Problem A. (Adopted from MIT EECS 6.S955)

Given a graph G = (V, E), $V = \{1, ..., n\}$, $E \subseteq V \times V$. A common problem is graph layout, where we choose positions of the vertices in V on the plane \mathbb{R}^2 respecting the connectivity of G. Assume that $(i, i) \notin E$ $\forall i \in V$ and (i, j) = (j, i) $\forall (i, j) \in E$.

Take $\vec{v}_1, \ldots, \vec{v}_n \in \mathbb{R}^2$ to be the positions of the vertices in V; these are the unknowns in graph layout. Suppose a finite positions of vertices are specified in $V_0 \subseteq V$ (labeled as $\vec{v}_k^0 \ \forall k \in V_0$). Derive two linear systems satisfied by (x_i, y_i) of \vec{v}_i 's to minimize the Dirichlet energy:

min
$$E(\vec{v}_1, \dots, \vec{v}_n) = \sum_{(i,j) \in E} ||\vec{v}_i - \vec{v}_j||_2^2$$

s.t $\vec{v}_k = \vec{v}_k^0 \quad \forall k \in V_0.$

Implement and compare both gradient descent and conjugate gradients for solving this system. Use a preconditioner of choice for conjugate gradients. Show improve of convergence.

Solution .

Since $\vec{v_i} \in \mathbb{R}^2$, let $\vec{v_i} = (x_i, y_i)$. This implies $||\vec{v_i} - \vec{v_j}||_2^2 = (x_i - x_j)^2 + (y_i - y_j)^2$. Then

$$E(\vec{v}_1, \dots, \vec{v}_n) = \sum_{(i,j)\in E} ||\vec{v}_i - \vec{v}_j||_2^2$$

$$= \sum_{(i,j)\in E} (x_i - x_j)^2 + (y_i - y_j)^2$$

$$= E_x + E_y \quad \text{where} \begin{cases} E_x \equiv \sum_{(i,j)\in E} (x_i - x_j)^2 \\ E_y \equiv \sum_{(i,j)\in E} (y_i - y_j)^2 \end{cases}$$

Rewriting our objective

$$\min E_x + E_y$$
 s.t $(x_k, y_k) = \vec{v}_k = \vec{v}_k^0 \quad \forall k \in V_0.$

We solve for E_x . Define $U = V \setminus V_0$. Then

$$\min \sum_{(i,j)\in E} (x_i - x_j)^2 \equiv \min \sum_{(i,u)\in E} (x_u - x_i)^2 \quad \forall u \in U$$

$$\frac{\partial E_x}{\partial x_u} = 2 \sum_{i \in N(u)} (x_u - x_i) = 0$$

$$\iff \sum_{i \in N(u)} (x_u - x_i) = 0$$

$$\iff \delta(u)x_u - \sum_{i \in N(u)\cap U} x_i = \sum_{i \in N(u)\cap V_0} x_j$$

Denote the cardinality of U as m = |U|. Let

1.
$$L \in \mathbb{R}^{m \times m}$$
; $L_{ii} = \delta(i)$; $L_{ij} = \begin{cases} -1 & \forall i \in N(j) \text{ for } j \in U \\ 0 & \text{otherwise} \end{cases}$ (the Laplacian matrix)

2.
$$\vec{x} \in \mathbb{R}^{m \times 1} = (x_{u_1}, \dots, x_{u_m})^{\top} \quad \forall u \in U$$

3.
$$\vec{b_x} \in \mathbb{R}^{m \times 1} = (x_{j_1}, \dots, x_{j_m})^{\top} \quad \forall j \in N(i) \cap V_0 \text{ with } i \in U$$

This defines a linear system $E_x \equiv L\vec{x} - \vec{b_x}$. The work for E_y is analogous.

With our assumption $(i, j) = (j, i) \quad \forall (i, j) \in E, L$ is intuitively symmetric. Also, the Laplacian graph matrix is always (strictly or not) diagonally dominant with non-negative diagonal entries. These combined guarantees positive semi-definiteness.

Complexity

The method of both steepest descent and conjugate gradients both depends on the matrix-vectors products operations. More often, the graph Laplacian matrix is sparse. Let t be the number of non-zero entries of L.[1] Then the complexity of:

• Gradient descent: $O(t\kappa)$

• Conjugate gradients: $O(t\sqrt{\kappa})$

with $\kappa \equiv \text{cond}(L)$.

Preconditioning

Let us decompose $L = D + A + A^{\top}$ where D is the diagonal matrix consists of the diagonal values of L, A is the strictly lower triangular matrix consists of corresponding part of L. Then, the SOR iterative scheme is

$$(\omega^{-1}D + A)\vec{x_{k+1}} = ((\omega^{-1} - 1)D - A^{\mathsf{T}})\vec{x_k} + \vec{b_x}$$

for some constant $\omega \in (0, 2)$.

Show of convergence.

By simply letting $M \equiv (\omega D + A)$ and $N \equiv ((\omega^{-1} - 1)D - A^{\top})$ yields the classic Jacobi's iterative scheme

$$M\vec{x_{k+1}} = N\vec{x_k} + \vec{b_x}$$

which converges to $L^{-1}\vec{b_x}$ when $\rho(M^{-1}N) < 1$.

Improve of convergence.

show later

References

[1] J. R. Shewchuk, "An introduction to the conjugate gradient method without the agonizing pain," tech. rep., USA, 1994.