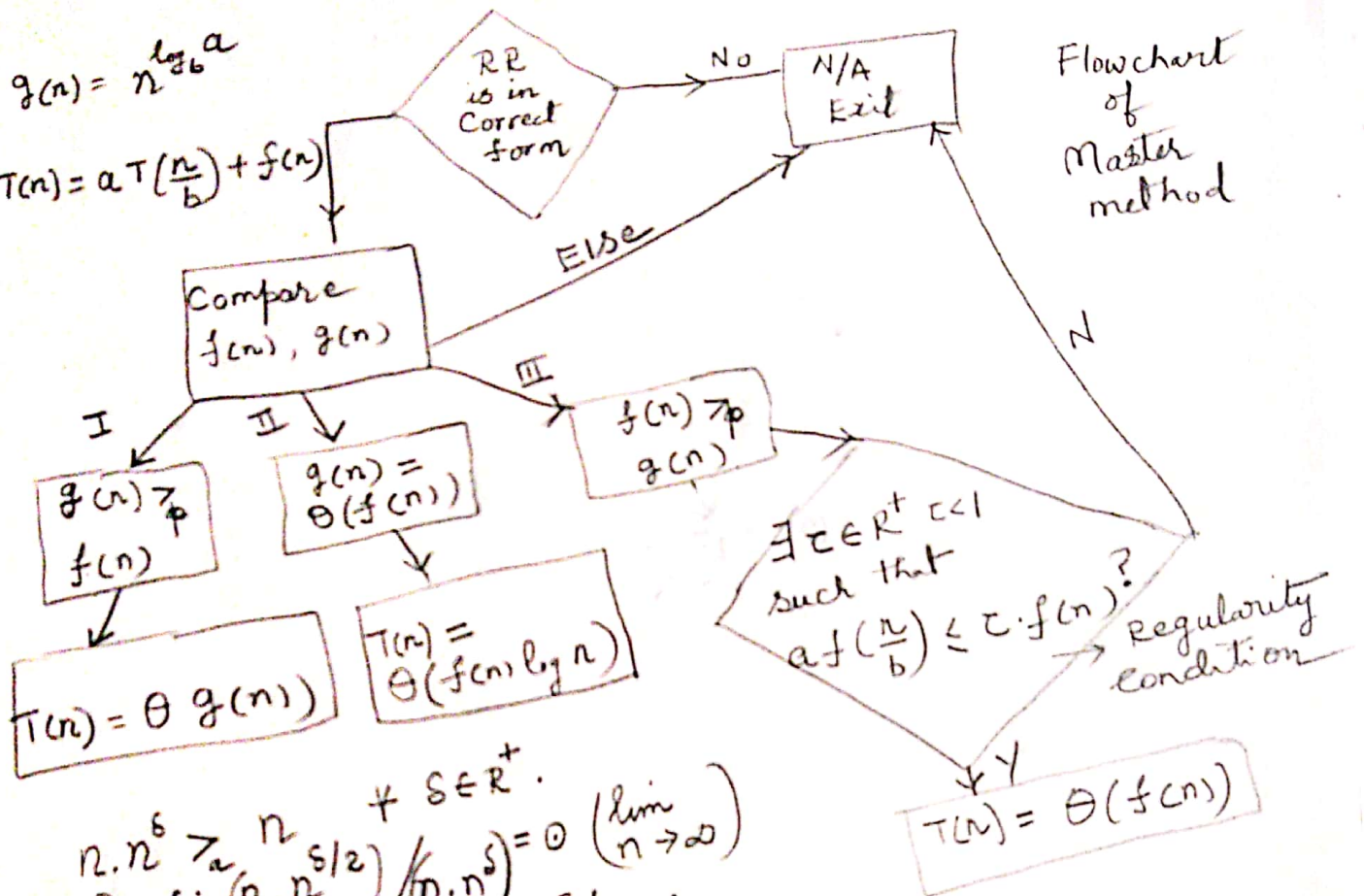


$$g(n) = n^{\log_b a}$$

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

Flowchart of Master method



Note: $n \cdot n^6 > n^7 \neq s \in \mathbb{R}^+$.
 Proof: $(n \cdot n^{s/2}) / (n \cdot n^s) = 0 \left(\lim_{n \rightarrow \infty} \right)$
 I.E. we choose $\epsilon = s/2$.
 $s > 0 \rightarrow \epsilon > 0$

Examples: (i) $T(n) = T(n-4) + n^2$. N/A wrong form
 (ii) $T(n) = 5T(n/2) + n^2$. $g(n) = n^{\log_2 5}$. Clearly $g(n) > n^2$. Case (I)
 $\therefore T(n) = \Theta(g(n))$
 (ii) $T(n) = 4T(n/2) + n^2$. $g(n) = f(n) = \Theta(n^2)$. $\therefore T(n) = \Theta(n^2 \log n)$ (Case II)
 (iii) $T(n) = 4T(n/2) + n^3$. $g(n) = n^{\log_2 4} = n^2$. $f(n) > g(n)$.
 Does $\exists 0 < c < 1$ s.t. $4\left(\frac{n}{2}\right)^3 \leq cn^3$
 $\Rightarrow \frac{4}{8}n^3 \leq cn^3 \Rightarrow \frac{1}{2} \leq c$. Any $c \in [\frac{1}{2}, 1)$
 half-open interval will do, e.g. $c = \frac{1}{2}$.
 Thus, $T(n) = \Theta(n^3)$ Case III.