

Some conservation results for higher order reverse mathematics

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ESI Reverse Mathematics: New Paradigms

Motivation for conservation results

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Either:

(1) The higher order theory has the same second order consequences as a familiar second order theory. This allows us to work in the expressive higher order theory and draw conclusions about second order arithmetic.

Or:

(2) The higher order theory does not have the same second order consequences as any familiar second order theory.

The base system RCA_0^ω

In his article *Higher order reverse mathematics* [5], Kohlenbach presents axioms for doing reverse mathematics using arithmetic in all finite types. The axioms for RCA_0^ω include:

Feferman's E-PRA^ω which includes axioms for arithmetic on numbers (type 0 objects), functions from numbers to numbers (type $0 \rightarrow 0$ objects, also called type 1), functions from type 1 objects to numbers, and so on. . .

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plus, the choice scheme $\text{QF} - \text{AC}^{1,0}$:

$$\forall x^1 \exists y^0 A(x, y) \rightarrow \exists Y^{1 \rightarrow 0} \forall x^1 A(x, Y(x))$$

where $A(x, y)$ is quantifier-free.

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Kohlenbach proves that RCA_0^ω has the same second order theorems as (a function based version of) RCA_0 .

The proof uses the type structure ECF (extensional hereditarily continuous functionals) of Troelstra.

Arithmetical comprehension

We can write ACA_0^ω for the higher order subsystem consisting of RCA_0^ω plus the axiom (\exists^2) .

The axiom (\exists^2) asserts the existence of a type 2 function \exists^2 with the following property:

If $f : \mathbb{N} \rightarrow \mathbb{N}$ is any function on the natural numbers, then $\exists^2(f) = 0$ if and only if 0 is in the range of f .

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The proof uses a re-axiomatization of ACA_0^ω that simplifies the verification that the higher order model satisfies the axioms.

A misleading question

$\text{RCA}_0^\omega + (\text{LPO})$ has the same second order theorems as ACA_0 .

Does $\text{RCA}_0^\omega + (\text{LLPO})$ have the same second order theorems as WKL_0 ?

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$\text{RCA}_0^\omega + (\text{LPO})$ has the same second order theorems as ACA_0 .

Does $\text{RCA}_0^\omega + (\text{LLPO})$ have the same second order theorems as WKL_0 ?

No. Absolutely not.

LPO and LLPO

(LLPO) is an axiom that asserts the existence of a type 2 function LLPO with the following property:

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The article [3] (joint with Carl Mummert) includes a proof of:

Theorem (RCA_0^ω) (LLPO) is equivalent to (\exists^2) (which is (LPO)).

The proof uses Kohlenbach's [5] equivalence of (\exists^2) with the existence of a sequentially discontinuous function. That result is based on Lemma 1 of Grilliot [2].

Weihrauch and higher order

In Weihrauch analysis, we have $\text{LLPO} <_W \text{LPO}$.

However, RCA_0^ω proves $(\text{LLPO}) \equiv (\text{LPO})$.

Why are these different?

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Why are these different?

The reverse math result (roughly) shows us that any realizer for LLPO can be used to compute a realizer for LPO. However, the computation is not uniform, so it doesn't show that LPO can be Weihrauch reduced to LLPO.

Weak König's Lemma

In *Higher order reverse mathematics* [5], Kohlenbach provides a conservation result for WKL_0 , based on a fan functional that can compute moduli of uniform continuity.

(MUC): There is a function Ω defined on functions from $2^{\mathbb{N}}$ to 2 , such that for any $\varphi : 2^{\mathbb{N}} \rightarrow 2$, and any $f_1, f_2 \in 2^{\mathbb{N}}$, if f_1 and f_2 agree on the first $\Omega(\varphi)$ entries then $\varphi(f_1) = \varphi(f_2)$.

(MUC) is not restricted to continuous functions, and so is inconsistent with (\exists^2) .

Prop. 3.15 [5] (paraphrased): $RCA_0^\omega + (MUC)$ is conservative over WKL_0 .

Note that Ω is a type 3 object.

More moduli of uniform continuity

Section 8 of *Banach's theorem in HORM* [3] (joint with Carl Mummert), discusses functions that compute moduli of uniform continuity.

$(M_{[0,1]})$: There is function M such that if $f : [0, 1] \rightarrow \mathbb{R}$ is continuous, then $M(f)$ is a modulus of uniform continuity for f .

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Lemma: RCA_0^ω plus (\exists^2) proves $(M_{[0,1]})$.

Prop: The second order theorems of $\text{RCA}_0^\omega + (M_{[0,1]})$ are exactly the same as those of WKL_0 .

Note: $M_{[0,1]}$ is a type 3 object.

Preliminary results

Viewing $2^{\mathbb{N}}$ and 2 as complete separable metric spaces, we can formulate restrictions of (MUC).

$(M_{2^{\mathbb{N}}})$: There is a function M such that if $f : 2^{\mathbb{N}} \rightarrow 2$ is continuous, then $M(f)$ is a modulus of uniform continuity for f .

Prop: RCA_0^ω plus (\exists^2) proves $(M_{2^{\mathbb{N}}})$. (See [3].)

Conj: The second order theorems of $\text{RCA}_0^\omega + (M_{2^{\mathbb{N}}})$ are exactly the same as those of WKL_0 .

More preliminary results

Further restrictions of $M_{2^{\mathbb{N}}}$ may be of interest.

$(M_{2^{\mathbb{N}},code})$: There is a function M such that if $f : 2^{\mathbb{N}} \rightarrow 2$ is a continuous function defined by a traditional RM code (set of quintuples), then $M(f)$ is a modulus of uniform continuity for f .

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Conj: RCA_0^ω plus (\exists^2) proves $(M_{2^{\mathbb{N}},code})$. (Section 8 of [3] plus extraction of functions from codes.)

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Conj: The second order theorems of $\text{RCA}_0^\omega + (M_{2^{\mathbb{N}},code})$ are exactly the same as those of WKL_0 .

Conj: RCA_0^ω plus $(M_{2^{\mathbb{N}},code})$ **does not prove** $(M_{2^{\mathbb{N}}})$.

Idea for the proof: $M_{2^{\mathbb{N}},code}$ is a type 2 object. Using an argument based on Grilliot's lemma, prove that if any principle asserting the existing of a type 2 object implies $(M_{2^{\mathbb{N}}})$, then it also implies (\exists^2) .

Stronger theories: Π_1^1 -CA₀

The theory Π_1^1 -CA₀ is related to the Suslin functional.

(SF): There is a function S mapping trees from $\mathbb{N}^{<\mathbb{N}}$ to 2 such that $S(T) = 0$ if and only if T is well-founded.

Theorem (Sakamoto and Yamizaki [6], based on work of Avigad and Feferman) RCA_0^ω plus (SF) is conservative over Π_1^1 -CA₀ for Π_3^1 sentences.

Question: Can Hunter's argument for ACA_0^ω be adapted? Does a Fefermanesque μ style reaxiomatization streamline the proof?

Stronger theories: ATR_0

In second order arithmetic, there are parallels between WKL_0 and ATR_0 .

$\text{WKL}_0 \equiv \Sigma_1^0$ -separation $\text{ATR}_0 \equiv \Sigma_1^1$ -separation

Questions:

- What are good candidates for functionals that have the same second order theorems as ATR_0 ?
- Are the ATR_0 level Weihrauch problems good candidates?
- Is the LPO/LLPO behavior repeated at this level?
- Is there a good candidate of type 2?

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