

Problem C. (*Adopted from CMU 21-387*)

Euclidean traveling salesman problem (TSP): Given a set of points $\vec{x}_1, \dots, \vec{x}_n \in \mathbb{R}^2$ representing the positions of cities on a map, we wish to visit each city exactly once while minimizing the total distance traveled. While Euclidean TSP is NP-hard, simulated annealing method provides a practical way to approximate its solution.

Solution

Let us define the permutation $\pi : I_n \rightarrow I_n$ where $I_n = \{1, 2, \dots, n\}$. A desired solution is a tour consisting of given $\vec{x}_i \in \mathbb{R}^2$, thus the feasible set of our problem should be in S_n which is the set of all permutations on I_n . This implies:

$$\pi \in S_n$$

For the weight, let us define $f : S_n \rightarrow \mathbb{R}$ that calculates the Euclidean distance for a given tour π along with the distance between the last point back to the initial point in the tour.

$$f(\pi) = \|x_{\pi_n} - x_{\pi_1}\|_2 + \sum_{i=1}^{n-1} \|x_{\pi_{i+1}} - x_{\pi_i}\|_2$$

Writing in optimization term, our objective now is:

$$\pi^* \equiv \arg \min_{\pi \in S_n} \|x_{\pi_n} - x_{\pi_1}\|_2 + \sum_{i=1}^{n-1} \|x_{\pi_{i+1}} - x_{\pi_i}\|_2$$

A simple shuffle function to generate a random tour that reach each city exactly once.

Algorithm 1 Generate a random tour

```

1: function RANDOMTOUR( $n$ ):
2:    $tour \leftarrow []$ 
3:   for  $i \leftarrow 1$  to  $n$  do
4:      $tour[i] \leftarrow i$ 
5:   for  $i \leftarrow n$  downto 1 do
6:      $j \leftarrow randInt(1, i)$ 
7:      $swap(tour[i], tour[j])$ 
8:   return  $tour$ 

```

Next up is a simple function to evaluate the total distance of a given tour.

Algorithm 2 Calculate total distance of tour

```

1: function TOURDISTANCE( $tour, cityCoordinates$ ):
2:    $totalDistance \leftarrow 0$ 
3:    $n \leftarrow \text{length}(tour)$ 
4:   for  $i \leftarrow 0$  to  $n - 2$  do
5:      $p_1 \leftarrow cityCoordinates[tour[i]]$ 
6:      $p_2 \leftarrow cityCoordinates[tour[i + 1]]$ 
7:      $totalDistance \leftarrow totalDistance + \text{distance}(p_1, p_2)$ 
8:    $p_1 \leftarrow cityCoordinates[tour[n - 1]]$ 
9:    $p_2 \leftarrow cityCoordinates[tour[0]]$ 
10:   $totalDistance \leftarrow totalDistance + \text{distance}(p_1, p_2)$ 
11:  return  $totalDistance$ 

```

In each step, we will generate a neighbor tour by performing a slight modification to the current tour, simply swapping two random cities from the current tour will be sufficient to implement the `GenerateNeighbor` function.

Algorithm 3 Generate neighbor tour

```

1: function GENERATE_NEIGHBOR_TOUR(tour):
2:   newTour  $\leftarrow$  tour
3:   n  $\leftarrow$  length(newTour)
4:   i  $\leftarrow$  randInt(0, n - 1)
5:   repeat
6:     j  $\leftarrow$  randInt(0, n - 1)
7:   until i  $\neq$  j
8:   swap(newTour[i], newTour[j])
9:   return newTour

```

Last but not least, we apply the simulated annealing method to approximate π^* .

Algorithm 4 TSP Simulated Annealing

```

1: function SIMULATED_ANNEALING(cities,  $T_0$ ,  $\alpha$ ):
2:    $T \leftarrow T_0$ 
3:   n  $\leftarrow$  length(cities)
4:   currentTour  $\leftarrow$  RandomTour(n)
5:   bestTour  $\leftarrow$  currentTour

6:   for i  $\leftarrow$  1, 2, 3, ... do
7:     neighborTour  $\leftarrow$  GenerateNeighbor(currentTour)
8:     currentEnergy  $\leftarrow$  TourDistance(currentTour, cities)
9:     neighborEnergy  $\leftarrow$  TourDistance(neighborTour, cities)

10:     $\Delta f \leftarrow$  neighborEnergy - currentEnergy
11:    if  $\Delta f < 0$  then ▷ Objective improved at neighborTour
12:      currentTour  $\leftarrow$  neighborTour
13:    else
14:      if random_uniform(0, 1)  $< e^{-\Delta f/T}$  then
15:        currentTour  $\leftarrow$  neighborTour ▷ True with probability  $e^{-\Delta f/T}$ 

16:     $d_1 \leftarrow$  TourDistance(currentTour, cities)
17:     $d_2 \leftarrow$  TourDistance(bestTour, cities)
18:    if  $d_1 < d_2$  then
19:      bestTour  $\leftarrow$  currentTour
20:     $T \leftarrow T \times \alpha$  ▷ Cool the temperature

21:   return bestTour

```

■