

# Notes on "Topoi" by Goldblatt

hisnobu-nakamura

<2019-06-27 Thu>

## Contents

<b>1 Chapter 3: Arrows instead of epsilons</b>	<b>1</b>
1.1 3.1. Monic arrows . . . . .	1
1.1.1 3.1. Exercise Solutions . . . . .	1
1.2 3.2. Epic arrows . . . . .	1
1.3 3.8 Products . . . . .	2

## 1 Chapter 3: Arrows instead of epsilons

### 1.1 3.1. Monic arrows

An arrow  $f : a \rightarrow b$  in a category  $\mathcal{C}$  is *monic* in  $\mathcal{C}$  if for any pair  $g, h : c \rightrightarrows a$  of  $\mathcal{C}$ -arrows, equality  $f \circ g = f \circ h$  implies that  $g = h$ . The symbolism  $f : a \rightarrowtail b$  is used to indicate that  $f$  is monic.

#### 1.1.1 3.1. Exercise Solutions

In any category (1)  $g \circ f$  is monic if both  $f$  and  $g$  are monic. (2) If  $g \circ f$  is monic, then so is  $f$ . **Solutions**

(1) If  $(g \circ f) \circ h_1 = (g \circ f) \circ h_2$ , then

$$\begin{aligned}(g \circ f) \circ h_1 &= (g \circ f) \circ h_2 \\ g \circ (f \circ h_1) &= g \circ (f \circ h_2) \\ f \circ h_1 &= f \circ h_2 \quad (\because g \text{ is monic}) \\ h_1 &= h_2 \quad (\because f \text{ is monic})\end{aligned}$$

Therefore,  $(g \circ f)$  is monic.

(2)

If  $f \circ h_1 = f \circ h_2$ , then

$$\begin{aligned}g \circ (f \circ h_1) &= g \circ (f \circ h_2) \\ (g \circ f) \circ h_1 &= (g \circ f) \circ h_2 \\ h_1 &= h_2 \quad (\because g \circ f \text{ is monic})\end{aligned}$$

Hence,  $f$  is monic.

### 1.2 3.2. Epic arrows

An arrow  $f : a \rightarrow b$  is *epic* (right-cancellable) in a category  $\mathcal{C}$  if for any pair  $g, h : b \rightrightarrows c$ , the equality  $g \circ f = h \circ f$  implies  $g = h$  i.e. whenever a diagram

$$\begin{array}{ccc} a & \xrightarrow{f} & b \\ f \downarrow & & \downarrow g \\ b & \xrightarrow{h} & c \end{array}$$

commutes, then  $g = h$ . The notation  $f : a \twoheadrightarrow b$  is used for epic arrows.

### 1.3 3.8 Products

**DEFINITION** A product in a category  $\mathcal{C}$  of two objects  $a$  and  $b$  is a  $\mathcal{C}$ -object  $a \times b$  together with a pair  $(pr_a : a \times b \rightarrow a, pr_b : a \times b \rightarrow b)$  of  $\mathcal{C}$ -arrows such that for any pair of  $\mathcal{C}$ -arrows of the form  $(f : c \rightarrow a, g : c \rightarrow b)$  there is exactly one arrow  $\langle f, g \rangle : c \rightarrow a \times b$  making

$$\begin{array}{ccccc}
 & & c & & \\
 & f \swarrow & \downarrow \langle f, g \rangle & \searrow g & \\
 a & \xleftarrow{pr_a} & a \times b & \xrightarrow{pr_b} & b
 \end{array}$$

commute, i.e. such that  $pr_a \circ \langle f, g \rangle = f$ ,  $pr_b \circ \langle f, g \rangle = g$ .  $\langle f, g \rangle$  is the product arrow of  $f$  and  $g$  with respect to the projections  $pr_a, pr_b$ .