

Double Turán Problem

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1 Introduction

Let $\text{ex}_2(n, m, F)$ be the maximum possible sum of the number of edges over m subgraphs G_1, \dots, G_m on the same vertex set $[n]$, with the constraint that $E(G_i) \cap E(G_j)$ does not contain graph F for $i \neq j$. Our goal is to determine $\text{ex}_2(n, m, F)$ for different forbidden graphs F . A trivial construction with $G_1 = K_n$ and G_2, \dots, G_m to be extremal graphs for F yields the lower bound $\binom{n}{2} + (m-1)\text{ex}(n, F)$. In this work, we use this bound as a benchmark to either show the tightness of $\text{ex}_2(n, m, F)$ or to provide a better bound.

Additionally, we are also interested in a more restrictive version where G_1, \dots, G_m are induced subgraphs of $G_1 \cup \dots \cup G_m$. We denote $\text{ex}_2^*(n, m, F)$ as the maximum possible sum of the number of edges over m induced subgraphs G_1, \dots, G_m on the same vertex set $[n]$ such that $E(G_i) \cap E(G_j)$ does not contain graph F for $i \neq j$. The trivial construction by taking G_1, \dots, G_m to be extremal graphs for F yields the lower bound $m \cdot \text{ex}(n, F)$. This is the benchmark we use to determine $\text{ex}_2^*(n, m, F)$.

In this work, we will first discuss the induced case, and then shift our focus to the general case. At the end, we will discuss the case where F is bipartite.

1.1 Definitions and Notation

Let $G = (V, E)$ be a graph. Let $V(G) = V$ denote the vertex set and $E(G) = E$ denote the edge set of G . We note by $v(G) = |V|$ the number of vertices and $e(G) = |E|$ the number of edges in G . For vertex $v \in V(G)$, we denote by $N_G(v) = \{u \in V(G) : \{u, v\} \in E(G)\}$ the neighborhood of v .

Given G_1, \dots, G_m subgraphs of G , we denote G_{i_1, \dots, i_k} as the subgraph of G with edge set $E(G_{i_1, \dots, i_k}) = \bigcap_{\alpha=1}^k E(G_{i_\alpha})$.

In this thesis, we reserve n to denote the number of vertices in a graph. Given a graph F , we denote $\text{ex}(n, F)$ to be the extremal number for F on a graph with n vertices, i.e. the maximum number of edges in a n -vertex graph that does not contain F as a subgraph.

We call a n -vertex complete graph K_n , and a complete bipartite graph $K_{a,b}$, where a, b are the size of its parts. We denote P_n as a path with n edges, and C_n as a cycle with n edges. Given graph G, H , define $G + H$ as the graph fully connecting G, H , i.e. $V(G+H) = V(G) \cup V(H)$ and $E(G+H) = E(G) \cup E(H) \cup \{\{u, v\} : u \in V(G), v \in V(H)\}$.

We also denote the set of first n positive integers as $[n] = \{1, 2, \dots, n\}$. Given a set X , we denote 2^X as the power set of X .

2 Induced Version

In this section, we assume that G_1, \dots, G_m are induced subgraphs of $G_1 \cup \dots \cup G_m$ and $E(G_i) \cap E(G_j)$ does not contain F for $i \neq j$. Here, we say that the extremal condition for m subgraphs is met if $\sum_{i=1}^m e(G_i) = \text{ex}_2^*(n, m, F)$.

TODO: add the condition for all G_i 's to be extremal graphs for F for all m , and generalize to hypergraph.

The following lemma shows that the problem can be reduced to only two graphs.

Lemma 2.1. *Let $n, m, k \geq 1$ such that $2 \leq k \leq m$, and let F be some graph. Then*

$$\text{ex}_2^*(n, m, F) \leq \frac{m}{k} \cdot \text{ex}_2^*(n, k, F).$$

Moreover, if the extremal condition for k induced subgraphs is met only when $G_1 = \dots = G_k$, then the above equality holds and the extremal condition for m induced subgraphs is met if and only if $G_1 = \dots = G_m$.

Not putting equality because I'm unsure if a construction for k subgraphs can always generalize to m subgraphs. For example, if $F = K_3$ and n is odd, the $G_1 = K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor}$ and $G_2 = K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1$ construction cannot be generalized to $m = n+1$ subgraphs.

Proof. Let G_1, \dots, G_m be induced subgraphs of $G_1 \cup \dots \cup G_m$ with $E(G_i) \cap E(G_j)$ not containing F for $i \neq j$. Put $G_{i+m} = G_i$ for all $i \in [m]$. Then

$$\sum_{i=1}^m e(G_i) = \frac{1}{k} \sum_{i=1}^m [e(G_i) + \dots + e(G_{i+k-1})] \leq \frac{1}{k} \sum_{i=1}^m \text{ex}_2^*(n, k, F) = \frac{m}{k} \cdot \text{ex}_2^*(n, k, F),$$

which establishes the upper bound.

Suppose $\sum_{i=1}^k e(G_i) = \text{ex}_2^*(n, k, F)$. By assumption $G_1 = \dots = G_k$, so $e(G_i) = \text{ex}_2^*(n, k, F)/k$ for $1 \leq i \leq k$. Hence, the construction $G_1 = \dots = G_m$ meets the upperbound. On the other hand, if $G_1 \neq G_2$ then $\sum_{i=1}^k e(G_i) < \text{ex}_2^*(n, k, F)$. Since $\sum_{i=1}^k e(G_{i+j}) \leq \text{ex}_2^*(n, k, F)$ for all $j \geq 1$, we have $\sum_{i=1}^m e(G_i) < \frac{m}{k} \cdot \text{ex}_2^*(n, k, F)$. Thus the extremal condition is met only when $G_1 = \dots = G_m$. \square

Lemma 2.1 allows us to reduce the problem to the case for two subgraphs G_1, G_2 . Let $C = V(G_1) \cap V(G_2)$, $a = |V(G_1) \setminus C|$, $b = |V(G_2) \setminus C|$, and $c = |C|$. Note that $a + b + c = n$. Since G_1, G_2 are induced subgraphs of $G_1 \cup G_2$, $G_1[C] = G_2[C] = G_{1,2}$. But then $G_{1,2}$ is F -free, so $e(G_1[C]) = e(G_2[C]) \leq \text{ex}(c, F)$. Given a, b, c , the optimal construction to maximize the number of edges over G_1, G_2 is to start with G_1, G_2 as complete graphs

and then remove edges from their intersection $G_{1,2}$ until it is F -free. This yields the inequality

$$e(G_1) + e(G_2) \leq \binom{a+c}{2} + \binom{b+c}{2} - 2 \left[\binom{c}{2} - \text{ex}(c, F) \right].$$

Define construction function $\mathcal{C}(a, b, c, F)$ as the expression on the right-hand-side of the above inequality. The maximum of \mathcal{C} tells us the extremal condition for two subgraphs. In particular, our benchmark construction of $G_1 = G_2$ being extremal graphs for F on n vertices meets the extremal condition if \mathcal{C} is maximized when $c = n$. Thus the problem is further reduced to an optimization problem over \mathcal{C} . For $0 \leq b \leq n - k$ and $k \leq c \leq n$, define

$$\begin{aligned} \Delta_k \mathcal{C}(a, b, c, F) &= \mathcal{C}(a, b, c, F) - \mathcal{C}(a, b+k, c-k, F) \\ &= \binom{a+c}{2} - \binom{a+c-k}{2} - 2 \left[\binom{c}{2} - \binom{c-k}{2} - \text{ex}(c, F) + \text{ex}(c-k, F) \right], \end{aligned}$$

and denote $\Delta \mathcal{C} = \Delta_1 \mathcal{C}$. Most of the work in this section will show that the maximum of \mathcal{C} happens when $c \geq n - k$ by proving that $\Delta_k \mathcal{C}(a, b, c, F) > 0$ for all $c \leq n - k$.

Lemma 2.2. *Let $n, c_0 \geq 1$, $m \geq 2$, and F be some graph. If $\mathcal{C}(a, b, c, F) < 2 \cdot \text{ex}(n, F)$ for $0 \leq c < c_0$ and $\text{ex}(c, F) - \text{ex}(c-1, F) > \frac{c-1}{2}$ for $c_0 \leq c \leq n$, then*

$$\text{ex}_2^*(n, m, F) = m \cdot \text{ex}(n, F)$$

and the extremal condition is met if and only if $G_1 = \dots = G_m$ are extremal graphs for F .

Not sure if this statement can be strengthened into if and only if. We would need to show there's no F such that $\text{ex}(c, F) - \text{ex}(c-1, F) \leq \frac{c-1}{2}$ for some c but \mathcal{C} has a unique maximum at $c = n$. In other words, there can not exist F such that its extremal number increases slowly at some small c but eventually catches up.

Proof. It suffices to show $\mathcal{C}(a, b, c, F)$ has a unique maximum of $2 \cdot \text{ex}(n, F)$ at $c = n$. We may assume $c \geq c_0$ by assumption. Suppose $b < n$. Since $\text{ex}(c, F) - \text{ex}(c-1, F) > \frac{c-1}{2}$,

$$\begin{aligned} \Delta \mathcal{C}(a, b, c, F) &= \binom{a+c}{2} - \binom{a+c-1}{2} \\ &\quad - 2 \left[\binom{c}{2} - \binom{c-1}{2} - \text{ex}(c, F) + \text{ex}(c-1, F) \right] \\ &= a - c + 1 + 2[\text{ex}(c, F) - \text{ex}(c-1, F)] > a \geq 0. \end{aligned}$$

Thus, \mathcal{C} is strictly increasing with respect to c for $c \geq c_0$. By symmetry of a and b , \mathcal{C} has a unique maximum of $2 \cdot \text{ex}(n, F)$ at $c = n$, which yields the unique extremal construction of $G_1 = G_2$ being extremal graphs for F on n vertices. \square

2.1 Complete Graph Case

Lemma 2.3. *For $n \geq 1$ and $r \geq 2$,*

$$\text{ex}(n, K_{r+1}) - \text{ex}(n-1, K_{r+1}) \geq \frac{n-1}{2},$$

with equality if and only if n is odd and $r = 2$.

Proof. By Turán's Theorem,

$$\text{ex}(n, K_{r+1}) - \text{ex}(n-1, K_{r+1}) = \delta(T_r(n)) = n - \left\lceil \frac{n}{r} \right\rceil \geq n - \left\lceil \frac{n}{2} \right\rceil.$$

The result now follows. \square

The following theorem for complete graphs with more than 3 vertices now follows directly from Lemma 2.2 and Lemma 2.3:

Theorem 2.4. *For $n \geq 1$, $m \geq 2$, and $r \geq 3$,*

$$\text{ex}_2^*(n, m, K_{r+1}) = m \cdot e(T_r(n)),$$

and the extremal condition is met if and only if $G_1 = \dots = G_m = T_r(n)$.

Surprisingly, the triangle case is more complicated than the case for larger complete graphs. As shown in Lemma 2.3, the condition given by Lemma 2.2 is not satisfied for all n in the triangle case, and there are indeed constructions of induced subgraphs G_1, G_2 that meet the extremal condition but are neither equal nor both complete bipartite graphs. For odd n , consider $G_1 = K_{\frac{n-1}{2}, \frac{n-1}{2}}$ and $G_2 = K_{\frac{n-1}{2}, \frac{n-1}{2}} + K_1$. The number of edges over G_1, G_2 is $\frac{(n-1)^2}{2} + n - 1 = \frac{n^2-1}{2} = 2 \left\lfloor \frac{n^2}{4} \right\rfloor$, which meets the benchmark construction of two complete bipartite graphs. We will show that this is the only deviant construction for the triangle case.

Theorem 2.5. *For $n \geq 1$ and $m \geq 2$,*

$$\text{ex}_2^*(n, m, K_3) = m \left\lfloor \frac{n^2}{4} \right\rfloor.$$

Moreover, if n is even or m is odd, then the extremal condition is met if and only if $G_1 = \dots = G_m$ are complete balanced bipartite graphs on n vertices. Otherwise, the extremal condition is met if and only if when either $G_1 = \dots = G_m = K_{\frac{n+1}{2}, \frac{n-1}{2}}$ or $G_{2k+1} = K_{\frac{n-1}{2}, \frac{n-1}{2}}$ and $G_{2k} = K_{\frac{n-1}{2}, \frac{n-1}{2}} + K_1$ for all $k \in [m/2]$.

Proof. We first show the following claim.

Claim 2.5.1. $\text{ex}_2^*(n, 2, K_3) = 2 \left\lfloor \frac{n^2}{4} \right\rfloor$, *and the extremal condition is met only when $G_1 = G_2 = K_{\lfloor \frac{n}{2} \rfloor, \lceil \frac{n}{2} \rceil}$, unless n is odd, $G_1 = K_{\frac{n-1}{2}, \frac{n-1}{2}}$, and $G_2 = K_{\frac{n-1}{2}, \frac{n-1}{2}} + K_1$.*

Proof. Consider $\Delta_2\mathcal{C}(a, b, c, K_3)$.

$$\begin{aligned} \Delta_2\mathcal{C}(a, b, c, K_3) &= \binom{a+c}{2} - \binom{a+c-2}{2} \\ &\quad - 2 \left[\binom{c}{2} - \binom{c-2}{2} - \left\lfloor \frac{c^2}{4} \right\rfloor + \left\lfloor \frac{(c-2)^2}{4} \right\rfloor \right] \\ &= 2(a+c) + 1 - 2[2c + 1 - (c+1)] \\ &= 2a + 1 > 0. \end{aligned}$$

This shows that $\mathcal{C}(a, b, c, K_3)$ has a maximum of $2 \left\lfloor \frac{n^2}{4} \right\rfloor$ when $c \geq n - 1$, which implies $\text{ex}_2^*(n, 2, K_3) = 2 \left\lfloor \frac{n^2}{4} \right\rfloor$. We are done if $c = n$, so assume that $c = n - 1$ and $a = 0$. Then in the extremal condition, $G_1 = G_{1,2} = K_{\lfloor \frac{n-1}{2} \rfloor, \lceil \frac{n-1}{2} \rceil}$ and

$$e(G_1) + e(G_2) = 2 \left\lfloor \frac{(n-1)^2}{4} \right\rfloor + \deg(v),$$

where v is the only vertex not in C . But then to meet the extremal condition,

$$\deg(v) = 2 \left\lfloor \frac{n^2}{4} \right\rfloor - 2 \left\lfloor \frac{(n-1)^2}{4} \right\rfloor = \begin{cases} n & \text{if } n \text{ is even,} \\ n-1 & \text{if } n \text{ is odd.} \end{cases}$$

Hence, n must be odd and G_2 must be a copy of G_1 with all vertices adjacent to the only remaining vertex, i.e. $G_2 = G_1 + K_1$. \square

By Lemma 2.1 and the above claim, it remains to show that for odd m , $G_1 = \dots = G_m = K_{\frac{n+1}{2}, \frac{n-1}{2}}$ when the extremal condition is met. Suppose not. the above claim then guarantees $G_i = K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1$ for some i . Put $G_{j+m} = G_j$ for all j . By applying the claim repeatedly,

$$\begin{aligned} G_i &= K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1 \\ G_{i+1} &= K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} \\ G_{i+2} &= K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1 \\ &\vdots \\ G_{i+m} &= K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor}, \end{aligned}$$

as m is odd. But then $G_{i+m} = G_i = K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1$, and this contradiction completes the proof. \square

Since n cannot be both even and odd, we have the following corollary:

Corollary 2.6. *For $n \geq 1$,*

$$\text{ex}_2^*(n, n, K_3) = n \left\lfloor \frac{n^2}{4} \right\rfloor,$$

and the extremal condition is met if and only if $G_1 = \dots = G_n$ are complete balanced bipartite graphs on n vertices.

Combining Theorem 2.4 and Theorem 2.5, we have the following theorem for complete graphs:

Theorem 2.7. *For $n \geq 1$ and $m, r \geq 2$,*

$$\text{ex}_2^*(n, m, K_{r+1}) = m \cdot \text{ex}(n, K_{r+1}).$$

Moreover, the extremal condition is met if and only if $G_1 = \dots = G_m = T_r(n)$, unless $r = 2$, n is odd, and m is even, in which case the extremal condition is met if and only if when either $G_1 = \dots = G_m = K_{\frac{n+1}{2}, \frac{n-1}{2}}$ or $G_{2k-1} = K_{\frac{n-1}{2}, \frac{n-1}{2}}$ and $G_{2k} = K_{\frac{n-1}{2}, \frac{n-1}{2}} + K_1$ for all $k \in [m/2]$.

2.2 Non-bipartite F

For non-bipartite F , it is hard to determine the extremal graphs for F in general, but their structures becomes more apparent when n is large. More specifically, the Erdős-Stone Theorem tells us that for large n , the extremal graph for F mimics the structure of the Turán graph. With this idea in mind, the following theorem is a generalization of Theorem 2.7 for large n .

Theorem 2.8. *For $n \geq 1$, and $m, r \geq 2$, if F is $(r+1)$ -colorable, then large enough n*

$$\text{ex}_2^*(n, m, F) = m \cdot \text{ex}(n, F),$$

and the extremal condition is met if and only if $G_1 = G_2 = \dots = G_n$ are extremal graphs for F , unless $r = 2$, n is odd, and m is even, in which case the extremal condition is met if and only if when either $G_1 = \dots = G_m = K_{\frac{n+1}{2}, \frac{n-1}{2}}$ or $G_{2k-1} = K_{\frac{n-1}{2}, \frac{n-1}{2}}$ and $G_{2k} = K_{\frac{n-1}{2}, \frac{n-1}{2}} + K_1$ for all $k \in [m/2]$.

Proof. It suffices to show for $m = 2$ by Lemma 2.1. We first show that $\mathcal{C}(a, b, c, F)$ fails to meet the desired bound for small c .

Claim 2.8.1. *If $c \leq \frac{n}{2}$, then $f(a, b, c) < 2 \cdot \text{ex}(n, F)$.*

Proof. Write $c = kn$ for some $k \in [0, 1/2]$. Since

$$f(a, b, kn) \leq 2 \binom{(1-k)n/2}{2} - 2 \left[\binom{kn}{2} - \text{ex}(kn, F) \right],$$

it suffices to show that

$$\text{ex}(n, F) - \text{ex}(kn, F) > \binom{(1-k)n/2}{2} - \binom{kn}{2}$$

for all $k \in [0, 1/2]$. By the Erdős-Stone theorem, $\text{ex}(n, F) = (1 - \frac{1}{r}) \frac{n^2}{2} + o(n^2)$ and so the left-hand-side is at least

$$\text{ex}(n, F) - \text{ex}(kn, F) \geq \text{ex}(n, F) - \text{ex}(n/2, F) \geq \left(1 - \frac{1}{r}\right) \left(\frac{n^2}{2} - \frac{n^2}{8}\right) - o(n^2) \geq \frac{3n^2}{16} - o(n^2).$$

On the right-hand-side,

$$\binom{(1-k)n/2}{2} - \binom{kn}{2} = (1 - 2k - 3k^2) \frac{n^2}{8} + o(n^2) \leq \frac{n^2}{8} + o(n^2).$$

Combining the above inequalities now yields the claim, as n is large. \square

Thus we may assume that $c > \frac{n}{2}$. A theorem of Simonovits states that for large enough c , $\text{ex}(c, F) = \text{ex}(c, K_{r+1}) + \text{ex}(c, \tilde{F})$, where \tilde{F} is the family of residue subgraphs of F after F is embedded into $T_r(c)$. This implies

$$\text{ex}(c, F) - \text{ex}(c-1, F) \geq \text{ex}(c, K_{r+1}) - \text{ex}(c-1, K_{r+1}),$$

as we assume n is sufficiently large. Thus by Lemma 2.2 and Lemma 2.3, we are done if $r > 3$. The above inequality also implies that for $r = 2$,

$$\Delta_2 \mathcal{C}(a, b, c, F) \geq \Delta_2 \mathcal{C}(a, b, c, K_3),$$

which is positive by the proof of Claim 2.5.1. Thus when c is $n - 1$ or n , $\mathcal{C}(a, b, c, F)$ attains its maximum, and plugging in the values yields

$$\mathcal{C}(a, b, c, F) \leq \max[2 \cdot \text{ex}(n, F), 2 \cdot \text{ex}(n - 1, F) + n - 1].$$

Assume that $a = 0$. By Lemma 2.3,

$$2 \cdot \text{ex}(n, F) - [2 \cdot \text{ex}(n - 1, F) + n - 1] \geq 2[\text{ex}(n, K_3) - \text{ex}(n - 1, K_3)] - n + 1 \geq 0,$$

with equality only if n is odd. Hence, $\mathcal{C}(a, b, c, F) \leq 2 \cdot \text{ex}(n, F)$. We may assume that n is odd and $c = n - 1$, otherwise we are done by Lemma 2.1. Then in the extremal condition, $G_1 = G_{1,2}$ is the extremal graph for F on $n - 1$ vertices, and G_2 must be $G_1 + K_1$. It remains to show that for odd m , $G_1 = \dots = G_m$ are extremal graphs for F when the extremal condition is met, and this follows from the parity argument in the proof of Theorem 2.5. \square

3 General Version

TODO: add introduction.

Theorem 3.1. *For all n and graph F ,*

$$\text{ex}_2(n, m, F) = m(1 + o(1))\text{ex}(n, F)$$

as $m \rightarrow \infty$.

Proof. Let $r = v(F)$. Pick $\epsilon > 0$. Reorder G_1, \dots, G_m so that $G_1, \dots, G_{m'}$ are all the G_i 's containing at least $(1 + \epsilon)\text{ex}(n, F)$ edges. A theorem of Simonovits states that G contains at least δn^r copies of F for some $\delta = \delta(\epsilon)$. Since there can be at most $\binom{n}{r}$ copies of F across all G_i 's,

$$m' \delta n^r \leq \binom{n}{r} \leq n^r \implies m' \leq \frac{1}{\delta}.$$

It now follows that

$$\begin{aligned} \sum_{i=1}^m e(G_i) &= \sum_{i=1}^{m'} e(G_i) + \sum_{i=m'+1}^m e(G_i) \\ &\leq \frac{1}{\delta} \binom{n}{2} + \left(m - \frac{1}{\delta}\right) (1 + \epsilon) \text{ex}(n, F) \\ &= m \left[1 + \epsilon + \frac{1}{m\delta} \left(\frac{\binom{n}{2}}{\text{ex}(n, F)} - (1 + \epsilon) \right) \right] \text{ex}(n, F). \end{aligned}$$

Since ϵ is arbitrary, the result follows. \square

Theorem 3.2. *For large enough n , suppose that G_1, \dots, G_m are graphs on common vertex set $[n]$ with no copy of F contained in any k of the G_i 's. If there exists extremal F -free subgraph H on n vertices such that $\binom{m}{k} \Delta(H) = o(n^{1/2})$, then*

$$\text{ex}_2(n, m, F) = (k - 1) \binom{n}{2} + \text{ex}(n, F) \binom{m}{k}.$$

Proof. For $S \subseteq [m]$, let E_S denote the set of edges that are contained in exactly $\{G_i\}_{i \in S}$. Then

$$\sum_{i=1}^m e(G_i) = \sum_{S \subseteq [m]} |S| |E_S| \leq (k - 1) \binom{n}{2} + \sum_{S \subseteq [m], |S| \geq k} (|S| - k + 1) |E_S|.$$

Let $A_S = \bigcup_{T \supseteq S} E_T$, i.e. the set of edges that are contained in all G_i with $i \in S$. When $|S| \geq k$, the edge set A_S is F -free and thus

$$\sum_{T \supseteq S} |E_T| \leq \text{ex}(n, F).$$

Hence,

$$\sum_{\substack{S \subseteq [m] \\ |S| \geq k}} (|S| - k + 1) |E_S| = \sum_{\substack{S \subseteq [m], T \subseteq S \\ |S|=k}} \sum_{\substack{|T| \geq k}} \frac{(|T| - k + 1) |E_T|}{\binom{|T|}{k}} \leq \sum_{\substack{S \subseteq [m], T \subseteq S \\ |S|=k}} \sum_{|T| \geq k} |E_T| \leq \binom{m}{k} \text{ex}(n, F),$$

as each $T \in [m]$ with $|T| \geq k$ is counted $\binom{|T|}{k}$ times in total and $|T| - k + 1 \leq \binom{|T|}{k}$. This proves the upper bound.

Now we show the bound is tight. In particular, we need to show there exists a construction such that the graph with edge set E_S is an extremal F -free graph, for all $S \subseteq [m]$ of size k . Let $M = \binom{m}{k}$ and H_1, \dots, H_M be copies of an extremal F -free graph on n vertices with $\Delta(H_i) = o(n^{1/2})$ for all i . It suffices to show that we can embed each H_i onto $[n]$ such that their edge sets are pairwise disjoint. We begin by an arbitrary embedding of each H_i and iteratively decrease the number of intersecting edges. Define a (u, v, i) -swap by swapping the embedding of vertex u and v of H_i , i.e. replacing each edge $\{u, w\} \in E(H_i)$ with the edge $\{v, w\}$ and each edge $\{v, w\} \in E(H_i)$ with the edge $\{u, w\}$. This preserves the type of isomorphism of H_i . Given a vertex v , let $N(v) = N_{H_1}(v) \cup \dots \cup N_{H_M}(v)$. Suppose there exists an intersecting edge $\{u, w\} \in E(H_i) \cap E(H_j)$. Since $|N(u)| \leq M \cdot \Delta(H_i) = o(n^{1/2})$, $|N(u) \cup N(N(u))| = o(n)$ so there exists a vertex $v \notin N(u) \cup N(N(u))$. Since $N(u) \cap N(v) = \emptyset$, performing a (u, v, i) -swap reduces the number of intersecting edges. The result now follows from iterating this process. \square

3.1 Triangle F

Consider F to be a triangle. Simply counting the number of triangles in each G_i shows the following:

Theorem 3.3. *For all n, m and $\epsilon > 0$,*

$$\text{ex}_2(n, m, K_3) < \left(m \cdot \frac{1 + \epsilon}{4} + \frac{1}{2\epsilon} - \frac{1}{2} \right) n^2.$$

Claim 3.3.1. *There are less than $\frac{2}{\epsilon}$ number of G_i 's with at least $(1 + \epsilon)\frac{n^2}{4}$ edges.*

Proof. Suppose $e(G_i) \geq (1 + \epsilon)\frac{n^2}{4}$ for $1 \leq i \leq k$. Let $K_3(G)$ denote the number of triangles in graph G . By the Moon-Moser inequality,

$$K_3(G_i) \geq \epsilon(1 + \epsilon)\frac{n^3}{12}.$$

Since there are no overlapping triangles from different G_i 's,

$$\binom{n}{3} \geq \sum_{i=1}^k K_3(G_i) \geq \frac{\epsilon(1 + \epsilon)}{12} kn^3.$$

Rearranging yields $k < \frac{2}{\epsilon}$. \square

By the claim,

$$\sum_{i=1}^m e(G_i) < \frac{2}{\epsilon} \binom{n}{2} + \left(m - \frac{2}{\epsilon}\right) (1 + \epsilon) \frac{n^2}{4} \leq m(1 + \epsilon) \frac{n^2}{4} + (1 - \epsilon) \frac{n^2}{2\epsilon},$$

which proves Theorem 3.3.

When $m = 2$,

$$e(G_1) + e(G_2) \leq \binom{n}{2} + e(G_{1,2}) \leq \binom{n}{2} + \left\lfloor \frac{n^2}{4} \right\rfloor,$$

which meets the benchmark bound and so $\text{ex}_2(n, 2, K_3) = \binom{n}{2} + \left\lfloor \frac{n^2}{4} \right\rfloor$.

This result is also true for $m = 3$:

Proposition 3.4. *For all n ,*

$$\text{ex}_2(n, 3, K_3) = \binom{n}{2} + \left\lfloor \frac{n^2}{2} \right\rfloor.$$

Proof. Define $H_k \subseteq G$ be the graph with edges contained in at least k number of G_i 's and note that $e(G_1) + e(G_2) + e(G_3) = e(H_1) + e(H_2) + e(H_3)$. Thus it suffices to show that $e(H_2) + e(H_3) \leq \frac{n^2}{2}$. Notice H_2 must not contain any triangles with two edges in H_3 , so

$$e(H_2) + e(H_3) \leq \binom{n}{2} + e(H_3) - |\{\{u, v\} : u \neq v, N_{H_3}(u) \cap N_{H_3}(v) \neq \emptyset\}|.$$

Let H'_3 be the graph with the same vertex set as H_3 and edge set $\{\{u, v\} : u \neq v, N_{H_3}(u) \cap N_{H_3}(v) \neq \emptyset\}$. It suffices to show that $\frac{n}{2} \geq e(H_3) - e(H'_3)$.

Let $d_1 \geq d_2 \geq \dots \geq d_n$ and $f_1 \geq f_2 \geq \dots \geq f_n$ each be the degree sequence of H_3 and H'_3 , respectively. We show that $f_i \geq d_i - 1$ for all i . Let v_i denote the vertex in H with degree d_i and u_i be the vertex in H with degree f_i . Let $S_i = |N_{H_3}(v_1) \cup \dots \cup N_{H_3}(v_i)|$. Since

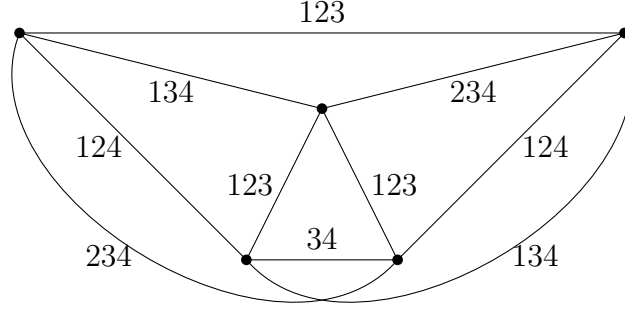
$$\sum_{u \in S_i} d_{H_3}(u) \geq d_1 + \dots + d_i,$$

we have that $|S_i| \geq i$. But then $S_i \setminus \{u_1, \dots, u_{i-1}\}$ is non-empty, and every $u \in S_i$ has degree $d_{H'_3}(u) \geq d_i - 1$. Hence, $f_i \geq d_i - 1$ for all i , which yields

$$e(H'_3) = \frac{1}{2} \sum_{i=1}^n f_i \geq \frac{1}{2} \sum_{i=1}^n (d_i - 1) = e(H_3) - \frac{n}{2}.$$

□

However, the bound in Proposition 3.1 is not tight for $m \geq 4$, as shown in the following graph:



The number on each edge denotes the set of G_i 's that contain the edge.

The above graph contains 29 edges, which exceeds the bound $\binom{5}{2} + 3\lfloor \frac{5^2}{4} \rfloor = 28$ by 1. By blowing up the above graph, we can construct a graph with $n \in 10\mathbb{Z}$ vertices that contains

$$5 \binom{n/5}{2} + 29 \cdot \frac{(n/5)^2}{4}$$

edges, which exceeds the bound $\binom{n}{2} + 3\lfloor \frac{n^2}{4} \rfloor$ by $n^2/100$.

3.2 Bipartite F

In this section, we discuss the case where F is bipartite. In particular, we focus on the cases where $F \subseteq K_{2,2}$ is P_2 , a path of length 2, or M_2 , a matching with two edges.

Theorem 3.5.

$$\text{ex}_2(n, m, P_2) \leq \left(\frac{1}{2} + o(1) \right) n^2 \sqrt{m}$$

as $n \rightarrow \infty$ or $m \rightarrow \infty$.

Proof. Since there are no overlapping P_2 's in different G_i 's,

$$\sum_{i=1}^m \#\{P_2 \subseteq G_i\} \leq \#\{P_2 \subseteq G\}$$

For each G_i , each vertex v in G_i and two of its neighbors form one unique P_2 , so

$$\#\{P_2 \subseteq G_i\} = \sum_{v \in V(G_i)} \binom{d_{G_i}(v)}{2}.$$

And by Jensen's inequality,

$$\sum_{v \in V(G_i)} \binom{d_{G_i}(v)}{2} \geq n \binom{d_{G_i}(v)/n}{2} = n \binom{2e(G_i)/n}{2} \geq \frac{2(e(G_i))^2}{n} - e(G_i).$$

On the other hand, since each three vertices in G can form at most three P_2 's,

$$\#\{P_2 \subseteq G\} \leq 3 \binom{n}{3} \leq \frac{n^3}{2}.$$

Combining the above inequalities yields

$$\frac{2m}{n} \left(\frac{1}{m} \sum_{i=1}^m e(G_i) \right)^2 - \sum_{i=1}^m e(G_i) \stackrel{\text{Jensen's}}{\leq} \sum_{i=1}^m \frac{2(e(G_i))^2}{n} - e(G_i) \leq \frac{n^3}{2},$$

and solving the quadratic equation gives

$$\sum_{i=1}^m e(G_i) \leq mn \cdot \frac{1 + \sqrt{4n^2/m + 1}}{4} = \left(\frac{1}{2} + o(1) \right) n^2 \sqrt{m},$$

as $n \rightarrow \infty$ or $m \rightarrow \infty$. □

When $m = n$, the following projective plane construction shows the above bound is tight asymptotically:

Theorem 3.6.

$$\text{ex}_2(n, n, P_2) = \left(\frac{1}{2} + o(1) \right) n^{5/2},$$

as $n \rightarrow \infty$.

Proof. It suffices to show the tightness of the bound in Theorem 3.5. Consider a finite projective plane of order q . The projective plane has $n = q^2 + q + 1$ points and n lines. Let $S_1, \dots, S_n \subseteq [n]$ be the n lines of the projective plane. Note that each line S_i contains $q + 1$ points, and the intersection of any two distinct lines S_i, S_j contains $|S_i \cap S_j| = 1$ point. Define G_1, \dots, G_n to be graphs on $[n]$, each with edge set $E(G_i) = \{\{j, k\} \subseteq [n] : j \neq k, j + k \in S_i \pmod n\}$. Note that the intersection of distinct G_i, G_j is P_2 free: since $|S_i \cap S_j| = 1$, if $\{a, b\}, \{a, c\} \in E(G_i) \cap E(G_j)$, then $a + b = a + c$ so $b = c$. We now count the number of edges in G_1, \dots, G_n . Since $|S_i| = q + 1$, for each point $j \in [n]$, there are $q + 1$ choices for $k \in [n]$ such that $j + k \in S_i$. But then we have to avoid counting the same edge twice and loops, so the number of edges in G_i is

$$e(G_i) = \frac{n(q + 1) - \#\text{loops counted for } G_i}{2}.$$

If $j \in [n]$ is even, then $k = j/2$ is the unique number in $[n]$ such that $k + k = j \pmod n$. If $j \in [n]$ is odd, then $k = (n + j)/2$ is the unique number in $[n]$ such that $k + k = j \pmod n$, as n is even. Hence, for each $j \in S_i$, there exists a unique $k \in [n]$ such that $k + k = j \pmod n$, and thus

$$\#\text{loops counted for } G_i = |S_i| = q + 1.$$

Since $q + 1 = (1 + o(1))n^{1/2}$, the number of edges in G_1, \dots, G_n is

$$\sum_{i=1}^n e(G_i) = n \cdot \frac{n(q + 1) - (q + 1)}{2} = \left(\frac{1}{2} + o(1) \right) n^{5/2},$$

as $n \rightarrow \infty$. □

Theorem 3.7. For all n, m ,

$$\text{ex}_2(n, m, M_2) \leq n^{5/2}.$$

Proof. Notice that $\#\{M_2 \subseteq G\} = \binom{e(G)}{2}$. On the other hand, each four vertices in G can form at most three M_2 's, so $\#\{M_2 \subseteq G\} \leq 3\binom{n}{4} \leq \frac{n^4}{8}$. By the same argument as in Theorem 3.4, we have

$$\sum_{i=1}^n \binom{e(G_i)}{2} = \sum_{i=1}^n \#\{M_2 \subseteq G_i\} \leq \#\{M_2 \subseteq G\} \leq \frac{n^4}{8}.$$

By Jensen's inequality,

$$\sum_{i=1}^n \binom{e(G_i)}{2} \geq n \binom{\sum_{i=1}^n e(G_i)/n}{2} = \frac{1}{2n} \left[\left(\sum_{i=1}^n e(G_i) \right)^2 - n \sum_{i=1}^n e(G_i) \right].$$

Combining the above inequalities yields

$$\left(\sum_{i=1}^n e(G_i) \right)^2 - n \sum_{i=1}^n e(G_i) \leq \frac{n^5}{4},$$

and solving the quadratic inequality gives

$$\sum_{i=1}^n e(G_i) \leq n^{5/2}.$$

□

We may obtain an exact result if we forbid both P_2 and M_2 at the same time:

Theorem 3.8. *For all n, m ,*

$$\text{ex}_2(n, m, \{P_2, M_2\}) = n^2 - n.$$

Proof. Denote the set of G_i 's as $\{G_i\} = \{G_1, \dots, G_n\}$, and the set of distinct pairs of G_i 's as $\{G_i\}^2 = \{\{G_j, G_k\} : j \neq k\}$. Consider the bipartite graph H with vertex set $V(H) = \{G_i\} \sqcup E(K_n)$ and edge set $E(H) = \{\{G_j, e\} \in \{G_i\} \times E(K_n) : e \in G_j\}$. Define $\phi : \{G_i\}^2 \rightarrow 2^{E(K_n)}$ by sending each $\{G_j, G_k\}$ to their common edge set $E(G_j) \cap E(G_k)$. Notice that each distinct G_j, G_k have at most one edge in common, so $|\phi(G_j, G_k)| \leq 1$. On the other hand, each edge $e \in E(G)$ can be obtained via ϕ by $\binom{d_H(e)}{2}$ possible distinct pairs (G_j, G_k) , and thus $|\phi^{-1}(e)| = \binom{d_H(e)}{2}$. But then

$$\binom{n}{2} \geq \sum_{(G_j, G_k) \in \{G_i\}^2} |\phi(G_j, G_k)| = \sum_{e \in E(K_n)} |\phi^{-1}(e)| = \sum_{e \in E(K_n)} \binom{d_H(e)}{2}.$$

By Jensen's inequality,

$$\sum_{e \in E(K_n)} \binom{d_H(e)}{2} \geq \binom{n}{2} \binom{\sum_{e \in E(K_n)} d_H(e)/n}{2} = \binom{n}{2} \binom{\sum_{i=1}^n e(G_i)/n}{2}.$$

Combining the above inequalities yields

$$2 \binom{n}{2}^2 \geq \left(\sum_{i=1}^n e(G_i) \right)^2 - \binom{n}{2} \sum_{i=1}^n e(G_i),$$

and the result now follows from solving the quadratic inequality.

To see that this bound is tight, consider the construction such that for each distinct $i, j \in [n]$, $E(G_i) \cap E(G_j)$ contains exactly one unique edge $e \in K_n$. The number of edges in this construction is $2\binom{n}{2} = n^2 - n$. \square