

# Double Turán Problem

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# 1 Introduction

Let  $\text{ex}_2(n, m, F)$  be the maximum possible sum of the number of edges over  $m$  subgraphs  $G_1, \dots, G_m$  on the same vertex set  $[n]$ , with the constraint that  $E(G_i) \cap E(G_j)$  does not contain graph  $F$  for  $i \neq j$ . Our goal is to determine  $\text{ex}_2(n, m, F)$  for different forbidden graphs  $F$ . A trivial construction with  $G_1 = K_n$  and  $G_2, \dots, G_m$  to be extremal graphs for  $F$  yields the lower bound  $\binom{n}{2} + (m-1)\text{ex}(n, F)$ . In this work, we use this bound as a benchmark to either show the tightness of  $\text{ex}_2(n, m, F)$  or to provide a better bound.

Additionally, we are also interested in a more restrictive version where  $G_1, \dots, G_m$  are induced subgraphs of  $G_1 \cup \dots \cup G_m$ . We denote  $\text{ex}_2^*(n, m, F)$  as the maximum possible sum of the number of edges over  $m$  induced subgraphs  $G_1, \dots, G_m$  on the same vertex set  $[n]$  such that  $E(G_i) \cap E(G_j)$  does not contain graph  $F$  for  $i \neq j$ . The trivial construction by taking  $G_1, \dots, G_m$  to be extremal graphs for  $F$  yields the lower bound  $m \cdot \text{ex}(n, F)$ . This is the benchmark we use to determine  $\text{ex}_2^*(n, m, F)$ .

In this work, we will first discuss the induced case, and then shift our focus to the general case. At the end, we will discuss the case where  $F$  is bipartite.

## 1.1 Definitions and Notation

Let  $G = (V, E)$  be a graph. Let  $V(G) = V$  denote the vertex set and  $E(G) = E$  denote the edge set of  $G$ . We note by  $v(G) = |V|$  the number of vertices and  $e(G) = |E|$  the number of edges in  $G$ . For vertex  $v \in V(G)$ , we denote by  $N_G(v) = \{u \in V(G) : \{u, v\} \in E(G)\}$  the neighborhood of  $v$ .

Given  $G_1, \dots, G_m$  subgraphs of  $G$ , we denote  $G_{i_1, \dots, i_k}$  as the subgraph of  $G$  with edge set  $E(G_{i_1, \dots, i_k}) = \bigcap_{\alpha=1}^k E(G_{i_\alpha})$ .

In this thesis, we reserve  $n$  to denote the number of vertices in a graph. Given a graph  $F$ , we denote  $\text{ex}(n, F)$  to be the extremal number for  $F$  on a graph with  $n$  vertices, i.e. the maximum number of edges in a  $n$ -vertex graph that does not contain  $F$  as a subgraph.

We call a  $n$ -vertex complete graph  $K_n$ , and a complete bipartite graph  $K_{a,b}$ , where  $a, b$  are the size of its parts. We denote  $P_n$  as a path with  $n$  edges, and  $C_n$  as a cycle with  $n$  edges. Given graph  $G, H$ , define  $G + H$  as the graph fully connecting  $G, H$ , i.e.  $V(G+H) = V(G) \cup V(H)$  and  $E(G+H) = E(G) \cup E(H) \cup \{\{u, v\} : u \in V(G), v \in V(H)\}$ .

We also denote the set of first  $n$  positive integers as  $[n] = \{1, 2, \dots, n\}$ . Given a set  $X$ , we denote  $2^X$  as the power set of  $X$ .

## 2 Induced Version

In this section, we assume that  $G_1, \dots, G_m$  are induced subgraphs of  $G_1 \cup \dots \cup G_m$  and  $E(G_i) \cap E(G_j)$  does not contain  $F$  for  $i \neq j$ . Here, we say that the extremal condition for  $m$  subgraphs is met if  $\sum_{i=1}^m e(G_i) = \text{ex}_2^*(n, m, F)$ .

**TODO:** add the condition for all  $G_i$ 's to be extremal graphs for  $F$  for all  $m$ , and generalize to hypergraph.

The following lemma shows that the problem can be reduced to only two graphs.

**Lemma 2.1.** *Let  $n, m, k \in \mathbb{N}$  such that  $2 \leq k \leq m$ , and let  $F$  be a graph. Then*

$$\text{ex}_2^*(n, m, F) \leq \frac{m}{k} \cdot \text{ex}_2^*(n, k, F).$$

Moreover, if the extremal condition for  $k$  induced subgraphs is met only when  $G_1 = \dots = G_k$ , then the above equality holds and the extremal condition for  $m$  induced subgraphs is met only when  $G_1 = \dots = G_m$ .

Not putting equality because I'm unsure if a construction for  $k$  subgraphs can always generalize to  $m$  subgraphs. For example, if  $F = K_3$  and  $n$  is odd, the  $G_1 = K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor}$  and  $G_2 = K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1$  construction cannot be generalized to  $m = n+1$  subgraphs.

*Proof.* Let  $G_1, \dots, G_m$  be induced subgraphs of  $G_1 \cup \dots \cup G_m$  with  $E(G_i) \cap E(G_j)$  not containing  $F$  for  $i \neq j$ . Put  $G_{i+m} = G_i$  for all  $i \in [m]$ . Then

$$\sum_{i=1}^m e(G_i) = \frac{1}{k} \sum_{i=1}^m [e(G_i) + \dots + e(G_{i+k-1})] \leq \frac{1}{k} \sum_{i=1}^m \text{ex}_2^*(n, k, F) = \frac{m}{k} \cdot \text{ex}_2^*(n, k, F),$$

which establishes the upper bound.

Suppose  $\sum_{i=1}^k e(G_i) = \text{ex}_2^*(n, k, F)$ . By assumption  $G_1 = \dots = G_k$ , so  $e(G_i) = \text{ex}_2^*(n, k, F)/k$  for  $1 \leq i \leq k$ . Hence, the construction  $G_1 = \dots = G_m$  meets the upperbound. On the other hand, if  $G_1 \neq G_2$  then  $\sum_{i=1}^k e(G_i) < \text{ex}_2^*(n, k, F)$ . Since  $\sum_{i=1}^k e(G_{i+j}) \leq \text{ex}_2^*(n, k, F)$  for all  $j \geq 1$ , we have  $\sum_{i=1}^m e(G_i) < \frac{m}{k} \cdot \text{ex}_2^*(n, k, F)$ . Thus the extremal condition is met only when  $G_1 = \dots = G_m$ .  $\square$

**Lemma 2.2.** *Let  $n \geq 1$ ,  $m \geq 2$  and  $F$  be some graph. If  $\text{ex}(c, F) - \text{ex}(c-1, F) > \frac{c-1}{2}$  for all  $1 \leq c \leq n$ , then*

$$\text{ex}_2^*(n, m, F) = m \cdot \text{ex}(n, F)$$

*and  $G_1 = \dots = G_m$  are extremal graphs for  $F$  when the extremal condition is met.*

Not sure if this statement can be strengthened into if and only if. We would need to show there's no  $F$  such that  $\text{ex}(c, F) - \text{ex}(c-1, F) \leq \frac{c-1}{2}$  for some  $c$  but  $f$  has a unique maximum at  $c = n$ . In other words, there can not exist  $F$  such that its extremal number increases slowly at some small  $c$  but eventually catches up.

*Proof.* By Lemma 2.1, it suffices to show the case for two subgraphs. Let  $G_1, G_2$  be induced subgraphs of  $G_1 \cup G_2$  with  $E(G_1) \cap E(G_2)$  not containing  $F$ . Moreover, let  $C = V(G_1) \cap V(G_2)$  and put  $a = |V(G_1) \setminus C|$ ,  $b = |V(G_2) \setminus C|$ , and  $c = |C|$ . Since  $G_1, G_2$  are induced subgraphs,  $G_1[C] = G_2[C] = G_{1,2}$ . But then  $G_{1,2}$  is  $F$ -free, so  $e(G_1[C]) = e(G_2[C]) \leq \text{ex}(c, F)$ . This yields the inequality

$$e(G_1) + e(G_2) \leq \binom{a+c}{2} + \binom{b+c}{2} - 2 \left[ \binom{c}{2} - \text{ex}(c, F) \right].$$

Let  $f(a, b, c)$  denote the expression on the right-hand-side of the above inequality.

**Claim 1.**  $f$  has a unique maximum at  $c = n$ .

*Proof.* Suppose  $b < n$ . Since  $\text{ex}(c, F) - \text{ex}(c-1, F) > \frac{c-1}{2}$ ,

$$\begin{aligned} f(a, b, c) - f(a, b+1, c-1) &= \binom{a+c}{2} - \binom{a+c-1}{2} \\ &\quad - 2 \left[ \binom{c}{2} - \binom{c-1}{2} - \text{ex}(c, F) + \text{ex}(c-1, F) \right] \\ &= a - c + 1 + 2[\text{ex}(c+1, F) - \text{ex}(c, F)] > a \geq 0. \end{aligned}$$

Thus,  $f$  is strictly increasing with respect to  $c$ . By symmetry,  $f$  has a unique maximum at  $c = n$ .  $\square$

By the claim,

$$e(G_1) + e(G_2) \leq 2 \binom{n}{2} - 2 \left[ \binom{n}{2} - \text{ex}(n, F) \right] = 2 \cdot \text{ex}(n, F),$$

and  $G_1 = G_2 = G_{1,2}$  are extremal graphs for  $F$  on  $n$  vertices.  $\square$

## 2.1 Complete Graph Case

**Lemma 2.3.** For  $n \geq 1$  and  $r \geq 2$ ,

$$\text{ex}(n, K_{r+1}) - \text{ex}(n-1, K_{r+1}) \geq \frac{n-1}{2},$$

with equality if and only if  $n$  is odd and  $r = 2$ .

*Proof.* By Turán's Theorem,

$$\text{ex}(n, K_{r+1}) - \text{ex}(n-1, K_{r+1}) = \delta(T_r(n)) = n - \left\lceil \frac{n}{r} \right\rceil \geq n - \left\lceil \frac{n}{2} \right\rceil.$$

The result now follows.  $\square$

The following theorem for complete graphs with more than 3 vertices now follows directly from Lemma 2.2 and Lemma 2.3:

**Theorem 2.4.** *For  $n \geq 1$ ,  $m \geq 2$ , and  $r \geq 3$ ,*

$$\text{ex}_2^*(n, m, K_{r+1}) = m \cdot e(T_r(n)),$$

*and  $G_1 = \dots = G_m = T_r(n)$  when the extremal condition is met.*

Surprisingly, the triangle case is more complicated than the case for larger complete graphs. As shown in Lemma 2.3, the condition given by Lemma 2.2 is not satisfied for all  $n$  in the triangle case, and there are indeed constructions of induced subgraphs  $G_1, G_2$  that meet the extremal condition but are neither equal nor both complete bipartite graphs. For odd  $n$ , consider  $G_1 = K_{\frac{n-1}{2}, \frac{n-1}{2}}$  and  $G_2 = K_{\frac{n-1}{2}, \frac{n-1}{2}} + K_1$ . The number of edges over  $G_1, G_2$  is  $\frac{(n-1)^2}{2} + n - 1 = \frac{n^2-1}{2} = 2 \left\lfloor \frac{n^2}{4} \right\rfloor$ , which meets the benchmark construction of two complete bipartite graphs. Hence, for the triangle case we have to make some compromises.

**Theorem 2.5.** *For  $n \geq 1$  and  $m \geq 2$ ,*

$$\text{ex}_2^*(n, m, K_3) = m \left\lfloor \frac{n^2}{4} \right\rfloor.$$

*Moreover, if  $n$  is even or  $m$  is odd, then the extremal condition is met only when  $G_1 = G_2$  are complete balanced bipartite graphs on  $n$  vertices. Otherwise, the extremal condition is met only when either  $G_1 = G_2 = K_{\frac{n+1}{2}, \frac{n-1}{2}}$  or  $G_1 = K_{\frac{n-1}{2}, \frac{n-1}{2}}$  and  $G_2 = K_{\frac{n-1}{2}, \frac{n-1}{2}} + K_1$ .*

*Proof.* By Lemma 2.2 and Lemma 2.3, we are done if  $n$  is even, so we may assume that  $n$  is odd. We first show the following claim.

**Claim 2.**  $\text{ex}_2^*(n, 2, K_3) = 2 \left\lfloor \frac{n^2}{4} \right\rfloor$ , *and the extremal condition is met only when either  $G_1 = G_2 = K_{\frac{n+1}{2}, \frac{n-1}{2}}$  or  $G_1 = K_{\frac{n-1}{2}, \frac{n-1}{2}}$  and  $G_2 = K_{\frac{n-1}{2}, \frac{n-1}{2}} + K_1$ .*

*Proof.* Let  $a, b, c$ , and  $f(a, b, c)$  be defined as in the proof of Lemma 2.2. Then

$$\begin{aligned} f(a, b, c) - f(a, b+2, c-2) &= \binom{a+c}{2} - \binom{a+c-2}{2} \\ &\quad - 2 \left[ \binom{c}{2} - \binom{c-2}{2} - \left\lfloor \frac{c^2}{4} \right\rfloor + \left\lfloor \frac{(c-2)^2}{4} \right\rfloor \right] \\ &= 2(a+c) + 1 - 2[2c+1 - (c+1)] \\ &= 2a + 1 > 0. \end{aligned}$$

Hence,  $f$  attains its maximum of  $2 \left\lfloor \frac{n^2}{4} \right\rfloor$  when  $c = n$  or  $n-1$ , that is,  $a+b \leq 1$ . Thus  $\text{ex}_2^*(n, 2, K_3) \leq 2 \left\lfloor \frac{n^2}{4} \right\rfloor$ . Assume that  $a = 0$ . When  $c = n$ ,  $G_1 = G_2 = G_{1,2}$  are extremal graphs for  $K_3$ , which is the complete balanced bipartite graph on  $n$  vertices. When  $c = n-1$ ,  $G_1 = G_{1,2} = K_{\frac{n-1}{2}, \frac{n-1}{2}}$  and  $G_2$  must be a copy of  $G_1$  with all vertices adjacent to the only remaining vertex, i.e.  $G_2 = G_1 + K_1$ , to meet the extremal number.  $\square$

By Lemma 2.1 and the above claim,  $\text{ex}_2^*(n, m, K_3) = m \left\lfloor \frac{n^2}{4} \right\rfloor$ . It remains to show that for odd  $m$ ,  $G_1 = \dots = G_m = K_{\frac{n+1}{2}, \frac{n-1}{2}}$  when the extremal condition is met. Suppose not. Then the claim guarantees  $G_i = K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1$  for some  $i$ . Put  $G_{m+i} = i$ . By applying the claim repeatedly,

$$\begin{aligned} G_i &= K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1 \\ G_{i+1} &= K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} \\ G_{i+2} &= K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor} + K_1 \\ &\vdots \\ G_{i+m} &= K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor}, \end{aligned}$$

as  $m$  is odd. But then  $G_{i+m} = G_i = K_{\lceil \frac{n-1}{2} \rceil, \lfloor \frac{n-1}{2} \rfloor}$ , and this contradiction completes the proof.  $\square$

## 2.2 Non-bipartite $F$

**Theorem 2.6.** *Suppose  $F$  is  $(r+1)$ -colorable, with  $r \geq 2$ . For large enough  $n$ ,*

$$\text{ex}_2^*(n, n, F) = n \cdot \text{ex}(n, F).$$

*In particular, the extremal number is reached only if  $G_1 = G_2 = \dots = G_n$  are  $n$ -vertex extremal graphs for  $F$ .*

By the same argument as in Theorem 2.1, it suffices to prove the following lemma:

**Lemma 2.7.** *Let  $F$  be  $(r+1)$ -colorable, with  $r \geq 2$ . Suppose  $E(G_1) \cap E(G_2)$  does not include  $F$ . For large enough  $n$ ,*

$$e(G_1) + e(G_2) \leq 2 \cdot \text{ex}(n, F),$$

*with equality if and only if  $G_1 = G_2$  are  $n$ -vertex extremal graphs for  $F$ , unless  $n$  is odd,  $G_1$  is an  $(n-1)$ -vertex extremal graph for  $F$ , and  $G_2 = G_1 + K_1$ .*

*Proof.* Let  $C = V(G_1) \cap V(G_2)$ ,  $A = V(G_1) \setminus C$ , and  $B = V(G_2) \setminus C$ . Put  $a = |A|$ ,  $b = |B|$ ,  $c = |C|$ . Since  $G_1, G_2$  are induced graphs,  $E(G_1[C]) = E(G_2[C]) = E(G[C]) = E(G_i) \cap E(G_j)$ , which is  $F$ -free. Hence,

$$e(G_1) + e(G_2) \leq \binom{a+c}{2} + \binom{b+c}{2} - 2 \left[ \binom{c}{2} - \text{ex}(c, F) \right]. \quad (2.1)$$

Define  $f(a, b, c)$  as the function on the right-hand-side of (2.2). We show that  $f(a, b, c)$  attains its maximum at  $a = b = 0$  and  $c = n$ .

**Claim 3.** *If  $c \leq \frac{n}{2}$ , then  $f(a, b, c) < 2 \cdot \text{ex}(n, F)$ .*

*Proof.* Write  $c = kn$  for some  $k \in [0, 1/2]$ . Since

$$f(a, b, kn) \leq 2 \binom{(1-k)n/2}{2} - 2 \left[ \binom{kn}{2} - \text{ex}(kn, F) \right],$$

it suffices to show that

$$\text{ex}(n, F) - \text{ex}(c, F) > \binom{(1-k)n/2}{2} - \binom{kn}{2}$$

for all  $k \in [0, 1/2]$ . By the Erdős-Stone theorem,  $\text{ex}(n, F) = (1 - \frac{1}{r}) \frac{n^2}{2} + o(n^2)$  and so the left-hand-side is at least

$$\text{ex}(n, F) - \text{ex}(c, F) \geq \text{ex}(n, F) - \text{ex}(n/2, F) \geq \left(1 - \frac{1}{r}\right) \left(\frac{n^2}{2} - \frac{n^2}{8}\right) - o(n^2) \geq \frac{3n^2}{16} - o(n^2).$$

On the right-hand-side,

$$\binom{(1-k)n/2}{2} - \binom{kn}{2} = (1 - 2k - 3k^2) \frac{n^2}{8} + o(n^2) \leq \frac{n^2}{8} + o(n^2).$$

Combining the above inequalities now yields the claim, as  $n$  is large.  $\square$

Thus we may assume that  $c > \frac{n}{2}$ . A theorem of Simonovits states that for large enough  $n$ ,  $\text{ex}(n, F) = \text{ex}(n, K_{r+1}) + \text{ex}(n, \tilde{F})$ , where  $\tilde{F}$  is the family of residue subgraphs of  $F$  after  $F$  is embedded into  $T_r(n)$ . This implies

$$\text{ex}(n+1, F) - \text{ex}(n, F) \geq \text{ex}(n, K_{r+1}) - \text{ex}(n+1, K_{r+1}),$$

and so

$$f(a, b-2, c+2) - f(a, b, c) \geq 2a - 2c - 1 + 2[\text{ex}(c+2, K_{r+1}) - \text{ex}(c, K_{r+1})].$$

Since  $\text{ex}(c+1, K_{r+1}) - \text{ex}(c, K_{r+1}) \geq c - \lfloor \frac{c}{r} \rfloor \geq c - \lfloor \frac{c}{2} \rfloor$ , we have

$$\text{ex}(c+2, K_{r+1}) - \text{ex}(c, K_{r+1}) \geq c+1 - \left\lfloor \frac{c+1}{2} \right\rfloor + c - \left\lfloor \frac{c}{2} \right\rfloor = c+1,$$

and thus

$$f(a, b-2, c+2) - f(a, b, c) \geq 2a+1 > 0.$$

By symmetry, we also have  $f(a-2, b, c+2) > f(a, b, c)$ . Thus,  $f$  attains its maximum when  $c$  is  $n-1$  or  $n$ . Equation (2.2) now yields,

$$e(G_1) + e(G_2) \leq \max[2 \cdot \text{ex}(n, F), 2 \cdot \text{ex}(n-1, F) + n - 1].$$

Assume that  $a = 0$ . Since

$$\begin{aligned} 2 \cdot \text{ex}(n, F) - [2 \cdot \text{ex}(n-1, F) + n - 1] & \\ & \geq 2[\text{ex}(n, K_{r+1}) - \text{ex}(n-1, K_{r+1})] - n + 1 \quad (2.2) \\ & = 2 \left( n - \left\lceil \frac{n}{r} \right\rceil \right) - n + 1 \\ & \geq n + 1 - 2 \left\lceil \frac{n}{2} \right\rceil \geq 0, \end{aligned}$$



we have

$$e(G_1) + e(G_2) \leq 2 \cdot \text{ex}(n, F).$$

If  $c = n$ , the equality holds only if  $G_1 = G_2$  are  $n$ -vertex extremal graphs for  $F$ . Suppose  $c = n - 1$  and the equality holds. Observe that equation (2.3) is equal to zero only when  $r = 2$  and  $n$  is odd. Thus if  $c = n - 1$ , then the equality could only be achieved when  $r = 2$ ,  $n$  is odd,  $G_1$  is an  $(n - 1)$ -vertex extremal graph for  $F$ , and  $G_2 = G_1 + K_1$ .  $\square$

### 3 General Version

TODO: add introduction.

**Theorem 3.1.** *For all  $n$  and graph  $F$ ,*

$$\text{ex}_2(n, m, F) = m(1 + o(1))\text{ex}(n, F)$$

as  $m \rightarrow \infty$ .

*Proof.* Let  $r = v(F)$ . Pick  $\epsilon > 0$ . Reorder  $G_1, \dots, G_m$  so that  $G_1, \dots, G_{m'}$  are all the  $G_i$ 's containing at least  $(1 + \epsilon)\text{ex}(n, F)$  edges. A theorem of Simonovits states that  $G$  contains at least  $\delta n^r$  copies of  $F$  for some  $\delta = \delta(\epsilon)$ . Since there can be at most  $\binom{n}{r}$  copies of  $F$  across all  $G_i$ 's,

$$m' \delta n^r \leq \binom{n}{r} \leq n^r \implies m' \leq \frac{1}{\delta}.$$

It now follows that

$$\begin{aligned} \sum_{i=1}^m e(G_i) &= \sum_{i=1}^{m'} e(G_i) + \sum_{i=m'+1}^m e(G_i) \\ &\leq \frac{1}{\delta} \binom{n}{2} + \left(m - \frac{1}{\delta}\right) (1 + \epsilon) \text{ex}(n, F) \\ &= m \left[ 1 + \epsilon + \frac{1}{m\delta} \left( \frac{\binom{n}{2}}{\text{ex}(n, F)} - (1 + \epsilon) \right) \right] \text{ex}(n, F). \end{aligned}$$

Since  $\epsilon$  is arbitrary, the result follows.  $\square$

**Theorem 3.2.** *For large enough  $n$ , suppose that  $G_1, \dots, G_m$  are graphs on common vertex set  $[n]$  with no copy of  $F$  contained in any  $k$  of the  $G_i$ 's. If there exists extremal  $F$ -free subgraph  $H$  on  $n$  vertices such that  $\binom{m}{k} \Delta(H) = o(n^{1/2})$ , then*

$$\text{ex}_2(n, m, F) = (k - 1) \binom{n}{2} + \text{ex}(n, F) \binom{m}{k}.$$

*Proof.* For  $S \subseteq [m]$ , let  $E_S$  denote the set of edges that are contained in exactly  $\{G_i\}_{i \in S}$ . Then

$$\sum_{i=1}^m e(G_i) = \sum_{S \subseteq [m]} |S| |E_S| \leq (k - 1) \binom{n}{2} + \sum_{S \subseteq [m], |S| \geq k} (|S| - k + 1) |E_S|.$$

Let  $A_S = \bigcup_{T \supseteq S} E_T$ , i.e. the set of edges that are contained in all  $G_i$  with  $i \in S$ . When  $|S| \geq k$ , the edge set  $A_S$  is  $F$ -free and thus

$$\sum_{T \supseteq S} |E_T| \leq \text{ex}(n, F).$$

Hence,

$$\sum_{\substack{S \subseteq [m] \\ |S| \geq k}} (|S| - k + 1) |E_S| = \sum_{\substack{S \subseteq [m], T \subseteq S \\ |S|=k}} \sum_{T \subseteq S} \frac{(|T| - k + 1) |E_T|}{\binom{|T|}{k}} \leq \sum_{\substack{S \subseteq [m], T \subseteq S \\ |S|=k}} \sum_{T \subseteq S} |E_T| \leq \binom{m}{k} \text{ex}(n, F),$$

as each  $T \in [m]$  with  $|T| \geq k$  is counted  $\binom{|T|}{k}$  times in total and  $|T| - k + 1 \leq \binom{|T|}{k}$ . This proves the upper bound.

Now we show the bound is tight. In particular, we need to show there exists a construction such that the graph with edge set  $E_S$  is an extremal  $F$ -free graph, for all  $S \subseteq [m]$  of size  $k$ . Let  $M = \binom{m}{k}$  and  $H_1, \dots, H_M$  be copies of an extremal  $F$ -free graph on  $n$  vertices with  $\Delta(H_i) = o(n^{1/2})$  for all  $i$ . It suffices to show that we can embed each  $H_i$  onto  $[n]$  such that their edge sets are pairwise disjoint. We begin by an arbitrary embedding of each  $H_i$  and iteratively decrease the number of intersecting edges. Define a  $(u, v, i)$ -swap by swapping the embedding of vertex  $u$  and  $v$  of  $H_i$ , i.e. replacing each edge  $\{u, w\} \in E(H_i)$  with the edge  $\{v, w\}$  and each edge  $\{v, w\} \in E(H_i)$  with the edge  $\{u, w\}$ . This preserves the type of isomorphism of  $H_i$ . Given a vertex  $v$ , let  $N(v) = N_{H_1}(v) \cup \dots \cup N_{H_M}(v)$ . Suppose there exists an intersecting edge  $\{u, w\} \in E(H_i) \cap E(H_j)$ . Since  $|N(u)| \leq M \cdot \Delta(H_i) = o(n^{1/2})$ ,  $|N(u) \cup N(N(u))| = o(n)$  so there exists a vertex  $v \notin N(u) \cup N(N(u))$ . Since  $N(u) \cap N(v) = \emptyset$ , performing a  $(u, v, i)$ -swap reduces the number of intersecting edges. The result now follows from iterating this process.  $\square$

### 3.1 Triangle $F$

Consider  $F$  to be a triangle. Simply counting the number of triangles in each  $G_i$  shows the following:

**Theorem 3.3.** *For all  $n, m$  and  $\epsilon > 0$ ,*

$$\text{ex}_2(n, m, K_3) < \left( m \cdot \frac{1 + \epsilon}{4} + \frac{1}{2\epsilon} - \frac{1}{2} \right) n^2.$$

**Claim 4.** *There are less than  $\frac{2}{\epsilon}$  number of  $G_i$ 's with at least  $(1 + \epsilon) \frac{n^2}{4}$  edges.*

*Proof.* Suppose  $e(G_i) \geq (1 + \epsilon) \frac{n^2}{4}$  for  $1 \leq i \leq k$ . Let  $K_3(G)$  denote the number of triangles in graph  $G$ . By the Moon-Moser inequality,

$$K_3(G_i) \geq \epsilon(1 + \epsilon) \frac{n^3}{12}.$$

Since there are no overlapping triangles from different  $G_i$ 's,

$$\binom{n}{3} \geq \sum_{i=1}^k K_3(G_i) \geq \frac{\epsilon(1 + \epsilon)}{12} k n^3.$$

Rearranging yields  $k < \frac{2}{\epsilon}$ .  $\square$

By the claim,

$$\sum_{i=1}^m e(G_i) < \frac{2}{\epsilon} \binom{n}{2} + \left(m - \frac{2}{\epsilon}\right) (1 + \epsilon) \frac{n^2}{4} \leq m(1 + \epsilon) \frac{n^2}{4} + (1 - \epsilon) \frac{n^2}{2\epsilon},$$

which proves Theorem 3.3.

When  $m = 2$ ,

$$e(G_1) + e(G_2) \leq \binom{n}{2} + e(G_{1,2}) \leq \binom{n}{2} + \left\lfloor \frac{n^2}{4} \right\rfloor,$$

which meets the benchmark bound and so  $\text{ex}_2(n, 2, K_3) = \binom{n}{2} + \left\lfloor \frac{n^2}{4} \right\rfloor$ .

This result is also true for  $m = 3$ :

**Proposition 3.4.** *For all  $n$ ,*

$$\text{ex}_2(n, 3, K_3) = \binom{n}{2} + \left\lfloor \frac{n^2}{2} \right\rfloor.$$

*Proof.* Define  $H_k \subseteq G$  be the graph with edges contained in at least  $k$  number of  $G_i$ 's and note that  $e(G_1) + e(G_2) + e(G_3) = e(H_1) + e(H_2) + e(H_3)$ . Thus it suffices to show that  $e(H_2) + e(H_3) \leq \frac{n^2}{2}$ . Notice  $H_2$  must not contain any triangles with two edges in  $H_3$ , so

$$e(H_2) + e(H_3) \leq \binom{n}{2} + e(H_3) - |\{\{u, v\} : u \neq v, N_{H_3}(u) \cap N_{H_3}(v) \neq \emptyset\}|.$$

Let  $H'_3$  be the graph with the same vertex set as  $H_3$  and edge set  $\{\{u, v\} : u \neq v, N_{H_3}(u) \cap N_{H_3}(v) \neq \emptyset\}$ . It suffices to show that  $\frac{n}{2} \geq e(H_3) - e(H'_3)$ .

Let  $d_1 \geq d_2 \geq \dots \geq d_n$  and  $f_1 \geq f_2 \geq \dots \geq f_n$  each be the degree sequence of  $H_3$  and  $H'_3$ , respectively. We show that  $f_i \geq d_i - 1$  for all  $i$ . Let  $v_i$  denote the vertex in  $H$  with degree  $d_i$  and  $u_i$  be the vertex in  $H$  with degree  $f_i$ . Let  $S_i = |N_{H_3}(v_1) \cup \dots \cup N_{H_3}(v_i)|$ . Since

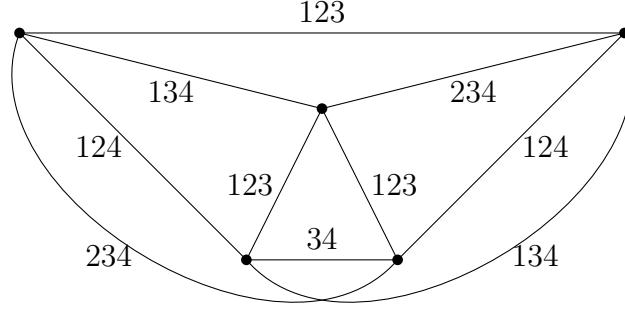
$$\sum_{u \in S_i} d_{H_3}(u) \geq d_1 + \dots + d_i,$$

we have that  $|S_i| \geq i$ . But then  $S_i \setminus \{u_1, \dots, u_{i-1}\}$  is non-empty, and every  $u \in S_i$  has degree  $d_{H'_3}(u) \geq d_i - 1$ . Hence,  $f_i \geq d_i - 1$  for all  $i$ , which yields

$$e(H'_3) = \frac{1}{2} \sum_{i=1}^n f_i \geq \frac{1}{2} \sum_{i=1}^n (d_i - 1) = e(H_3) - \frac{n}{2}.$$

□

However, the bound in Proposition 3.1 is not tight for  $m \geq 4$ , as shown in the following graph:



The number on each edge denotes the set of  $G_i$ 's that contain the edge.

The above graph contains 29 edges, which exceeds the bound  $\binom{5}{2} + 3\lfloor \frac{5^2}{4} \rfloor = 28$  by 1. By blowing up the above graph, we can construct a graph with  $n \in 10\mathbb{Z}$  vertices that contains

$$5 \binom{n/5}{2} + 29 \cdot \frac{(n/5)^2}{4}$$

edges, which exceeds the bound  $\binom{n}{2} + 3\lfloor \frac{n^2}{4} \rfloor$  by  $n^2/100$ .

## 3.2 Bipartite $F$

In this section, we discuss the case where  $F$  is bipartite. In particular, we focus on the cases where  $F \subseteq K_{2,2}$  is  $P_2$ , a path of length 2, or  $M_2$ , a matching with two edges.

**Theorem 3.5.**

$$\text{ex}_2(n, m, P_2) \leq \left( \frac{1}{2} + o(1) \right) n^2 \sqrt{m}$$

as  $n \rightarrow \infty$  or  $m \rightarrow \infty$ .

*Proof.* Since there are no overlapping  $P_2$ 's in different  $G_i$ 's,

$$\sum_{i=1}^m \#\{P_2 \subseteq G_i\} \leq \#\{P_2 \subseteq G\}$$

For each  $G_i$ , each vertex  $v$  in  $G_i$  and two of its neighbors form one unique  $P_2$ , so

$$\#\{P_2 \subseteq G_i\} = \sum_{v \in V(G_i)} \binom{d_{G_i}(v)}{2}.$$

And by Jensen's inequality,

$$\sum_{v \in V(G_i)} \binom{d_{G_i}(v)}{2} \geq n \binom{d_{G_i}(v)/n}{2} = n \binom{2e(G_i)/n}{2} \geq \frac{2(e(G_i))^2}{n} - e(G_i).$$

On the other hand, since each three vertices in  $G$  can form at most three  $P_2$ 's,

$$\#\{P_2 \subseteq G\} \leq 3 \binom{n}{3} \leq \frac{n^3}{2}.$$

Combining the above inequalities yields

$$\frac{2m}{n} \left( \frac{1}{m} \sum_{i=1}^m e(G_i) \right)^2 - \sum_{i=1}^m e(G_i) \stackrel{\text{Jensen's}}{\leq} \sum_{i=1}^m \frac{2(e(G_i))^2}{n} - e(G_i) \leq \frac{n^3}{2},$$

and solving the quadratic equation gives

$$\sum_{i=1}^m e(G_i) \leq mn \cdot \frac{1 + \sqrt{4n^2/m + 1}}{4} = \left( \frac{1}{2} + o(1) \right) n^2 \sqrt{m},$$

as  $n \rightarrow \infty$  or  $m \rightarrow \infty$ . □

When  $m = n$ , the following projective plane construction shows the above bound is tight asymptotically:

**Theorem 3.6.**

$$\text{ex}_2(n, n, P_2) = \left( \frac{1}{2} + o(1) \right) n^{5/2},$$

as  $n \rightarrow \infty$ .

*Proof.* It suffices to show the tightness of the bound in Theorem 3.5. Consider a finite projective plane of order  $q$ . The projective plane has  $n = q^2 + q + 1$  points and  $n$  lines. Let  $S_1, \dots, S_n \subseteq [n]$  be the  $n$  lines of the projective plane. Note that each line  $S_i$  contains  $q + 1$  points, and the intersection of any two distinct lines  $S_i, S_j$  contains  $|S_i \cap S_j| = 1$  point. Define  $G_1, \dots, G_n$  to be graphs on  $[n]$ , each with edge set  $E(G_i) = \{\{j, k\} \subseteq [n] : j \neq k, j + k \in S_i \pmod n\}$ . Note that the intersection of distinct  $G_i, G_j$  is  $P_2$  free: since  $|S_i \cap S_j| = 1$ , if  $\{a, b\}, \{a, c\} \in E(G_i) \cap E(G_j)$ , then  $a + b = a + c$  so  $b = c$ . We now count the number of edges in  $G_1, \dots, G_n$ . Since  $|S_i| = q + 1$ , for each point  $j \in [n]$ , there are  $q + 1$  choices for  $k \in [n]$  such that  $j + k \in S_i$ . But then we have to avoid counting the same edge twice and loops, so the number of edges in  $G_i$  is

$$e(G_i) = \frac{n(q + 1) - \#\text{loops counted for } G_i}{2}.$$

If  $j \in [n]$  is even, then  $k = j/2$  is the unique number in  $[n]$  such that  $k + k = j \pmod n$ . If  $j \in [n]$  is odd, then  $k = (n + j)/2$  is the unique number in  $[n]$  such that  $k + k = j \pmod n$ , as  $n$  is even. Hence, for each  $j \in S_i$ , there exists a unique  $k \in [n]$  such that  $k + k = j \pmod n$ , and thus

$$\#\text{loops counted for } G_i = |S_i| = q + 1.$$

Since  $q + 1 = (1 + o(1))n^{1/2}$ , the number of edges in  $G_1, \dots, G_n$  is

$$\sum_{i=1}^n e(G_i) = n \cdot \frac{n(q + 1) - (q + 1)}{2} = \left( \frac{1}{2} + o(1) \right) n^{5/2},$$

as  $n \rightarrow \infty$ . □

**Theorem 3.7.** For all  $n, m$ ,

$$\text{ex}_2(n, m, M_2) \leq n^{5/2}.$$

*Proof.* Notice that  $\#\{M_2 \subseteq G\} = \binom{e(G)}{2}$ . On the other hand, each four vertices in  $G$  can form at most three  $M_2$ 's, so  $\#\{M_2 \subseteq G\} \leq 3\binom{n}{4} \leq \frac{n^4}{8}$ . By the same argument as in Theorem 3.4, we have

$$\sum_{i=1}^n \binom{e(G_i)}{2} = \sum_{i=1}^n \#\{M_2 \subseteq G_i\} \leq \#\{M_2 \subseteq G\} \leq \frac{n^4}{8}.$$

By Jensen's inequality,

$$\sum_{i=1}^n \binom{e(G_i)}{2} \geq n \binom{\sum_{i=1}^n e(G_i)/n}{2} = \frac{1}{2n} \left[ \left( \sum_{i=1}^n e(G_i) \right)^2 - n \sum_{i=1}^n e(G_i) \right].$$

Combining the above inequalities yields

$$\left( \sum_{i=1}^n e(G_i) \right)^2 - n \sum_{i=1}^n e(G_i) \leq \frac{n^5}{4},$$

and solving the quadratic inequality gives

$$\sum_{i=1}^n e(G_i) \leq n^{5/2}.$$

□

We may obtain an exact result if we forbid both  $P_2$  and  $M_2$  at the same time:

**Theorem 3.8.** *For all  $n, m$ ,*

$$\text{ex}_2(n, m, \{P_2, M_2\}) = n^2 - n.$$

*Proof.* Denote the set of  $G_i$ 's as  $\{G_i\} = \{G_1, \dots, G_n\}$ , and the set of distinct pairs of  $G_i$ 's as  $\{G_i\}^2 = \{\{G_j, G_k\} : j \neq k\}$ . Consider the bipartite graph  $H$  with vertex set  $V(H) = \{G_i\} \sqcup E(K_n)$  and edge set  $E(H) = \{\{G_j, e\} \in \{G_i\} \times E(K_n) : e \in G_j\}$ . Define  $\phi : \{G_i\}^2 \rightarrow 2^{E(K_n)}$  by sending each  $\{G_j, G_k\}$  to their common edge set  $E(G_j) \cap E(G_k)$ . Notice that each distinct  $G_j, G_k$  have at most one edge in common, so  $|\phi(G_j, G_k)| \leq 1$ . On the other hand, each edge  $e \in E(G)$  can be obtained via  $\phi$  by  $\binom{d_H(e)}{2}$  possible distinct pairs  $(G_j, G_k)$ , and thus  $|\phi^{-1}(e)| = \binom{d_H(e)}{2}$ . But then

$$\binom{n}{2} \geq \sum_{(G_j, G_k) \in \{G_i\}^2} |\phi(G_j, G_k)| = \sum_{e \in E(K_n)} |\phi^{-1}(e)| = \sum_{e \in E(K_n)} \binom{d_H(e)}{2}.$$

By Jensen's inequality,

$$\sum_{e \in E(K_n)} \binom{d_H(e)}{2} \geq \binom{n}{2} \binom{\sum_{e \in E(K_n)} d_H(e)/n}{2} = \binom{n}{2} \binom{\sum_{i=1}^n e(G_i)/n}{2}.$$

Combining the above inequalities yields

$$2 \binom{n}{2}^2 \geq \left( \sum_{i=1}^n e(G_i) \right)^2 - \binom{n}{2} \sum_{i=1}^n e(G_i),$$

and the result now follows from solving the quadratic inequality.

To see that this bound is tight, consider the construction such that for each distinct  $i, j \in [n]$ ,  $E(G_i) \cap E(G_j)$  contains exactly one unique edge  $e \in K_n$ . The number of edges in this construction is  $2\binom{n}{2} = n^2 - n$ .  $\square$