CSC110 Lecture 16: Greatest Common Divisor

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1 Ex 1: Property of the greatest common divisor

For your reference, here is the definition of greatest common divisor.

Let $x, y, d \in \mathbb{Z}$. We say that d is a **common divisor** of x and y when d divides x and d divides y. We say that d is the **greatest common divisor** of x and y when it the largest number that is a common divisor of x and y, or 0 when x and y are both 0.

We can define the function $gcd : \mathbb{Z} \times \mathbb{Z} \to \mathbb{N}$ as the function which takes numbers x and y, and returns their greatest common divisor.

And here is the Quotient-Remainder Theorem, slightly modified to handle negative integers.

Quotient-Remainder Theorem. Let $n \in \mathbb{Z}$ and $d \in \mathbb{Z}$. If $d \neq 0$ then there exist $q \in \mathbb{Z}$ and $r \in \mathbb{N}$ such that n = qd + r and $0 \leq r < d$. Moreover, these q and r are unique for a given n and d.

We say that q is the **quotient** when n is divided by d, and that r is the **remainder** when n is divided by d. We use n % d to denote the remainder of n divided by d.

And finally, here is the Divisibility of Linear Combinations Theorem that we covered in lecture today.

Divisibility of Linear Combinations. Let $m, n, d \in \mathbb{Z}$. If d divides m and d divides n, then d divides every linear combination of m and n.

1. Prove the following statement: For all $a, b, c \in \mathbb{Z}$, $b \neq 0$ and $c \mid a$ and $c \mid b$, then $c \mid a \% b$.

$$\forall a, b, c \in \mathbb{Z}, b \neq 0 \land c \mid a \land c \mid b \implies c \mid a \% b$$

Tip: introduce variables q and r using the Quotient-Remainder theorem in your proof.

Let $a, b, c \in \mathbb{Z}$

Assume $b \neq 0$, $c \mid a$, $c \mid b$

By QRT,
$$\exists q \in \mathbb{Z}, a = qb + r \implies r = a - qb$$

r is a linear combination of a and b

Since $c \mid a$ and $c \mid b$, then by DLC, $c \mid r$ which means $c \mid a \% b \blacksquare$

2. Prove the following statement: For all $a, b \in \mathbb{Z}$, if $b \neq 0$ then $\gcd(b, a \% b) \mid a$.

$$\forall a, b \in \mathbb{Z}, b \neq 0 \implies \gcd(b, a \% b) \mid a$$

Tip: introduce variables q and r using the Quotient-Remainder theorem in your proof.

Let $a, b \in \mathbb{Z}$

Assume $b \neq 0$

Let r = a % b, $d = \gcd(b, r)$

By QRT,
$$\exists q \in \mathbb{Z}, a = qb + r$$

Since $d = \gcd(b, r)$, then by definition, $d \mid b \land d \mid r$

Since $d \mid b \wedge d \mid r$, then by DLC $d \mid qb + r$

Therefore $d \mid a$ as needed.

- 3. Suppose we have three numbers $a, b, c \in \mathbb{Z}$, and we know that $b \neq 0$. If we want to prove that $c = \gcd(a, b)$, what would we need to prove, using the above definition of gcd?
 - (a) $c \mid a$
 - (b) c | b
 - (c) $\forall e \in \mathbb{N}, (e \mid a \land e \mid b) \implies e \leq c$

2 Ex 2: The Extended Euclidean Algorithm

We left off our discussion of the Extended Euclidean Algorithm in lecture with the following code:

```
def gcd_extended(a: int, b: int) -> Tuple[int, int, int]:
1
           """Return the gcd of a and b, and integers p and q such that
2
3
          gcd(a, b) == p * a + b * q.
 4
5
          >>> gcd_extended(10, 3)
6
 7
           (1, 1, -3)
8
          x, y = a, b
9
10
11
           # Initialize px, qx, py, and qy
           px, qx = 1, 0
12
           py, qy = 0, 1
13
14
           while y != 0:
15
               # Loop invariants
16
               assert math.gcd(a, b) == math.gcd(x, y) \# L.I.1
17
               assert x == px * a + qx * b
                                                         # L.I. 2
18
                                                         # L.I. 3
19
               assert y == py * a + qy * b
20
               q, r = div mod(x, y) # quotient and remainder when a is divided by b
21
22
23
               # Update x and y
24
               x, y = y, r
25
26
               # Update px, qx, py, and qy
27
               px, qx, py, qy = py, qy, px - q * py, qx - q * qy
```

return (x, px, qx)

1. Recall that the loop invariants must hold for the initial values of the loop variables. Given this, how should we initialize px, qx, py, and qy? (*Hint*: use the simplest possible numbers.)

Write your answer below, or directly in the gcd_extended code.

```
1 px, qx = 1, 0
2 py, qy = 0, 1
```

2. Inside the loop body, we need to figure out how to update the new variables px, qx, py, and qy. Now in order to derive the updates for px, qx, py, and qy, we need to do a bit of math.

For an arbitrary iteration of the while loop:

- $Let x_0, y_0, p x_0, q x_0, p x_0, p y_0$ be the values of the variables x, y, px, qx, py, qx, and qy at the *start* of the iteration.
- $x_1, y_1, px_1, qx_1, py_1, qy_1$ be the values of these variables at the end of the iteration.
- q and r be the values of the variables q and r during the iteration.
- (a) What do the loop invariants 2 and 3 tell you about the relationships between x_0 , y_0 , px_0 , qx_0 , py_0 , qy_0 ? (This is what we can assume to be True at the start of the loop iteration.)

$$x_0 = (px_0)(a) + (qx_0)(b)$$

 $y_0 = (py_0)(a) + (qy_0)(b)$

(b) What do the loop invariants tell you about the desired relationships between x_1 , y_1 , px_1 , qx_1 , py_1 , qy_1 ? (This is what we want to be True at the end of the loop iteration.)

```
x_1 = (px_1)(a) + (qx_1)(b)

y_1 = (py_1)(a) + (qy_1)(b)
```

(c) From the Quotient-Remainder Theorem, what is the relationship between x_0 , y_0 , q, and r? $x_0 = q \cdot y_0 + r$

```
(d) What are the values of x_1 and y_1 in terms of x_0, y_0, q, and/or r? x_1 = y_0 y_1 = r
```

(e) What should px_1 and qx_1 equal to satisfy the invariant for x_1 , using only the "0" variables?

```
px_1 = py_1qx_1 = qy_1
```

(f) (This is the hardest part) Given your answers from earlier parts, what should py_1 and qy_1 equal in order to satisfy the invariant for y_1 , using only the "0" variables?

Hint: You'll need to do a calculation to start with your invariant in (b) and replace all of the "1" variables with "0" variables.

```
py_1 = px_0 - q \cdot py_0qy_1 = qx_0 - p \cdot qy_0
```

3. Using your answers to 2(e) and 2(f), complete the code for the Euclidean algorithm to update the variables px, qx, py, and qy inside the loop body. You should then be able to run the code with all of the loop invariants passing.

```
1 px, qx, py, qy = py, qy, px - q * py, qx -q * qy
```

Congratulations, you've just completed a derivation of the Extended Euclidean Algorithm!