

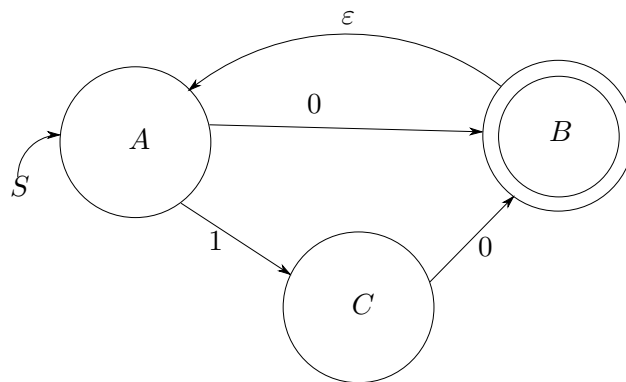
CSC236 Week 08: Machines, Expressions: Equivalence

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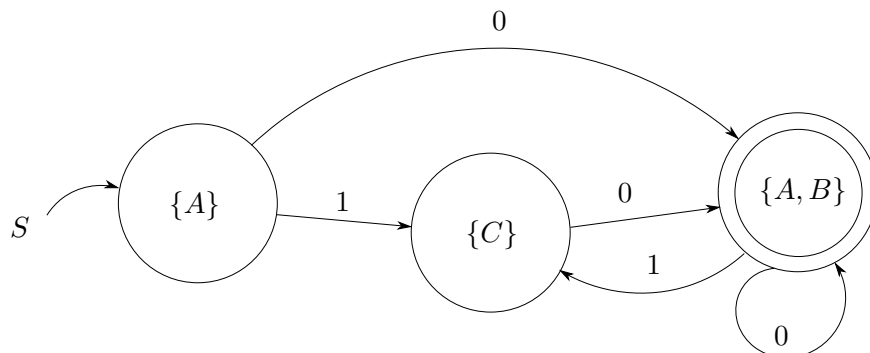
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1	NFSA that accepts $L((0 + 10)(0 + 10)^*)$	

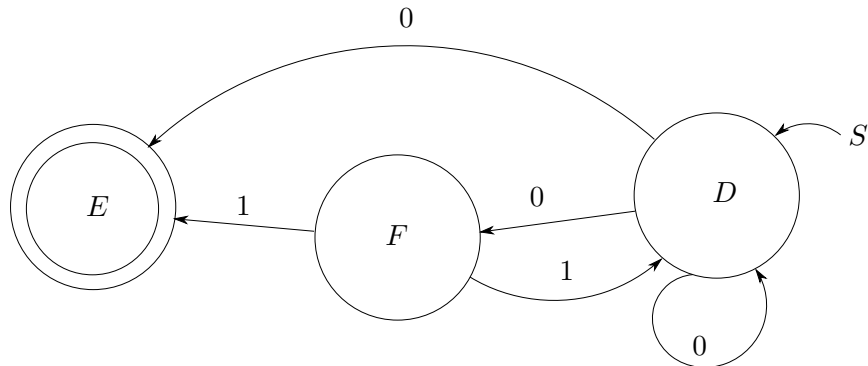


The ϵ transition makes it non deterministic. $A \xrightarrow{0} A \cup B$ and $C \xrightarrow{0} A \cup B$.

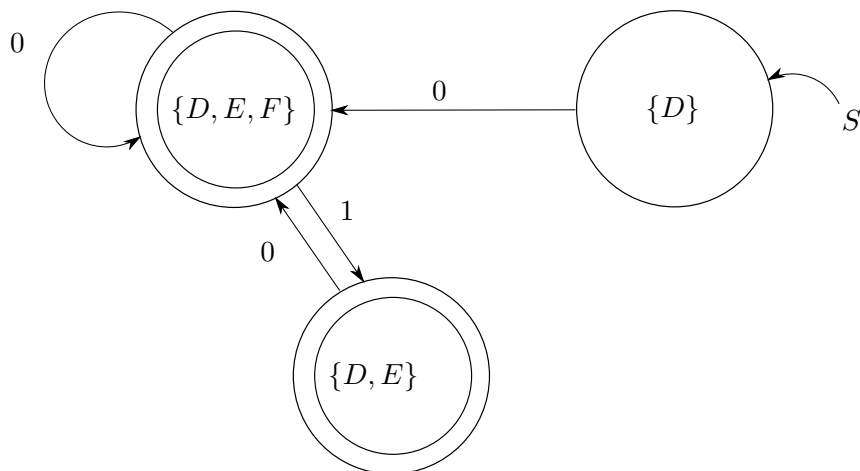
The corresponding DFSA is as follows:



2 NFSA that accepts $\text{Rev}(L((0 + 10)(0 + 10)^*))$



The corresponding DFSA is as follows:



3 FSAs and regexes are equivalent.

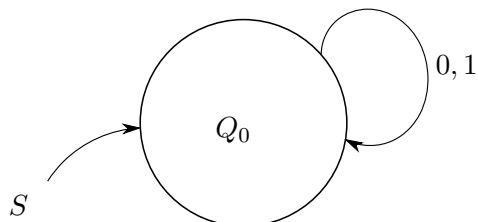
$L = L(M)$ for some DFSA $M \iff L = L(M')$ for some NFSA $M' \iff L = L(R)$ for some regular expression R .

Step 1.0: convert $L(R)$ to $L(M')$.

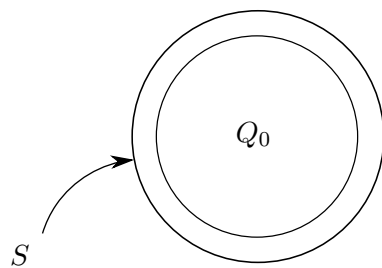
Start with $\emptyset, \varepsilon, a \in \Sigma$.

- Base case: Let s in $\{\emptyset, \varepsilon, a\}$ for some $a \in \Sigma$.

$L(\emptyset) = L(M)$, where M is:



$L(\varepsilon) = L(M)$, where M is:



For correctness, assume $a = 0$

$L(0) = L(M)$, where M is:

