

# SCoP Detection: A Fast Algorithm for Industrial Compilers

Sebastian Pop and Aditya Kumar

SARC: Samsung Austin R&D Center

Jan 19, 2016

# Polyhedral compilation in industrial compilers

- ▶ Goal: enable isl scheduler in GCC at -O3

# Polyhedral compilation in industrial compilers

- ▶ Goal: enable isl scheduler in GCC at -O3
- ▶ search loops that can benefit from polyhedral compilation
- ▶ minimal overhead: search as fast as possible
- ▶ only use existing analysis information
- ▶ use the right abstract representation

# What is a SCoP?

Regions of code that can be represented in the Polyhedral Model.

- ▶ SCoPs = Static Control Parts

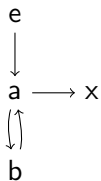
# What is a SCoP?

Regions of code that can be represented in the Polyhedral Model.

- ▶ SCoPs = Static Control Parts
- ▶ ACLs = Affine Control Loops
- ▶ PWACs = Parts With Affine Control

# Step 1: accept natural loops

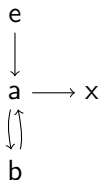
Natural loop



maybe SCoP

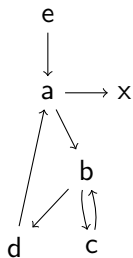
## Step 1: accept natural loops

Natural loop



maybe SCoP

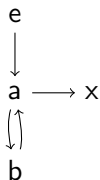
Nested loops



maybe SCoP

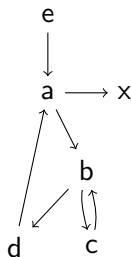
## Step 1: accept natural loops

Natural loop



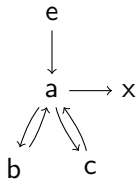
maybe SCoP

Nested loops



maybe SCoP

Irreducible

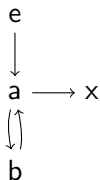


not a SCoP:  
ambiguous  
iteration order



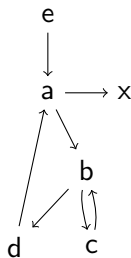
## Step 1: accept natural loops

Natural loop



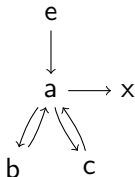
maybe SCoP

Nested loops



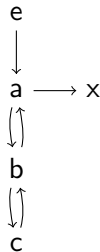
maybe SCoP

Irreducible



not a SCoP:  
ambiguous  
iteration order

Irreducible



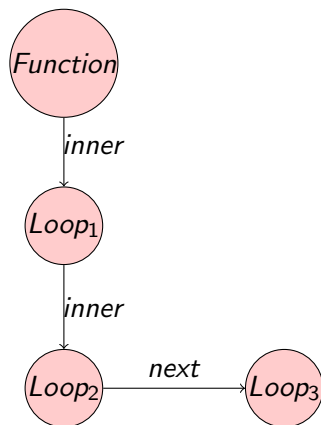
not a SCoP:  
ambiguous  
iteration order

# Natural Loop Tree

```
int foo(int N)
{
    int i, j, k;
    for(i=0; i<N; ++i){//Loop1
        stmt1;
        for (j=0; j<N; ++j)//Loop2
            stmt2;
        for (k=0; k<N; ++k)//Loop3
            stmt3;
    }
}
```

# Natural Loop Tree

```
int foo(int N)
{
    int i, j, k;
    for(i=0; i<N; ++i){ //Loop1
        stmt1;
        for (j=0; j<N; ++j) //Loop2
            stmt2;
        for (k=0; k<N; ++k) //Loop3
            stmt3;
    }
}
```



## Step 2: check for side-effects

- ▶ function calls
- ▶ inline assembly
- ▶ volatile operations

## Step 3: affine scalar evolutions

Linear

```
i0 = phi_l1(0, i1)
// i0={0,+,1}_l1
i1 = i0 + 1
// i1={1,+,1}_l1
```

maybe SCoP

## Step 3: affine scalar evolutions

### Linear

```
i0 = phi_l1(0, i1)
// i0={0,+,1}_l1
i1 = i0 + 1
// i1={1,+,1}_l1
```

maybe SCoP

### Non-linear

```
j2 = phi_l1(3, j3)
j3 = j2 + i1
// j2={3,+,{1,+,1}_l1}_l1
```

not an ACL: polynomial of degree 2

## Step 3: affine scalar evolutions

### Linear

```
i0 = phi_l1(0, i1)
// i0={0,+,1}_l1
i1 = i0 + 1
// i1={1,+,1}_l1
```

maybe SCoP

### Non-linear

```
j2 = phi_l1(3, j3)
j3 = j2 + i1
// j2={3,+,{1,+,1}_l1}_l1
```

not an ACL: polynomial of degree 2

### Non-linear

```
k4 = phi_l2(4, k5)
k5 = k4 * 2
// k4={4,*,2}_l2
```

not an ACL: exponential

## Step 3: affine scalar evolutions

### Linear

```
i0 = phi_l1(0, i1)
// i0={0,+,1}_l1
i1 = i0 + 1
// i1={1,+,1}_l1
```

maybe SCoP

### Non-linear

```
j2 = phi_l1(3, j3)
j3 = j2 + i1
// j2={3,+,{1,+,1}_l1}_l1
```

not an ACL: polynomial of degree 2

### Non-linear

```
k4 = phi_l2(4, k5)
k5 = k4 * 2
// k4={4,*,2}_l2
```

not an ACL: exponential

analyzed expressions

- ▶ branch conditions
- ▶ memory accesses



## Step 4: delinearize memory access functions

Linear access functions

$A[100*i + 400*j]$

$B[i][j]$

can represent in isl

## Step 4: delinearize memory access functions

### Linear access functions

$A[100*i + 400*j]$

$B[i][j]$

can represent in isl

### Non-linear access functions

$C[i*i]$

$D[4*N*M*i + 4*M*j + 4*k]$

$E[4*i*N + 4*j]$

cannot represent in isl

## Step 4: delinearize memory access functions

### Linear access functions

$A[100*i + 400*j]$   
 $B[i][j]$

can represent in isl

### delinearization

- ▶ recognize array multi-dimensions
- ▶ compute linear access functions

### Non-linear access functions

$C[i*i]$   
 $D[4*N*M*i + 4*M*j + 4*k]$   
 $E[4*i*N + 4*j]$

cannot represent in isl

## Step 4: delinearize memory access functions

### Linear access functions

```
A[100*i + 400*j]  
B[i][j]
```

can represent in isl

### Non-linear access functions

```
C[i*i]  
D[4*N*M*i + 4*M*j + 4*k]  
E[4*i*N + 4*j]
```

cannot represent in isl

### delinearization

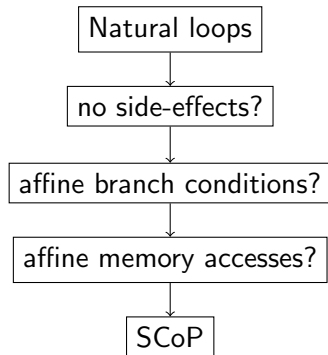
- ▶ recognize array multi-dimensions
- ▶ compute linear access functions

### delinearized access functions

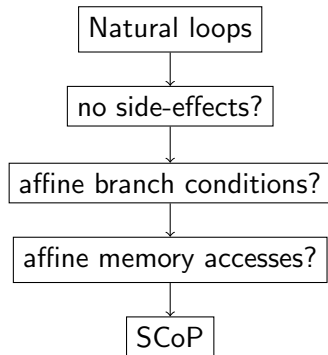
```
int D[][N][M];  
D[i][j][k]  
  
int E[][N];  
E[i][j]
```

can represent in isl

# Overall picture: SCoP detection



# Overall picture: SCoP detection



Required analyses:

- ▶ natural loops tree
- ▶ (post-)dominators tree
- ▶ alias analysis
- ▶ scalar evolution analysis

# Detecting SCoPs by induction on Natural Loops Tree

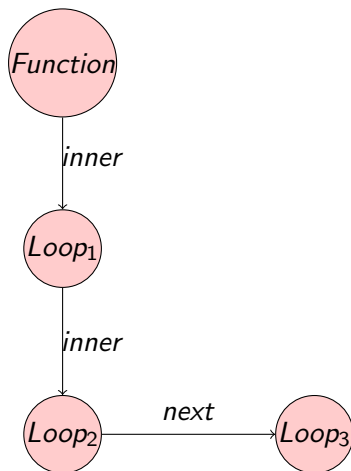
- ▶ Start with a loop in the natural loops tree rather than the root of the CFG

# Detecting SCoPs by induction on Natural Loops Tree

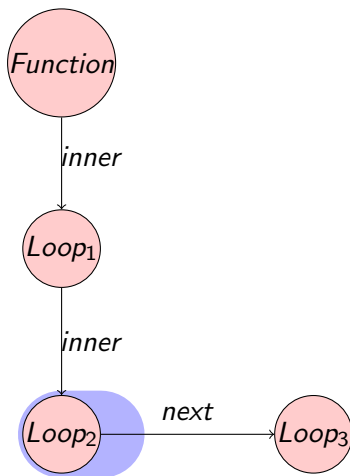
- ▶ Start with a loop in the natural loops tree rather than the root of the CFG
- ▶ Focus on structure of natural loops before the validity of each statement



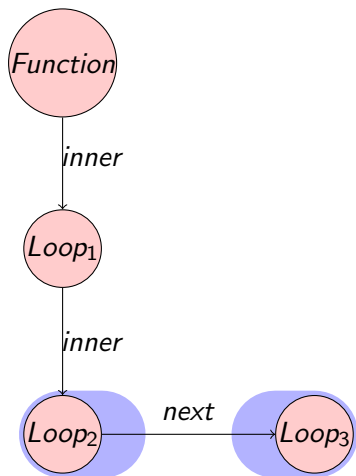
## Example: Induction on Natural Loops Tree



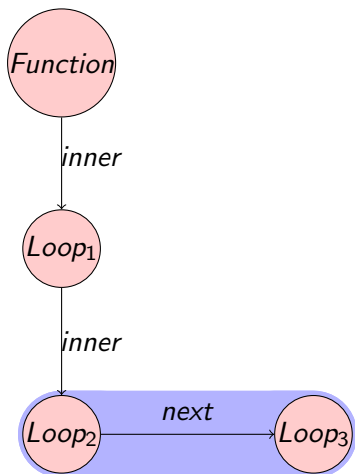
## Example: Induction on Natural Loops Tree



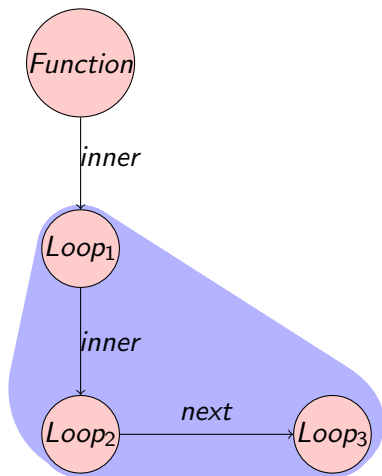
## Example: Induction on Natural Loops Tree



## Example: Induction on Natural Loops Tree



## Example: Induction on Natural Loops Tree



## Other implementations of SCoP Detection

- ▶ Previous graphite SCoP detection based on CFG and DOM (misses the structure of loops)

# Other implementations of SCoP Detection

- ▶ Previous graphite SCoP detection based on CFG and DOM (misses the structure of loops)
- ▶ Polly's SCoP detection based on structure of SESE regions (full function body analysis even without interesting loops)

# Other implementations of SCoP Detection

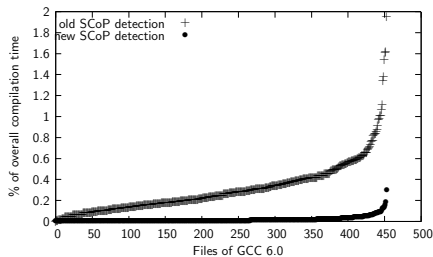
- ▶ Previous graphite SCoP detection based on CFG and DOM (misses the structure of loops)
- ▶ Polly's SCoP detection based on structure of SESE regions (full function body analysis even without interesting loops)
- ▶ Pet, Rose, other source-to-source compilers: SCoP detection based on the AST of a specific programming language



# Experimental Results

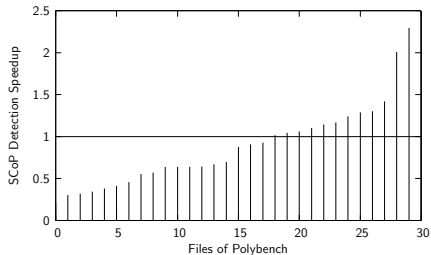
## Compilation time overhead

Benchmark	Old %	New %
Polybench	1.4	1.9
Tramp3d-v4	7.0	0.3
GCC 6.0	0.24	0.01



## SCoP Metrics on Polybench

SCoP Metric	Old	New	Polly
Loops/SCoP	2.59	6.09	5.17



# Conclusion and Future work

## Conclusion

- ▶ New faster algorithm for SCoP detection
- ▶ Enable polyhedral optimization in industrial compilers

## Future Work

- ▶ SCoP detection to drive polyhedral optimization (avoid maximal SCoPs)
- ▶ Use profile data to guide and select polyhedral transforms