EM314 - NUMERICAL METHODS ASSIGNMENT - 1

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E/15/131

QI. use taylor series sinx,

$$f(n) = f(a) + (n-a) f'(a) + (n-a)^2 f''(a) + (n-a)^3 f''(a) + \dots$$

:
$$\sin x = 0 + x \cdot \cos(0) + \frac{\chi^2}{2!} \cdot 0 + \frac{\chi^3}{3!} (-\cos 0) + \cdots$$

$$sin x = 0 + x + 0 + E_3(f, \xi)$$

$$|\sin \chi - \chi| = |E_3(f, \xi)|$$

$$= |\cos(\xi) \cdot \chi^3|$$
31

if approximation has to give correct result rounded to six decimal place then.

$$\frac{\text{error} < 10^6}{31} < 10^6$$

$$|x^3| < 6 \times 10^6$$

: |cos & | < 1

Q2. any $\chi \in \mathbb{F}$ can be represented as, $\chi = \pm (d_1, \dots d_t) \cdot \beta^E$

the sign bit can only assume 2 values. d, cannot assume 0. therefore it can only assume $\beta-1$ values. but each of digits $d_2, d_3, ..., d_k$ can assume β different values. therefore mantissa assumes $(\beta-1) \times \beta^{k-1}$ values.

Range of exponent is [L,u] therefore exponent can assume (U-L+1) values

number of elements, = $2\times(\beta-1)\beta^{t-1}\times(u-L+1)$ set IF contains $2(\beta-1)\beta^{t-1}(u-L+1)$ elements.

Q3. Taylor series., $f(x) = \sum_{k=0}^{N} \frac{1}{k!} \cdot f^{(k)}(\chi_{o})(\chi - \chi_{o})^{k} + E_{N+1} \cdot (f, \xi)$ where $E_{N+1}(f, \xi) = \frac{1}{(N+1)!} \cdot f^{(N+1)}(\xi)(\chi - \chi_{o})^{N+1}$

$$f(x+h) = f(x) + h f'(x) + \frac{h^2}{2} f''(\xi) \quad \text{where} \\ x < \xi < x + h$$

$$f'(x) = f(x+h) - f(x) - \frac{h^2}{2} f''(\xi)$$

$$h = \frac{f(x+h) - f(x)}{h} - \frac{h f''(\xi)}{2} = f'_h(x) - \frac{h f''(\xi)}{2}$$

$$\vdots \quad E_h(x) = \left[f'(x) - f'_h(x)\right] = \left[h \cdot \frac{f''(\xi)}{2}\right]$$

 $\lim_{h\to 0} \frac{|E_n(x)|}{|h|} = \lim_{h\to 0} \frac{|f''(\xi)|}{2} = \frac{|f''(\xi)|}{2} = \text{constant}.$

.. Eh = O(h)

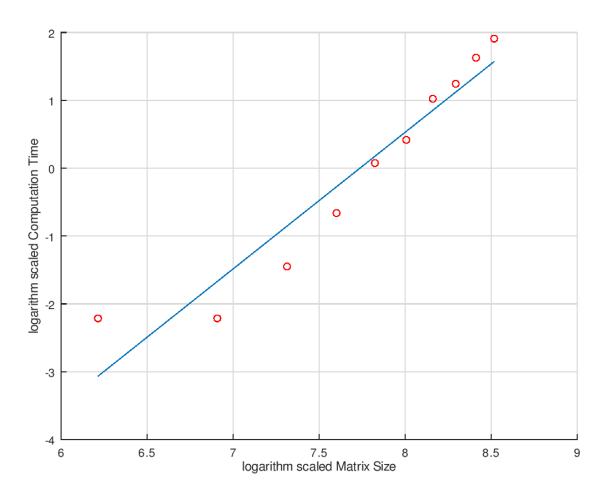
Computer Experiments

Q4. OCTAVE code (q4.m) (a), (b), (c), (d)

Output of the code

```
>> q4
co =
2.3552 -18.2233
```

Output Graph (log t vs log n)



- **(e)** Hence here $\alpha = 2.3552$
- (f) Theoretical value > computing value

 The theoretical value was 3. Here alpha is 2. 3552. It's because the time taken for all arithmetic was assumed to be constant. But real computation time will vary.

Q5. When n = 10

OCTAVE code (q5.m)

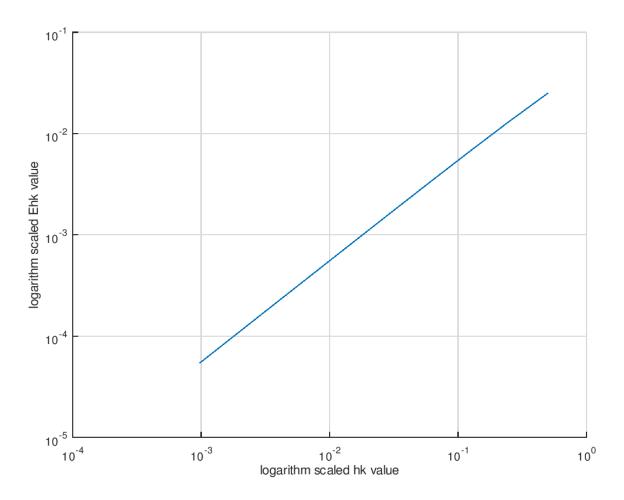
(a) Output of the code

>> q5

k	hk	f'hk	Ehk
1	0.500000	0.308301	0.025032
2	0.250000	0.320171	0.0131625
3	0.125000	0.326576	0.00675738
4	0.062500	0.329909	0.00342474
5	0.031250	0.331609	0.00172415
6	0.015625	0.332468	0.000865053
7	0.007812	0.332900	0.000433276
8	0.003906	0.333117	0.000216826
9	0.001953	0.333225	0.00010846
10	0.000977	0.333279	5.42417e-05

0.98706 -2.96094

Output Graph (log Ehk vs log hk)



(b) Ehk tends to zero (Ehk \rightarrow 0) Hence according to the graph Υ = 0. 98706 Eh/h is tents to a constant while h tents to infinity. Hence Eh = O(h)

(c) When n = 40

Mathlab code (q6.m)

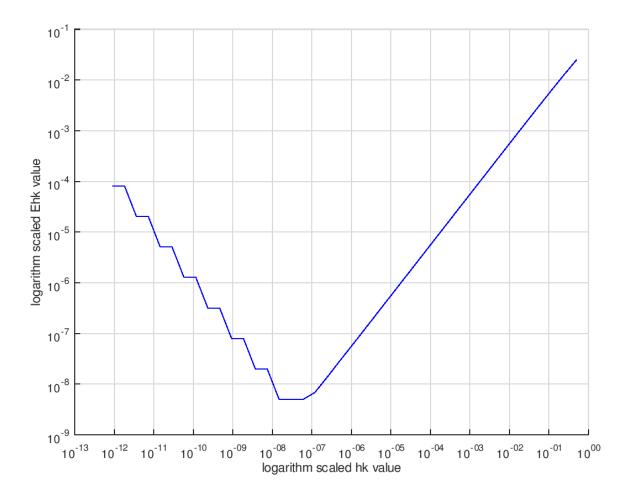
Output of the code

>> q5

k	hk	f'hk	Ehk
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10	0.000977	0.333279	5.42417e-05

0.000488	0.333306	2.71238e-05
0.000244	0.333320	1.35626e-05
0.000122	0.333327	6.7815e-06
0.000061	0.333330	3.3908e-06
0.000031	0.333332	1.69541e-06
0.000015	0.333332	8.47712e-07
0.000008	0.333333	4.23858e-07
0.000004	0.333333	2.11953e-07
0.000002	0.333333	1.06016e-07
0.000001	0.333333	5.3163e-08
0.000000	0.333333	2.68531e-08
0.000000	0.333333	1.3349e-08
0.000000	0.333333	6.8297e-09
0.000000	0.333333	4.96705e-09
0.000000	0.333333	4.96705e-09
0.000000	0.333333	4.96705e-09
0.000000	0.333333	1.98682e-08
0.000000	0.333333	1.98682e-08
0.000000	0.333333	7.94729e-08
0.000000	0.333333	7.94729e-08
0.000000	0.333333	3.17891e-07
0.000000	0.333333	3.17891e-07
0.000000	0.333332	1.27157e-06
0.000000	0.333332	1.27157e-06
0.000000	0.333328	5.08626e-06
0.000000	0.333328	5.08626e-06
0.000000	0.333313	2.03451e-05
0.000000	0.333313	2.03451e-05
0.000000	0.333252	8.13802e-05
0.000000	0.333252	8.13802e-05
	0.000244 0.000122 0.000061 0.000031 0.000008 0.000000 0.000000 0.000000 0.000000 0.000000	0.0002440.3333200.0001220.3333270.0000610.3333300.0000150.3333320.0000080.3333330.0000040.3333330.0000010.3333330.0000010.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333330.0000000.3333320.0000000.3333320.0000000.3333320.0000000.3333330.0000000.3333320.0000000.3333320.0000000.3333330.0000000.3333330.0000000.333332

Output Graph (log Ehk vs log hk)



- (d) when N increases Ehk decreases and reach it's minimum value and then increases. Therefore, Ehk not tends to zero when N → 0. It happens because machine can handle only up to certain precision value. When Ehk values too small machine will round off. Therefore after some point roundoff constant dominates In that time Ehk value increases.
- (e) From the graph Ehk minimum is 4.96705e-09, when k = 26 (hk = 0.000000). Therefore $k_{min} = 26$ minimizes the Ehk.