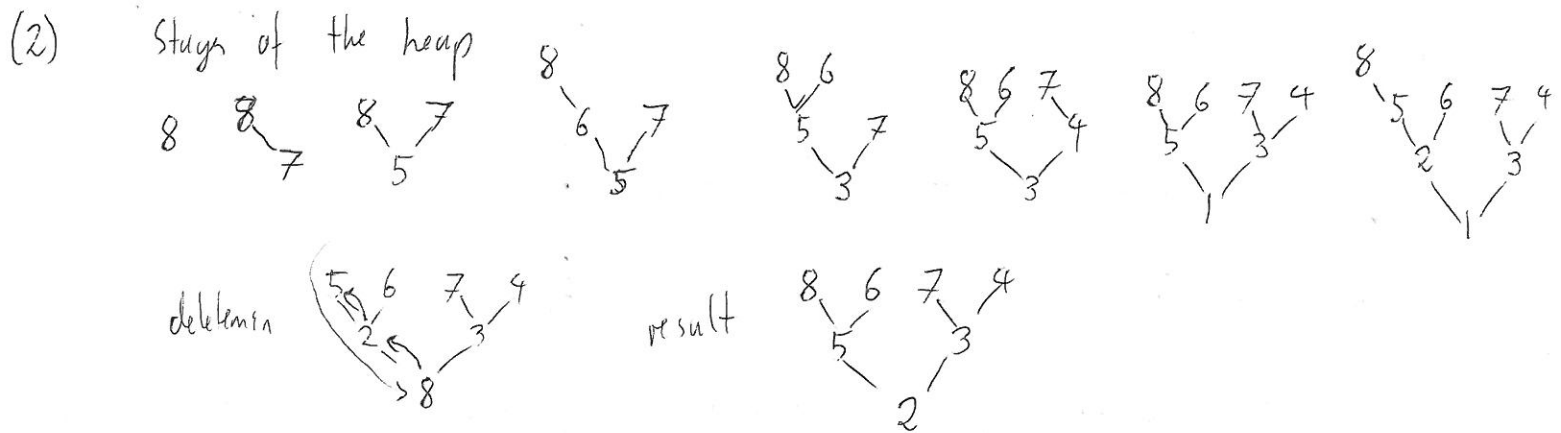


CSc 220 Algorithms Midterm Exam Sample Solution 10/11/2017

(1) $f(n) = n!$ does not satisfy $f(n+1) \leq C f(n)$ for any C , so $f(n+1) \neq O(f(n))$



(3) Have: $f(n) \leq \sum_{i=1}^{n-1} \frac{1}{i} f(i) + C$ for some given C

Want: $f(n) \leq A \cdot n$ for some A (we choose)

Inductive Assumption: $f(k) \leq A \cdot k$ for $k < n$, so $f(i) \leq A \cdot i$ for $i = 1 \dots n-1$

$$\text{so } f(n) \leq \sum_{i=1}^{n-1} \frac{1}{i} f(i) + C \leq \sum_{i=1}^{n-1} \frac{1}{i} A i + C = (n-1) \cdot A + C \leq n \cdot A \text{ for } A \geq C \checkmark$$

(4) Have: $f(n) \leq \frac{1}{2} n^2 f(\frac{1}{2}n)$

Want: $f(n) \leq C \cdot 2^{(\log n)^2}$ for some C

Inductive Assumption: $f(k) \leq C \cdot 2^{(\log k)^2}$ for all $k < n$,

$$\text{so } f(\frac{1}{2}n) \leq C \cdot 2^{(\log \frac{1}{2}n)^2} = C \cdot 2^{(\log n - 1)^2} = C \cdot 2^{(\log n)^2 - 2(\log n) + 1} = C \cdot 2^{(\log n)^2} \cdot \frac{1}{n^2} \cdot 2$$

$$\text{thus } f(n) \leq \frac{1}{2} n^2 f(\frac{1}{2}n) \leq \frac{1}{2} n^2 \cdot C \cdot 2^{(\log n)^2} \cdot \frac{1}{n^2} \cdot 2 = C \cdot 2^{(\log n)^2} \checkmark$$

(5) list of comparisons $(1,2), (5,2), (5,3), (5,4), (5,6), (7,6), (7,10), (8,10), (9,10), (12,10)$

(6)

0	1	2	3	4	5	6	7	8	9
90	91	42	13	44	15		07	98	89
00	51	92	73				37		09
10	01		23						29
	21								

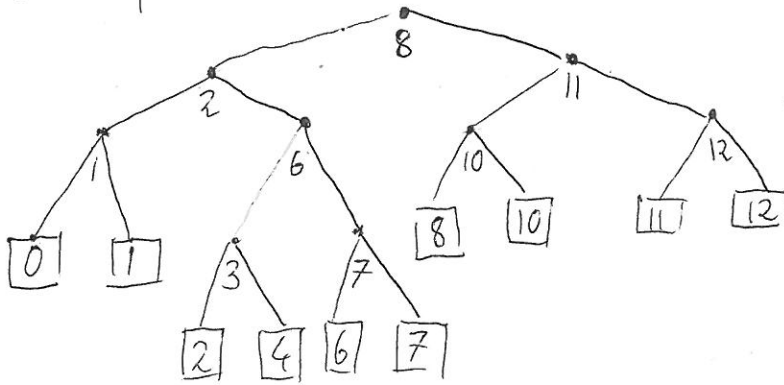
sequence after first round 90 00 10 91 51 01 21 42 92 13 73 23 44 15 07 37 98 89 09 29

0	1	2	3	4	5	6	7	8	9
00	10	21	37	42	51		73	89	90
01	13	23		44					91
07	15	29							92
09									98

sequence after second round 00 01 07 09 10 13 15 21 23 29 37 42 44 51 73 89 90 91 92 98

(7) $\text{int lowerdigit}(\text{int } i, \text{int } n) \{ \text{return } (i \% n); \}$
 $\text{int upperdigit}(\text{int } i, \text{int } n) \{ \text{return } (i / n); \}$

(8) after insert 10 and rebalance (needs double rotation)



after delete 1 and rebalance (needs single rotation)

