theorem

Let  $F$  be any graph of chromatic number  $r + 1 \geq 3$ . Then  
 $\text{ex}(n, F) = (1 + o(1)) T_r(n)$  as  $n \rightarrow \infty$ .

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## Question

What is the value of  $\phi(m, n, F) = \max \sum_{i=1}^m |G_i|$ ?

# Double Turán Problem

Double Turán problems are closely related to Turán problems for 3-uniform hypergraphs  $H$  through *link graphs*.

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## Definition

For  $i \in V(H)$ , define graph  $H_i$  with

$$V(H_i) = V(H) \setminus \{i\} \quad \text{and} \quad E(H_i) = \{\{j, k\} : \{i, j, k\} \in E(H)\}.$$



# Double Turán Problem

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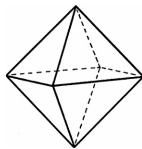


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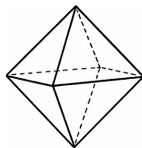


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$H$  is octahedron-free  $\implies H_1, H_2, \dots, H_n$  are double  $C_4$ -free.

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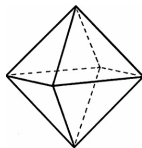


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$H$  is octahedron-free  $\implies H_1, H_2, \dots, H_n$  are double  $C_4$ -free.

$$\boxed{\text{ex}(n, O) \leq \phi(n, n, C_4)}$$

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We call graphs  $G_1, G_2, \dots, G_m$  *induced* if each  $G_i$  is an induced subgraph of  $\bigcup_{j=1}^m G_j$ .

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We call graphs  $G_1, G_2, \dots, G_m$  *induced* if each  $G_i$  is an induced subgraph of  $\bigcup_{j=1}^m G_j$ .

Let graphs  $G_1, G_2, \dots, G_m$  be induced and double  $F$ -free.

## Question

What is the value of  $\phi^*(m, n, F) = \max \sum_{i=1}^m e(G_i)$ ?

# Induced Double Turán Problem

## Generalized Turán problem

What is the maximum number  $ex(n, F, K_3)$  of triangles in a graph  $H$  on  $[n]$  with no copy of  $F$  as a subgraph?



# Induced Double Turán Problem

## Generalized Turán problem

What is the maximum number  $\text{ex}(n, F, K_3)$  of triangles in a graph  $H$  on  $[n]$  with no copy of  $F$  as a subgraph?

For  $i \in V(G)$ , define  $G_i$  with

$$V(G_i) = V(G) \quad \text{and} \quad E(G_i) = \{\{j, k\} : \{i, j\}, \{j, k\}, \{i, k\} \in E(G)\}$$

# Induced Double Turán Problem

Ex. Octahedron-free graph  $G$ .

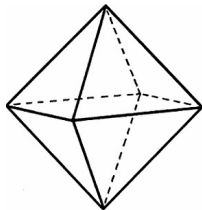


Figure: Octahedron

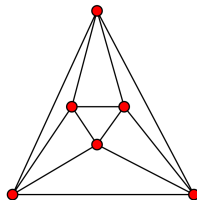


Figure: Octahedron Graph  $K_{2,2,2}$

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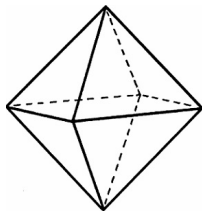


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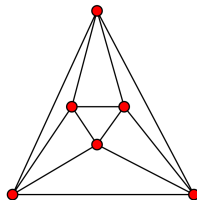


Figure: Octahedron Graph  $K_{2,2,2}$

$G$  is  $K_{2,2,2}$ -free  $\implies G_1, G_2, \dots, G_n$  are double  $K_{2,2}$ -free.

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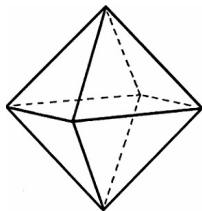


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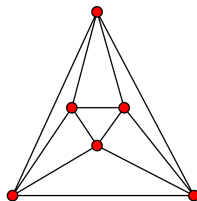


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$G$  is  $K_{2,2,2}$ -free  $\implies G_1, G_2, \dots, G_n$  are double  $K_{2,2}$ -free.

$$\text{ex}(n, K_{2,2,2}, K_3) \leq \phi(n, n, K_{2,2})$$

## Theorem A

For  $m \geq 3$  and non-bipartite  $F$ , if  $n$  is large enough, then

$$\phi^*(m, n, F) = m \cdot \text{ex}(n, F),$$

with equality only for identical extremal  $n$ -vertex  $F$ -free graphs.

# Main Results

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For  $m \geq 3$  and non-bipartite  $F$ , if  $n$  is large enough, then

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## Theorem B

For  $m, n, r \geq 3$ ,

$$\phi^*(m, n, K_r) = m \cdot e(T_{r-1}(n)),$$

with equality for induced  $K_r$ -free graphs  $G_1, G_2, \dots, G_m$  only if  $G_1 = G_2 = \dots = G_m = T_{r-1}(n)$ .

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We need the idea of  $(m, n, k)$ -blowup to state the result for  $\phi(m, n, K_r)$ .

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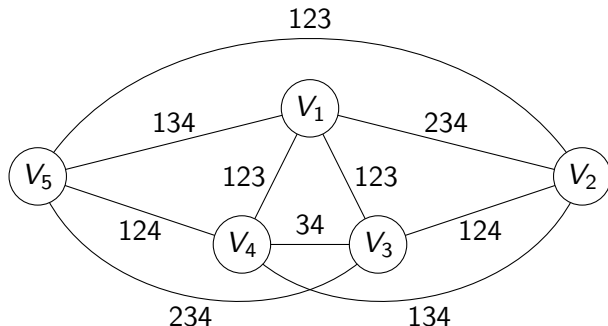


Figure: Example of an  $(4, n, 5)$ -blowup not containing a double  $K_3$ .



# Main Results

Let  $f(m, n, r)$  denote the maximum possible sum of edges in an double  $K_r$ -free  $(m, n, k)$ -blowup with  $k \leq \binom{m}{2}$ .

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## Theorem C

For  $n \geq 1$ ,

①

$$\phi(3, n, K_3) = \binom{n}{2} + \left\lfloor \frac{n^2}{2} \right\rfloor.$$

② if  $r \geq 2$  and  $m \geq 1$ ,

$$\phi(m, n, K_r) = f(m, n, r).$$

In particular,

$$\lim_{n \rightarrow \infty} \frac{f(4, n, 3)}{\binom{n}{2} + \left\lfloor \frac{n^2}{2} \right\rfloor} > 1$$

## Conjecture

Let  $F$  be any non-empty graph and  $m, n \geq 1$ . Then

$$\phi^*(m, n, F) = \Theta(m \cdot \text{ex}(n, F) + n^2).$$

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$$\phi^*(m, n, F) \geq \max\left\{\binom{n}{2}, m \cdot \text{ex}(n, F)\right\}.$$

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By Theorem 1, the conjecture is true when  $F$  is non-bipartite.

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$$\phi^*(m, n, F) = \Theta(m \cdot \text{ex}(n, F) + n^2).$$

Since

$$\text{ex}(n, K_{2,2,2}, K_3) \leq \phi^*(n, n, K_{2,2})$$

the conjecture implies

$$\text{ex}(n, K_{2,2,2}, K_3) \leq O(n^2)$$

## Theorem D

Let  $F$  be a graph. If there exists an extremal  $F$ -free  $n$ -vertex graph with maximum degree at most  $\sqrt{n}/m^2$ , then

$$\phi(m, n, F) = \binom{n}{2} + \binom{m}{2} \text{ex}(n, F).$$

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$$\phi(m, n, F) = \binom{n}{2} + \binom{m}{2} \text{ex}(n, F).$$

If  $P$  is a path of length 2 and  $m = o(n^{1/4})$ ,

$$\binom{n}{2} + m - 1 \leq \phi^*(m, n, P) \leq \phi(m, n, P) = \binom{n}{2} + \binom{m}{2} \left\lfloor \frac{n}{2} \right\rfloor.$$



## Theorem E

Let  $P$  be a path with two edges. Then  $\phi(n, n, P) = \Omega(n^{5/2})$ , whereas  $\phi^*(n, n, P) = o(n^{5/2})$ , as  $n \rightarrow \infty$ . In particular,

$$\lim_{n \rightarrow \infty} \frac{\phi^*(n, n, P)}{\phi(n, n, P)} = 0.$$

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$$\lim_{n \rightarrow \infty} \frac{\phi^*(n, n, P)}{\phi(n, n, P)} = 0.$$

This shows that  $\phi(n, n, P)$  and  $\phi^*(n, n, P)$  are very different problems.

## Theorem B

Let  $m, n, r \geq 3$ . Then

$$\phi^*(m, n, K_r) = m \cdot e(T_{r-1}(n)),$$

with equality for induced  $K_r$ -free graphs  $G_1, G_2, \dots, G_m$  only if  $G_1 = G_2 = \dots = G_m = T_{r-1}(n)$ .

## Proof Roadmap

- Step 1: Reduce to the case of smaller  $m$
- Step 2: Further reduce to an optimization problem
- Step 3: Solve the optimization problem

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## Step 1: Reduce to the case of smaller $m$

Put  $G_{i+m} = G_i$  for all  $i \in [m]$ .

$$\sum_{i=1}^m e(G_i) = \frac{1}{k} \sum_{i=1}^m [e(G_i) + \cdots + e(G_{i+k-1})]$$

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Since  $e(G_i) + \cdots + e(G_{i+k-1}) \leq \phi^*(k, n, F)$ ,

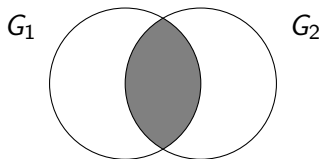
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## Step 2: Further reduce to an optimization problem

### Proof Roadmap

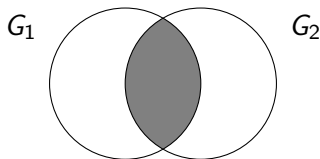
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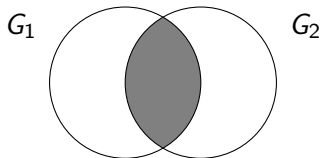


### Observation

If  $G_1, G_2$  intersects in  $t$  vertices, Then

$$e(G_1) + e(G_2) \leq \binom{n-t}{2} + (n-t)t + 2\text{ex}(t, F)$$

## Step 2: Further reduce to an optimization problem



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$\implies$  we are done if the unique maximum is at  $t = n$ .

# Step 3: Solve the optimization problem

## Proof Roadmap

- Step 1: Reduce to the case of smaller  $m$
- Step 2: Further reduce to an optimization problem
- **Step 3:** Solve the optimization problem

# Extra Ingredients to Prove Theorem A

## Theorem A

For  $m \geq 3$  and non-bipartite  $F$ , if  $n$  is large enough, then

$$\phi^*(m, n, F) = m \cdot \text{ex}(n, F),$$

with equality only for identical extremal  $n$ -vertex  $F$ -free graphs.

- First show that  $t = |V(G_1 \cap G_2)| \geq \sqrt{n}$ .
- For large enough  $t$ , any extremal  $t$ -vertex  $F$ -free graph contains a spanning  $T_{r-1}(t)$ .

# Proof of Theorem C: Upper Bound

We need to show

$$\phi(m, n, F) \leq \binom{n}{2} + \text{ex}(n, F) \binom{m}{2}.$$

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Let  $E_S$  be the set of edges in exactly  $\{G_i\}_{i \in S}$ .

$$\implies \sum_{i=1}^m e(G_i) = \sum_{S \subseteq [m]} |S| |E_S| \leq \binom{n}{2} + \sum_{S \subseteq [m], |S| \geq 2} (|S| - 1) |E_S|.$$

# Proof of Theorem C: Upper Bound

$$\sum_{\substack{S \subseteq [m] \\ |S| \geq 2}} (|S| - 1) |E_S| = \sum_{\substack{S \subseteq [m], \\ |S|=2}} \sum_{T \supseteq S} \frac{(|T| - 1) |E_T|}{\binom{|T|}{2}} \leq \sum_{\substack{S \subseteq [m], \\ |S|=2}} \sum_{T \supseteq S} |E_T|,$$

as each  $T \in [m]$  with  $|T| \geq 2$  is counted  $\binom{|T|}{2}$  times.



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Observation: If  $|S| \geq 2$ , the edge set  $\bigcup_{T \supseteq S} E_T$  is  $F$ -free.

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$$\left| \bigcup_{T \supseteq S} E_T \right| = \sum_{T \supseteq S} |E_T| \leq \text{ex}(n, F)$$

$$\Rightarrow \sum_{\substack{S \subseteq [m] \\ |S| \geq 2}} (|S| - 1) |E_S| \leq \sum_{\substack{S \subseteq [m], \\ |S|=2}} \sum_{T \supseteq S} |E_T| \leq \binom{m}{2} \text{ex}(n, F)$$

# Proof of Theorem C: Lower Bound

Let  $H_1, \dots, H_{\binom{m}{2}}$  be extremal  $F$ -free graphs on  $[n]$  with  $\Delta(H_i) \leq \sqrt{n}/m^2$

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IDEA: start with any embedding and iteratively decrease overlapping edges

# Proof of Theorem C: Lower Bound

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Resulting graph remains isomorphic