Calculus 3 for Computer Science Project

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1 Introduction

Much of the code used within this assignment is Ruby version 1.9.x. Since it is, for most intents and purposes, a lesser known language than say, Java, we felt it necessary to explain bits of the language which might surface in the code we wrote for this project.

1.1 Ruby

First and foremost, Ruby shares many similarites with, and is inspired by, Lisp, Smalltalk, and Perl (most likely in that order). It contains many of the niceties of functional programming that Lisp does, it is truly object oriented in the same way Smalltalk is, and it is a terse scripting language with powerful features like Perl is.

1.1.1 Syntax, Language Quirks, Etc.

Notations that might seem confusing unless one has a background in all three of these languages are such:

Basic Syntax Comments are preceded by a #. Strings are wrapped in either single- or double-quotes. Indentation is two spaces. Rather than using indentation-based block delimiters or curly braces, Ruby simply uses the end keyword.

Blocks Blocks are notations for anonymous functions in Ruby (like lambda in Lisp). They can be written a number of ways, such as:

```
# Assigning a lambda/proc to a variable, and then calling it
f = ->(x) { x + 1 }
f[2] #=> 3

# Array#map is one of several functions which takes a block as an argument
a = [1,2,3]
a.map { | i | i * i } #=> [1,4,9]

# This is the expanded form of the block, used for creating multi-line anonymous
    functions
a.each do | item |
    puts item
end
# would print 1, 2, and 3 on separate lines
```

Ranges Ranges are represented in Ruby with one of two operators, .. or 0..10 is an inclusive range ([0, 10]) in a more mathematical notation, whereas 0...5 is exclusive ([0, 5)). Ranges can be iterated across.

Method Invocation In Ruby, method invocation has optional parens. Rather than using the form instance.method(arg1, arg2), one can use the form instance.method

arg1, arg2. In the case where an invocation doesn't have arguments, the parentheses are still optional.

- **Object Oriented** Ruby, like Smalltalk, is object oriented down to the primitives of the language. This means that all things in Ruby are objects, and thus have methods that can operate on them. This library was written to make use of this, monkey-patching functionality into the existing Matrix and Vector classes in Ruby.
- Notation Ruby has a common nomenclature for expressing its classes and their methods. Object#method is the de facto standard among Rubyists, hence, that's the form we'll use here. Similarly, #=> is used to denote return values.
- Further Notes Ruby's Matrix and Vector classes lack #[] = methods, therefore, we often convert these two datatypes to arrays and back again to perform matrix or vector arithmetic or other operations.

Hopefully, that should clear up any misconceptions or confusion before addressing the actual code at hand. That said, all three parts of this report do make use of both some standard libraries in Ruby, as well as extensions upon them.

- http://www.ruby-doc.org/core/classes/Array.html
- http://www.ruby-doc.org/core/classes/Matrix.html
- http://www.ruby-doc.org/core/classes/Vector.html

1.1.2 Common Code

Additionally, we wrote an abstraction layer into some of these classes via monkeypatching in order to add some common functionality:

Listing 1: Common Code for All Three Parts

```
class Vector
  private
  def sign(x)
    return 1 if x > 0
    return -1 if x < 0
    return 0
  end
end

class Matrix
  def pretty_print
    str = ""
    self.to_a.each do |row|
    row.each do |i|
        if i.to_i >= 0
        str << ""</pre>
```

```
end
          if ("\%.3f"\% i).to_f == i.to_i
            str << "#{i.to_i}
          else
             \operatorname{str} << "%.3 f " % i
          end
       end
       str << "\n"
     \mathbf{end}
     puts str
  end
  \mathbf{def} \ \mathsf{inf\_norm}
     self.to_a.map do |a|
       a.map do | ar |
          ar.abs
       end. inject(&:+)
     end.sort[0]
  def is_lower_triangular?
     triangular (self.column_vectors)
  def is_upper_triangular?
     triangular (self.row_vectors)
  end
  private
  def triangular (vecs)
      \  \, \textbf{for} \  \, \textbf{i} \  \, \textbf{in} \  \, 0 \ldots \textbf{vecs.length} \\
       vec = vecs[i].to_a
       unless i <= 1
          return false unless vec[0...i]. all? { |n| n == 0 } and vec[i..-1]. all? { |n| n != 0
       end
    end
     return true
  end
end
```

1.2 Java

Part Two is written in Java. The code therein extends the Matrix class with a few methods and constructors. We assume that Java's syntax is understood by the reader given its prominence, especially here at Georgia Tech.

1.3 About This Document

This document was typeset in LaTeX. It uses the *color* and *listings* packages for the code formatting. It additionally uses *hyperref* for in-document links in the compiled PDF.

The source code for this document is available, along with the compiled PDF form, at http://github.com/wfarr/calc3-for-cs/blob/master/report/report.tex. All of the

actual code used in the report, along with data output from the code can be viewed at http://github.com/wfarr/calc3-for-cs/tree/master/code. The unique advantage of viewing the code this way is that the user can walk through individual commits and see the code evolve into its final form (the same can be said of the report, since it is in fact just code and content).

2 Part One

The purpose of Part One of the project is to solve the typical $A\vec{x} = \vec{b}$ equation, with A being a Hilbert matrix. A Hilbert matrix is a square matrix whose elements follow the form

$$H_{ij} = \frac{1}{i+j-1}$$

Here's an implementation in Ruby:

Listing 2: Hilbert Matrix Implementation

Often times, simplifying a single matrix A into two or more "nicer" matrices (in the case of these algorithms, LU or QR) can make solving the equation $A\vec{x} = \vec{b}$ easier. Such algorithms introduce the potential for error, namely because they are modified forms of the original matrix.

2.1 LU Decomposition

LU Decomposition uses matrix multiplication to reduce a matrix A into two matrices, L (a lower triangular matrix) and U (an upper triangular matrix).

2.1.1 Explanation of the Algorithm

The algorithm for doing so is fairly simple in and of itself:

- 1. Starting with the first column, find the first non-zero entry below the diagonal. Let this entry be considered x. Let that column's diagonal element be y.
- 2. Multiply an Identity matrix, with the location of the entry x set to the value $-\frac{x}{y}$. This matrix is L_n .
- 3. The resulting matrix is the new A for the next iteration.
- 4. Repeat these steps until the resulting A is upper triangular. At this point, A becomes U.

- 5. To find L, multiply $L_1^{-1}L_2^{-1}...L_n^{-1}$.
- 6. Substitute A with LU in the equation $A\vec{x} = \vec{b}$ and solve.

This is commonly known as the Doolittle algorithm.

2.1.2 Implementation of the Algorithm

Listing 3: LU Decomposition

```
class Matrix
  def lu_decomposition
     return nil unless self.square?
     n = self.row_size
     a = self
     l_n = []
     cvs = a.column_vectors.map \{ |v| v.to_a \}
     for k in 0...cvs.length
         \  \, \textbf{for} \  \, \textbf{j} \  \, \textbf{in} \  \, 0 \ldots cvs.length \\
           l_new = Matrix.identity(n).to_a
           if l_new[j][k] == 1 || j < k
           end
           l_{-}new\,[\,j\,\,]\,[\,k\,] \;=\; -\;\,(\,cvs\,[\,k\,]\,[\,j\,\,]\;\;/\;\;cvs\,[\,k\,]\,[\,k\,]\,)
           l_n \ll l_new
           a = Matrix[*l_new] * Matrix[*cvs.transpose]
           cvs = a.column_vectors.map \{ |v| v.to_a \}
     end
     l_{\text{final}} = l_{\text{n}} \cdot \text{map} \{ |m| \text{Matrix}[*m] \cdot \text{inverse} \} \cdot \text{inject}(\&:*)
      u_final = a
     return l_final, u_final
  end
end
```

The algorithm first begins with an essential check: the method self.square? determines if the matrix is a square matrix, and returns true if it is. LU Decomposition can only be done on square matrices, thus, the method returns nil when given a non-square matrix. Next, the algorithm defines A (written as a in the code because A would've been a Constant rather than a variable) to be the instance of self. To iterate across the columns efficiently, we use Matrix#column_vectors, which returns an array of column vectors. This array is then mapped over to convert the vectors into arrays. The end result is that cvs is an array of arrays representing the columns of self.

The actual computation lies in the nested for loops. For each iteration, an l_new matrix is created and converted to an array. If the current values of j and k are above the diagonal, then the algorithm skips to the next iteration. Next, l_new[j][k] is set to $-\frac{x}{y}$, as above in the algorithm's description. A new a is made as the product of l_new and cvs.transpose (the same matrix as A). The last step of each iteration is rebuilding cvs based off of the newest a.

Finally, L and U are assigned and returned. While u_final is straight-forward, l_final is a bit more complicated. l_n.map { |m| Matrix[*m].inverse } returns an array of inverted matrices from the original array of arrays (of arrays). The one bit of syntactic sugar in that line is the use of Matrix[*m]. In this case, * is acting as the glob operator, essentially inserting all the content of the array it's called on rather than simply inserting the array itself.

This is necessary because Matrix[...] takes a list of rows (in the form of arrays) as its argument. Finally, this new array is passed Array#inject, which applies a given block to all elements of an array and returns the result. In this case, the injection is making use of a feature in Ruby 1.9 called symbol_to_proc, which allows for passing the method the :* symbol and automatically converting it into a proc/lambda. Thus, the result of the injection is to multiply all the results of the map together, in order.

2.1.3 Results and Analysis

Below are the results from the solving $LU\vec{x} = \vec{b}$ for \vec{x} . sol is the solution vector. Errors are not listed for these computations because the resulting errors are all 0. This is because the LU algorithm used here doesn't introduce any additional error. However, there is still technically error inherent in the problem itself, which can be found by evaluating cond(LU) because $cond(A) = cond(LU) \leq cond(L)cond(U)$.

```
\begin{split} N &= 2 \\ sol &= Vector[-2, 6] \\ N &= 3 \\ sol &= Vector[3, -24, 30] \\ N &= 4 \\ sol &= Vector[-4, 60, -180, 140] \\ N &= 5 \\ sol &= Vector[5, -120, 630, -1120, 630] \\ N &= 6 \\ sol &= Vector[-6, 210, -1680, 5040, -6300, 2772] \\ N &= 7 \\ sol &= Vector[7, -336, 3780, -16800, 34650, -33264, 12012] \\ N &= 8 \\ sol &= Vector[-8, 504, -7560, 46200, -138600, 216216, -168168, 51480] \\ \end{split}
```

N = 9

sol = Vector[9, -720, 13860, -110880, 450450, -1009008, 1261260, -823680, 218790]

N = 10

sol = Vector[-10, 990, -23760, 240240, -1261260, 3783780, -6726720, 7001280, -3938220, 923780]

N = 11

sol = Vector[11, -1320, 38610, -480480, 3153150, -12108096, 28588560, -42007680, 37413090, -18475600, 3879876]

N = 12

sol = Vector[-12, 1716, -60060, 900900, -7207200, 34306272, -102918816, 199536480, -249420600, 193993800, -85357272, 16224936]

N = 13

sol = Vector[13, -2184, 90090, -1601600, 15315300, -88216128, 325909584, -798145920, 1309458150, -1422621200, 981608628, -389398464, 67603900]

N = 14

 $sol = Vector[-14,\ 2730,\ -131040,\ 2722720,\ -30630600,\ 209513304,\ -931170240,\ 2793510720,\ -5761615860,\ 8180071900,\ -7852869024,\ 4867480800,\ -1757701400,\ 280816200]$

N = 15

sol = Vector[15, -3360, 185640, -4455360, 58198140, -465585120, 2444321880, -8779605120, 22086194130, -39264345120, 49080431400, -42184833600, 23728968900, -7862853600, 1163381400]

N = 16

 $sol = Vector[-16, \, 4080, \, -257040, \, 7054320, \, -105814800, \, 977728752, \, -5975009040, \\ 25241364720, \, -75724094160, \, 163601438000, \, -255218243280, \, 284747626800, \, -221470376400, \\ 114011377200, \, -34901442000, \, 4808643120]$

N = 17

 $\begin{aligned} & \text{sol} = \text{Vector}[17, -4896, 348840, -10852800, 185175900, -1955457504, 13742520792, \\ -67310305920, \ 236637794250, \ -607662484000, \ 1148482094760, \ -1594586710080, \\ 1605660228900, \ -1140113772000, \ 540972351000, \ -153876579840, \ 19835652870] \end{aligned}$

 $\begin{array}{l} N=18\\ sol=Vector[-18,5814,-465120,16279200,-313374600,3747960216,-29983681728,\\ 168275764800,-683620294500,2050860883500,-4593928379040,7707169098720,\\ -9633961373400,8835881733000,-5770371744000,2538963567360,-674412197580,\\ 81676217700] \end{array}$

 $\begin{array}{l} N=19\\ sol=Vector[19,-6840,610470,-23876160,514829700,-6919311168,62466003600,\\ -397742716800,1845774795150,-6380456082000,16652990374020,-33030724708800,\\ 49775467095900,-56549643091200,47605566888000,-28774920430080,11802213457650,\\ -2940343837200,335780006100] \end{array}$

 $\begin{array}{l} N=20\\ sol=Vector[-20,7980,-790020,34321980,-823727520,12355912800,-124932007200,\\ 894921112800,-4698335842200,18503322637800,-55509967913400,127994058246600,\\ -227544992438400,311023037001600,-323717854838400,251780553763200,-141626561491800,\\ 54396360988200,-12759640231800,1378465288200] \end{array}$

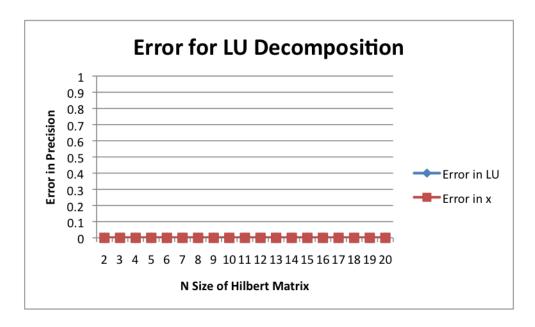


Figure 1: Error for LU Decomposition.

2.2 Householder Reflections

The Householder Reflection is, as Wikipedia might say: "a linear transformation that describes a reflection in a hyperplane (containing the origin)".

2.2.1 Description of the Algorithm

- 1. Let \vec{a} be the first column of the given matrix A.
- 2. Let $\vec{u} = (a_1 + sign(a_1) * ||\vec{a}||, a_2, a_3, ..., a_n)^t$.
- 3. Let

$$H(\vec{a}) = I - 2\frac{\vec{u}\vec{u}^t}{||\vec{u}||^2}$$

- 4. If $H(\vec{a})$ is the same size as A, $H(\vec{a}) = H_1$. Else, let H_1 be an Identity matrix, where the bottom right $n \times n$ entries are the contents of $H(\vec{a})$.
- 5. Let $A_1 = H_1 A$.
- 6. Let $\vec{a2}$ be the second column of A_1 , using only the items including and below the diagonal.
- 7. Let $\vec{u2} = (a2_1 + sign(a2_1) * ||\vec{a2}||, a2_2, a2_3, ..., a2_n)^t$.
- 8. Find $H(\vec{a2})$ as described above.
- 9. Find H_2 as described above.
- 10. Let $A_2 = H_2 A_1$.
- 11. Repeat the above as necessary until A_n is upper triangular. $A_n = R$.
- 12. $Q = H_1 H_2 H_{(n-1)} H_n$.

2.2.2 Implementation of the Algorithm

Listing 4: QR Decomposition via Householder Reflections

```
h_list << h
    current_iteration = h * current_iteration
    for i in 0...self.row_size
      cv = current_iteration.get_column_vector(i+1)
      break if cv.size < 2 | current_iteration.is_upper_triangular?
      h = (cv.find_householder_reflection - Matrix.identity(cv.size)).expand_to_dimensions(
           init_dim , init_dim ) + Matrix.identity(init_dim)
      h_list \ll h
      current_iteration = h * current_iteration
    q,r = h_list.inject(&:*), current_iteration
    return q,r
  end
  def expand_to_dimensions(x,y)
    {\tt curr\_x} \;,\;\; {\tt curr\_y} \;,\;\; {\tt a} \;=\; {\tt self.row\_size} \;,\;\; {\tt self.column\_size} \;,\;\; {\tt self.to\_a}
    a.each_index do |row|
      for i in 0...(y - curr_y)
        a[row] = a[row].insert(0,0)
      end
    end
    for i in 0...(x - curr_x)
      a = a.insert(0, Array.new(y)\{0\})
    return Matrix.rows(a)
  end
  def get_column_vector(x)
    return Vector . elements (self . column(x)[x..-1])
  end
end
class Vector
  def find_householder_reflection
    a = self.to_a
    a = a[0] if a[0]. is a? (Array)
    a[0] = a[0] + sign(a[0]) * self.r
    u = Vector[*a]
    norm_u\_sqrd = u.r**2
    uut = u.covector.transpose * u.covector
    h = Matrix.identity(uut.row_size) - (uut * (2 / norm_u_sqrd))
    return h
  end
end
```

Matrix#householder begins by asserting that the instance of self is a square matrix (which, since it ought to be a Hilbert matrix, it should be). current_iteration is assigned to self, for now, but in the algorithm itself, it is really A_n . init_dim is n for this $n \times n$ matrix, and is used for Matrix#expand_to_dimensions, which serves to form an $n \times n$ matrix the same size as self, where $H(\vec{an})$ is the bottom-right $j \times j$ entries of the matrix and all other entries are 0, for iterations of the Householder Reflection past the first (as described in step 4). Because the Householder Reflection works by essentially iterating across the column vectors of a matrix, we store the current column vector in cv. h is found by first calling the Vector#find_householder_reflection method, which takes self (an instance of Vector) and finds $(I - 2\vec{u}\vec{u}^t/||\vec{u}||^2)$; then, we subtract an Identity matrix of cv.size from it, so that,

when the matrix is expanded to the right dimensions, we can simply add an Identity matrix of size $init_dim$ to it, and thus determine H_n . h is then appended to h_list, which stores all given H_n matrices so Q can be found at the end of the process. $current_iteration$ then becomes h * $current_iteration$, as it is again A_n . This process is repeated in the for loop until A_n is upper triangular.

Once the for loop is done, q is found by injecting: * to the array (finding the product of all its elements), and then the method returns both q and r.

2.2.3 Results and Analysis

Below are the solutions for n=2,3,...,20 and their respective errors. sol is the solution to the equation $QR\vec{x}=\vec{b}$, err1 is $||QR-H||_{\infty}$ and err2 is $||H\vec{x}-\vec{b}||_{\infty}$. It is important to note that regardless of the other error measurements, there remains inherent error in the original problem; however, it is more stable than other decompositions as Q is orthogonal, and its condition number becomes 1, making cond(A) = cond(QR) = cond(R).

```
N=2
sol = Vector[-2.0, 6.0]
err1 = 1.66533453693773e-16
err2 = 1.60118641699469e-15
N = 3
sol = Vector[3.00000000000003, -24.00000000001, 30.000000000001]
err1 = 1.94289029309402e-16
err2 = 1.52807993228568e-14
N = 4
sol = Vector[-3.999999999912, 59.999999999991, -179.99999999973, 139.999999999983]
err1 = 8.32667268468867e-17
err2 = 9.99822253751169e-14
N = 5
sol = Vector[5.000000000005, -120.00000000098, 630.000000000437, -1120.0000000007,
630.00000000032]
err1 = 2.77555756156289e-17
err2 = 1.52010783481848e-11
N = 6
-6300.00000054552, 2772.00000021677]
err1 = 1.66533453693773e-16
```

```
err2 = 4.50239798778978e-11
N = 7
34650.0000600293, -33264.0000519156, 12012.0000170525]
err1 = 5.55111512312578e-17
err2 = 1.06241414529746e-08
N = 8
sol = Vector[-7.99999865279824, 503.999925690703, -7559.99901077151, 46199.9945707619,
-138599.985226035, 216215.978917003, -168167.984892845, 51479.9957131147
err1 = 1.52655665885959e-16
err2 = 2.09425467726484e-07
N = 9
450447.11486721, -1009002.0608654, 1261253.1328125, -823675.828304291, 218788.96420002
err1 = 1.17961196366423e-16
err2 = 1.21468854032131e-06
N = 10
sol = Vector[-9.9992164587602, 989.931160077453, -23758.5153179169, 240226.370779037,
-1261194.48214722,\ 3783598.73800659,\ -6726421.01654053,\ 7000989.76269531,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -6726421.01654053,\ -67264
3938067.03869629, 923746.248901367
err1 = 3.81639164714898e-16
err2 = 8.29218156290765e-05
N = 11
3138670.80151367, -12057010.7783203, 28476986.6474609, -41855156.7421875, 37286086.8505859,
-18416710.2607422, 3868220.01245117]
err1 = 2.28983498828939e-16
err2 = 0.00185948780040064
N = 12
-7113479.57226562, 33916613.15625, -101885819.21875, 197749121.4375, -247409658.1875,
192575582.5, -84787765.75, 16125571.2402344]
```

err1 = 8.32667268468867e-17err2 = 0.0585635613263169

N = 13

 $\begin{aligned} & \operatorname{sol} = \operatorname{Vector}[\text{-}17.629118386656, 2490.83803939819, \text{-}86599.147064209, 1297060.88525391,} \\ & -10418336.40625, 50111503.28125, \text{-}153161826.6875, 306042975.75, \text{-}401453162.5,} \\ & 338343685.75, \text{-}172737279.125, 46734931.15625, \text{-}4675286.6953125}] \end{aligned}$

err1 = 2.56739074444567e-16

err2 = 1.56440400784855

N = 14

 $\begin{aligned} & \text{sol} = \text{Vector}[11.0518277585506, -1764.99642848969, 68637.5826416016, -1136988.75390625, \\ & 9953654.8203125, -51021878.3125, 159625676.25, -299363132.0, 285460295.0, 19638141.5, \\ & -388452240.125, 449289740.5625, -231062811.6875, 47002828.21875] \end{aligned}$

err1 = 1.11022302462516e-16

err2 = 1.33430942738243

N = 15

 $\begin{aligned} & \text{sol} = \text{Vector}[3.29442912340164, -380.200534820557, 7961.71643066406, 4005.939453125, \\ & -1492676.53125, 17134611.625, -93579339.75, 288969014.0, -501815422.0, 360213708.0, \\ & 326874200.0, -1012593684.0, 1003466117.0, -480407496.0, 93219587.1875] \end{aligned}$

err1 = 1.52655665885959e-16err2 = 2.91871352864347

N = 16

 $\begin{aligned} & \text{sol} = \text{Vector}[3.48647512495518, -349.450300216675, 5386.19491577148, 54487.87109375, \\ & -1854269.609375, 17502754.9375, -84299294.625, 228281192.75, -323287753.5, 95025811.625, \\ & 424792122.0, -671701552.0, 339574854.0, 80873993.5, -147689482.0, 42722294.125] \end{aligned}$

err1 = 1.2490009027033e-16

err2 = 0.32930532191147

N = 17

 $\begin{aligned} & \operatorname{sol} = \operatorname{Vector}[-3.22108361124992, 728.225301265717, -36159.6916503906, 724344.018554688, \\ & -7402839.390625, \ 43077150.21875, \ -148688391.75, \ 295754646.5, \ -275331067.0, \ -47381104.0, \ 241467401.0, \ 230313529.5, \ -896512947.0, \ 851724014.0, \ -286636087.0, \\ & -30766117.5, \ 29693161.5] \end{aligned}$

err1 = 1.38777878078145e-16

err2 = 9.96697505353901

N = 18

sol = Vector[-13.0393237173557, 2013.47146224976, -75089.7684326172, 1170257.80664062, -9283656.0390625, 39911290.9375, -85617516.25, 27606976.125, 287936783.125, -

 $636646629.0,\ 483502426.0,\ -37884332.0,\ 181593149.0,\ -677921490.0,\ 485359060.0,\ 129016824.0,\ -273697056.0,\ 85027202.375]$

err1 = 1.38777878078145e-16

err2 = 8.97655984917988

N = 19

 $\begin{aligned} & \text{sol} = \text{Vector}[\text{-}10.575250312686, \ 1318.33017349243, \ \text{-}38301.4317016602, \ 435838.275390625, \\ & -2257109.34375, \ 5587472.8125, \ \text{-}14746395.5, \ 107632294.5, \ \text{-}505401583.0, \ 1146077540.0, \\ & -970212002.0, \ \text{-}862750557.0, \ 2432907180.0, \ \text{-}1225124416.0, \ \text{-}950342148.0, \ 820311713.0, \\ & 549310000.0, \ \text{-}743750960.5, \ 212360400.75] \end{aligned}$

err1 = 2.25514051876985e-16

err2 = 2.63674345603329

N = 20

 $sol = Vector[2.00276151299477, -391.758289337158, 18499.1511230469, -364428.8125, \\ 3674418.6875, -20653072.0625, 65033263.25, -100826814.25, 24646409.0, 77025627.0, \\ 218001766.0, -800348840.0, 587811592.0, 426437661.0, -533096001.0, -313934860.0, \\ 292470652.0, 454225316.5, -534529997.0, 154409446.5]$

err1 = 1.04083408558608e-16

err2 = 1.33007968616477

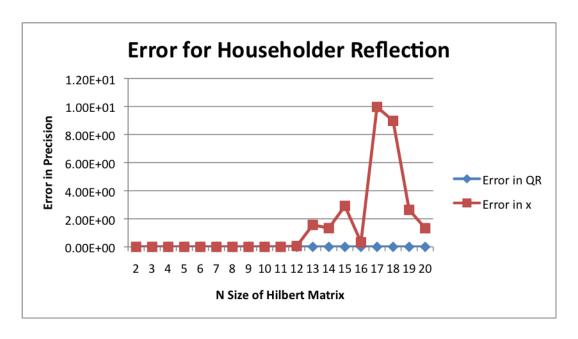


Figure 2: Error for Householder Reflection.

There are a few trends to examine within these results. First and foremost, it is clear that for the first handful of iterations, the Householder Reflections are extremely accurate. However, it is important to note that inaccuracy present at, and the huge increases in err2 after N=13. Since the errors measured above indicated this, but didn't offer any conclusive reason why (other than error amplification in \vec{x} , further investigations were necessary. Invoking the same property above, and using the property $||A|| = |\max \lambda_i|$, I used the Power method to find the absolute value of the maximum eigenvalue of A to compute the cond(R) to determine the inherent error amplification in the problem.

Listing 5: Determining cond(R) for Householder

```
for i in 2..20
  m = Matrix.hilbert(i)
  q,r = m.householder
  norm_r = r.power_method.abs
  norm_r_inv = r.inverse.power_method.abs
  puts "cond(N=#{i}): #{norm_r * norm_r_inv}"
end
```

As expected, this code snippet printed out condition numbers for each matrix... up until the 13th iteration.

```
\begin{array}{l} \operatorname{cond}(N=2)\colon 15.0000000633676\\ \operatorname{cond}(N=3)\colon 299.040132219316\\ \operatorname{cond}(N=4)\colon 6349.76484639525\\ \operatorname{cond}(N=5)\colon 138457.647018696\\ \operatorname{cond}(N=6)\colon 3063599.23761458\\ \operatorname{cond}(N=7)\colon 68420347.6630276\\ \operatorname{cond}(N=8)\colon 1537937583.62974\\ \operatorname{cond}(N=9)\colon 34733563292.8343\\ \operatorname{cond}(N=10)\colon 787269502031.433\\ \operatorname{cond}(N=11)\colon 17829448576474.5\\ \operatorname{cond}(N=12)\colon 4.02515205148727e+14\\ \end{array}
```

Having let the program run for twenty minutes, fully utilizing one 2.4Ghz core and 8MB RAM, we determined that the solution could not be computed. This was because the Power Method did not converge to an eigenvalue for $||R^{-1}||$. This is most likely because the matrix itself is ill-conditioned, having incredibly small eigenvalues. This is more than likely a contributing factor of the inaccuracy of the results after the 13th iteration.

2.3 Givens Rotations

The Givens Rotation, like the Householder Reflection, is a more efficient and stable decomposition method than Gram-Schmidt. Unlike the Householder Reflection, however, the Givens

Rotation uses counterclockwise rotations in order to zero out items below the diagonal. This is done based off of the principle of a standard rotation matrix in 2 dimensions:

$$\left[\begin{array}{cc} c & s \\ -s & c \end{array}\right] \left[\begin{array}{c} a \\ b \end{array}\right] = \left[\begin{array}{c} r \\ 0 \end{array}\right]$$

where

$$c = \frac{a_{ii}}{\sqrt{a_{ii}^2 + a_{ij}^2}}$$

and

$$s = -\frac{a_{ij}}{\sqrt{a_{ii}^2 + a_{ij}^2}}$$

and c and s represent $\cos \theta$ and $\sin \theta$ respectively.

2.3.1 Explanation of the Algorithm

- 1. Iterate through the columns of A, looking for non-zero entries below the diagonal. When such an entry is found, let it be a_{ij} .
- 2. Find c and s using the formulas above.
- 3. Create a new Identity matrix of size n, called G_n .
- 4. Replace g_{ii} and g_{jj} with c. Then, replace g_{ij} with -s and g_{ji} with s.
- 5. Let $A_n = G_n A_{(n-1)}$.
- 6. Repeat the above process until A_n is upper triangular.
- 7. $Q = G_1^t G_2^t ... G_n^t$ and $R = A_n$.

2.3.2 Implementation of the Algorithm

Listing 6: QR Decomposition via Givens Rotations

```
class Matrix
  def givens
    return nil unless self.square?
    n = self.row_size
    a = self
    g_n = []
    cvs = a.column_vectors.map { |v| v.to_a }
    for i in 0...cvs.length
```

```
for j in 0...cvs.length
    next unless j > i
    g = Matrix.identity(n).to_a
    c = cvs[i][i] / Math.sqrt(cvs[i][i]**2 + cvs[i][j]**2)
    s = -cvs[i][j] / Math.sqrt(cvs[i][i]**2 + cvs[i][j]**2)
    g[i][i], g[j][j] = c, c
    g[j][i], g[i][j] = s, -s
    g = Matrix[*g]
    g_n << g
    a = g * a
    cvs = a.column_vectors.map { |v| v.to_a }
    end
    end
    q,r = g_n.map { |m| m.t }.inject(&:*), a
    return q,r
    end
end
end</pre>
```

Matrix#givens, like Matrix#householder, iterates across the column vectors of A_n . The nested for loops serve to track i and j for the iterations. The inner loop automatically skips to the next iteration unless the current j is below the diagonal. Next, c and s are assigned, and inserted into g. The Array \rightarrow Matrix conversion is again used due to the lack of an Matrix#[]= method.

Once the loops are finished, q and r are computed, with a map to transpose the G_n and an injection to find the product of all G_n in the case of the former. Finally, they are returned.

2.3.3 Results and Analysis

```
\begin{array}{l} N=2\\ sol=Vector[-2.0,\ 6.0]\\ err1=5.55111512312578e-17\\ err2=9.93013661298909e-16\\ \\ N=3\\ sol=Vector[3.00000000000007,\ -24.000000000004,\ 30.00000000000004]\\ err1=1.94289029309402e-16\\ err2=1.03578512786223e-14\\ \\ N=4\\ sol=Vector[-3.9999999999983,\ 59.99999999975,\ -179.99999999994,\ 139.99999999999]\\ err1=5.55111512312578e-17\\ err2=3.11850560999409e-13\\ \\ N=5\\ sol=Vector[4.99999999999852,\ -119.99999999996,\ 630.0000000000051,\ -1120.00000000016,\ 630.0000000000073] \end{array}
```

```
err1 = 1.2490009027033e-16
err2 = 1.36606012668301e-11
N = 6
-6300.00000049779, 2772.00000019732]
err1 = 1.80411241501588e-16
err2 = 2.77657346835272e-10
N = 7
sol = Vector[7.0000000199434, -336.000000780041, 3780.00000739936, -16800.0000284091,
34650.0000515878, -33264.0000442229, 12012.0000144187]
err1 = 6.93889390390723e-17
err2 = 1.55557238094569e-09
N = 8
sol = Vector[-7.99999896803638, 503.999942819588, -7559.99923548102, 46199.9957879186,
-138599.988499641, 216215.983540893, -168167.988176227, 51479.9966375828
err1 = 8.32667268468867e-17
err2 = 7.94069060736149e-08
N = 9
sol = Vector[8.99992426042445, -719.994748193771, 13859.9107559323, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818863, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.360818860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -110879.008860, -1108799.008860, -1108799.008860, -1108799.008860, -1108799.008860, -1108799.008860, -1108799.008860, -1108799.008860, -1108799.008860, -1108799.008860, -1108799.008860,
450447.649126053, -1009003.18890953, 1261254.46387482, -823676.650474548, 218789.171199322
err1 = 7.63278329429795e-17
err2 = 4.66503786431221e-06
N = 10
-1261083.16352844, 3783295.4758606, -6725927.26672363, 7000515.76196289, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.26672363, -6725927.2667236, -6725927.2667236, -6725927.2667236, -6725927.2667236, -6725927.2667236, -6725927.2667236, -6725927.2667236, -6725927.2667236, -6725927.266729, -6725927.266729, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -6725927, -672592, -672592, -672592, -672592, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259, -67259
3937819.6257019, 923692.114944458]
err1 = 5.55111512312578e-17
err2 = 0.000118021401875031
N = 11
3145750.00866699, -12082177.5825195, 28532307.1542969, -41931200.3964844, 37349712.1035156,
-18446337.2919922, 3874106.02575684]
err1 = 6.93889390390723e-17
```

err2 = 0.00195217793208091

N = 12

sol = Vector[-11.7637949250638, 1689.72214841843, -59325.9497375488, 891934.624389648, -7147851.09570312, 34069381.0234375, -102316196.78125, 198536336.8125, -248341644.9375, 193264411.125, -85076605.21875, 16178025.9941406]

err1 = 4.16333634234434e-17

err2 = 0.115012096881196

N = 13

 $\begin{aligned} & sol = Vector[-61.8027538955212, 9494.16192626953, -359111.035888672, 5871992.68164062, \\ & -51801904.90625, 276049891.0, -945991737.0, 2154025841.0, -3292805052.0, 3339661792.0, \\ & -2154944515.0, \, 800648371.0, \, -130364935.875] \end{aligned}$

err1 = 6.24500451351651e-17err2 = 8.90425266073271

N = 14

sol = Vector[7.45895887166262, -1295.68168830872, 54003.9275512695, -952275.251953125, 8870901.421875, -48738288.0, 166811664.25, -360888921.25, 476149459.5, -321513091.25, -8250174.0, 188159185.5625, -129611016.0, 29910008.9082031]

err1 = 9.0205620750794e-17err2 = 0.62963504942535

N = 15

 $\begin{aligned} & \text{sol} = \text{Vector}[12.3366819694638, -1821.77809238434, 65374.663269043, -992064.502929688, \\ & 7841929.44140625, -35339854.515625, 92208062.125, -126442974.375, 41272087.75, \\ & 109135805.75, -80595176.125, -157957808.5, 290534658.25, -180913882.0, 41185843.65625] \\ & \text{err1} = 6.24500451351651\text{e}\text{-}17 \\ & \text{err2} = 1.07872075743074 \end{aligned}$

N = 16

 $\begin{aligned} & \operatorname{sol} = \operatorname{Vector}[70.3090637922287, -11236.7543334961, \, 440947.072753906, -7426134.3671875, \\ & 66566162.375, -351925447.5, \, 1145727382.0, \, -2272672160.0, \, 2441160344.0, \, -556420716.0, \\ & -1655614908.0, \, 1192563112.0, \, 1264526288.0, \, -2290058780.0, \, 1285152142.0, \, -262007026.0] \\ & \operatorname{err1} = 4.85722573273506\text{e}\text{-}17 \end{aligned}$

err2 = 30.713229091634

N = 17

 $sol = Vector[-3.61188149452209, 725.89476776123, -33253.7346801758, 621643.462890625, \\ -5928474.9375, 31899825.4375, -99654787.25, 170816763.0, -115289738.0, -55190619.125, \\ 15787368.75, 275831749.625, -188990795.5, -450670627.5, 804762953.25, -494987422.6875, \\ -640670627.5, -64067067.5, -64067067.5, -64067067.5, -64067067.5, -64067067.5, -64067067.5, -64067067.5, -6406706$

```
111024901.21875
err1 = 9.0205620750794e-17
err2 = 1.22101337499855
N = 18
-5161550.10546875, 34689883.0859375, -131933774.5, 287567094.5, -336914135.0,
214197943.0, -338028503.0, 819475402.0, -438687772.0, -1515573252.0, 3161195832.0,
-2670750844.0, 1102708686.0, -183166515.5
err1 = 1.07552855510562e-16
err2 = 5.73895360472133
N = 19
sol = Vector[-5.86808938533068, 761.973731994629, -22226.7529602051, 229556.625488281,
-614158.8046875, -5032363.90625, 43757105.6875, -130368521.375, 118144546.0,
259113075.25, -781019987.5, 660212292.0, 74774548.5, -430512041.5, 483743559.0,
-992752915.0, 1328748118.0, -817227380.5, 188826241.8125]
err1 = 8.32667268468867e-17
err2 = 10.3948622428631
N = 20
sol = Vector[-8.68420749902725, 1128.25052070618, -33044.6605834961, 343626.84375,
-916979.609375, -8018512.4375, 71636052.25, -228503793.0, 268607218.0, 276694216.5,
-1191775110.0, 1161745104.0, 109201064.0, -942929664.0, 881892328.0, -1341010140.0,
2050000240.0, -1638338480.0, 616628050.0, -85223057.0
err1 = 6.59194920871187e-17
err2 = 3.44515872161303
```

Again, like the Householder Reflection method, the Givens Rotation suffers odd errors once it reaches the 13×13 Hilbert matrix. And again, the problem arises because it is an ill-conditioned matrix.

Listing 7: Determining cond(R) for Givens

```
\label{eq:formula} \begin{array}{ll} \textbf{for i in } 2..20 \\ m = Matrix.hilbert(i) \\ q,r = m.givens \\ r\_norm = r.power\_method.abs \\ r\_inv\_norm = r.inverse.power\_method.abs \\ puts "cond(N=\#\{i\}): \#\{r\_norm * r\_inv\_norm\}" \\ \textbf{end} \\ Results were: \\ \\ cond(N=2): 15.0000011125923 \end{array}
```

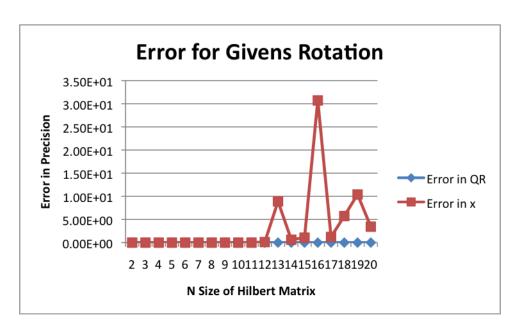


Figure 3: Error for Givens Rotation.

 $\begin{array}{l} \operatorname{cond}(N=3)\colon 299.040144310689\\ \operatorname{cond}(N=4)\colon 6349.7651396892\\ \operatorname{cond}(N=5)\colon 138457.652248714\\ \operatorname{cond}(N=6)\colon 3063599.34834093\\ \operatorname{cond}(N=7)\colon 68420350.0720479\\ \operatorname{cond}(N=8)\colon 1537937656.71739\\ \operatorname{cond}(N=9)\colon 34733602598.1239\\ \operatorname{cond}(N=10)\colon 787216766308.864\\ \operatorname{cond}(N=11)\colon 17863296168299.8\\ \operatorname{cond}(N=12)\colon 4.05102981075616e+14\\ \operatorname{cond}(N=13)\colon 2.11135930217763e+16\\ \end{array}$

The Givens Rotation's cond(R) ceases to calculate at N=14, one dimension after the Householder Reflection.

3 Part Two

Listing 8: Convolutional Code

```
import java.util.Random;
public class ConvolutionalCode
    public static void main (String [] args)
        Matrix input, output, Y0, Y1, A, B;
        //Encode specific input sequence
        double[] vector = \{1, 0, 1, 1, 1, 0, 1, 1, 1, 0, 0, 0, 1, 0\};
        input = new Matrix(vector);
        System.out.print(input);
        System.out.println("Encoding...");
        //Encodes the Y0 output stream for the input
        A = new Matrix(input.length, input.length);
        A = A. fillMatrixZero();
        Y0 = A. matrixMultiply(input);
        System.out.print(Y0);
        //Encode the Y1 output stream for the input
        \stackrel{'}{B} = new Matrix(input.length, input.length);
        B = B. fillMatrixOne();
        Y1 = B. matrixMultiply(input);
        System.out.print(Y1);
        //Generates a random input stream and encodes both output streams
        input = Matrix.generateInput();
        A = new Matrix(input.length, input.length);
        A = A. fillMatrixZero();
        B = new Matrix(input.length, input.length);
        B = B. fill Matrix One ();
        Y0 = A. matrix Multiply (input);
        Y1 = B. matrixMultiply(input);
        System.out.println("Generating input...");
        System.out.print(input);
        System.out.print("Encoding...");
        System.out.print(Y0);
        System.out.print(Y1);
        //Generates two random output stream and decodes their input streams using
        //Jacobi and Gauss-Seidel iterations. The first random stream is decoded as
        //a Y0 stream and the second as a Y1 stream. Both iterative methods are fun 8 times.
        int maxit = 8;
        System.out.println("Generating random output...");
        output = Matrix.generateOutput();
        A = new Matrix (output.length, output.length);
        A = A. fillMatrixZero();
        B = new Matrix(output.length, output.length);
        B = B. fill Matrix One ();
        System.out.print(output);
        System.out.println("Decoding...");
        input = output.jacobiIteration(A, maxit);
        System.out.print(input);
```

```
input = output.gaussseidelIteration(A, maxit);
       System.out.print(input);
       System.out.println("Generating random output...");
       output = Matrix.generateOutput();
       A = new Matrix (output.length, output.length);
       A = A. fillMatrixZero();
       B = new Matrix(output.length, output.length);
       B = B. fillMatrixOne();
       System.out.print(output);
       System.out.println("Decoding...");
        input = output.jacobiIteration(B, maxit);
       System.out.print(input);
       input = output.gaussseidelIteration(B, maxit);
       System.out.print(input);
       System.out.print("Done");
   }
}
                                   Listing 9: Algorithms
import java.util.Random;
public class ConvolutionalCode
    public static void main (String[] args)
        Matrix input, output, Y0, Y1, A, B;
        //Encode specific input sequence
       input = new Matrix (vector);
       System.out.print(input);
       System.out.println("Encoding...");
        //Encodes\ the\ YO\ output\ stream\ for\ the\ input
       A = new Matrix(input.length, input.length);
       A = A. fillMatrixZero();
       Y0 = A. matrixMultiply(input);
       System.out.print\left( Y0\right) ;
        //Encode the Y1 output stream for the input
       B = new Matrix(input.length, input.length);
       B = B. fillMatrixOne();
       Y1 = B. matrixMultiply(input);
       System.out.print(Y1);
        //Generates a random input stream and encodes both output streams
       input = Matrix.generateInput();
       A = new Matrix(input.length, input.length);
       A = A. fillMatrixZero();
       B = new Matrix(input.length, input.length);
       B = B. fillMatrixOne();
       Y0 = A. matrixMultiply(input);
       Y1 = B. matrix Multiply (input);
       System.out.println("Generating input...");
       System.out.print(input);
       System.out.print("Encoding...");
```

```
System.out.print(Y0);
       System.out.print(Y1);
       //Jacobi and Gauss-Seidel iterations. The first random stream is decoded as
       \dot{I}/a Y0 stream and the second as a Y1 stream. Both iterative methods are fun 8 times.
       int maxit = 8;
       System.out.println("Generating random output...");
       output = Matrix.generateOutput();
       A = new Matrix (output.length, output.length);
       A = A. fill Matrix Zero ();
       B = new Matrix (output.length, output.length);
       B = B. fillMatrixOne();
       System.out.print(output);
       System.out.println("Decoding...");
       input = output.jacobiIteration(A, maxit);
       System.out.print(input);
       input = output.gaussseidelIteration(A, maxit);
       System.out.print(input);
       System.out.println("Generating random output...");
       output = Matrix.generateOutput();
       A = new Matrix(output.length, output.length);
       A = A. fillMatrixZero();
       B = new Matrix (output.length, output.length);
       B = B. fillMatrixOne();
       System.out.print(output);
       System.out.println("Decoding...");
       input = output.jacobiIteration(B, maxit);
       System.out.print(input);
       input = output.gaussseidelIteration(B, maxit);
       System.out.print(input);
       System.out.print("Done");
   }
}
```

4 Part Three

The purpose of Part Three is to implement the Power Method in order to find the largest eigenvalue of an $n \times n$ matrix, in particular a Leslie matrix that models the population of a city. Using the Leslie matrix we can study the growth of populations by modeling the population distribution in future years and exploring the meaning of the Leslie matrix eigenvalue.

4.1 Leslie Matrices

Leslie Matrices are used to describe population growth in a city. As shown in the problem, the first row of the matrix represents the per capita average number of female offsprings of that age group, or the fecundity. The following numbers in each column are the fractions of people surviving to the next age class. Each entry in the given population distribution matrix is the part of the population in that age group, each group being of ten years (0-9,10-19, etc). The first population distribution matrix given is for the year 2000 and is shown later in listing 7.

The given Leslie Matrix of the city shown in Ruby is shown later on in listing 7 as well.

4.2 Power Method

The Power Method is used to calculate eigenvalues of $n \times n$ matrices. In the class notes it was shown that after defining

$$\vec{u}(n+1) = A\vec{u}(n)$$

it follows that

$$\lim_{n \to \infty} \frac{\vec{w} \cdot \vec{u}(n+1)}{\vec{w} \cdot \vec{u}(n)} = \lambda$$

The following is an implementation of the Power Method in Ruby. Its explanation is found in the solutions below.

4.3 Implementation of the Power Method

Listing 10: Power Method

```
class Matrix
  def power_method
    a = self
    u_prev = Array.new(a.column_size) { 1 }
    w = Array.new(a.column_size) { | i | i == 0 ? 1 : 0 }
```

```
lambda, innerProd1, innerProd2 = 0, 0, 0
    u_new_matrix = a * Vector[*u_prev].covector.transpose
    u_new_array = u_new_matrix.to_a
    u_new = []
     \begin{tabular}{lll} \textbf{for} & i & \textbf{in} & 0 \dots u\_new\_array.size \\ \end{tabular} 
      u_new.push(u_new_array[i][0])
    innerProd1 = innerProd1 + Vector[*w].inner_product(Vector[*u_new])
    innerProd2 = innerProd2 + Vector[*w].inner_product(Vector[*u_prev])
    lambda_prev = lambda
    lambda = (innerProd1/innerProd2).to_f
    u_prev = u_new
    u_new_matrix = a * Vector[*u_prev].covector.transpose
    u_new_array = u_new_matrix.to_a
    u_new = []
    for i in 0...u_new_array.size
      u_new.push(u_new_array[i][0])
    while true
      innerProd1 = innerProd1 + Vector[*w].inner_product(Vector[*u_new])
      innerProd2 = innerProd2 + Vector[*w].inner-product(Vector[*u-prev])
      lambda_prev = lambda
      lambda = (innerProd1/innerProd2).to_f
      u_prev = u_new
      u_new_matrix = a * Vector[*u_prev].covector.transpose
      u_new_array = u_new_matrix.to_a
      u_new = []
      for i in 0...u_new_array.size
        u_new.push(u_new_array[i][0])
      break if (lambda - lambda_prev).abs <= 0.00000001
    return lambda
  end
end
```

4.4 Answering the Question of Population Trends

4.4.1 Social Factors and Influences of the Leslie Matrix

The original problem asks to interpret the data in the matrix and discuss the social factors that influence those numbers. The solution was derived from an analysis of the given Leslie matrix.

In the given Leslie matrix the survival rate in the different age groups and the fecundity both vary. The fecundity can be seen when analyzing the first row of the matrix. In the age groups 0-9 and 50 and over, the average birth of females from that age class is zero because women are not physically fertile. In the ranges 10-19 and 20-29 the birth rate is low because women generally have the most children in their thirties, which is apparent in the .9 per capita average in that age group.

Looking at the data in each of the columns of the matrix, the fraction of surviving

individuals varies due to several factors. The number in each column represents the surviving fraction that makes it to the next group. Newborns of age 0-9 have a survival rate of .7 possibly due to infant deaths which occur due to birth complications and medical problems. The survival rates increase in the groups of 10-19 and 20-29, particularly because individuals are the healthiest during this "prime" of their life. The rate slowly decreases after 30-39 and over the next few age groups until it drops .4 in the 70-79 age group in which health complications due to old age cause more deaths.

4.4.2 Projected Population Distributions

In this implementation, the output is written to a file using the **content** variable and the new matrices for the population distributions are shown in the code.

Listing 11: Populations output = File.new("data3.txt", "w+") content = "" leslie = Matrix[[0, 1.2, 1.1, 0.9, 0.1]][0.7,0, 0, 0, 0, 0, 0,,0], [0, 0.85, 0, 0, 0, 0,0 ,0 ,0,0], $\begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$,0 0.9, 0, 0,,0]0, 0 ,00, 0.9, 0,0 ,0 ,0,0],0 0, 0, 0.88, 0,0], [0, 0,[0, 0,0, 0, 0, 0,0.8,0,0 ĺ0, 0, 0, 0, 0, 0,0, 0.77, 0,00, 0, 0, # largest eigenvalue of leslie content << "largest lambda: #{leslie.power_method}\n"</pre> # population distributions [0.924], [0.36]] $x_{-2} = leslie * x_{-1} \# > Matrix [[5.1875], [4.445], [1.2495], [1.4535], [1.458], [1.6632],$ [1.408], [1.0472], [0.3696]] $x_{-3} = leslie * x_{-2} \# Matrix[[8.1624], [3.63125], [3.77825], [1.12455], [1.30815],$ [1.28304], [1.33056], [1.08416], [0.41888] $x_4 = leslie * x_3 \not \# > \mathit{Matrix}[[9.656485], [5.71368], [3.0865625], [3.400425], [1.012095],$ [1.151172], [1.026432], [1.0245312], [0.433664] $x_{-5} = leslie * x_{-4} # > Matrix[[13.41322675], [6.7595395], [4.856628], [2.77790625],$ [3.0603825], [0.8906436], [0.9209376], [0.79035264], [0.40981248]# total populations populations = [$x_0. to_a. flatten. inject(\&:+) * 10000,$ $x_1 \cdot to_a \cdot flatten \cdot inject(\&:+) * 10000,$ $x_2 \cdot to_a \cdot flatten \cdot inject(\&:+) * 10000,$ $x_3 \cdot to_a \cdot flatten \cdot inject(\&:+) * 10000,$

 $x_4.to_a.flatten.inject(\&:+) * 10000,$

```
x_{-5}. to_a. flatten.inject(&:+) * 10000
content << "2010: #{populations[1]}
2020: #{populations[2]}
2030: #{populations[3]}
2040: #{populations [4]}
2050: #{populations [5]}
# percentage of population change
content << "2000 -> 2010: %.2 f\%
2010 \rightarrow 2020: \%.2 f\
2020 \rightarrow 2030: \%.2 \text{ f}\
2030 \rightarrow 2040: \%.2 \text{ f}
2040 -> 2050: %.2 f\%" % [((populations [1] - populations [0]) / populations [0] * 100),
                            ((populations[2] - populations[1]) / populations[1] * 100),
                            ((populations[3] - populations[2]) / populations[2] * 100),
                            ((populations [4] - populations [3]) / populations [3] * 100),
                            ((populations [5] - populations [4]) / populations [4] * 100)]
```

The next section requires a calculation of the population distribution, total population, and fraction of change of the total population in 2010, 2020, 2030, 2040, and 2050.

The approach to this problem was to use the Leslie Matrix Model and multiply the Leslie Matrix by each previous population distribution, starting in 2010 and using the given distribution for 2000.

The population distribution is shown in the code as matrices and written is

```
 \vec{x}(1) = (6.35, 1.47, 1.615, 1.62, 1.89, 1.76, 1.36, 0.924, 0.36)(10^5) 
 \vec{x}(2) = (5.1875, 4.445, 1.2495, 1.4535, 1.458, 1.6632, 1.408, 1.0472, 0.3696)(10^5) 
 \vec{x}(2) = (8.1624, 3.63125, 3.77825, 1.12455, 1.30815, 1.28304, 1.33056, 1.08416, 0.41888)(10^5) 
 \vec{x}(4) = (9.656485, 5.71368, 3.0865625, 3.400425, 1.0121, 1.1512, 1.02643, 1.02453, 0.4337)(10^5) 
 \vec{x}(5) = (13.413227, 6.75954, 4.8566, 2.777906, 3.06038, 0.89064, 0.92094, 0.790353, 0.409812)(10^5)
```

The total population of the city in question for each year was as follows:

```
2010: 173490.0
2020: 182815.0
2030: 221212.4
2040: 265050.467
2050: 338794.2932
```

With the fractional change of the total population:

 $2000 \rightarrow 2010$: 22.18% $2010 \rightarrow 2020$: 5.37% $2020 \rightarrow 2030$: 21.00% $2030 \rightarrow 2040$: 19.82% $2040 \rightarrow 2050$: 27.82%

By implementing the formula given in the problem

$$\vec{x}(k+1) = A\vec{x}(k)$$

where k = 0, 1, 2, 3, 4, 5 and $\vec{x}(0)$ is the given population distribution, the population distribution could be calculated for each decade increase. This, more simply put, is a multiplication of the Leslie matrix with the previous population distribution. This is seen in the code as $x_1 = leslie * x_0$ where x_0 is the $\vec{x}(0)$ population distribution. The next decades are calculated in the same manner. In order to find the total population of one year, a simple summation of the elements of each of the population distribution matrices (multiplied by 10^5) shown in the code as population gives the total population. The percent increase of the population was calculated by taking the difference in total population divided by the previous population, for each decade as seen in the last portion of the code in listing 7.

4.4.3 Obtaining the Eigenvalue and its Meaning

In order to evaluate these populations it was required to write a program implementing the Power method to calculate the largest eigenvalue of an $n \times n$ matrix where the method stops after 8 digits of accuracy. The program could be used to calculate the largest eigenvalue of the given Leslie matrix.

The bulk of this problem was in developing a Power method algorithm in Ruby that took in an $n \times n$ matrix that iterated until 8 digit decimal precision was achieved. After testing to see if the program returned correct eigenvalues for any square matrix, it was passed the Leslie matrix in order to find the largest eigenvalue.

After running the program using the Leslie matrix, the method returns the eigenvalue 1.28865626235758.

Implementing the power method in Ruby only required the creation of one method because of some shared coding with part 1 including the Vector#inner_product method to calculate the dot product, used in part of the algorithm. The general formula for the power method can be found in the web notes by Laszlo Erdos for *Linear Algebra: Numerical Methods in Numerical Computation of Eigenvalues*(p 62).

$$\lim_{n \to \infty} \frac{\vec{w} \cdot \vec{u}(n+1)}{\vec{w} \cdot \vec{u}(n)} = \lambda$$

where \vec{w} and \vec{u}_0 are randomly chosen nonzero vectors. As shown in the declarations of \mathbf{u} -prev and \mathbf{w} in the power method program implemented uses a \vec{w} vector of $[1,0,0...,0_n]^t$ and a \vec{u}_0 vector of $[1,1,1...,1_n]^t$ both with the number of entries equal to n of the $n \times n$ Leslie matrix. These are the vectors we used in class to compute eigenvalues with the Power Method but they could be any vectors as long as they are not zero. The coding of the algorithm for the actual iterations of the method can be seen in the while loop. There are two general parts to calculating the eigenvalue. First, \vec{u}_1 is calculated and later $\vec{u}_{(n+1)}$ with the same code but over many iterations in the while loop, represented in the code, after certain conversions, as \mathbf{u} -new. This is modeled by $\vec{u}_{(n+1)} = A\vec{u}_{(n)}$ on page 62 of the web notes by Erdos, where in the code \mathbf{u} -new is $\vec{u}_{(n+1)}$, \mathbf{u} -prev is \vec{u}_n and \mathbf{a} is A, the given matrix.

The second part involves the eigenvalue calculated using the inner products according to the limit exampled earlier. The code is first implemented when using innerProd1 and innerProd2 and later on over multiple iterations, where lambda is the approximated eigenvalue and lambda_prev is saved for precision calculation. The break statement terminates the loop after the difference in lambdas is less than 10^{-8} , guaranteeing that precision. The first iteration does not have a lambda to compare so it occurs outside and before the loop, creating some repetition of code. When converting matrices to arrays in order to multiply the inner product we encounter a small problem in Ruby. When using the Matrix#.to_a the given matrix is converted to an array of arrays containing numbers instead of just an array of numbers. To take care of this, a new array is created and filled with the elements of the old array shown as u_new.push(u_new_array[i][0]). In the code the first calculation of $\vec{u}_{(n+1)}$ is u_new_matrix and in Ruby is of the type Matrix. The array conversion is denoted as u_new_array and finally the array of numbers which can actually be used is u_new. Once lambda is calculated as a float point of precision to 8 decimals it is returned by Matrix#power_method.

The next part of the problem asked to analyze the eigenvalue obtained for the Leslie matrix and discuss what this means about the growth of the population in the long run.

Obtaining the eigenvalue of 1.28865626 reveals the eventual behavior of the population of the given city. The eigenvalue is greater than 1 which means the population will grow. If it was less than 1, the population would over time become extinct and if it was equal to 1 then the population would remain the same, with equal birth and death. This can be shown using the theorem covered in chapter 4 of Erdos's notes (p 49):

Theoreom 4.1

For a square matrix A, $||A^k|| \to 0$ if and only if, all eigenvalues of A satisfy $|\lambda| < 1$

When looking at the iterative model given in the question

$$\vec{x}(k) = A^k \vec{x}(0)$$

it is possible to see that a norm that converges to zero with a large enough k would eventually give a population of zero when multiplied by the distribution vector. This would happen if the largest eigenvalue was less than 1, and it can be drawn that because our eigenvalue computed for the Leslie matrix is greater than 1 the population must grow, getting closer to that rate as k gets larger.

Therefore we arrive at the equation

$$\vec{x}(n+1) = \lambda \vec{x}(n)$$

for large values of n, where λ is the dominant or largest eigenvalue.

This shows that as n gets larger $\vec{x}(n+1)$ is a scalar multiple of $\vec{x}(n)$ and if that scalar, the computed eigenvalue, is positive the population will eventually grow by that growth rate, as shown in our given case of the population of a city.

4.4.4 Supposing a Change in Birth Rate

In this implementation, the output is written to a file using the content variable

Listing 12: Populations

```
leslie = Matrix[[0, 0.6, 1.1, 0.9, 0.1]]
                    [0.7,0,0,0,0,0,0,
                    [0, 0.85, 0, 0,
                                                0
                                                                 ,0],
                                          0,
                                                    ,0
                                                           ,0
                     \begin{bmatrix} 0 & 0 & 0.9 & 0 & 0, \\ 0 & 0 & 0 & 0.9 & 0, \\ 0 & 0 & 0 & 0.9 & 0, \\ 0 & 0 & 0 & 0 & 0.8 \\ \end{bmatrix} 
                                                    ,0
                                                0
                                                           ,0
                                                                  ,0]
                                                0
                                                    ,0
                                                           ,0
                                                                  ,0],
                                                    ,0
                                0, 0, 0.88, 0
                                                                 ,0]
                    [0, 0, 0]
                                                                 ,0],
                                0, 0, 0, 0, 0.8, 0
                                                0, 0.77, 0
                              0\,,\quad 0\,,\quad 0\,,
                                                                  ,0],
                                0, 0, 0,
                                              0, 0, 0.40, 0]
x_3 = leslie * x_2
x_4 new = leslie * x_3 new
x_5new = leslie * x_4new
populations_new = [
                        x_2 \cdot to_a \cdot flatten \cdot inject(\&:+) * 10000,
                       x_3new.to_a.flatten.inject(&:+) * 10000,
                       x_4new.to_a.flatten.inject(&:+) * 10000,
                       x_5 new.to_a.flatten.inject(&:+) * 10000
content << "\n\nWith 2020's second set modified.
content << ~"2020: ~\#\{populations\_new\,[\,0\,]\,\}
2030: #{populations_new[1]}
2040: \#\{\text{populations\_new}[2]\}
2050: #{populations_new[3]}
content << "largest lambda: #{leslie.power_method}\n"</pre>
```

The problem asks to decrease the birth rate of the second age group by half in 2020 and then predict the population for 2030, 2040 and 2050, discussing again the meaning of the eigenvalue for this new matrix over a long period of time.

This was approached by changing the original Leslie matrix to reflect the new change in birth rate and then applying the previous method of calculating using the total populations and fractional change for the years of 2030, 2040, and 2050.

The newly computed eigenvalue is 1.16790279. The new populations of the following years are:

```
2020: 182815.0
2030: 194542.4
2040: 224593.967
2050: 262190.9132
```

and the percentage of change in populations are

```
2020 \rightarrow 2030: 6.41%

2030 \rightarrow 2040: 15.45%

2040 \rightarrow 2050: 16.74%
```

These values were arrived at by first altering the Leslie matrix as seen above, where the birth rate of the second age group, which was previous at 1.2 was decreased to 0.6. This birth rate changes in 2020 so the population distribution in 2010 and 2020 remains the same as presented in the data. In calculating the distribution in 2030, the calculation used before still applies but includes the new Leslie matrix. The population distribution in 2030 is equal to the new Leslie matrix times the population distribution in 2020. In Ruby this looks like x_3new = leslie * x_2 and the code following it, where x_3new, x_4new, x_5new are all of the new distribution matrices. Once these are calculated the total populations can be found by using the same method earlier, which was to add each of the numbers in the population distribution in populations_new. Finding the percent change was also the same as previously done.

The eigenvalue is computed by using the Power Method on the new Leslie matrix, using a simple call to the method leslie.power_method

This eigenvalue is still a positive number greater than 1, showing a population that will still grow positively and eventually level off at a factor of 1.167902791647756 for the new birth rate, which is somewhat lower than the previous rate. The population totals show the growing populations diminish due to the halved birthrate of the second age group in 2020. By 2050 where the population would have been 338794.2932 it is now 262190.9132. However, in the long run the population will not diminish because it's dominant eigenvalue is still greater than 1.