

CSC236 Week 02: Complete Induction

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1 Complete Induction

- Every natural number greater than 1 has a prime factorization

$$2 = 2$$

$$3 = 3$$

$$4 = 2 \times 2$$

$$5 = 5$$

$$6 = 2 \times 3$$

$$7 = 7$$

$$8 = 2 \times 2 \times 2$$

$$9 = 3 \times 3$$

$$10 = 2 \times 5$$

- How does the factorization of 8 help with the factorization of 9?

The fact that 8 can be expressed as a product of primes has nothing to do with 9 being a product of primes.

2 Notational Convenience

Sometimes you will see the following:

$$\bigwedge_{k=0}^{k=n-1} P(k)$$

... as equivalent to

$$\forall k \in \mathbb{N}, k < n \implies P(k)$$

3 More dominoes

$$\left(\forall n \in \mathbb{N}, \left[\bigwedge_{k=0}^{k=n-1} P(k) \right] \implies P(n) \right) \implies \forall n \in \mathbb{N}, P(n)$$

If all the previous cases always imply the current case then all cases are true.

4 Complete induction outline

- **Inductive step:** introduce n and state inductive hypothesis $H(n)$
 - **Derive conclusion** $C(n)$: show that $C(n)$ follows from $H(n)$, indicating where you use $H(n)$ and why that is valid.
- **Verify base case(s):** verify that the claim is true for any cases not covered in the inductive step

This is the same outline as simple induction but we modify the inductive hypothesis, $H(n)$ so that it assumes the main claim for every natural number from the starting point up to $n - 1$, and the conclusion, $C(n)$ is now the main claim for n .

5 Watch the base cases, part 1

$$f(n) = \begin{cases} 1 & n \leq 1 \\ [f(\lfloor \sqrt{n} \rfloor)]^2 + 2f(\lfloor \sqrt{n} \rfloor) & n > 1 \end{cases}$$

Check a few cases, and make a conjecture:

$$\begin{aligned}f(0) &= 1 \\f(1) &= 1 \\f(2) &= 3 \\f(3) &= 3 \\f(4) &= 15 \\f(5 \dots 15) &= 15 \\f(16) &= 255\end{aligned}$$

All of these things are divisible by 3. The square of something that is divisible by 3 is still divisible by 3 and the double of something that is divisible by 3 is still divisible by 3.

6 For all natural numbers $n > 1$, $f(n)$ is a multiple of 3?

For natural numbers n define $P(n) : f(n)$ is a multiple of 3.

I will prove, using complete induction, that $\forall n > 1, P(n)$.

i. Induction on n

Let $n \in \mathbb{N}$. Assume $n > 1$. Also assume that $P(k)$ is true for all natural numbers k less than n , and greater than 1.

Notice that the floor of the square root of n is greater than 1. Also, the square root of n is less than n (since $n > 1 \implies n^2 > n \implies n > \sqrt{n}$).

Thus by the induction hypothesis, I have $P(\lfloor \sqrt{n} \rfloor)$, this number is a multiple of 3.

Let $k \in \mathbb{N}$ s.t. $\lfloor \sqrt{n} \rfloor = 3k$, so $f(n) = (3k)^2 + 2(3k) = 3(3k^2 + 2k)$, a multiple of 3.

So $P(n)$ follows in both possible cases.

ii. Base Cases

$P(2)$ claims that $f(2) = 3$ is a multiple of 3, which is true.

$P(3)$ claims that $f(3) = 3$ is a multiple of 3, which is true.

7 Zero pair-free binary strings

Deont by $zpfbs(n)$ the number of binary strings of length n That contain no paris of adjacent zeros. What is $zpfbs(n)$ for the first few natural numbers n ?

$$\begin{aligned}
zpbs(0) &= 1 \\
zpbs(1) &= 2 \\
zpbs(2) &= 3 \\
zpbs(3) &= 5 \\
zpbs(4) &= 8 \\
zpbs(5) &= 13 \\
&\dots \\
zpbs(n) &= zpbs(n-1) + zpbs(n-2)
\end{aligned}$$

$$f(n) = \begin{cases} 1, & n = 0 \\ 2, & n = 1 \\ f(n-1) + f(n-2), & n > 1 \end{cases}$$

$\forall n \in \mathbb{N}$, defined predicate $P(n)$ as: $f(n) = zpbs(n)$

Prove by complete induction that for all natural numbers n , $P(n)$.

Let $n \in \mathbb{N}$. Assume that P is true for $0, \dots, n-1$. I will show that $P(n)$ follows.

For the case $n \geq 2$: Partition the zero-pair-free binary strings of length n into those that end in 1 and those that end in 0. Those that end in 1 are simply those of length $n-1$ with a 1 appended, and by $P(n-1)$ (since $n-1 < n$ and $n-1 \geq 0$, $n \geq 1$), there are $f(n-1)$ of these. Those that end in 0 must actually end in 10 (otherwise they are a zero-pair), and by $P(n-2)$ (since $n-2 \leq n$ and $n-2 \geq 0$, $n \geq 2$), there are $f(n-2)$ of these. Altogether there are $f(n-1) + f(n-2)$ zero-pair-free binary strings of length n when $n \geq 2$, which is $P(n)$.

For the base case $n = 0$: There is one binary string (the empty one) of length 0, and it is zero-pair-free, and $f(0) = 1$ and $P(0)$ is true.

For the base case $n = 1$: There are two binary strings of length 1, and neither have pairs of zeros, and $f(1) = 2$ so $P(1)$ is true.

Thus in all possible cases, $P(n)$ follows.

8 Every natural number greater than 1 has a prime factorization

Each natural number n , let predicate $P(n)$ be: n can be expressed as a product of primes.

Prime factorization: represent as product of 1 or more primes.

Prove by complete induction that for all natural numbers n , $P(n)$.

Let $n \in \mathbb{N}$ s.t. $n > 1$. Assume P is true for $2, \dots, n-1$. I will show that $P(n)$ follows.

Case n is composite: By definition, n has a natural number factor f_1 such that $1 < f_1 < n$. By $P(f_1)$ (since $1 < f_1 < n$) we know f_1 can be expressed as a product of primes. Let $f_2 = \frac{n}{f_1}$, since $f_1 > 1$, we know that $\frac{n}{f_1} < n$, and also since $f_1 < n$, we know that $\frac{n}{f_1} > \frac{f_1}{f_1} = 1$. Since $f_1 > 1$, then $f_1 = \frac{n}{f_2} > 1$, so $n > f_2$. So, by $P(f_2)$, we know that f_2 can be expressed as a product of primes.

Therefore since f_1 and f_2 are both products of primes, $n = f_1 \times f_2$ is a product of primes and $P(n)$ follows.

Case n is prime: Then n is its own primes factorization, and $P(n)$ follows.

In all possible cases, $P(n)$ follows.