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week 2 - Assignment - Game Theory

Q1) a)  $\rightarrow$  If each  $s_i^*$  is weakly dominant, then for any deviation  $s_i$ , for all  $s_{-i}$ .

$$u_i(s_i^*, s_{-i}) \geq u_i(s_i, s_{-i})$$

$\rightarrow$  So no player can strictly gain by deviating  $\rightarrow s^*$  is a Nash Equilibrium.

b) If each  $s_i^*$  is strictly dominant, then  
 $u_i(s_i^*, s_{-i}) > u_i(s_i, s_{-i}) \quad \forall s_i \neq s_i^*$ .

So each player will strictly prefer  $s_i^*$ , making  $s^*$  the only NE.

c) A dominant strategy maximises  $u_i(s_i, s_{-i}) \quad \forall s_{-i}$ .  
Maximin strategy maximizes  $\min_{s_{-i}} u_i(s_i, s_{-i})$ .

So dominant strategy also gives best worst case  $\rightarrow$  it is a maximin strategy.

Q2) a) Payoff matrix

William/Henry	A (10K)	B (give 40K)
A	(10K, 10K)	(10K, 90K)
B	(40K, 10K)	(0, 0)

b) Both prefer 40K over 10K and 10K over 0.

each wants the other to choose A.

So (B, B) is the only NE - no incentive to deviate if the other picks B.



3) a) let  $x$  be number of others who clean.

If  $i$  chooses  $C$ :  $u_i = 10 - 1 = 9$

If  $x=0$  and  $i$  chooses  $NC$ :  $u_i = 0$ .

Else (at least one cleans) :  $u_i = 10$

So: If someone else cleans :  $C \rightarrow 9, NC \rightarrow 10$

. If no one else cleans :  $C \rightarrow 9, NC \rightarrow 10$

b) PSNE: only when exactly one student chooses  $C$ .

then : • Cleaner gets 9

• others get 10

→ no one wants to change :  $C$  person would get

0 if they switch, others drop to 9 if they switch to

$C$ . So 10 PSNEs, each with 1 cleaner and 9 non-cleaners.

Q.4)

a)

	$C1$	$C2$
$R1$	$(2, 1)$	$(0, 0)$
$R2$	$(0, 0)$	$(1, 2)$
NEs : $(R1, C1), (R2, C2)$		

b) 1.  $C2$  weakly dominated by  $C1$  for columns → remove  $C2$ .

2. Then  $R2$  weakly dominated by  $R1$  → remove  $R2$ .  
→ left with  $(R1, C1)$

c) 1.  $C1$  weakly dominated by  $C2$  → remove  $C1$

2.  $R1$  weakly dominated by  $R2$  → remove  $R1$   
→ left with  $(R2, C2)$



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Q-5) Suppose Player  $i$  bids  $b_i \neq v_i$ .

- If  $b_i < v_i \rightarrow$  may lose  $\rightarrow$  even when they value more than price  $\rightarrow$  regret.
- If  $b_i > v_i$ , may win and overpay (price  $>$  value)  $\rightarrow$  negative utility.

$\rightarrow$  Truthful bid

$b_i = v_i$  :

- wins only if  $v_i$  highest
  - pays second-highest bid ( $< v_i$ )  
 $\rightarrow$  utility  $v_i - \text{2nd bid} \geq 0$ , else gets 0
- $\rightarrow$  So truthful bidding is weakly dominant. if lower.