CSC111 Lecture 17: Iterative Sorting Algorithms, Part 1

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March 15, 2021

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1 Exercise 1: Implementing selection sort

Here is the skeleton of a selection sort algorithm we developed in lecture:

```
def selection_sort(lst: list) -> None:
1
2
         """Sort the given list using the selection sort algorithm.
3
4
        Note that this is a *mutating* function.
5
        >>> 1st = [3, 7, 2, 5]
6
7
        >>> selection_sort(lst)
        >>> lst
8
9
        [2, 3, 5, 7]
10
11
        for i in range(0, len(lst)):
            # Loop invariants
12
            # - lst[:i] is sorted
13
            \# - if i > 0, lst[i - 1] is less than all items in lst[i:]
14
15
            # Find the index of the smallest item in lst[i:] and swap that
16
            # item with the item at index i.
17
            index_of_smallest = _min_index(lst, i)
18
            lst[index_of_smallest], lst[i] = lst[i], lst[index_of_smallest]
19
20
21
    def _min_index(lst: list, i: int) -> int:
22
        """Return the index of the smallest item in lst[i:].
23
24
        In the case of ties, return the smaller index (i.e., the index that appears first).
25
```

```
26
27
         Preconditions:
              - 0 \le i \le len(lst) - 1
28
29
30
         >>> _min_index([2, 7, 3, 5], 1)
         2
31
         11 11 11
32
         min_index_so_far = i
33
34
         for x in range(i + 1, len(lst)):
              if lst[x] < lst[min_index_so_far]:</pre>
35
                  min_index_so_far = x
36
37
38
         return min_index_so_far
```

Complete this implementation by implementing the helper function _min_index. Hint: this is similar to one of the functions you implemented on this week's prep!

2 Exercise 2: Running-time analysis

1. Analyse the running time of the helper function $_min_index$ in terms of n, the length of the input 1st, and/or i, the second argument.

We will assume that the smallest item is at the end of the list for a worst case running time analysis.

```
def _min_index(lst: list, i: int) -> int:
1
2
       min_index_so_far = i
                                                 # 1 step
        for x in range(i + 1, len(lst)):
                                                 \# iterates n-i-1 times
3
            if lst[x] < lst[min_index_so_far]: # 1 step for whole if blcok</pre>
4
5
                min_index_so_far = x
6
7
        return min_index_so_far
                                                 # 1 step
```

Therefore an upper bound for the running time is $2 + n - i - 1 = 1 + n - i \in \mathcal{O}(n - i)$.

Next, we will can find an input family that proves a lower bound, but we will ommit it in this example.

Therefore, since we have found both an upper and lower bound that match, $RT_{\mathtt{min_index}} \in \Theta(n-i)$.

2. Analyse the running time of selection_sort.

We will assume that the list to be sorted is in reverse order, thus slection sort will take as long as possible.

```
def selection_sort(lst: list) -> None:
    for i in range(0, len(lst)):  # n steps
        index_of_smallest = _min_index(lst, i) # n - i steps
        lst[index_of_smallest], lst[i] =\  # 1 step
        lst[i], lst[index_of_smallest]
```

The loop runs n times for $i = 0, 1, \dots, n-1$

The iteration i takes n-i steps (because of $min_index(lst, i)$) plus one step for constant time operations.

Therefore
$$RT_{\texttt{selection_sort}} = \sum_{i=0}^{n-1} (n-i+1) \in \Theta(n^2).$$

3 Additional exercises

1. Translate the two loop invariants in selection_sort into Python assert statements. You can use is_sorted=/=is_sorted_sublist from this week's prep.

(One version is included in the Course Notes, but it's a good exercise for you to try it yourself without looking there first!)