

# Factor Oracle for Machine Improvisation

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August 2016



## Preliminaries

# Preliminaries

## Word

A **word**  $s$  is a finite sequence  $s = s_1 s_2 \dots s_m$  of length  $|s| = m$  on a finite alphabet  $\Sigma$ .

$s =$ 

a	b	b	c	a	b	c	d	a	b	c
---	---	---	---	---	---	---	---	---	---	---

## Factor

A word  $x \in \Sigma^*$  is a **factor** of  $s$  if and only if  $s$  can be written  $s = uxv$  with  $u, v \in \Sigma^*$ . Given integers  $i, j$  where  $1 \leq i \leq j \leq m$ , we denote a *factor* of  $s$  as  $s[i \dots j] = s_i s_{i+1} \dots s_j$ .

$s =$ 

a	b	b	c	a	b	c	d	a	b	c
---	---	---	---	---	---	---	---	---	---	---

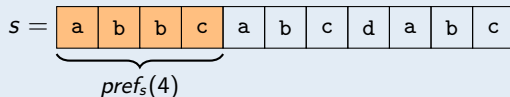
  

$\underbrace{\hspace{10em}}$   
 $s[3, 5]$

# Preliminaries

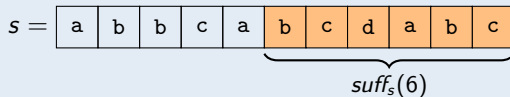
## Prefix

A factor  $x$  of  $s$  is a **prefix** of  $s$  if  $s = xu$  with  $u \in \Sigma^*$ . The  $i$ th *prefix* of  $s$ , denoted  $\text{pref}_s(i)$ , is the prefix  $s[1 \dots i]$ .



## Suffix

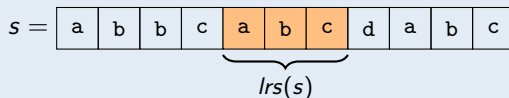
A factor  $x$  of  $s$  is a **suffix** of  $s$  if  $s = ux$  with  $u \in \Sigma^*$ . The  $i$ th *suffix* of  $s$ , denoted  $\text{suff}_s(i)$ , is the suffix  $s[i \dots m]$ .



# Preliminaries

## Longest Repeated Suffix (LRS)

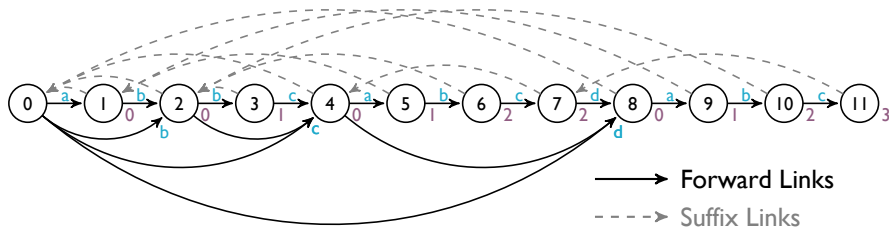
A factor  $x$  of  $s$  is the **longest repeated suffix** of  $s$  if  $x$  is a suffix of  $s$  and  $|x|$  is maximal.



Factor Oracle

# Factor Oracle

## Overview

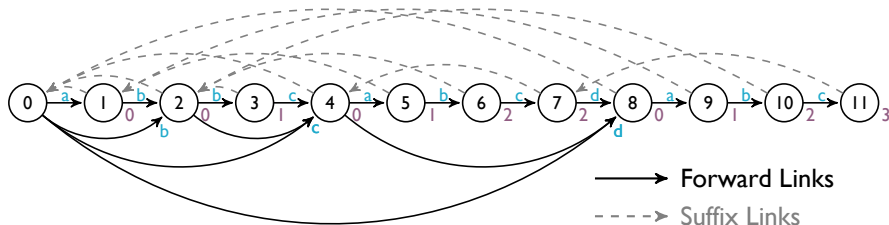


## Factor Oracle

The **factor oracle** of a word  $s$  of length  $m$  is a *deterministic finite automaton*  $(Q, q_0, F, \delta)$  where  $Q = \{0, 1, \dots, m\}$  is the set of states,  $q_0 = 0$  is the starting state,  $F = Q$  is the set of terminal states and  $\delta$  is the transition function.

# Factor Oracle

## Overview



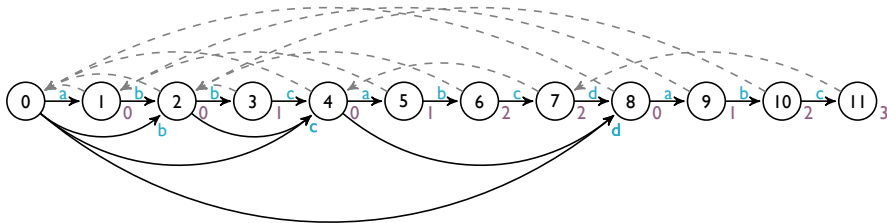
## Suffix Link

The **suffix link** of a state  $i$  of the factor oracle of a word  $s$ , is equal to the state in which the *longest repeated suffix* ( $lrs$ ) of  $s[1 \dots i]$  is recognized.



# Factor Oracle

## Overview

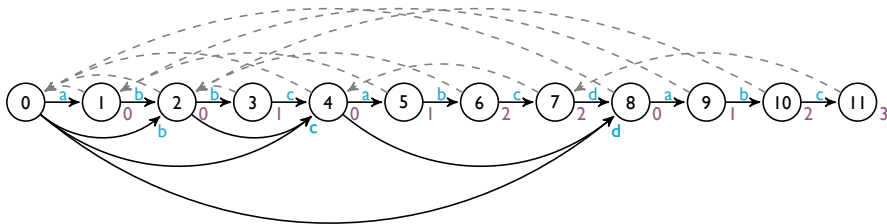


## Suffix Links

- $s = \text{abbcabcbcdabc}$

# Factor Oracle

## Overview

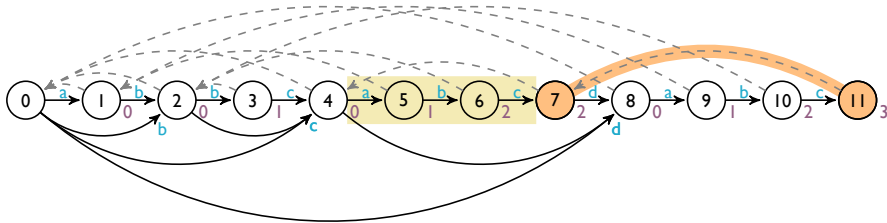


## Suffix Links

- $s = \text{abbc}\text{abcdabc}$
- $\text{lrs}(s) = \text{abc}$

# Factor Oracle

## Overview



## Suffix Links

- $s = \text{abbcababcdabc}$
- $\text{lrs}(s) = \text{abc}$
- $S(11) = 7$

# Factor Oracle

## Algorithm

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### Algorithm 1 Construction of a Factor Oracle

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```
1: function FactorOracle( $p = p_1 p_2 \dots p_m$ )
2:   Create a new oracle  $P$  with an initial state 0
3:    $S_P(0) \leftarrow -1$ 
4:   for  $i \leftarrow 1, m$  do
5:      $Oracle(p = p_1 p_2 \dots p_i) \leftarrow \text{AddLetter}(Oracle(p = p_1 p_2 \dots p_{i-1}), p_i)$ 
6:   end for
7:   return  $Oracle(p = p_1 p_2 \dots p_m)$ 
8: end function
```

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# Factor Oracle

## Algorithm

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### Algorithm 2 Incremental update of Factor Oracle

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```
1: function AddLetter( $Oracle(p = p_1, p_2 \dots p_m), \sigma$ )
2:   Create state  $m + 1$ 
3:   Create a new transition from  $m$  to  $m + 1$  labeled by  $\sigma$   $\triangleright \delta(m, \sigma) = m + 1$ 
4:    $k \leftarrow S_p(m)$ 
5:    $i \leftarrow m$ 
6:   while  $k > -1$  and there is no transition from  $k$  by  $\sigma$  do
7:     Create a new transition from  $k$  to  $m + 1$  by  $\sigma$   $\triangleright \delta(k, \sigma) = m + 1$ 
8:      $i \leftarrow k$ 
9:      $k \leftarrow S_p(k)$ 
10:  end while
11:  if  $k = -1$  then
12:     $S_{p\sigma} \leftarrow 0$ 
13:     $lrs_{p\sigma} \leftarrow 0$ 
14:  else
15:     $S_{p\sigma} \leftarrow$  state that leads the transition from  $k$  by  $\sigma$ 
16:     $lrs_{p\sigma} \leftarrow \text{LengthCommonSuffix}(i, S(m + 1) - 1) + 1$ 
17:  end if
```

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# Factor Oracle

## Algorithm

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### Algorithm 3 Incremental update of Factor Oracle

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```
18:    $k \leftarrow \text{FindBetter}(m + 1, p[m + 1 - \text{lrs}(m + 1)])$ 
19:   if  $k \neq 0$  then
20:      $\text{lrs}_{p\sigma} \leftarrow \text{lrs}(m + 1) + 1$ 
21:      $S_{p\sigma} \leftarrow k$ 
22:   end if
23:    $T(S_{p\sigma}) \leftarrow T(S(m + 1) \cup \{i\})$ 
24:   return  $\text{Oracle}(p = p_1 p_2 \dots p_m \sigma)$ 
25: end function
```

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# Factor Oracle

## Algorithm

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### Algorithm 4 Find Better Algorithm

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```
1: function FindBetter( $i, a$ )
2:   for all the elements  $j$  of  $T(i)$  in increasing order do
3:     if  $lrs(j) = lrs(i)$  and  $p[j - lrs(i)] = a$  then
4:       return  $j$ 
5:     end if
6:   end for
7:   return 0
8: end function
```

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# Factor Oracle

## Algorithm

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### Algorithm 5 Length Common Suffix Algorithm

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```
1: function LengthCommonSuffix( $i, j$ )  
2:   if  $S(i) = j$  then  
3:     return  $lrs(i)$   
4:   else  
5:     while  $S(i) \neq S(j)$  do  
6:        $j \leftarrow S(j)$   
7:     end while  
8:   end if  
9:   return  $\min(lrs(i), lrs(j))$   
10: end function
```

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Thank you for your attention! 😊

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