

CSC236 Week 05: Languages: Definitions

Hisbaan Noorani

October 7 – October 13, 2021

Contents

1	Some definitions	1
2	More notation — string operations	2
3	Language operations	2
4	Another way to define languages	2
5	Regular expression to languages:	3
6	Regex examples	3

1 Some definitions

- **alphabet:** Finite, non-empty set of symbols, e.g. $\{a, b\}$ or $\{0, 1, -1\}$. Conventionally denotes Σ .
- **string:** Finite (including empty) sequence of symbols over an alphabet: abba is a string over $\{a, b\}$. Convention: ε is the empty string, never an allowed symbol, Σ^* is set of all strings over Σ .
- **language:** Subset of Σ^* for some alphabet Σ . Possibly empty, possibly empty, possibly infinite subset. E.e. $\{\}, \{aa, aaa, aaaa, \dots\}$.

N.B.: $\{\} \neq \{\varepsilon\}$. $|\{\}| = 0 \neq 1 = |\{\varepsilon\}|$

Many problems can be reduced to languages: logical formulas, identifiers fro compilation, natural language processing. Key question is recognition:

Given language L ans string s , is $s \in L$?

2 More notation — string operations

- **string length:** denotes $|s|$, is the number of symbols in s , e.g. $|bba| = 3$.
- $s = t$: if and only if $|s| = |t|$, and $s_i = s_t$, for $0 \leq i < |s|$.
- s^R : reversal of s is obtained by reversing symbols of s , e.g. $1011^R = 1101$.
- st **or** $s \circ t$: concatenation of s and t — all characters of s followed by all those of t , e.g. $bba \circ bb = bbabb$.
- s^k : denotes s concatenated with itself k times, e.g. $ab^3 = ababab$, $101^0 = \varepsilon$.
- Σ^n : all strings of length n over Σ , Σ^* denotes all strings over Σ .

3 Language operations

- \bar{L} : Complement of L , i.e. $\Sigma^* - L$. If L is a language of strings over $\{0, 1\}$ that start with 0, then \bar{L} is the language of strings that begin with 1 plus the empty string.
- $L \cup L'$: Union.
- $L \cap L'$: Intersection.
- $L - L'$: Difference.
- $\text{Rev}(L)$: $= cbrs^R : s \in L$
- **concatenation:** LL' or $L \circ L' = \{rt : r \in L, r \in L'\}$. Special cases $L\{\varepsilon\} = L = \{\varepsilon\}L$, and $L\{\} = \{\} = \{\}L$.
- **exponentiation:** L^k is concatenation of L , k times. Special case, $L^0 = \{\varepsilon\}$, including $L = \{\}$.
- **Kleene star:** $L^* = L^0 \cup L^1 \cup L^2 \cup \dots$

4 Another way to define languages

In addition to the set description $L = \{\dots\}$.

Definition: The regular expressions (regexprs or REs) over alphabet Σ is the smallest set such that

- \emptyset, ε , and x , for every $x \in \Sigma$ are REs over Σ .
- If T and S are REs over Σ , then so are:
 - $(T + S)$ (union) — lowest precedence operator
 - (TS) (concatenation) — middle precedence operator
 - T^* (star) — highest precedence

5 Regular expression to languages:

The $L(R)$, the language denoted (or described) by R is defined by structural induction.

- Basis; If R is a regular expression by the basis of the definition of regular expressions, then define $L(R)$:
 - $L(\emptyset) = \emptyset$ (the empty language – no strings!)
 - $L(\varepsilon) = \{\varepsilon\}$ (the language consisting of just the empty string)
 - $L(x) = \{x\}$ (the language consisting of the one-symbol string)
- Induction step: If R is a regular expression by the induction step of the definition, then define $L(R)$:
 - $L((T + S)) = L(S) \cup L(T)$
 - $L((TS)) = L(S)L(T)$
 - $L(T^*) = L(T)^*$

We are assuming about $(S + T)$ and (ST) above?

6 Regexp examples

- $L(0 + 1) = \{0, 1\}$
- $L(br0 + 1^*)$ All binary strings over $\{0, 1\}$
- $L((01)^*) = \{\varepsilon, 01, 0101, 010101, \dots\}$
- $L(0^*1^*)$ 0 or more 0s followed by 0 or more 1s
- $L(0^* + 1^*)$ 0 or more 0s or 0 or more 1s
- $L((0 + 1)(0 + 1)^*)$ Non-empty binary strings over $\{0, 1\}$